

Operating Systems Design

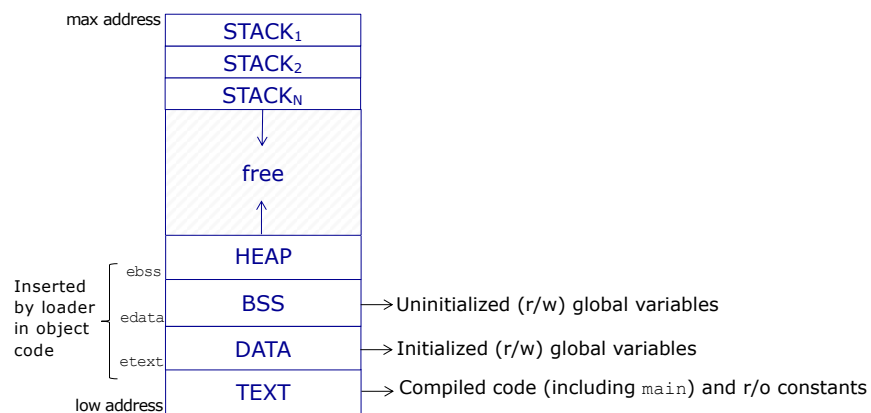
Xinu Memory Management

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1

Memory management in Xinu

- A single address space shared by all processes
 - Each process has its own stack
 - Xinu processes are “threads” (**lightweight processes**)

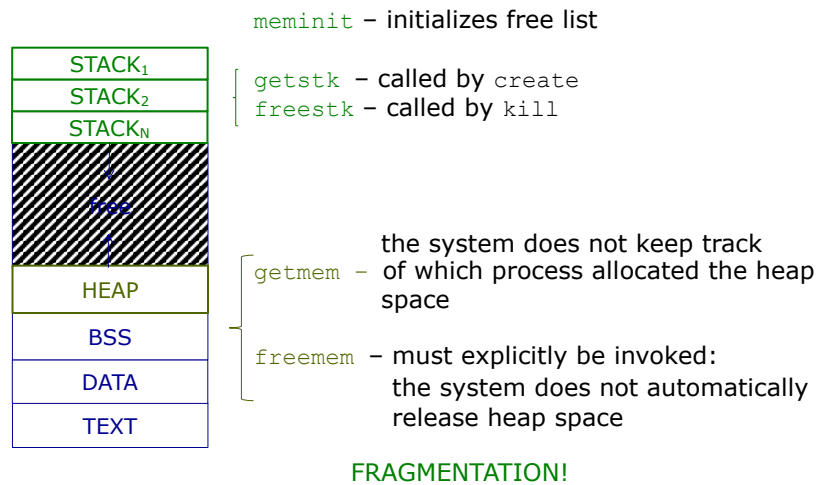


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2

2

Dynamic memory allocation in Xinu



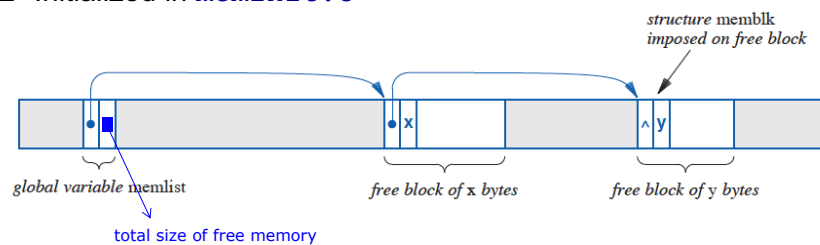
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3

3

Free list in Xinu

- Linked list of **free blocks** ordered by increasing address
- Stored in free space
- `memlist` = pointer to first free block
- Each block (`memblk`) contains:
 - Pointer to next block
 - Size of the block (except for `memlist`)
- Initialized in `meminit.c`



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4

4

Xinu data structures for memory management

■ In `memory.h`

```
/* Block of free list */
struct memblk {
    struct memblk *mnext;      /* Ptr to next free memory blk */
    uint32 mlength;           /* Size of blk (includes memblk header) */
};

extern struct memblk memlist; /* Head of free memory list */

extern void *minheap;         /* Start of heap */
extern void *maxheap;         /* Highest valid heap address */

/* Added by linker */
extern int text;              /* Start of text segment */
extern int etext;             /* End of text segment */
extern int data;              /* Start of data segment */
extern int edata;             /* End of data segment */
extern int bss;               /* Start of bss segment */
extern int ebss;              /* End of bss segment */
```

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5

5

Xinu – memory requests rounding

- `memblk` must contain at least 8 bytes
- memory allocation requests are rounded to multiple of `memblk` size (8 bytes)
- see in `memory.h`
 - `roundmb(x)`
 - `truncmb(x)` – only used at startup on initial free block size – see `meminit.c`

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6

6

Xinu - heap space allocation & release

- `getmem`
 - uses **first-fit allocation policy**
 - **splits** the block if necessary
- `freemem`
 - uses address to locate block in free list
 - tries to **coalesce** (to limit fragmentation)
 - with previous free block, next free block, or both

Xinu – stack space allocation & release

- `getstk`
 - allocates stack from highest block in free list that fits the request
 - visits whole free list to find suitable block
 - **splits** block if necessary
 - returns the *highest* address in the block
- `freestk`
 - **uses** `freemem`
 - converts the address to be passed to the `freemem` from highest address in the block (returned by `getstk` and passed as its argument) to lowest address in the block