• Given the earlier conditions, it holds that for any $\delta > 0$,

$$P_{8}(x > \mu(1+8)) \leq \left(\frac{e^{\delta}}{(1+8)^{1+8}}\right)^{\mu}$$
• So, $P_{8}(y > e^{t}\mu(1+8)) = e^{-\mu(1-e^{t})} - t\mu(1+8)$

- Since t is a free parameter in the above, we can find a t that minimizes the right hand side.
- To simplify, let f(t) = lne $= -\mu(1-e^{t}) t\mu(1+\delta)$ $= -\mu(1-e^{t}) t\mu(1+\delta)$

• Given the earlier conditions, it holds that for any $\delta > 0$,

$$P_{s}(\times \geq \mu(1+\delta)) \leq \left(\frac{e^{\delta}}{(1+\delta)^{1+\delta}}\right)^{\mu}$$

$$-\mu(1-e^{t}) - t\mu(1+\delta)$$

- To simplify, let $f(t) = lne^{-\mu(1-e^t)} t\mu(1+\delta)$ = - m (1-e2) - t m (1+8)
- Differentiating f(t) wrt t, we get f'(t) = met _ m(1+8)
- So, f'(t) = 0 at $f = l_n (1+\delta)$
- Verify that the above t corresponds to a minima. OFFLINE.

• Given the earlier conditions, it holds that for any $\delta > 0$,

$$P_{\delta}(\times \geq M(1+\delta)) \leq \left(\frac{e^{\delta}}{(1+\delta)^{1+\delta}}\right)^{M}$$

• With $t = \ln(1+\delta)$, we get that

$$P_8(x > \mu(1+8)) \leq \frac{e^{-\mu(1-(1+8))}}{(1+8)^{\mu(1+8)}} = \frac{e^{\mu s}}{(1+s)^{1+s}}$$

$$= \left(\frac{e^{S}}{(1+S)^{1+S}}\right)^{\gamma}$$

completing the proof.

• Given the earlier conditions, it holds that for any $\delta > 0$,

$$P_{s}(\times \geq M(1+\delta)) \leq \left(\frac{e^{\delta}}{(1+\delta)^{1+\delta}}\right)^{M}$$

• With $t = \ln(1+\delta)$, we get that

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A simplification of the RHS gives

$$P_{8}(x > \mu(148)) \leq \begin{cases} e^{-\mu 8/4} & \text{if } 8 \leq 1 \\ e^{-\mu 8/n8} & \text{if } 8 > 1 \end{cases}$$

A Simple Example

- Here is one more practical application of the Chernoff bounds.
- Consider once again counting the number of heads out of tossing n fair coins independently.
 - So, $p = \frac{1}{2}$.
- Let X_i denote the random variable that takes 1 if the ith coin toss results in a head, and 0 otherwise.
 - $E[X_i] = \frac{1}{2}$.
- Let $X = \Sigma_i X_i$
- X counts the total number of heads over the n tosses.
 - E[X] = n/2.
 - With n = 100, say, we expect 50 heads over 100 coin tosses.

A Simple Example

- Markov inequality tells that the probability that X takes a value beyond 70 is Pr (X ≥ 70) = Pr (X ≥ 70/50 x 50)
 <= 5/7 = 0.7 (approx.)
- To apply Chbychev's inequality, we need to do some extra work.
- Var(X_i) for any i is computed as E[X_i²] E[X_i]².
- $E[X_i^2] = 0x(1/2) + 1x(1/2) = \frac{1}{2}$.
- $Var(X_i) = \frac{1}{2} (\frac{1}{2})^2 = \frac{1}{4}$.
- Var(X) = 100x Var(X₁) due to independence.
- So, Var(X) = 100/4 = 25, and $\sigma_X = 5$.
- Now, $\Pr(X \ge 70)$ can be rewritten as $\Pr(|X 50| \ge 20)$ and further as $\Pr(|X 50| \ge (20x 1/\sigma_X) \sigma_X)$) which is now at most 25/400 = 1/16 = 0.0625.