

CS 32 Week 8

Discussion 1I

Srinath

Outline

- Trees
- Binary Search Trees
- Worksheet 8

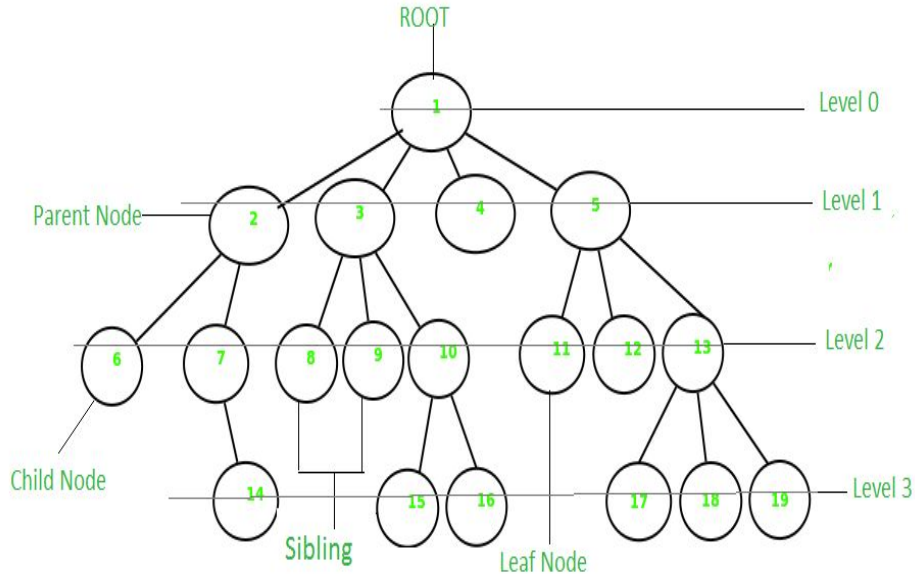
Trees



Trees :

Kind of data structure that holds hierarchy. Can be thought of as extension to linked lists

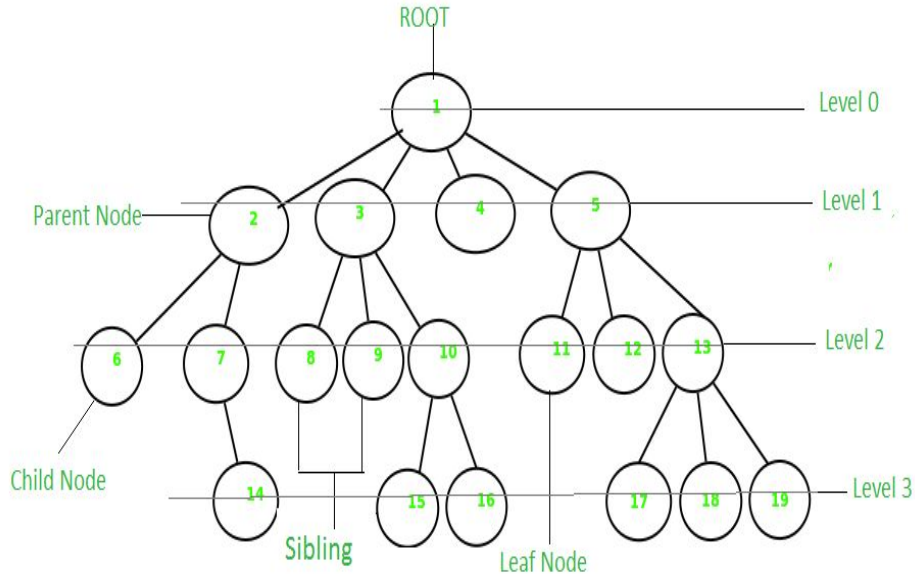
Terms : parent, child, root, leaf, sibling, height, subtree



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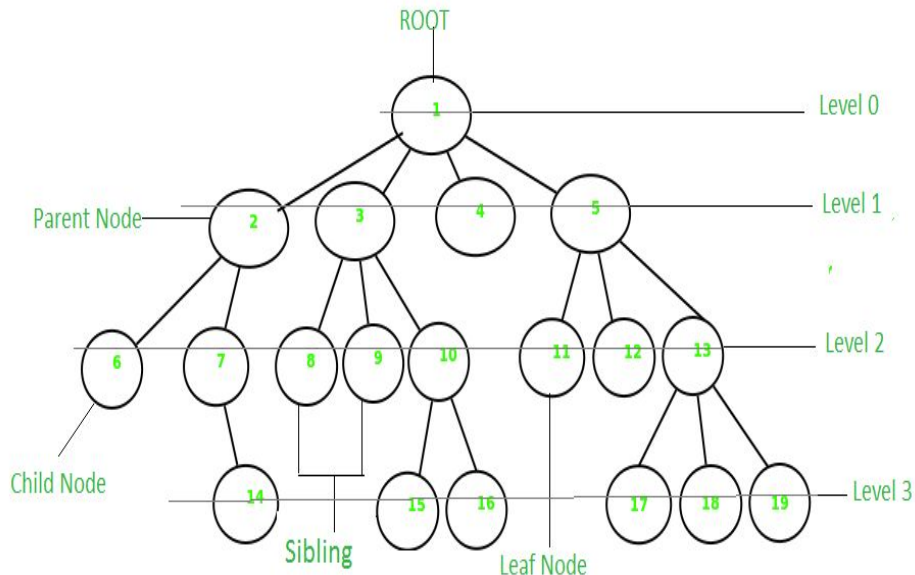
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Code??

Trees :

Kind of data structure that holds hierarchy. Can be thought of as extension to linked lists



Terms : parent, child, root, leaf, sibling, height, subtree

Code??

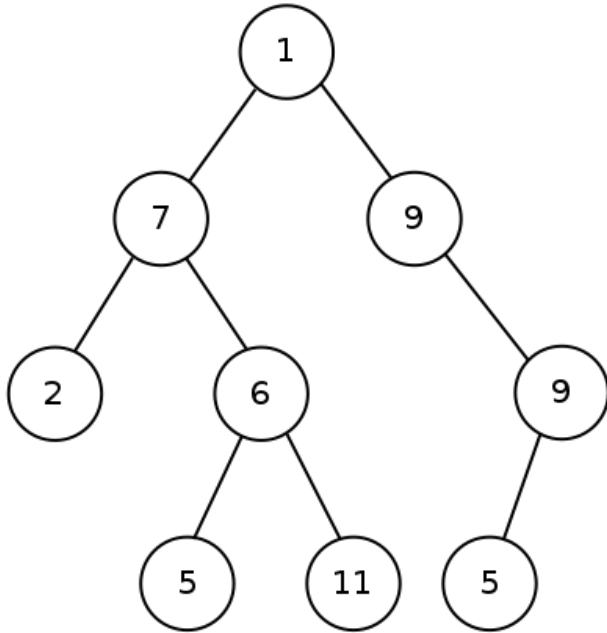
```
struct Node{  
    int val;  
    vector<Node*> children;  
};
```

```
Node* root = new Node();  
root->val = 100;
```

```
Node* root = NULL; // empty tree
```

Trees : Binary Tree

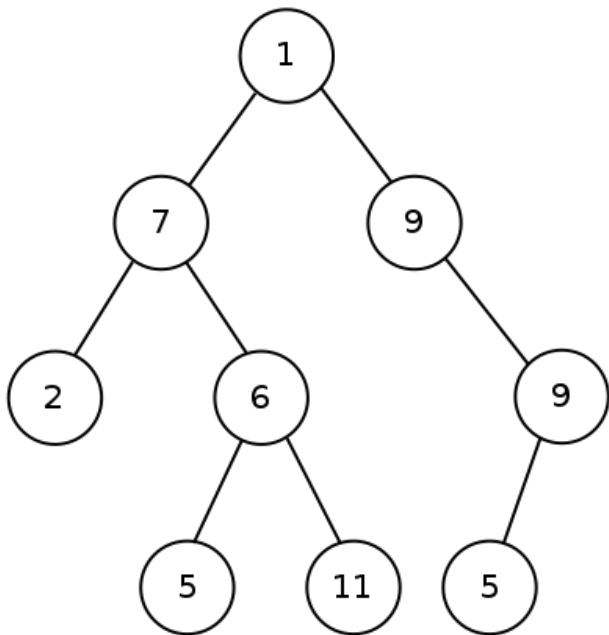
A simplified tree, each node has **at most 2** children.



Code??

Trees : Binary Tree

A simplified tree, each node has **at most 2** children.



Code??

```
struct Node{  
    int val;  
    Node* left;  
    Node* right;  
};
```

```
Node* root = new Node();  
root->val = 100;
```

```
Node* root = NULL; // empty tree
```


Trees : Insertion

Insert a new value (n) at given Node* p

```
void insert(Node* p, int n){  
  
}
```

Trees : Insertion

Insert a new value (n) at given Node* p

```
void insert(Node* p, int n){  
    // create a new node  
    Node* myNode = new Node();  
    myNode->val = n;  
    if (p==NULL){ // empty tree  
        p = myNode;  
    }else{  
        (p->children).push_back(myNode);  
    }  
    return;  
}
```

```
Node * root = NULL;  
insert(root, 190);
```

Trees : Insertion

Insert a new value (n) at given Node* p

```
void insert(Node* p, int n){  
    // create a new node  
    Node* myNode = new Node();  
    myNode->val = n;  
    if (p==NULL){ // empty tree  
        p = myNode;  
    }else{  
        (p->children).push_back(myNode);  
    }  
    return;  
}
```

```
Node * root = NULL;  
insert(root, 190);
```

Are we missing anything?

Trees : Insertion

Insert a new value (n) at given Node* p

```
void insert(Node* & p, int n){  
    // create a new node  
    Node* myNode = new Node();  
    myNode->val = n;  
    if (p==NULL){ // empty tree  
        p = myNode;  
    }else{  
        (p->children).push_back(myNode);  
    }  
    return;  
}
```

```
Node * root = NULL;  
insert(root, 190);
```

Are we missing anything?

- Yes, 'root' won't be modified as we are passing by value

Trees : Traversal

To go through all nodes of a Tree.

Pre-Order : parent, then children

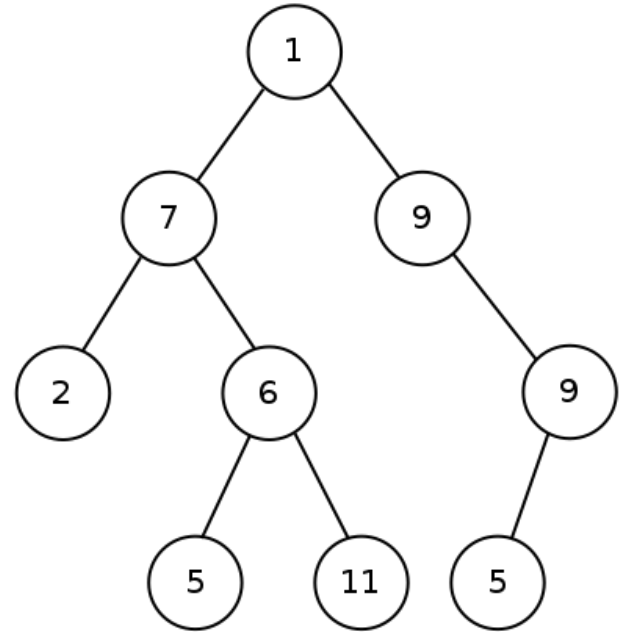
Post-Order : children, then parent

In-Order : left subtree, parent, right subtree (defined for binary trees)

Pre-Order :

Post-Order :

In-Order :



Trees : Traversal

To go through all nodes of a Tree.

Pre-Order : parent, then children

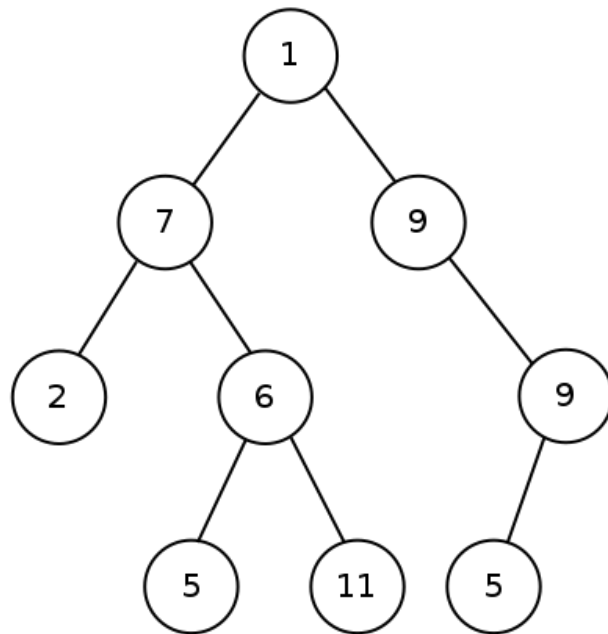
Post-Order : children, then parent

In-Order : left subtree, parent, right subtree (defined for binary trees)

Pre-Order : 1, 7, 2, 6, 5, 11, 9, 9, 5

Post-Order : 2, 5, 11, 6, 7, 5, 9, 9, 1

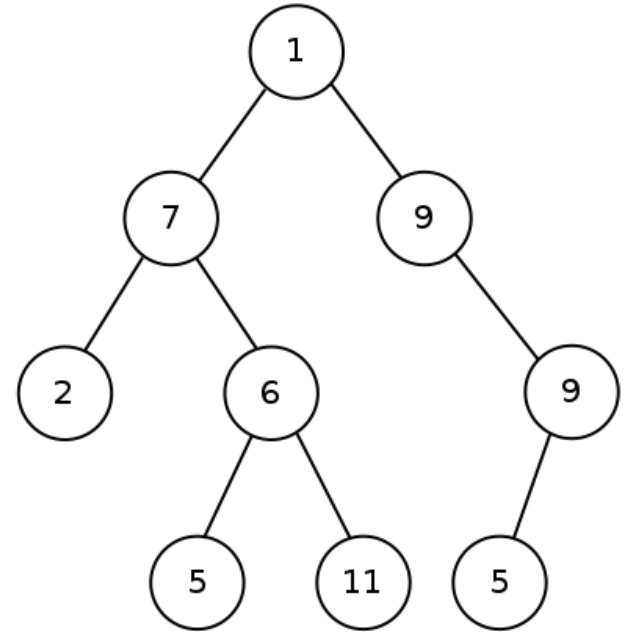
In-Order : 2, 7, 5, 6, 11, 1, 9, 5, 9



Trees : Traversal

Code for pre-order traversal?

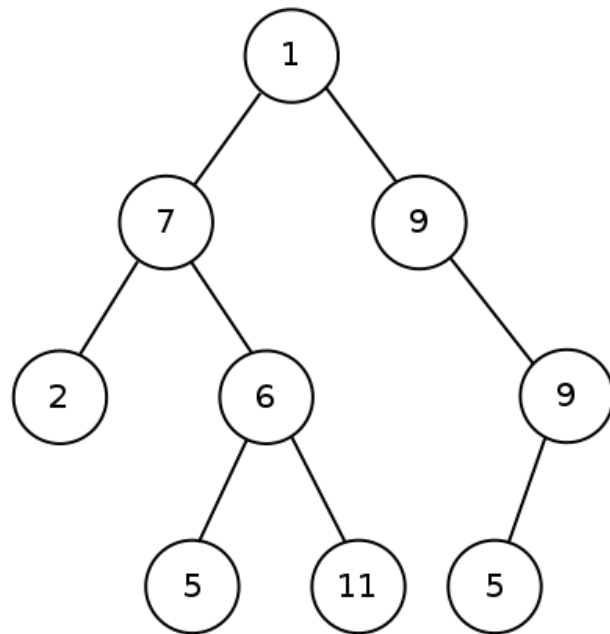
```
void pre-order-traversal(Node* root){  
}
```



Trees : Traversal

Code for pre-order traversal?

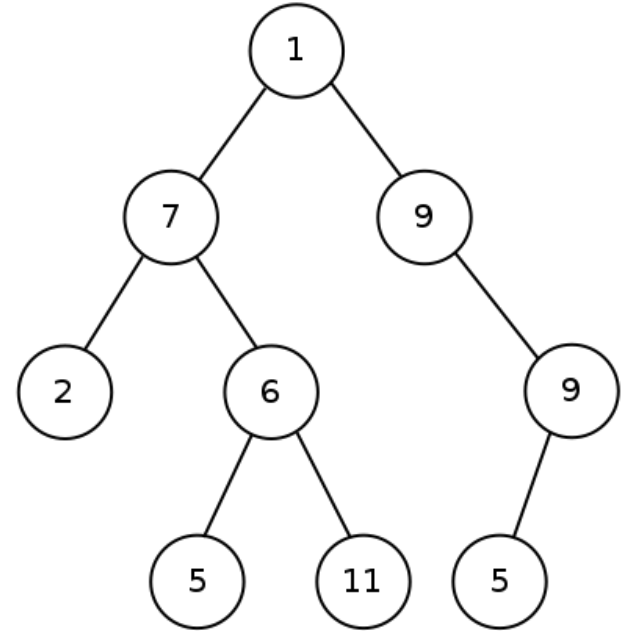
```
void pre-order-traversal(Node* root){  
    // base case  
    if(root == NULL) return;  
    cout << root->val << " "; // curr node  
    // then children  
    for(int i=0; i<(root->children).size(); i++){  
        pre-order-traversal((root->children)[i]);  
    }  
    return;  
}
```



Trees : Traversal

Code for post-order traversal?

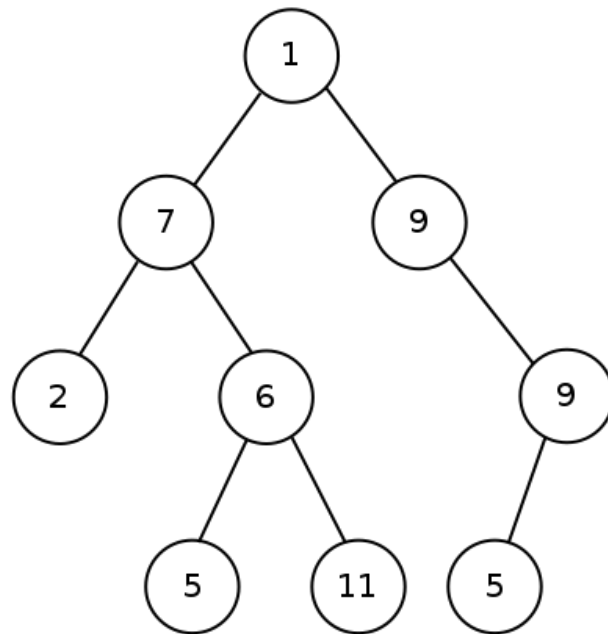
```
void post-order-traversal(Node* root){  
}
```



Trees : Traversal

Code for post-order traversal?

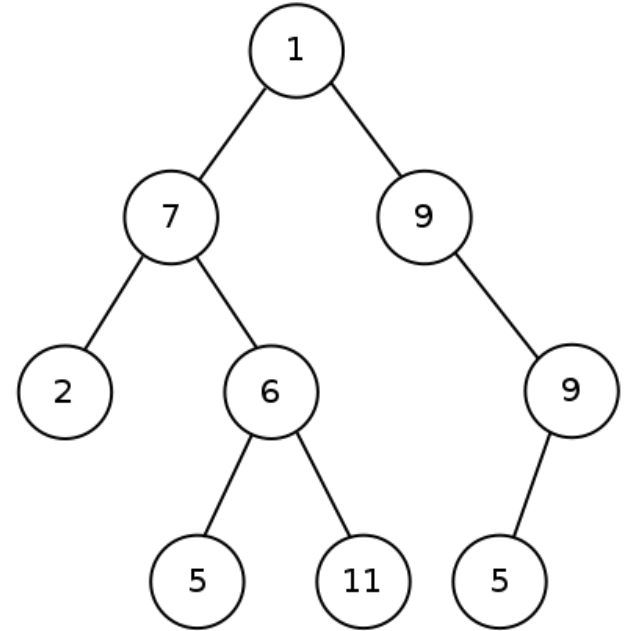
```
void post-order-traversal(Node* root){  
    // base case  
    if(root == NULL) return;  
    // first children  
    for(int i=0; i<(root->children).size(); i++){  
        post-order-traversal((root->children)[i]);  
    }  
    // curr node  
    cout << root->val << " ";  
    return;  
}
```



Trees : Traversal

Code for in-order traversal? (for a binary tree)

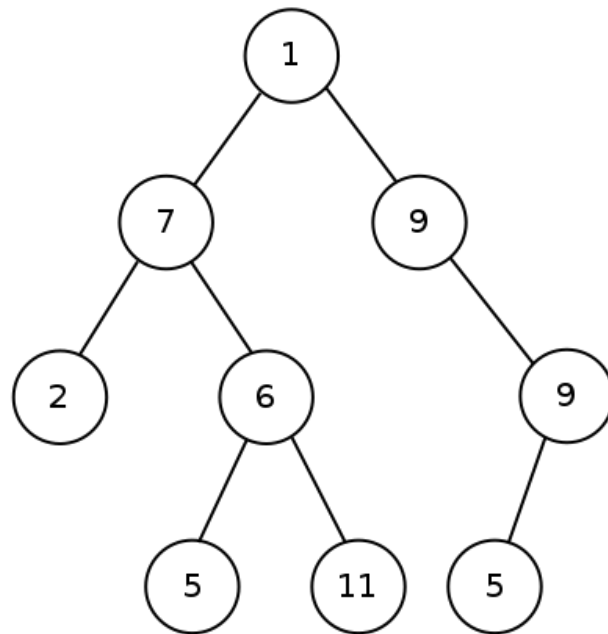
```
void in-order-traversal(Node* root){  
  
}
```



Trees : Traversal

Code for in-order traversal? (for a binary tree)

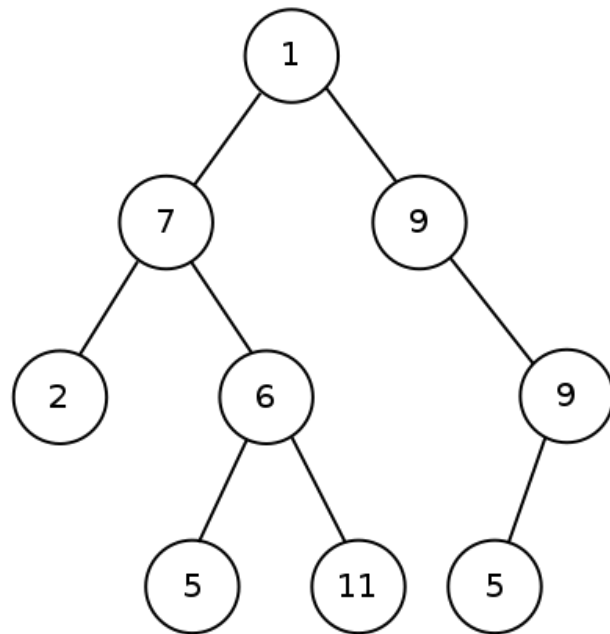
```
void in-order-traversal(Node* root){  
    if(root == NULL) return;  
    in-order-traversal(root→left);  
    cout << root->val << " ";  
    in-order-traversal(root→right);  
}
```



Trees : Height

How to find height of a Tree?

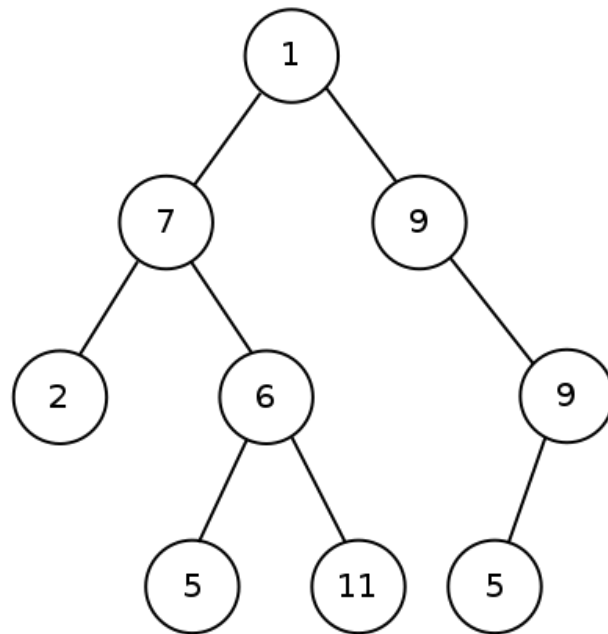
```
int height(Node* root){  
}
```



Trees : Height

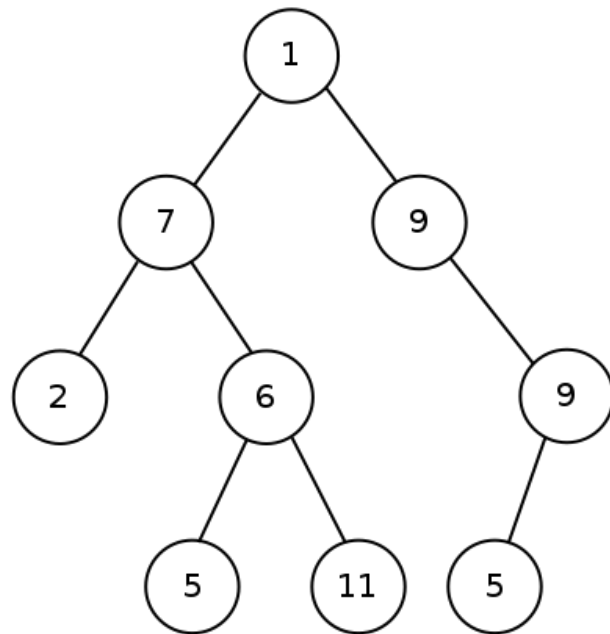
How to find height of a Tree?

```
int height(Node* root){  
    if(root == NULL || (root->children).empty()) return 0;  
    int maxChildHt = INT_MIN;  
    for(int i=0; i<(root->children).size(); i++){  
        maxChildHt = max(maxChildHt, height((root->children)[i]));  
    }  
    return 1+maxChildHt;  
}
```



Trees : Deletion

How to delete a particular value (k) from tree.



Binary Search Trees (BST)

BST : Definition

Organised binary tree such that the following property holds

Given any node **p**,

Let **x** be any node in **left-subtree** of **p**,

Let **y** be any node in **right-subtree** of **p**,

then

$x \rightarrow \text{val} \leq p \rightarrow \text{val} \leq y \rightarrow \text{val}$

BST : Definition

Organised binary tree such that the following property holds

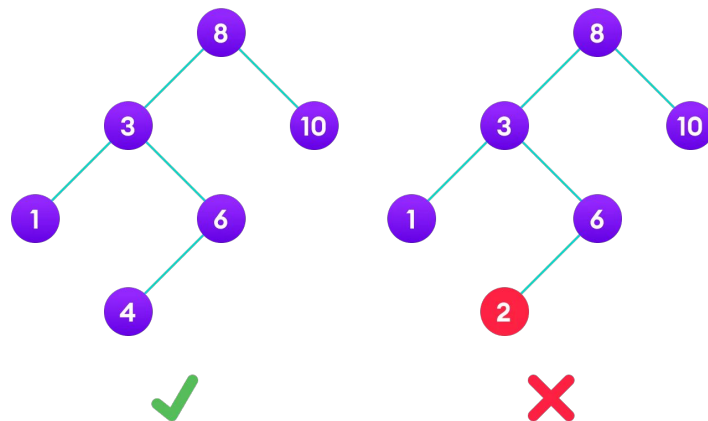
Given any node p ,

Let x be any node in **left-subtree** of p ,

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$x \rightarrow \text{val} \leq p \rightarrow \text{val} \leq y \rightarrow \text{val}$

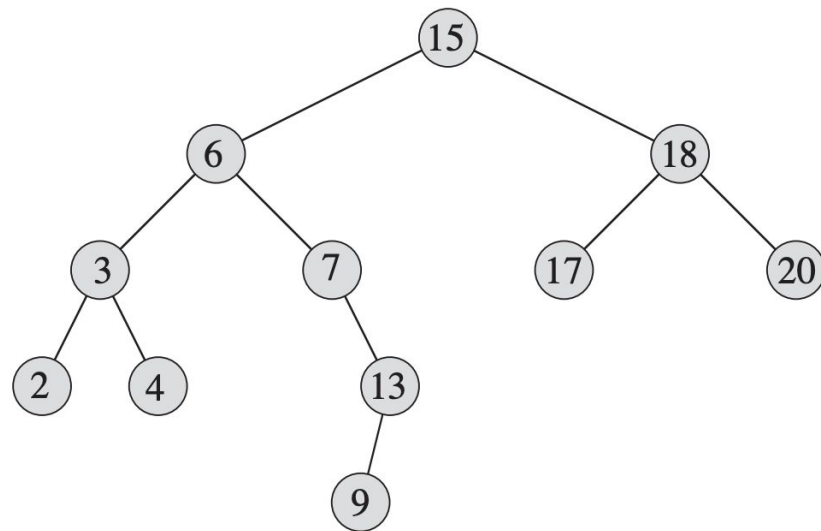


Src : <https://cdn.programiz.com/sites/tutorial2program/files/bst-vs-not-bst.png>

BST : Insertion

Insert a new value (n) in BST rooted at 'root'

Insert should be in such a way that BST property is still valid after insertion.

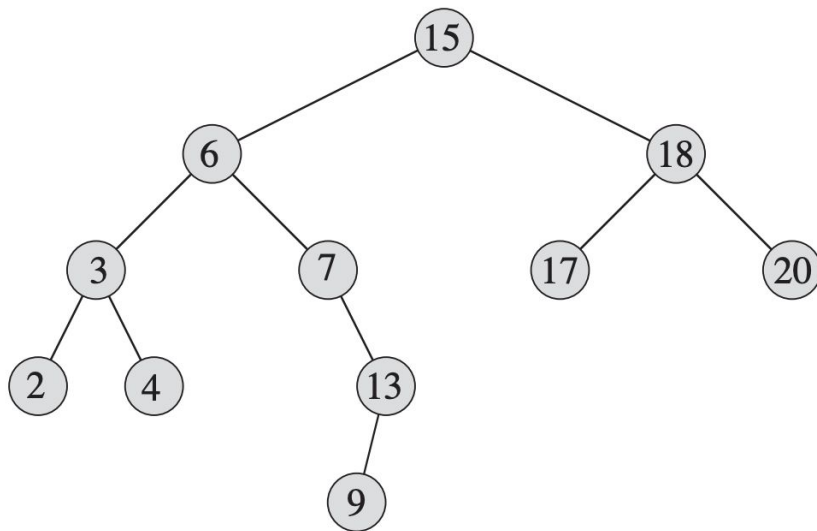


BST : Insertion

Insert a new value (n) in BST rooted at 'root'

Insert should be in such a way that BST property is still valid after insertion.

```
void insert(Node* & root, int n){  
    if(root == NULL){  
        root = new Node();  
        root->val = n;  
        root->left = NULL; root->right = NULL;  
        return;  
    }  
    if(n < root->val){ // insert left  
        insert(root->left, n);  
    }else if (n > root->val){ // insert right  
        insert(root->right, n);  
    }  
    else{ // equal case (already exists)  
        return;  
    }  
}
```



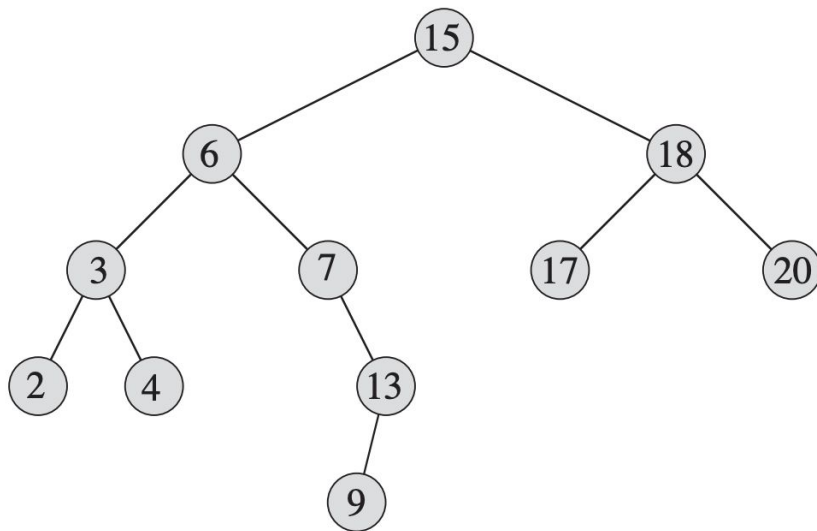
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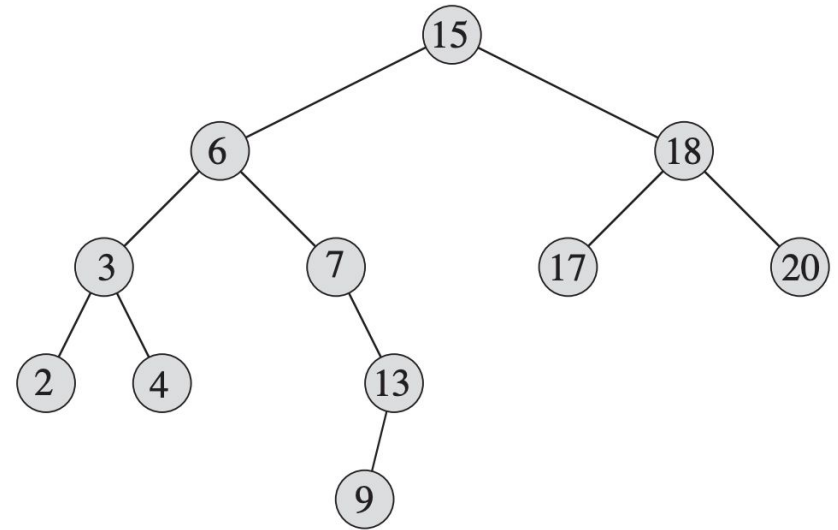
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        return;  
    }  
}
```

Time Complexity : $O(\log N)$ - Average case



BST : Lookup | Search

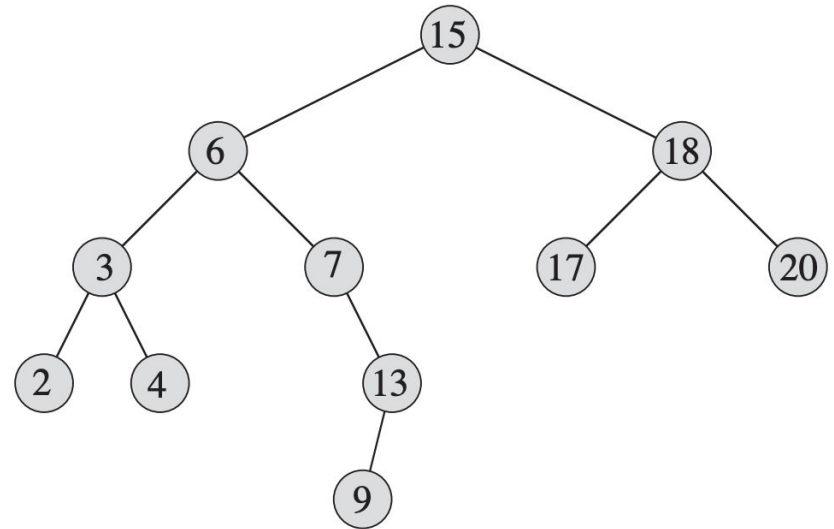
Does a given value (n) exist in BST rooted at 'root'?



BST : Lookup | Search

Does a given value (n) exist in BST rooted at 'root'?

```
bool search(Node* root, int n){  
    if(root == NULL) return false;  
    if(root->val == n) return true;  
    if(n < root->val) return search(root->left, n);  
    else return search(root->right, n);  
}
```

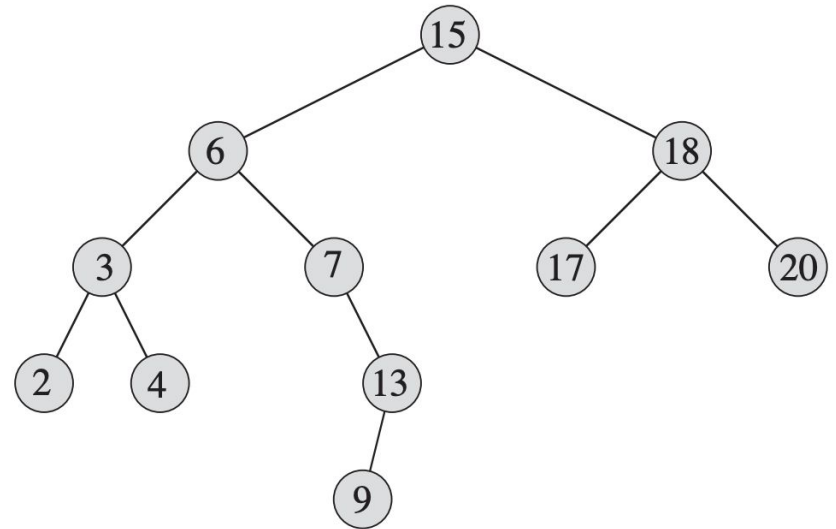


BST : Lookup | Search

Does a given value (n) exist in BST rooted at 'root'?

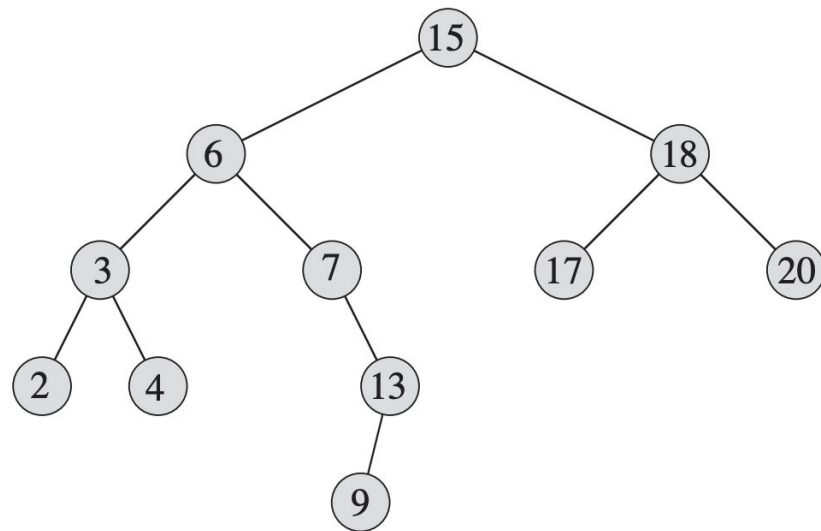
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    if(root == NULL) return false;  
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    else return search(root->right, n);  
}
```

Time Complexity : $O(\log N)$ - Average case



BST : Traversal

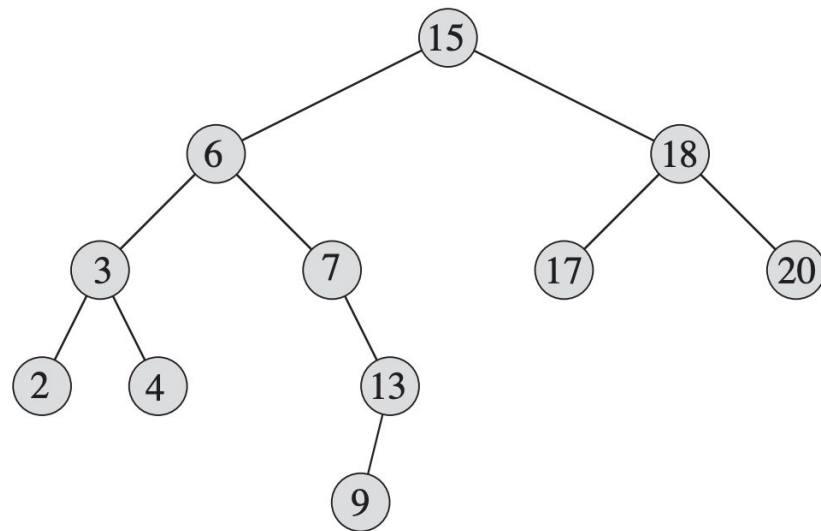
Pre-Order, Post-Order and In-Order
Similar to Binary Tree



BST : Traversal

Pre-Order, Post-Order and In-Order
Similar to Binary Tree

```
void in-order(Node* root){  
    if(root == NULL) return;  
    in-order(root→left);  
    cout << root→val << " ";  
    in-order(root→right);  
}
```

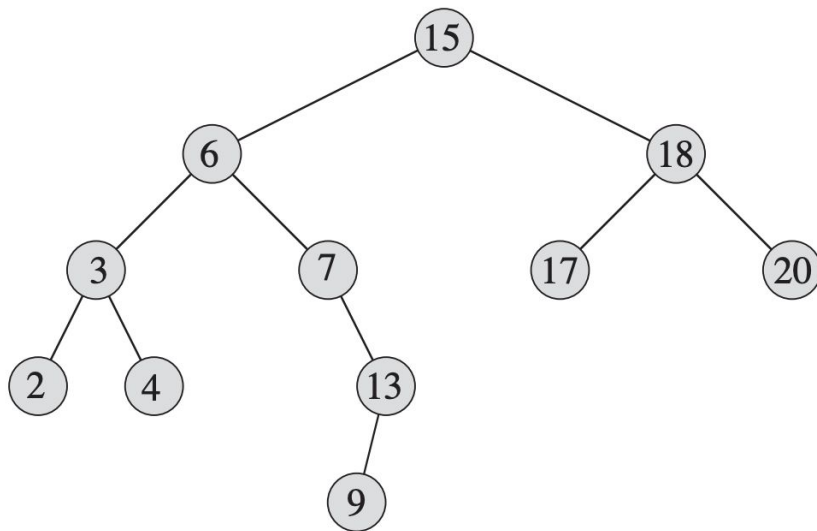


BST : Traversal

Pre-Order, Post-Order and In-Order
Similar to Binary Tree

```
void in-order(Node* root){  
    if(root == NULL) return;  
    in-order(root→left);  
    cout << root→val << " ";  
    in-order(root→right);  
}
```

Time Complexity : $O(N)$



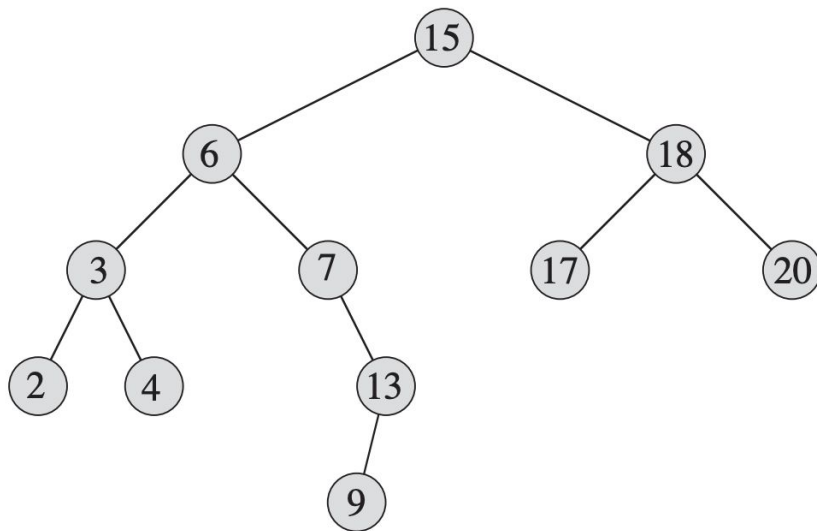
BST : Traversal

Pre-Order, Post-Order and In-Order
Similar to Binary Tree

What does in-order traversal print??

```
void in-order(Node* root){  
    if(root == NULL) return;  
    in-order(root->left);  
    cout << root->val << " ";  
    in-order(root->right);  
}
```

Time Complexity : $O(N)$



BST : Traversal

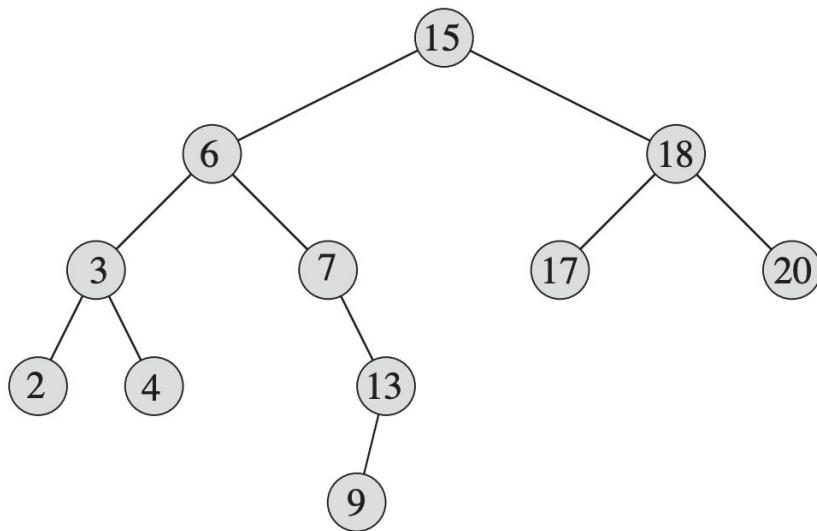
Pre-Order, Post-Order and In-Order
Similar to Binary Tree

What does in-order traversal print??

- Sorted order of elements

```
void in-order(Node* root){  
    if(root == NULL) return;  
    in-order(root->left);  
    cout << root->val << " ";  
    in-order(root->right);  
}
```

Time Complexity : $O(N)$

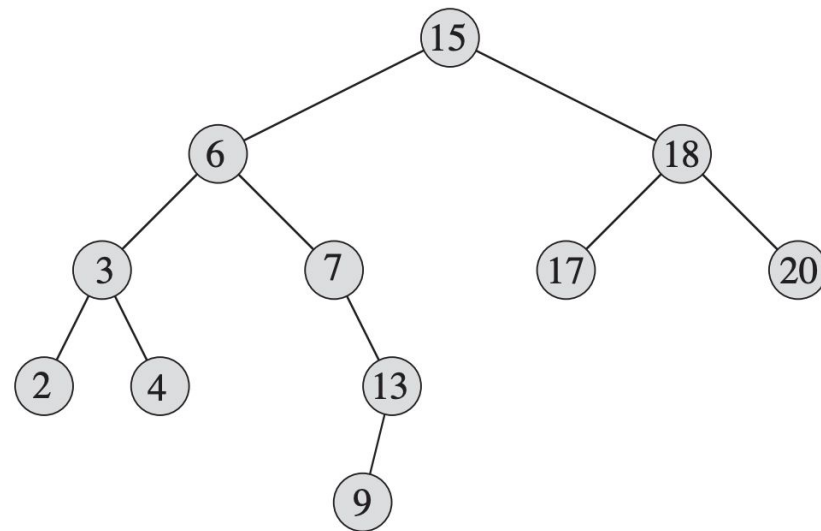


BST : Deletion

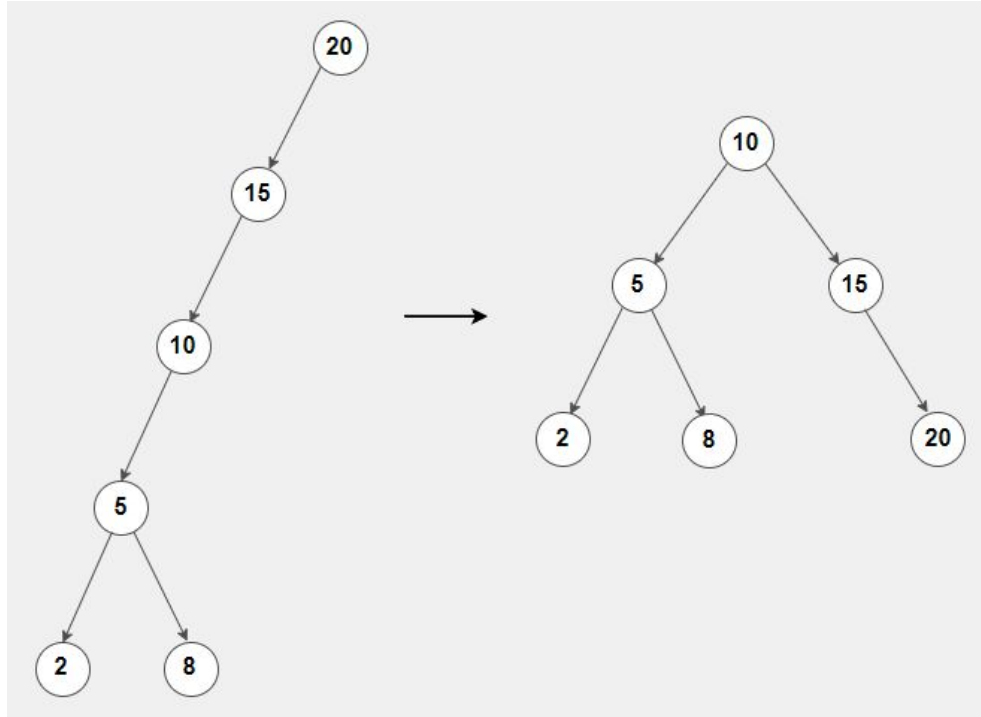
Delete a given value(k) from BST??

```
void delete(Node* root, int k){  
}
```

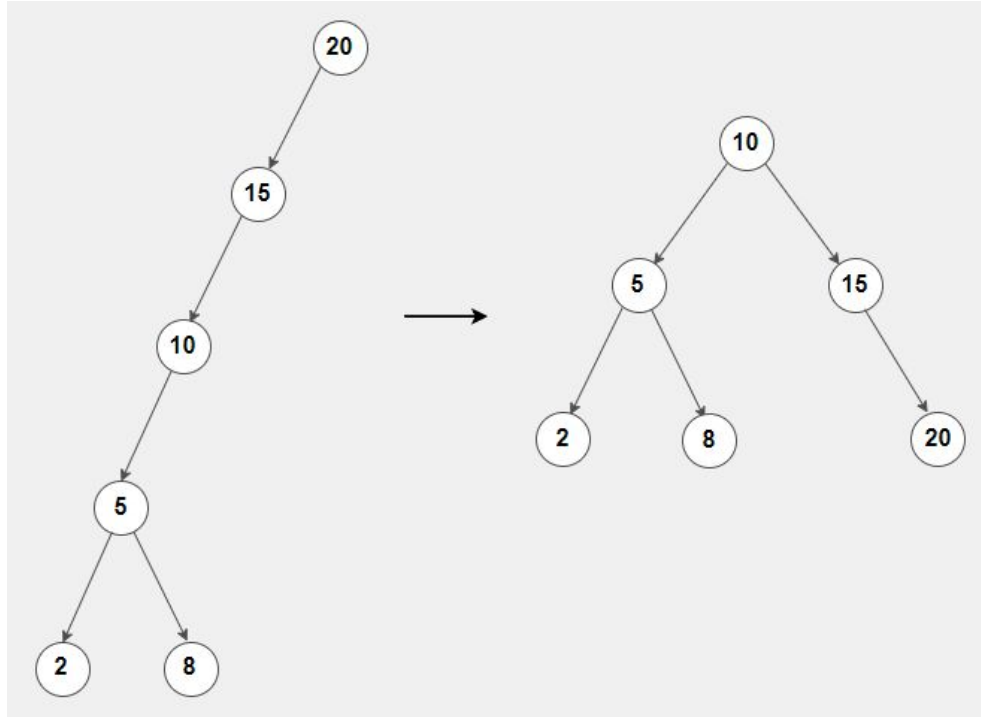
Time Complexity :



BST : Balanced vs Unbalanced



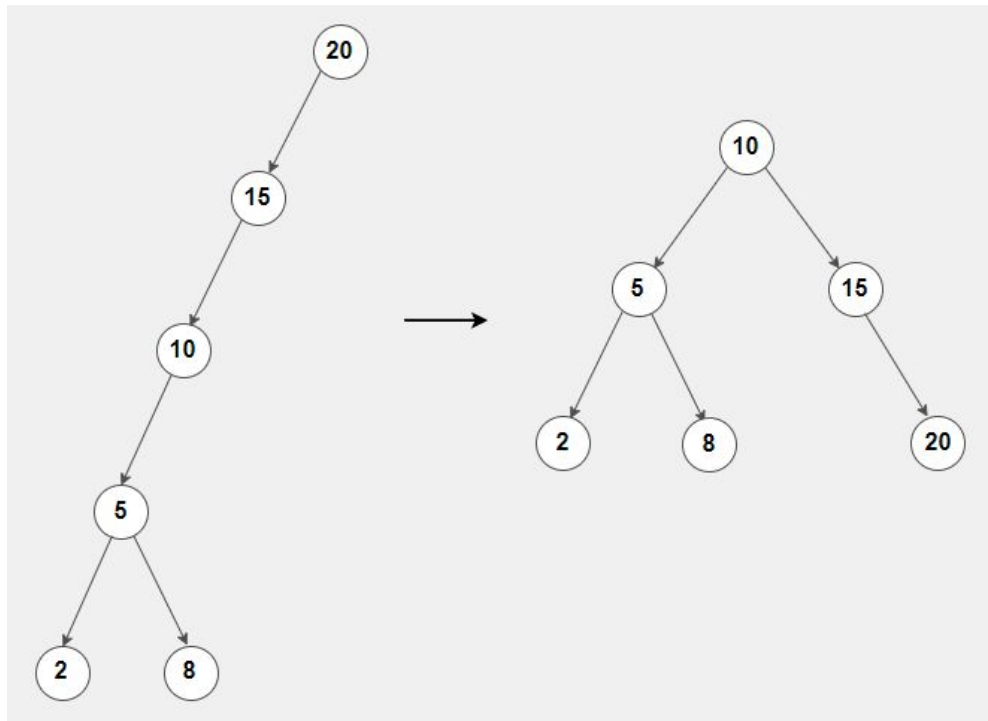
BST : Balanced vs Unbalanced



Balancing schemes

- AVL Tree
- Red-Black Tree
- 2-3 Tree
- etc....

BST : Balanced vs Unbalanced



Balancing schemes

- AVL Tree
- Red-Black Tree
- 2-3 Tree
- etc.....

Basic operations on balanced ones has **$O(\log N)$** complexity.

Unbalanced might go to worst case of $O(N)$

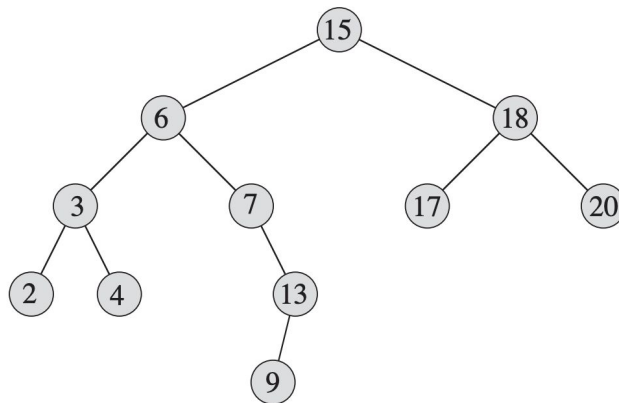
Keep in mind the **cost of balancing**.

BST : Problems

Find k'th smallest element in a BST??

BST Node is defined as

```
struct Node{  
    int val;  
    Node* left;  
    Node* right;  
    int m_size; //number of nodes in subtree.  
}
```



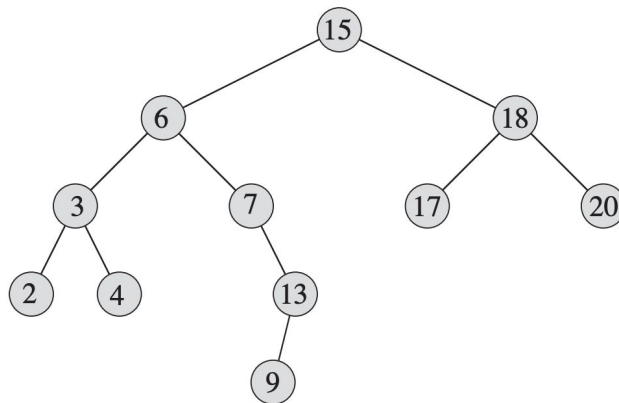
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}
```

```
int Ksmallest(Node* root, int k){  
    int l_size = 0;  
    if(root->left != NULL) l_size = root->left->m_size;  
    if(l_size==k-1) return root->val;  
    if(l_size < k-1) return Ksmallest(root->right, k-1-l_size);  
    else return Ksmallest(root->left, k);  
}
```



BST : Problems

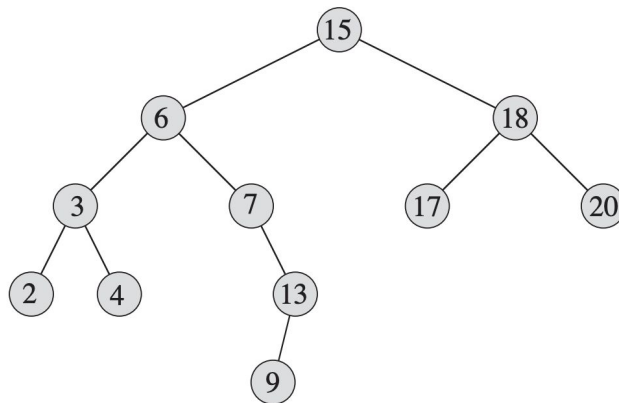
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}
```

Time Complexity : $O(\log N)$ - Average case



BST : Problems

Given N elements, what is time-complexity of converting them to BST structure ??

BST : Problems

Given N elements, what is time-complexity of converting them to BST structure ??

$O(N \log N)$

How to find median of an input data stream ??

BST : Usage

STL **Set**, **Multiset** and **Map** use self-balancing BST.

```
#include <set> // includes both set and multiset
```

```
// declaration
```

```
set<type> s; // a set do not keep duplicates
```

```
multiset<type> ms; // a multiset can have duplicates
```

```
#include <map> // for map usage
```

```
// declaration
```

```
map<keyType, valueType> mp;
```

```
// stores paired data, (key, value)
```

```
// auto sorted by keys
```

```
// do not allow duplicate keys
```

refer www.cplusplus.com for detailed usage

References

BST

<https://www.programiz.com/dsa/binary-search-tree>

Content in some of the slides is taken from Yiyou Chen.