- We study three ways to estimate the tail probabilities of random variables.
- It will be noted that, the more information we know about the random variable the better the estimate we can derive about a given tail probability.

Markov Inequality: If X is a non-negative valued random variable with an expectation of μ, then Pr[X ≥ cμ] ≤ 1/c.

Proof of Markov inequality:

$$\mu = \sum_{\alpha} \alpha P_{\gamma}(X=\alpha)$$

$$= \sum_{\alpha} \alpha P_{\gamma}(X=\alpha) + \sum_{\alpha} \alpha p_{\gamma}(X=\alpha)$$

$$= \sum_{\alpha < \gamma} P_{\gamma}(X=\alpha)$$

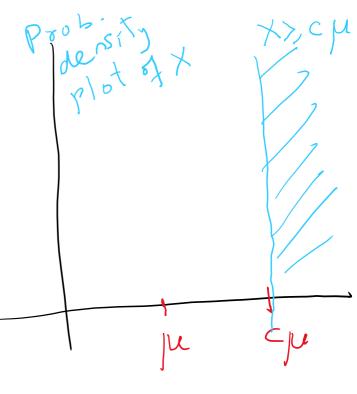
$$= \sum_{\alpha < \gamma} P_{\gamma}(X=\alpha)$$

$$= \sum_{\alpha > \gamma < \gamma} P_{\gamma}(X=\alpha)$$

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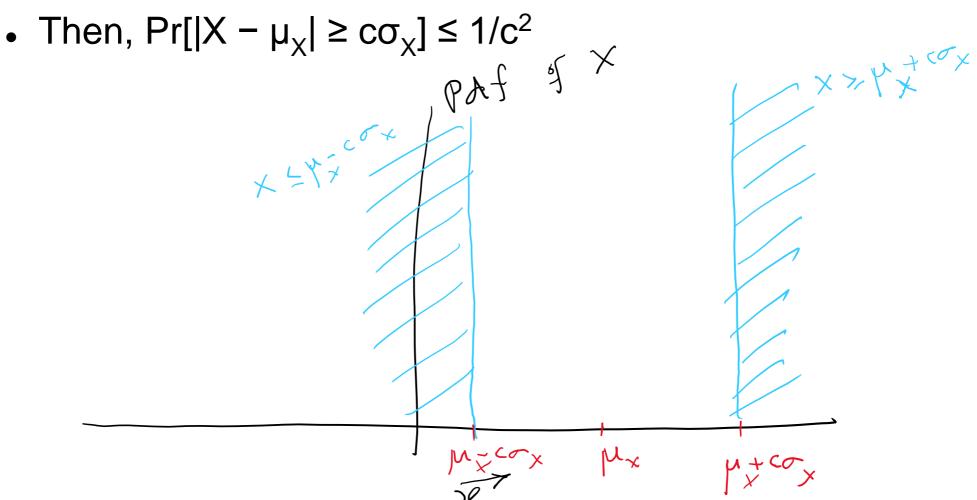
$$= \sum_{\alpha > \gamma} P_{\gamma}(X=\alpha)$$



- Markov Inequality: If X is a non-negative valued random variable with an expectation of µ, then Pr[X ≥ cµ] ≤ 1.
- Applying this inequality tells us that the randomized quick sort algorithm has a run time of more than twice its expectation with a probability of 1/2.
- The run time is n<sup>2</sup> with probability of nearly log n /n.

- Chebychev Inequality: We first define the terms standard deviation and variance of a random variable X.
- Let X be a random variable with an expectation of μ. The variance of X, denoted by var(X), is defined as var(X) = E[(X μ)²]. The standard deviation of X, denoted by σ<sub>X</sub>, is defined as σ<sub>X</sub> = var(X).
- Note that by definition,  $var(x) = E[(X-\mu)^2] = E[X^2-2X\mu+\mu^2] = E[X^2]-\mu^2$ .
- The second equality follows from the linearity of expectations.

 Chebychev inequality: Let X be a random variable with expectation μ<sub>X</sub> and standard deviation σ<sub>X</sub>.



- Chebychev inequality: Let X be a random variable with expectation  $\mu_X$  and standard deviation  $\sigma_X$ .
- Then,  $Pr[|X \mu_X| \ge c\sigma_X] \le 1$
- Proof. Let random variable  $Y = (X \mu_X)^2$ . Then,

$$E[Y] = E[(X - \mu_X)^2] = \sigma_X^2$$
 by definition

Now, 
$$Pr[|X - \mu_X| \ge c\sigma_X] = Pr[(X - \mu_X)^2 \ge c^2\sigma_X^2]$$

$$= Pr[Y \ge c^2 \sigma_X^2].$$

Applying Markov Inequality to the random variable Y

$$Pr[Y \ge c^2 \sigma_X^2] = Pr[Y \ge c^2 \mu_Y] \le \frac{1}{c^2}$$

- Better tail inequalitites can be obtained by the powerful technique called Chernoff bounds.
- However, applicability is a little restricted too.
- Let us study the most popular version to start with.

- Let X be a random variable defined as the sum of n independent and identically distributed random variables X<sub>1</sub>, X<sub>2</sub>, ..., X<sub>n</sub>.
  - In other words,  $X = \Sigma_i X_i$
  - Short form i.i.d.
- Let us assume that each X<sub>i</sub> is a Bernoulli random variable.
  - In other words, each X<sub>i</sub> takes values in {0, 1}.
- Let  $Pr(X_i = 1) = p$  and hence  $Pr(X_i = 0) = 1 p$ .
- Finally, let  $E[X] = \mu$ .
  - Notice that  $E[X] = \Sigma_i E[X_i] = n. (1.p + (1-p).0)) = np.$

- Several settings relate to the above statements.
- Consider throwing a (biased) coin over n trials.
- Each trial, the probability of Heads is p.
- So, each X<sub>i</sub> corresponds to the fact that the ith trial results in a Heads.
- Let us count the number of Heads over the n trails.
   Indeed, X = Σ<sub>i</sub> X<sub>i</sub> captures this count as a random variable.
- Note that the expected number of Heads over n trails is exactly np.

- Finally, to the theorem.
- Given the earlier conditions, it holds that for any  $\delta > 0$ ,

$$P_{\delta}(X \ge M(1+\delta)) \le \left(\frac{e^{\delta}}{(1+\delta)^{1+\delta}}\right)^{M}$$

- The normal strategy employed to prove tail estimates of sums of independent random variables is to make use of exponential moments.
- While proving Chebychev inequality, we made use of the second-order moment. It can be observed that using higher order moments would generally improve the bound on the tail inequality.
- But using exponential moments would result in a vast improvement.

$$P_{\delta}(\times > \mu(1+\delta)) \leq \left(\frac{e^{\delta}}{(1+\delta)^{1+\delta}}\right)^{\mu}$$

- While proving Chebychev inequality, we made use of the second-order moment of a random variable.
- It can be observed that using higher order moments would generally improve the bound on the tail inequality
- But using exponential moments would result in a vast improvement in the bound.
  - An exponential moment of a random variable X is the expectation of functions of X such as e<sup>X</sup>.

$$P_{s}(x \geq \mu(1+\delta)) \leq \left(\frac{e^{\delta}}{(1+\delta)^{1+\delta}}\right)^{\mu}$$

- For each i, we define the random variable Y<sub>i</sub> = e<sup>tXi</sup> for a real number t > 0 that will be chosen later.
- Notice that
  - 1) Y<sub>i</sub> is a positive valued random variable.

2) 
$$E[Y_i] = E[e^{t \times i}] = p_i \cdot e^t + (1 - p_i) \cdot e^t$$
  
=  $p_i \cdot e^t + 1 - p_i$ 

• Given the earlier conditions, it holds that for any  $\delta > 0$ ,

$$P_{\delta}(X \ge \mu(1+\delta)) \le \left(\frac{e^{\delta}}{(1+\delta)^{1+\delta}}\right)^{\mu}$$

- Now define another random variable Y = Y<sub>1</sub>.Y<sub>2</sub>....Y<sub>n</sub>.
- Now, we can note that

$$E[Y] = E[Y_1, Y_2, \dots, Y_n]$$

$$= \int_{i=1}^{n} f(Y_i) = (P_i - e^t + 1 - P_i)$$

Why does the above calculation hold?

$$P_{s}(x > \mu(1+\delta)) \leq \left(\frac{e^{\delta}}{(1+\delta)^{1+\delta}}\right)^{\mu}$$

- The next step we do is do apply Markov inequality on Y as follows.
- First, notice that  $Y = e^{tX}$  as

$$y = y_1 \cdot y_2 \cdot y_n = e^{t \times 1} \cdot e^{t \times 2}$$

$$= e^{t \times 1 + x_2 + \dots + x_m} = e^{t \times 1}$$

$$= e^{t \times 1 + x_2 + \dots + x_m} = e^{t \times 1}$$

• Given the earlier conditions, it holds that for any  $\delta > 0$ ,

$$P_{\delta}(X \ge \mu(1+\delta)) \le \left(\frac{e^{\delta}}{(1+\delta)^{1+\delta}}\right)^{\mu}$$

• First, notice that  $Y = e^{tX}$ . And,

• Further, 
$$X > \mu(1+\delta) \Leftrightarrow e^{t} > e^{t} \mu(1+\delta)$$

$$\Rightarrow Y > e^{t} \mu(1+\delta)$$

• So, we are interested in the event  $Y \ge e^{t\mu(1+\delta)}$ .

• Given the earlier conditions, it holds that for any  $\delta > 0$ ,

$$P_{s}(\times \geq M(1+\delta)) \leq \left(\frac{e^{\delta}}{(1+\delta)^{1+\delta}}\right)^{M}$$

• So, we are interested in the event  $Y \ge t\mu(1+\delta)$ . We proceed as:

$$Pr(Y \ge e^{t\mu(1+8)}) \le \frac{E[Y]}{e^{t}\mu(1+8)} = \frac{TT(1-k_1+k_2e^t)}{e^{t}\mu(1+8)}$$

$$\le \frac{Tt e^{-k_1+k_2e^t}}{e^{t}\mu(1+8)}$$

$$= \frac{e^{-\mu(1+e^t)}}{e^{t}\mu(1+8)}$$

• Given the earlier conditions, it holds that for any  $\delta > 0$ ,

$$P_{s}(\times \geq M(1+\delta)) \leq \left(\frac{e^{\delta}}{(1+\delta)^{1+\delta}}\right)^{M}$$

So, we are interested in the event Y ≥ tµ(1+δ). We proceed as:

$$Pr(Y \ge e^{t\mu(1+8)}) \le \frac{E[Y]}{e^{t\mu(1+8)}} = \frac{\pir(1-k_1+k_1e^t)}{e^{t\mu(1+8)}}$$

$$\le \frac{T+e^{-k_1+k_1e^t}}{e^{t\mu(1+8)}} = e^{-\mu(1-e^t)} - t\mu(1+8)$$

$$= \frac{e^{-\mu(1-e^t)}}{e^{t\mu(1+8)}} = e^{-\mu(1-e^t)} - t\mu(1+8)$$

$$P_{8}(x > \mu(1+8)) \leq \left(\frac{e^{\delta}}{(1+8)^{1+8}}\right)^{\mu}$$
• So,  $P_{8}(y > e^{t}\mu(1+8)) = e^{-\mu(1-e^{t})} - t\mu(1+8)$ 

- Since t is a free parameter in the above, we can find a t that minimizes the right hand side.
- To simplify, let f(t) = lne  $= -\mu(1-e^{t}) t\mu(1+\delta)$   $= -\mu(1-e^{t}) t\mu(1+\delta)$

$$P_{s}(\times \geq \mu(1+\delta)) \leq \left(\frac{e^{\delta}}{(1+\delta)^{1+\delta}}\right)^{\mu}$$

$$-\mu(1-e^{t}) - t\mu(1+\delta)$$

- To simplify, let  $f(t) = lne^{-\mu(1-e^t)} t\mu(1+\delta)$ = - m (1-e2) - t m (1+8)
- Differentiating f(t) wrt t, we get f'(t) = met \_ m(1+8)
- So, f'(t) = 0 at  $f = l_n (1+\delta)$
- Verify that the above t corresponds to a minima. OFFLINE.

• Given the earlier conditions, it holds that for any  $\delta > 0$ ,

$$P_{\delta}(\times \geq M(1+\delta)) \leq \left(\frac{e^{\delta}}{(1+\delta)^{1+\delta}}\right)^{M}$$

• With  $t = \ln(1+\delta)$ , we get that

$$P_8(x > \mu(1+8)) \leq \frac{e^{-\mu(1-(1+8))}}{(1+8)^{\mu(1+8)}} = \frac{e^{\mu s}}{(1+s)^{1+s}}$$

$$= \left(\frac{e^{S}}{(1+S)^{1+S}}\right)^{\gamma}$$

completing the proof.