

Iterative Solution to Nonlinear Equations (1D)

Given a function $f: \mathbb{R} \rightarrow \mathbb{R}$, consider the root-finding problem

(P) Find $\xi \in \mathbb{R}$ such that $f(\xi) = 0$.

We've all seen linear equations: $ax + b = 0 \Rightarrow x = -b/a$

Quadratic equations: $ax^2 + bx + c = 0 \Rightarrow x_{\pm} = \frac{-b \pm (b^2 - 4ac)^{1/2}}{2a}$

Q: What about higher-order polynomials?

A: No! For all $n \geq 5$, there exists a polynomial with integer coefficients with solution that cannot be written in terms of radicals.

So: No general formula for ξ can exist! Instead, we construct an approximation to ξ via iteration.

Simple Iteration

Let $f: [a, b] \rightarrow \mathbb{R}$, $a < b$, $f \in C^0[a, b]$

We want to solve the problem (P).

Ex: $f(x) = x^2 + 1$. $\xi = \pm i$, so $\xi \notin \mathbb{R}$!

Thm: Let $f \in C^0[a, b]$, $f: [a, b] \rightarrow \mathbb{R}$. Assume $f(a) \cdot f(b) \leq 0$.

Then $\exists \xi \in [a, b]$ with $f(\xi) = 0$.

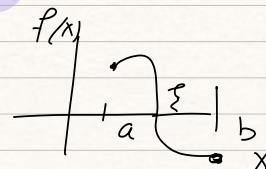
Pf: Assume $f(a) \cdot f(b) \neq 0$, otherwise we are done.

Then $f(a) \cdot f(b) < 0$, so $0 \in (f(a), f(b))$. By the intermediate value theorem, done. \square

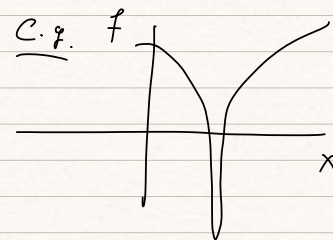
Note: Can be hard to use the above theorem because such an

" f is a continuous fn. on $[a, b]$ "

" $f(a)$ and $f(b)$ have opp. sign"



interval $[a, b]$ could be hard to find!
(see book for example)



Instead, we often transform the problem to be a **fixed-point problem** by writing

$$x - g(x) = 0 \Leftrightarrow f(x) = 0$$

(for example, $g = f + x$)

Def: We say that ξ is a **fixed-point** of g if $g(\xi) = \xi$.

Thm (Brouwer): Let $g: [a, b] \rightarrow \mathbb{R}$, $g \in C([a, b])$. Further assume that $g(x) \in [a, b]$ for all $x \in [a, b]$. Then $\exists \xi \in [a, b]$ such that $\xi = g(\xi)$.

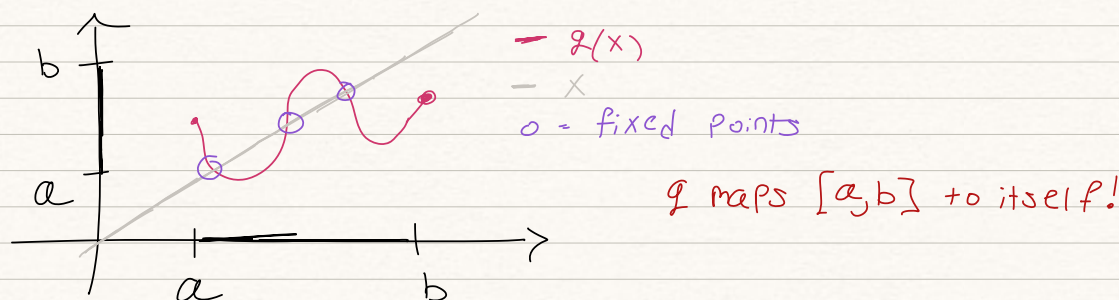
Pf: Let $f(x) = x - g(x)$. Then:

$$(i) \quad f(a) = a - g(a) \leq 0 \quad (g(x) \geq a)$$

$$(ii) \quad f(b) = b - g(b) \geq 0 \quad (g(x) \leq b)$$

So $f(a)f(b) \leq 0$, $f \in C([a, b])$. Hence, by the

previous theorem, $\exists \xi$ with $f(\xi) = 0 \Rightarrow g(\xi) = \xi$. \square



Note: Many ways to transform root-finding problem to a fixed-point problem, only need to find one with $g(x) \in [a, b]$.

Note: So far we have verified the existence of solutions to the problem $f(x) = 0$. The next def. moves towards algorithms.

Def: Let $g \in C^0([a, b])$ and assume $g(x) \in [a, b]$ for $x \in [a, b]$.

Let $x_0 \in [a, b]$. We call the recursion

$$(1) \quad x_{k+1} = g(x_k), \quad k = 0, 1, \dots,$$

a simple iteration. The $\{x_k\}$ are iterates.

Claim: If the iteration (1) converges to some ξ , then ξ is a fixed point of g .

PF: By continuity,

$$\xi = \lim_{k \rightarrow \infty} x_{k+1} = \lim_{k \rightarrow \infty} g(x_k) = g\left(\lim_{k \rightarrow \infty} x_k\right) = g(\xi) \quad \square$$

We now give sufficient conditions for convergence of (1).

Def: Let $g \in C^0([a, b], [a, b])$. $g(x) \in [a, b]$ for $x \in [a, b]$

We say g is a contraction if $\exists 0 < L < 1$ with

$$|g(x) - g(y)| \leq L |x - y| \quad \forall x, y \in [a, b]$$