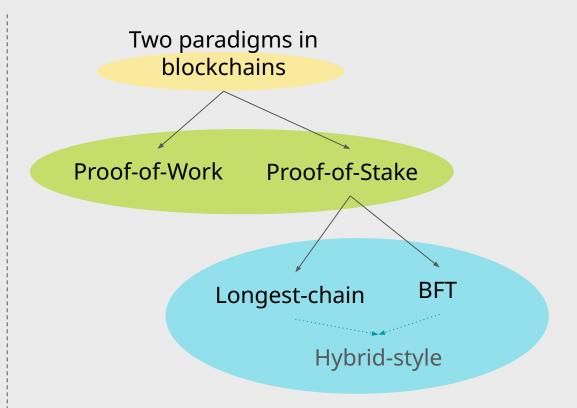
An overview of proof-of-work and proof-of-stake consensus

Valeria (Lera) Nikolaenko
The 13th BIU Winter School on cryptography



Outline

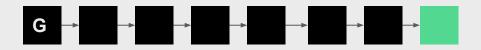
Def a blockchainPoW vs PoSLongest-chain vs BFT



0: What is a blockchain?

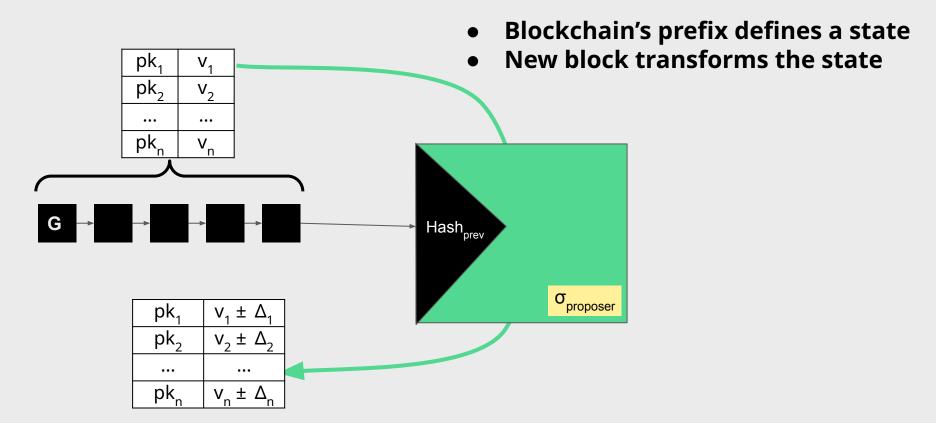
Blockchain: what is it?

Blockchain is a distributed ledger - a log of transactions*.

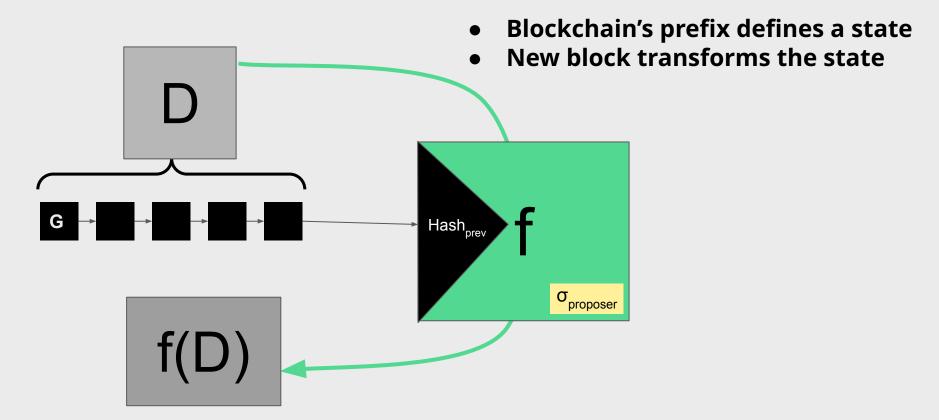


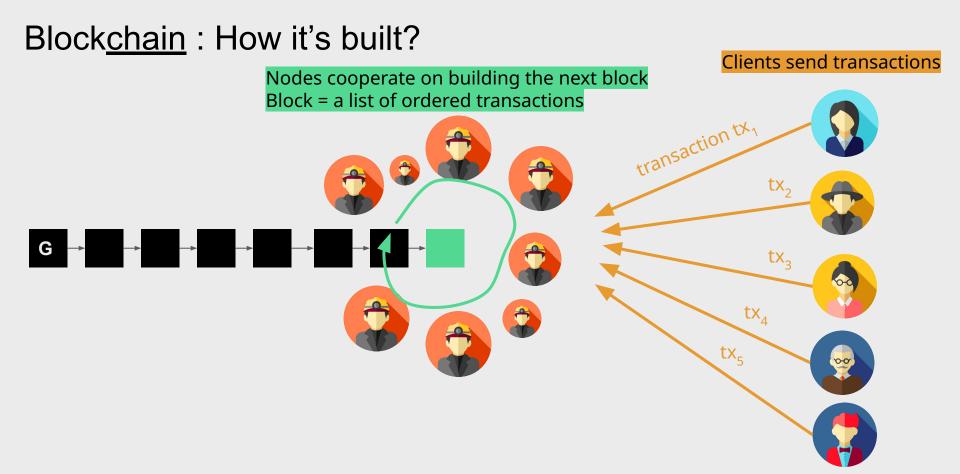
- Blockchain is a chain of blocks.
- Each block is a list of transactions + hash of the previous block.
- <u>Safety</u>: everybody agrees on history.
- <u>Liveness</u>: every valid transaction eventually gets added
 not censoring

Blockchain: what is it?

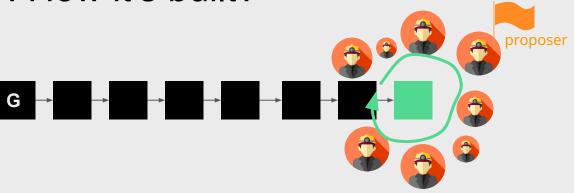


Blockchain: what is it?





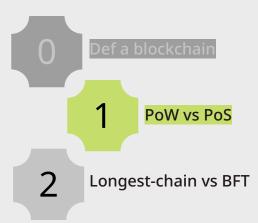
Blockchain: How it's built?

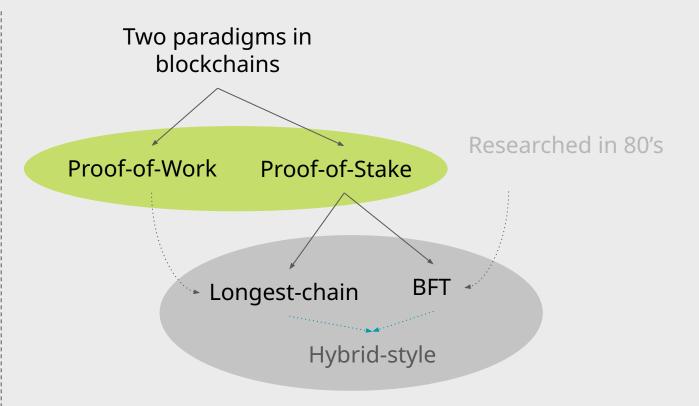


- A proposer(s) is elected to announce a new block
 Pre-electing a proposer limits the options for Step 2
- 2. Everybody decides whether to accept this block or ignore

Consensus mechanism

Outline





1: Proof-of-Work vs. Proof-of-Stake for proposer election

Proof-of-Work for proposer election

- The proposer searches for a block B, s.t. Hash(B) < target.
 - For a hash function Hash: $\{0,1\}^* \rightarrow \{0,1\}^{256}$
 - Find B, s.t. Hash(B) $< 2^{256}$ / D (takes time O(D) to solve)
 - D "difficulty"
 - Currently for Bitcoin D ≈ 2^{45} a block is found once in 10 min

Proof-of-Work for proposer election

- The proposer searches for a block B, s.t. Hash(B) < target.
- Once found, announce B it to peers.
- Easy to verify the proposer.

Blocks are hard to find => small number of possible proposers



This proposer election mechanism:

- Live: blocks are proposed regularly
- Fair: each miner is elected proportional to its computational power
- Permissionless: open participation

Proof-of-Work electricity loss

Bitcoin: <u>85 TWh/yr</u> (June 2022)

• Ethereum: <u>50-100 TWh/yr</u> (June 2022)

Argentina: 130 TWh/yr

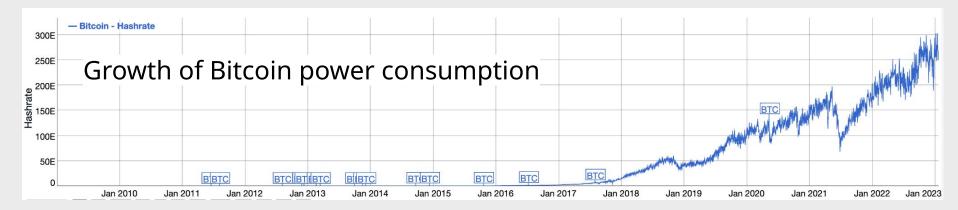
Mining gold: 240 TWh/yr

Market-cap of coins:



Jan 2022, https://coin360.com/

|Bitcoin + Ethereum| < Mining gold

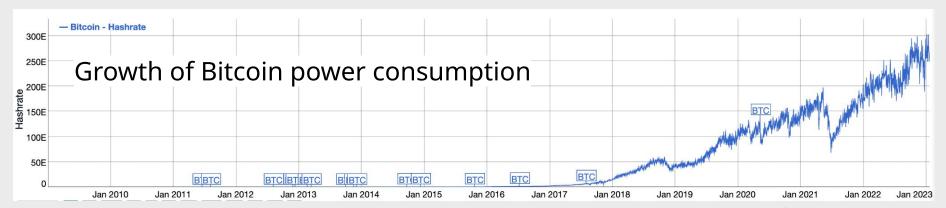


Useful number to keep in mind!

Bitcoin hash-rate:

- 200 EH/s = 2^{76} hashes per 10 min
- 6.25 BTC per block/per 10 min ≈ \$150,000

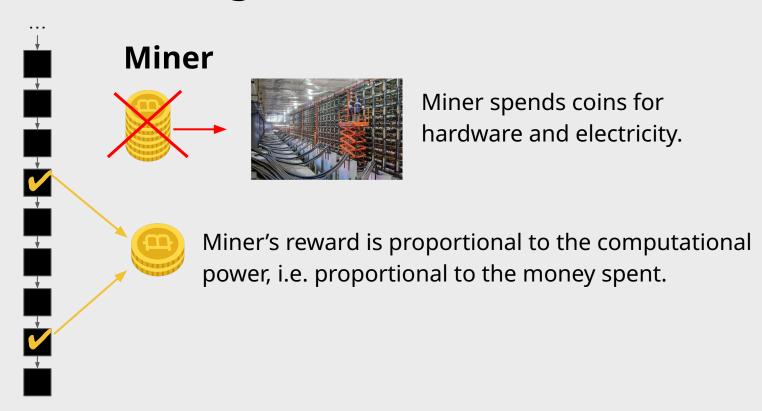




Can we make PoW better?

- Increase the block size
 - Delays in block propagation => chain will fork more => decreases reliability
- Increase the frequency of the blocks
 - o Gervais et al. (CCS'16): Bitcoin block frequency can be reduced to 1 min
- Lightning network: off-chain transactions
- Change the proof-of-work puzzle
 - o proof-of-useful-work
- Bitcoin-NG:
 - use PoW sporadically to elect the leader for a longer duration of time (to propose multiple blocks).

Proof-of-Work mining



Proof-of-Stake mining



Proof-of-Stake for proposer election

- The validator (miner) locks* T coins
- N number of validators
- Simplest approach: round-robin
- The validator is <u>elected</u> as a proposer with probability 1/N

Desired properties:

- Live
- Fair 🗸
- Permissionless: open participation
 - Validator needs to be able to acquire coins.

^{* &}quot;locking" coins requires the blockchain to have basic smart-contract capabilities

Staking design choices

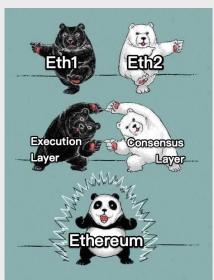
- Required coins to be locked can be
 - o fixed (e.g. Ethereum 2.0: 32 ETH) or
 - variable
- Delegated-proof-of-stake : allows stakeholders delegate their stake to validators.
- Some protocols want to limit the number of validators:
 - o Random selection of a subcommittee
 - Delegate stake to active nodes

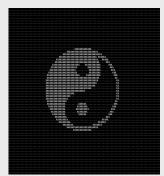
This talk: fixed staking, no delegation.

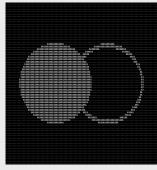
Ethereum switched from PoW to PoS

- The Merge Sep 15, 2022
- Ethereum reduced electricity use by 30,000x
- World electricity consumption → ~99.8%

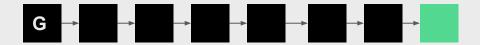








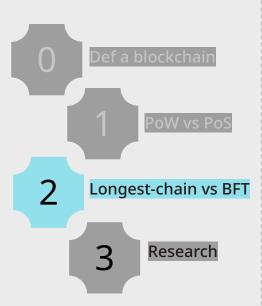
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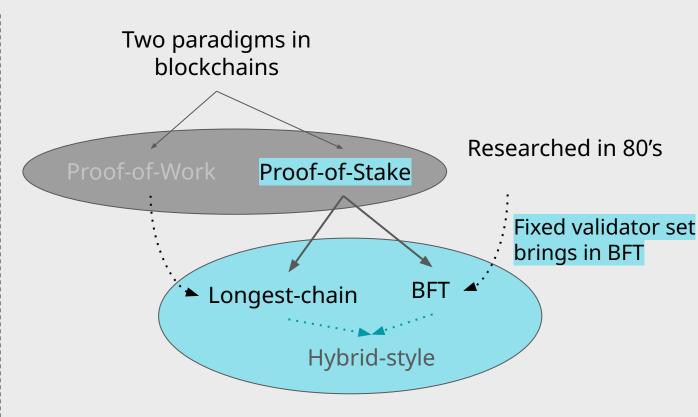


- 1. A proposer(s) is elected to announce a new block
 - a. Proof-of-Work or
 - b. Proof-of-Stake
- 2. Everybody decides whether to accept this block or not
 - a. Longest-chain or
 - b. BFT

Consensus mechanism

Outline





Consensus: agree on proposed block

- Safety: everybody agrees on history
- Liveness: every valid transaction eventually gets added
- Resilience against byzantine validators
- Minimal network assumption
 - o **Synchronous**: shared global clock + known bound Δ on network delay
 - \circ **Partially-synchronous**: shared global clock + exists unknown bound Δ on network delay (equivalent: synchronous after unknown GST)
 - Asynchronous: no bounds on network delays (subsumes network partitions)

Possibility-Impossibility Results are well known for <u>fixed</u> validator set

	Liveness	Safety	Byzantine threshold	
Synchrony	YES	YES	ANY* (<50%)	OUR FOCUS
Partial-synchrony	YES	YES	< 33%	
Asynchrony*	NO*	NO*	≤ 1	
(Partial)-Synchrony with network partitions	YES	NO	1	
	NO	YES	1	

More Resources

- Youtube Lectures by Tim Roughgarden
- Consensus cannon from a16z: https://a16zcrypto.com/consensus-canon/
- Ittai Abraham:
 https://decentralizedthoughts.github.io/







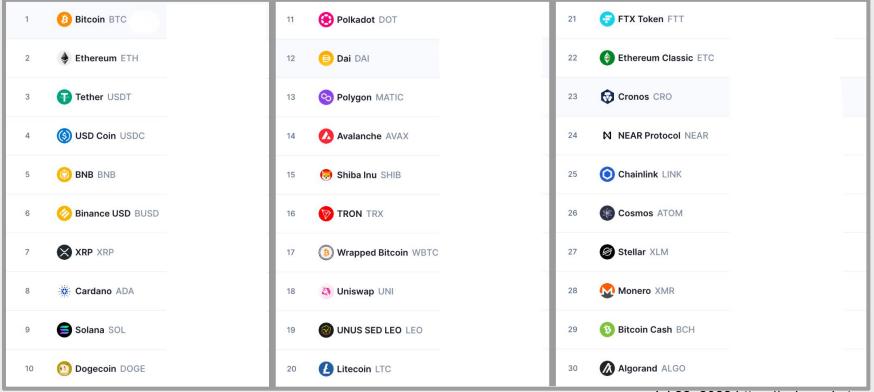
Andrew Lewis-Pye

Desired consensus properties

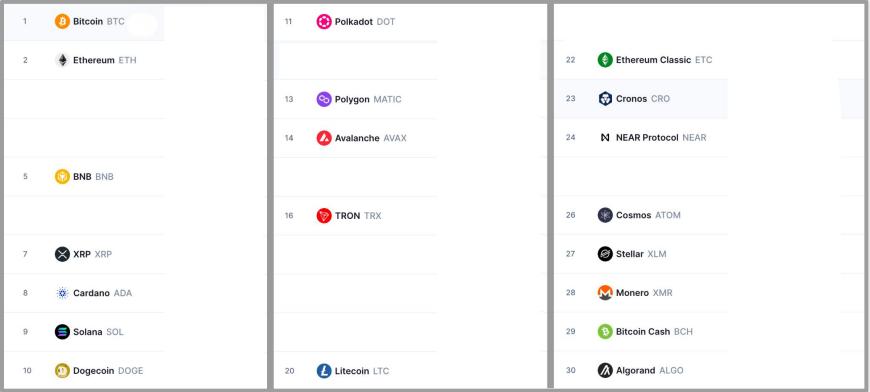
- Safety: validators agree on history
- Liveness: every valid transaction eventually gets added
- Resilience against ½ byzantine validators
- Partially-synchronous: exists unknown bound Δ on network delay
 - + shared global clock

2: Longest-chain vs. BFT consensus

Blockchains by market cap: 30

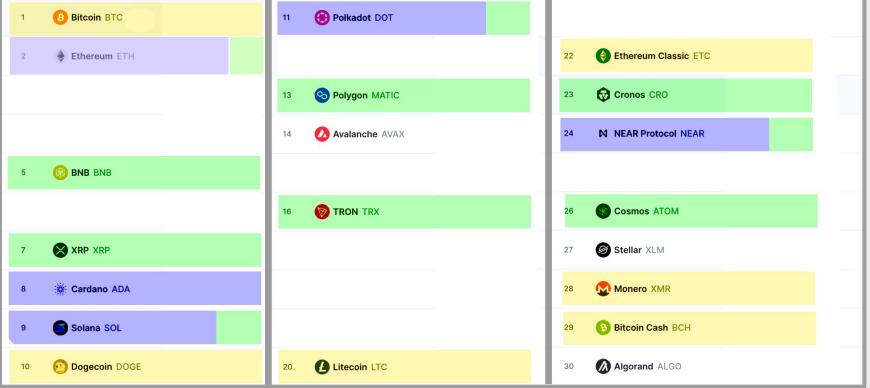


Blockchains by market cap: 20 non-ERC-20



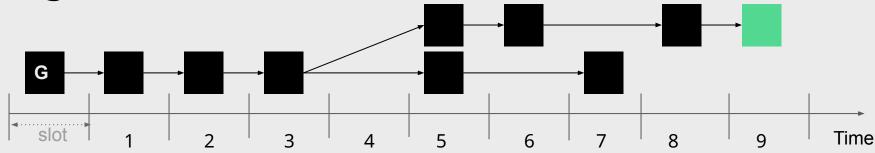
Blockchains by market cap

Longest-chain PoW ,
BFT PoS ,
Longest-chain PoS .



2: <u>Longest-chain</u> vs. <u>BFT consensus</u>

Longest-chain consensus

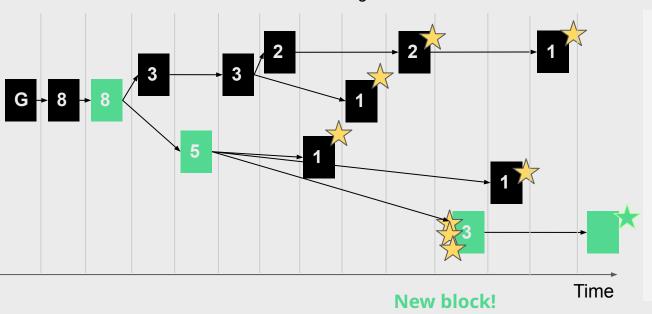


- Each validator stores the tree of blocks.
- 2. Proposer is elected per slot from the validator set (can have multiple proposers).
- 3. Proposer picks a chain to extend according to "fork-choice rule".
- 4. Creates a block on top, publishes it to the network.

Fork-choice rules: usually greedy, e.g. longest (more blocks in it) Parameters are tuned so that with high probability only one chain **eventually** wins.

Ethereum's fork-choice rule "LMD GHOST":

grows the blockchain on sub-branches with the "most activity".



LMD GHOST:

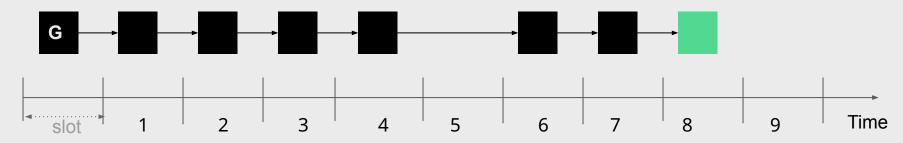
- One leader is selected per slot to propose a block.
- Each validator votes once in an epoch (32 slots) at its designated slot.
- Latest vote per validator is counted.
- The number on each block is the number of validators' votes on the subtree.
- The greedy algorithms walk the tree choosing the "heaviest" blocks.

Most popular PoS approaches

	Longest-chain consensus
Examples:	Ethereum,Peercoin,Ouroboros
Finality:	Eventual = Probabilistic
Safety under partition:	NO
Liveness under partition:	YES
#validators:	unlimited
Typical network topologies:	Flexible:

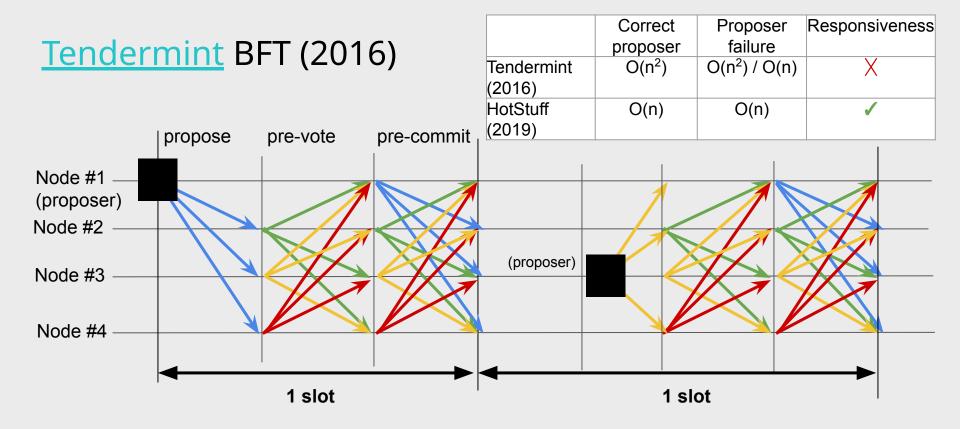
2: Longest-chain vs. <u>BFT consensus</u>

BFT-consensus

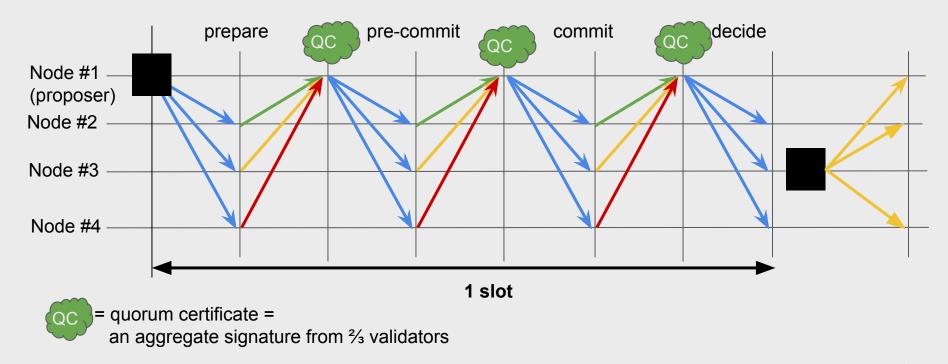


Each slot:

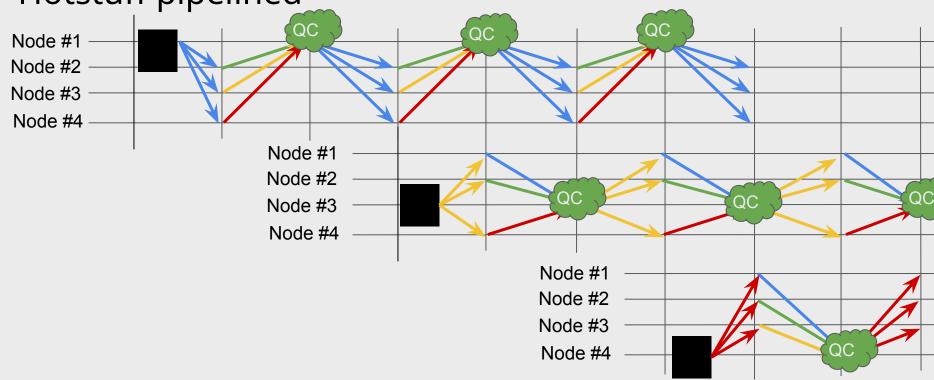
- 1. A proposer is elected from the validator set.
- 2. The proposer proposes a new block.
- 3. Validators work on finalizing the block, it either gets finalized or not.



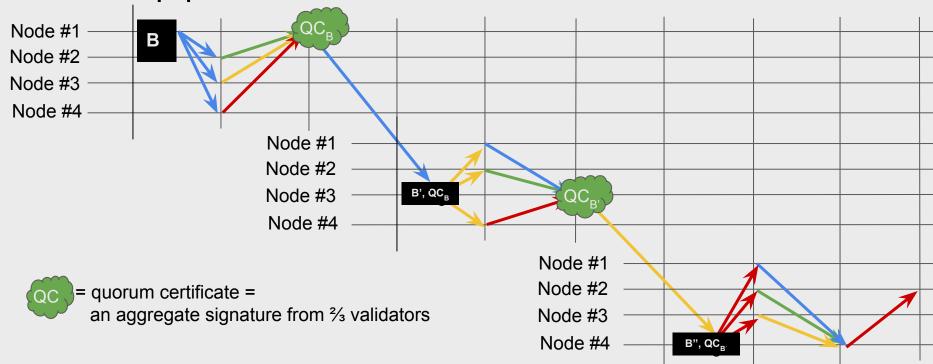
Hotstuff: rerouting through the proposer and use aggregatable signatures



Hotstuff pipelined



Hotstuff pipelined



HotStuff: 3-chain commit rule



Hotstuff: B₁ gest committed after there are three blocks with QCs on top of it.

https://malkhi.com/posts/2019/08/hotstuff-three-chain-rules/

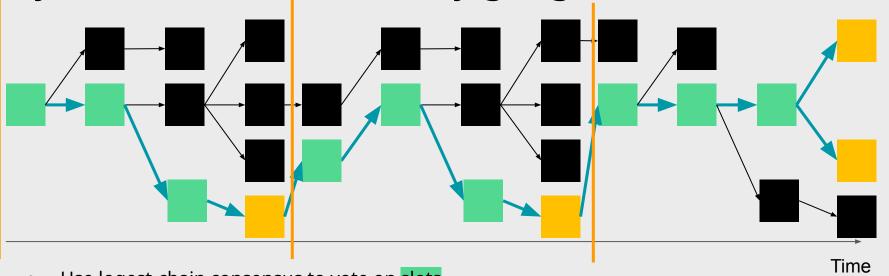
Most popular PoS approaches

	Longest-chain consensus	BFT consensus	
Examples:	Peercoin, Ouroboros-*	Tendermint, Hotstuff	
Finality:	Eventual = Probabilistic	Deterministic Immediate	
Safety under partition:	NO	YES	
Liveness under partition:	YES	NO	
#validators:	unlimited	4-100 slower with more validators	
Typical network topologies:	Flexible:	Complete:	

3: Hybrid consensus: Longest-chain + BFT

a16z crypto

Hybrid-consensus / finality gadgets



- Use logest-chain consensus to vote on slots
- Use BFT consensus to vote on epochs
- Epoch is a fixed number of slots
 - Ethereum: 32 slots per epoch, 12 second per slot, 6.4 min per epoch
- Typically: 2-stage BFT voting, hence 2 epochs to finality

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Most popular PoS approaches

	Longest-chain consensus	BFT consensus	Hybrid
Examples:	Peercoin, Ouroboros-*	Tendermint, Hotstuff	
Finality:	Eventual = Probabilistic	Deterministic Immediate	Deterministic Slow
Safety under partition:	NO	YES	YES
Liveness under partition:	YES	NO	TEMPORAL LOSS
#validators:	unlimited	4-100 slower with more validators	unlimited
Typical network topologies:	Flexible:	Complete:	Flexible:

Algorand: multiple leaders per slot

- No locked tokens, no slashing. A user registers a participation key for a certain number of rounds.
- Each node runs a VRF privately twice:
 - o to determine whether it is elected as a leader, and
 - to determine whether it is elected in a committee (new committee is selected for every step).
- There could be multiple leaders per round
- Similar to longest-chain: the lowest block (by hash) is propagated.
- Committee then votes twice on the block (in BFT-style).
- Forward-secure: nodes discard old keys disabling attacks on the past.
- Network-partition-resilient: halts during the partition, resumes when the network is restored.

2: Longest-chain vs. BFT consensus more...

Avalanche: leaderless consensus

- Snowball/Snowman <u>leaderless</u> consensus protocol!
- UTXO model, thus txs form a DAG.
- Establishes a <u>partial order</u>.
- Each node accepts a valid transaction tx_1 and sends it to other nodes, until it sees a conflicting transaction, tx_2 . If it hears about tx_2 from a > threshold of peers, it switches tx_1 to tx_2 and sends that to its peers.
- Finality: probabilistic

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