

# **CS 32 Week 10**

## **Discussion 1I**

Srinath

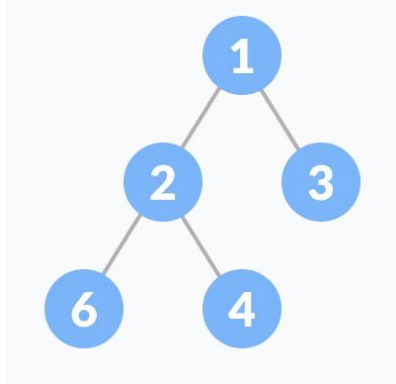
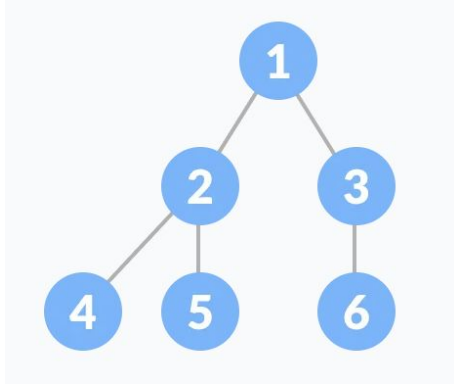
# Outline

- Complete Binary Tree
- Heaps
- Priority Queue : STL
- Worksheet 9 - Heaps

# **Complete Binary Tree**

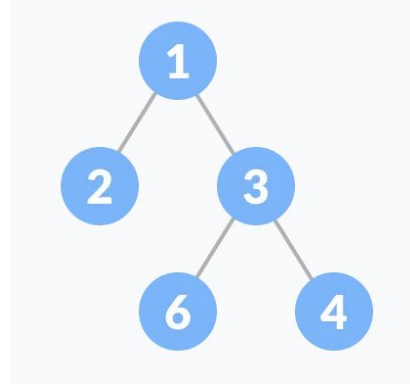
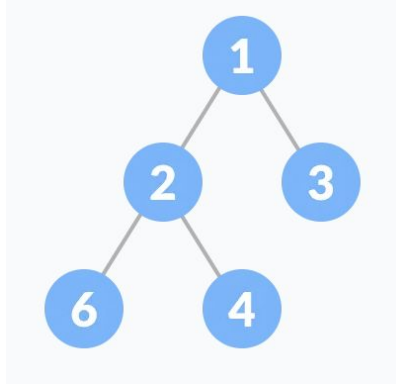
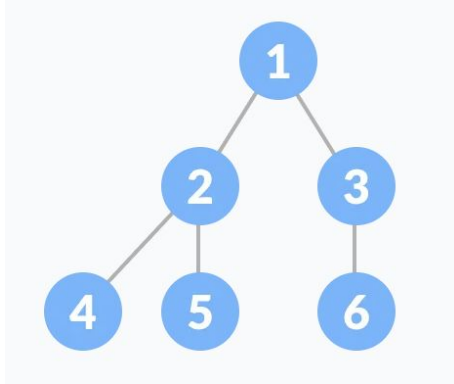
# Complete Binary Tree : Definition

A binary tree with **all levels** completely **filled** except the lowest one, Lowest one is filled from left



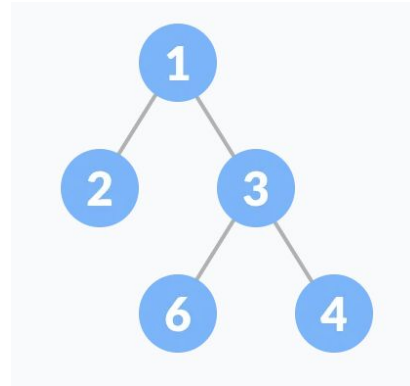
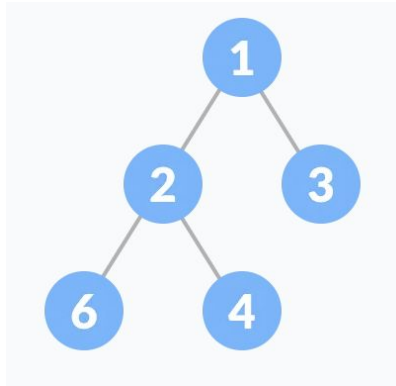
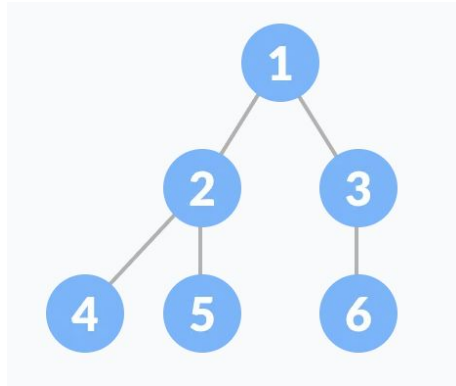
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# Complete Binary Tree : Definition

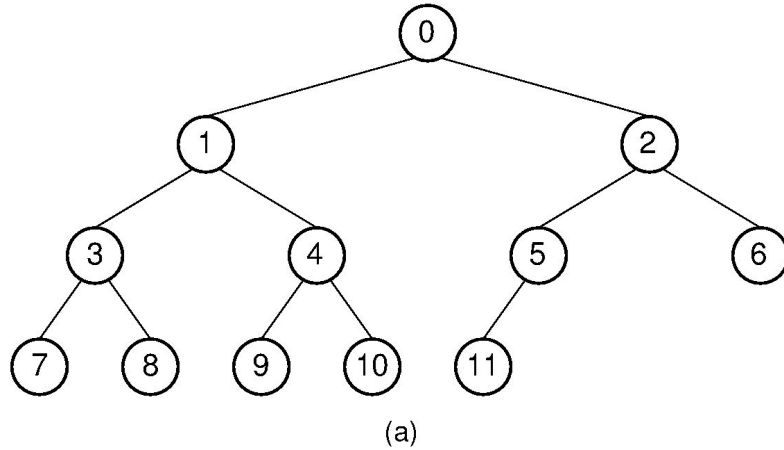
A binary tree with **all levels** completely **filled** except the lowest one, Lowest one is filled from left



Not a CBT

# Complete Binary Tree : Definition

A binary tree with **all levels** completely **filled** except the lowest one, Lowest one is filled from left



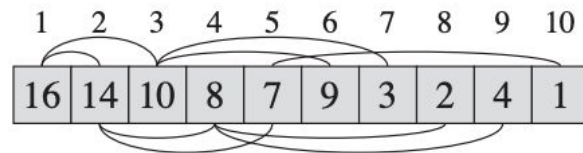
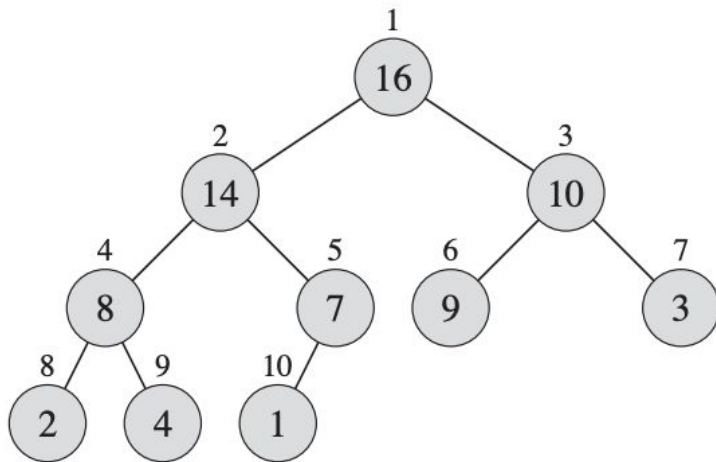
# Complete Binary Tree : Representation

An **N** Node CBT can be represented using an array

For each node at **i**, its **left** child is  **$2*i$** , **right** child is  **$2*i+1$** , **parent** is  **$\lfloor i/2 \rfloor$**

Or..

For each node at **i**, its **left** child is  **$2*i+1$** , **right** child is  **$2*i+2$** , **parent** is  **$\lfloor (i-1)/2 \rfloor$**  (if you prefer **0** - indexing)





# Heaps

# Heaps : Definition

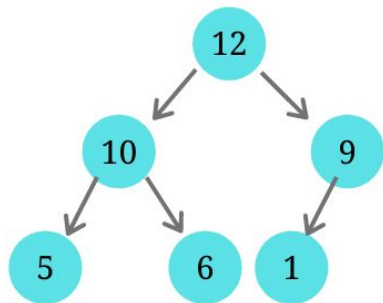
Heap is a CBT, satisfying **heap-property**.

For a **max-heap**, all parents have greater value than their children.

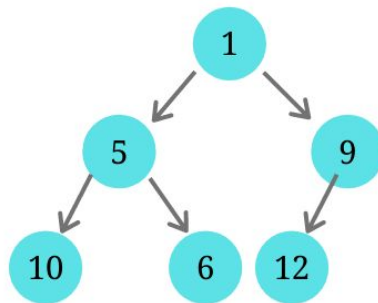
For a **min-heap**, all parents have smaller value than their children.

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Max-Heap



Min-Heap

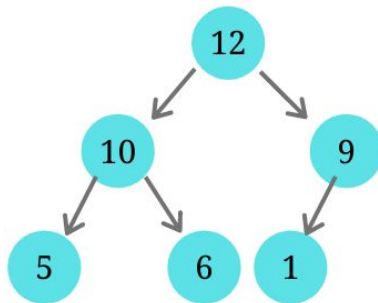


# Heaps : Definition

As it is a CBT, you can use an array or vector to represent it.

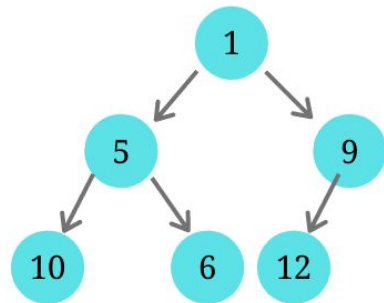
```
//heap of valuetype double  
//may also use fixed-size array for heap  
vector<double> heap;
```

Max-Heap



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Min-Heap



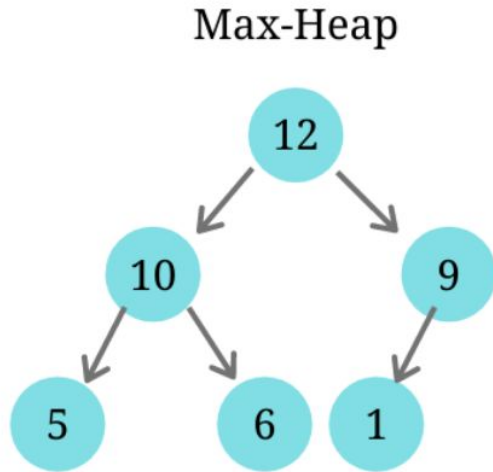
# Heaps : Insertion

How to insert a new element into existing heap?

```
void insert(vector<double>& heap, const double& val) {  
    //insert a value val to heap  
}
```

# Heaps : Insertion

How to insert a new element(say 25) into existing heap?



# Heaps : Insertion

```
void insert(vector<double>& heap, const double& val) {  
    heap.push_back(val);  
    int cur_ind = heap.size() - 1;  
    while(cur_ind != 0  
        && heap[cur_ind] > heap[(cur_ind-1)/2]) {  
        swap(heap[cur_ind], heap[(cur_ind-1)/2]);  
        cur_ind = (cur_ind-1)/2;  
    }  
}
```

# Heaps : Deletion

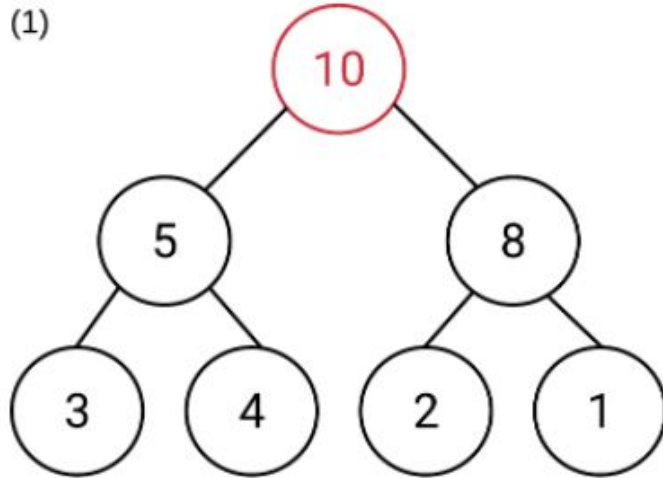
We are interested in deleting the root element of heap in general.

Why?

# Heaps : Deletion

We are interested in deleting the root element of heap in general.

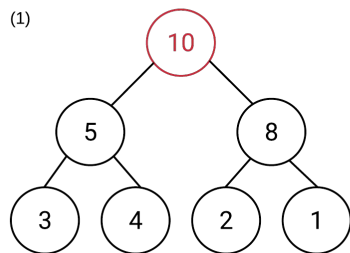
**Why?** - We use up(consume) the max/min element in general



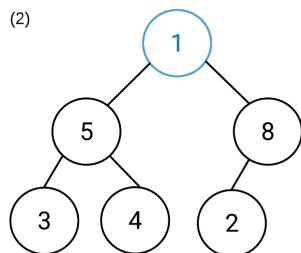


# Heaps : Deletion

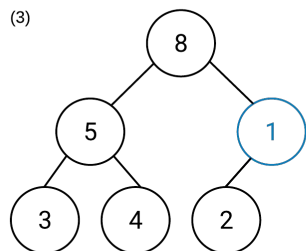
## Deleting from this heap



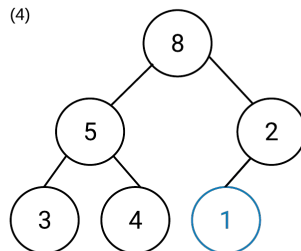
Starting with this max heap



Step 1: the bottom most, left most node, the 1 node, gets placed at the root



Step 2: Because 1 is less than both of its children, it swaps with the larger element, the 8 node



Step 3: Once again, 7 is bigger than its parent, the 6 node, so it gets swapped

# Heaps : Deletion

Recursive

Swap **root** and **last** elements  
**pop** the last element.  
call MAX-HEAPIFY(A, 0);

MAX-HEAPIFY( $A, i$ )

```
1   $l = \text{LEFT}(i)$ 
2   $r = \text{RIGHT}(i)$ 
3  if  $l \leq A.\text{heap-size}$  and  $A[l] > A[i]$ 
4       $\text{largest} = l$ 
5  else  $\text{largest} = i$ 
6  if  $r \leq A.\text{heap-size}$  and  $A[r] > A[\text{largest}]$ 
7       $\text{largest} = r$ 
8  if  $\text{largest} \neq i$ 
9      exchange  $A[i]$  with  $A[\text{largest}]$ 
10     MAX-HEAPIFY( $A, \text{largest}$ )
```

# Heaps : Deletion

## Non-Recursive

```
void remove(vector<double>& heap) {
    swap(heap[0], heap[heap.size() - 1]); //swap root to leaf
    heap.pop_back();
    int sz = heap.size();
    int cur = 0;
    while (2 * cur + 1 < sz) {
        if (2 * cur + 2 >= sz) { //only left child exists
            if (heap[2 * cur + 1] > heap[cur]) {
                swap(heap[2 * cur + 1], heap[cur]);
                cur = 2 * cur + 1;
            }
            else break;
        }
        else {
            //larger than both left and right
            if (heap[cur] > heap[2 * cur + 1] && heap[cur] > heap[2 * cur + 2]) {
                break;
            }
            //pick the larger element of left and right
            if (heap[2 * cur + 1] > heap[2 * cur + 2]) {
                swap(heap[cur], heap[2 * cur + 1]);
                cur = 2 * cur + 1;
            }
            else {
                swap(heap[cur], heap[2 * cur + 2]);
                cur = 2 * cur + 2;
            }
        }
    }
}
```

# Heaps : Max or Min

How to get Max or Min element?

```
double get_max(const vector<double>& heap) {  
    //return maximal element of the heap  
}
```

# Heaps : Max or Min

How to get Max or Min element?

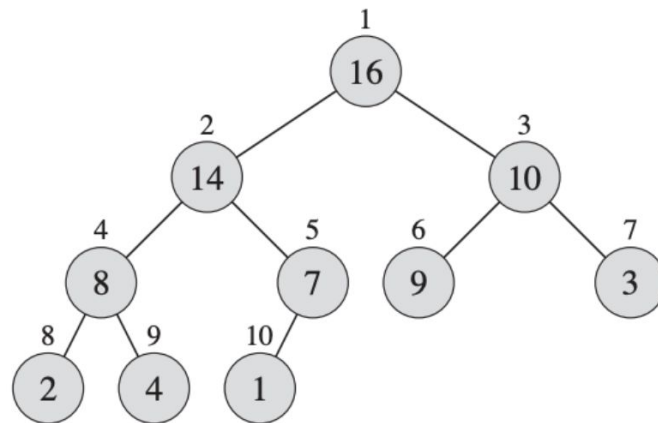
```
double get_max(const vector<double>& heap) {  
    return heap[0];  
}
```

# Heaps : HeapSort

Given a heap(that means it satisfies the heap property)

How to get the elements sorted?

```
void heap_sort(vector<double>& heap) {  
    //sort the heap  
}
```

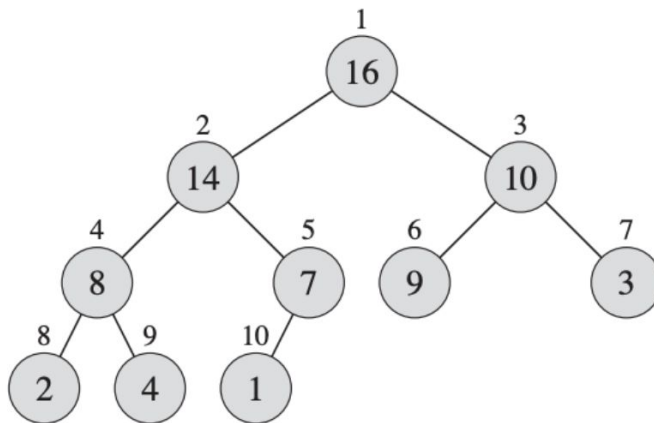


# Heaps : HeapSort

Given a heap(that means it satisfies the heap property)

How to get the elements sorted?

```
void heap_sort(vector<double>& heap) {  
    if (heap.size() <= 1) return;  
    double val = heap[0]; //save the largest  
    remove(heap); //remove the largest  
    heap_sort(heap); //sort the rest  
    heap.push_back(val); //add largest back  
}
```



# Heaps : Complexity

	Average	worst
Insertion:		
Deletion:		
Get_max for max heap:		
Heap_sort:		



# Heaps : Complexity

	Average	worst
Insertion:	$O(\log N)$	$O(\log N)$
Deletion:	$O(\log N)$	$O(\log N)$
Get_max for max heap:	$O(1)$	$O(1)$
Heap_sort:	$O(N \log N)$	$O(N \log N)$

# **Priority Queue : STL**

# Priority Queue :

A linear data structure.

Looks like a queue, but totally different. (queue uses linked list, **priority\_queue** uses **heap**). For standard types, the priority is **larger** values (max heap), but like **set** and **map**, one can **overload** the **< operator** or define a priority comparator.

Like a heap, a `priority_queue` is not totally sorted. But **its top element is guaranteed** to have the **highest priority** among all elements. It **automatically adjust** the heap after each **pop** and **push**.

## *fx* Member functions

<b>(constructor)</b>	Construct priority queue (public member function )
<b>empty</b>	Test whether container is empty (public member function )
<b>size</b>	Return size (public member function )
<b>top</b>	Access top element (public member function )
<b>push</b>	Insert element (public member function )
<b>emplace</b> <small>C++11</small>	Construct and insert element (public member function )
<b>pop</b>	Remove top element (public member function )
<b>swap</b> <small>C++11</small>	Swap contents (public member function )

# Priority Queue : Custom Comparator

```
struct LessThanByAge
{
    bool operator()(const Person& lhs, const Person& rhs) const
    {
        return lhs.age < rhs.age;
    }
};
```

then instantiate the queue like this:

```
std::priority_queue<Person, std::vector<Person>, LessThanByAge> pq;
```

If you just need min-priority queue instead of max

```
priority_queue <int, vector<int>, greater<int>> g
```

# Priority Queue : Example

```
priority_queue<int> g1;
priority_queue<int, vector<int>, greater<int>> g2;
int b[5] = {3, 2, 6, 1, 8};
for (int i = 0; i < 5; ++i) {
    g1.push(b[i]);
    g2.push(b[i]);
}
while(!g1.empty()) {
    cout << g1.top() << endl;
    g1.pop();
}
while(!g2.empty()) {
    cout << g2.top() << endl;
    g2.pop();
}
```

Output:

# Priority Queue : Example

```
priority_queue<int> g1;  
priority_queue<int, vector<int>, greater<int>> g2;  
int b[5] = {3, 2, 6, 1, 8};  
for (int i = 0; i < 5; ++i) {  
    g1.push(b[i]);  
    g2.push(b[i]);  
}  
while(!g1.empty()) {  
    cout << g1.top() << endl;  
    g1.pop();  
}  
while(!g2.empty()) {  
    cout << g2.top() << endl;  
    g2.pop();  
}
```

Output:

8  
6  
3  
2  
1  
1  
2  
3  
6  
8

# Priority Queue : Complexity

	Average	worst
push:	$O(\log N)$	$O(\log N)$
pop:	$O(\log N)$	$O(\log N)$
top:	$O(1)$	$O(1)$

# Priority Queue : Sample Problem

How to use priority\_queues(heaps) to keep track of the median of a data stream?



# Priority Queue : Sample Problem

How to use priority\_queues(heaps) to keep track of the median of a data stream?

Hint :

- Use 2 priority queues(PQ's)

- 1 for left half (max-priority)

- 1 for right half (min-priority)

- For every new element

- push into left PQ.

- get top element of left PQ, push it into right PQ.

- pop top element of left PQ.

# STL Data Structures : Summary

**Unordered\_set (Hash):** sorted/unordered?. insertion, deletion, look-up

**Set (BST):** sorted/unordered? insertion, deletion, look-up.

**Unordered\_map (Hash):** sorted/unordered? insertion, deletion, look-up.

**Map (BST):** for mapping, sorted/unordered?. insertion, deletion, look-up.

**Priority\_queue (heap):** for knowing extreme values, sorted/unordered?. knowing the max(min) from max(min) heap. insertion, deletion, look-up.

# STL Data Structures : Summary

**Unordered\_set (Hash):** fast for look-up, **unsorted**.  $O(1)$  for insertion, deletion, look-up.

**Set (BST):** for look-up, **sorted**.  $O(\log N)$  for insertion, deletion, look-up.

**Unordered\_map (Hash):** fast for mapping, **unsorted**.  $O(1)$  for insertion, deletion, map by key.

**Map (BST):** for mapping, **sorted**.  $O(\log N)$  for insertion, deletion, map by key.

**Priority\_queue (heap):** for knowing extreme values, **unsorted**.  $O(1)$  for knowing the max(min) from max(min) heap.  $O(\log N)$  for insertion, deletion.  $O(N)$  for look-up.

# References

## Chapter 6 - HeapSort

- Introduction to Algorithms, T.H Cormen
- [https://web.cs.ucla.edu/~srinath/static/pdfs/DataStructures&Algorithms\\_Cormen.pdf](https://web.cs.ucla.edu/~srinath/static/pdfs/DataStructures&Algorithms_Cormen.pdf)

Most content of the slides is taken from Yiyou Chen.

# **Thank You!!**

**Good luck for your finals!**

# Graphs

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A data structure to store key-value pairs.  
Something like a dictionary.

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