01>

Booting Process

1. System Steefup

Blos (Boot monitor

2. Staget both Loader)

MBR (master boot record)

3. Stuge 2 boot Loader

LILO, GRUB elc

4. Kernel

Linux

5. (Init)

User-space

1. System Sterrfup -> when a System is first booted freset the processor encuefes tode at a well known location. In PC-BIOS Storla on a flash memory on mother board.

1-) Bios is a Small piece of Coch (S12 bytes) - 1 st proglos -> Bios identifies HW (boot device) which is bootable and active -> From this HW, Linux is booted where MBR Contains prime boot Loader MBR Loaded into RAM, Bios yields Contral to it.

9. Stage 1 Boot Loader -> job is to find and Load Secondar, boot Loader

->MBR Loads this Stage 1 Bootherds

- Stage 2 bood Loader) Aptly Called askernel Loader it's task is to Load the line & Keand.
- .> 15t & 2 hd Stage Loot Loader Con sine of are Called
 - Loader)
- still System is Consulted, default keenel image and initial image are Loaded into memory
- h. Keenel -> Af the head of the Keenel image, a gontine does Bame & W schep and decompress the Keenel image and place it in high memory.

 Skeenel is then called and Keenel boot begins

 -> During this sinited is Copied into RAM and monufed This acts as temporary root file System in AAM
- 5 Enit S keenel storts the first Closer Space application -> Compiled into E Library

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is The to money or the principle to with the

-) first pros labinlinit

62 Function of 05

01) Scurify

- -s Protects User data.
- -> Et prevents unauthorized aless to preserams and US or data.

02) 306 accounting

Used by Various tabiks.

03) Earor dectecting aids

S Constantly monitors the System to detect elevors and avoid malfemetioning of System.

Oh) Coordination between software and Users.
-> s es Co-ordinates and assigns interpreters
Compilers, assembles.

05) Memory management

- -> The OS manages the primary memory.
- if keeps trock of memory, which bytes are used by which page,
- -> Allocate memory for the process, deallocate when the process has terminated.

66) Processor management

sos decides the order in which the Rownes have access to the processor

> Keeps track of the status of processes

7) File management

-> A file system is organised into directories for efficient navigation and Usage

(3) Monolithic and Miceo Kernel difference with diagram

(Micro Kernel)

-Suser services and kernel. are kept in seperate address space son

-sos is Complex to design

-) Mi uo Kernel is smell in size

- Deasier of add new functionalities

does not effect the Working Of micro kernel.

-) En ecution speed is Low

Monolithic Kelnel

Both User services and Kernel Services are Keptin Same address space.

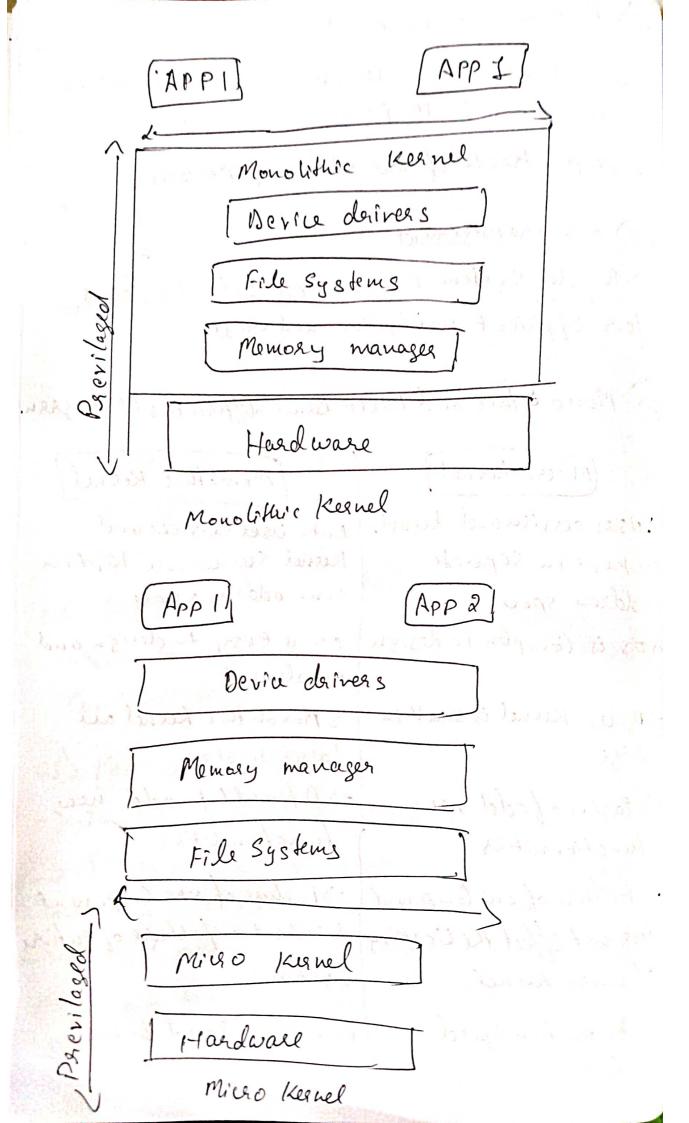
os is easy to design and i'm plement

s monohithic keenel are Larger Insize.

-> Difficult to add new functional fies

- > Failure of one Component -> Failure of one Component heads to failure of entire System.

Enecution speed is high.



-> UEFI Stavols for unified entensible firmware interface. Most new mothe boards Consist of this type It has more advantage than cesing 1208.

most impostantly, it provides uses friendly graphical uses interface (GUS)

BIOS Provides blue screen BIOS Cannot regognise harge Storage derivers, VET-I Provider a good alternative.

-> En algular BIOS that uses the keybaard to select the option OEFI allows Controls via mouse.

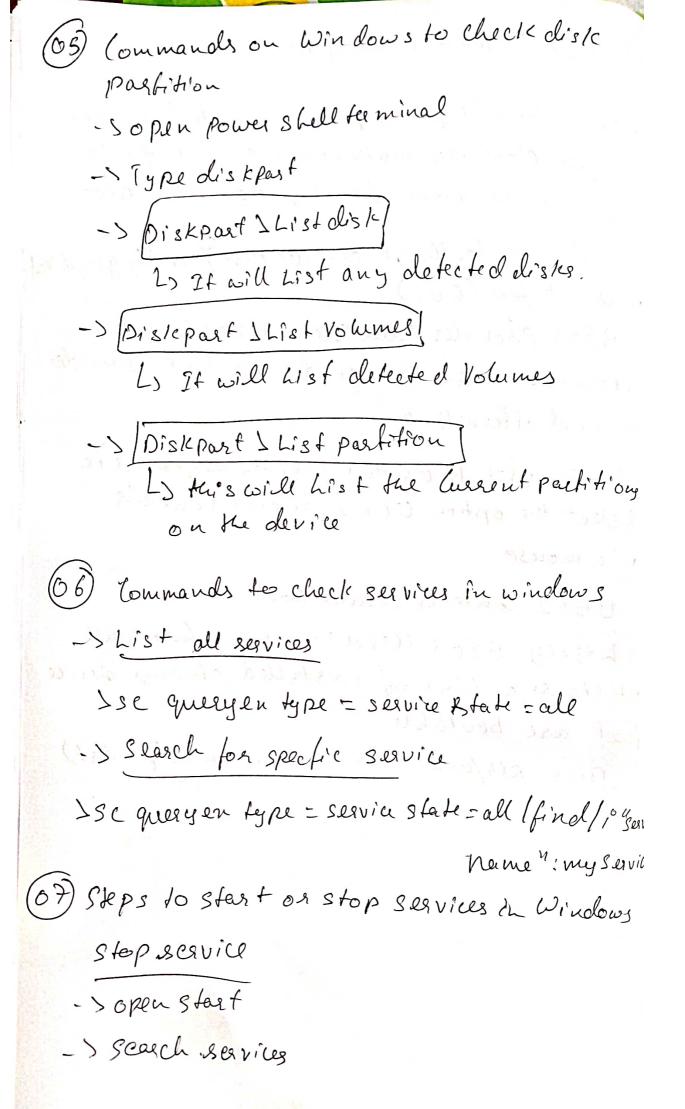
VEFI Contains Securce boot.

-) Lesaly BIOS Used hay BIOS firmward. It stores a hist of installed storage device that are bootable

Bios perform Posi (Powenon Self-test)

Se grosgen Apr = service shart = old of notifice

(0) is to start or stop seavices in children.



-> Click on the service you would to start -> click Start button -> Apply button

Stop 3 la vice

-> open Start

-> search services

-> eliek on the service you want to stop

-schick stop

s elick apply