

Intermediate Code Generation - Part 2

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NPTEL Course on Principles of Compiler Design

Outline of the Lecture

- Introduction (covered in part 1)
- Different types of intermediate code (covered in part 1)
- Intermediate code generation for various constructs

Code Template for *Function* Declaration and Call

Assumption: No nesting of functions

```
result foo(parameter list){ variable declarations; Statement list; }
```

```
func begin foo
```

```
/* creates activation record for foo - */
```

```
/* - space for local variables and temporaries */
```

```
code for Statement list
```

```
func end /* releases activation record and return */
```

```
x = bar(p1,p2,p3);
```

```
code for evaluation of p1, p2, p3 (result in T1, T2, T3)
```

```
/* result is supposed to be returned in T4 */
```

```
param T1; param T2; param T3; refparam T4;
```

```
call bar, 4
```

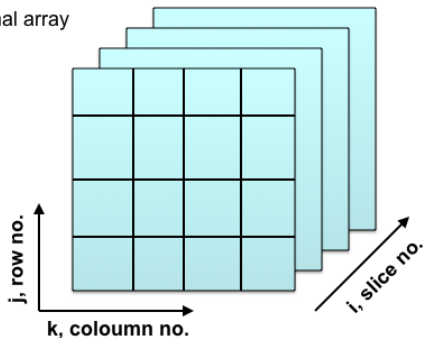
```
/* creates appropriate access links, pushes return address */
```

```
/* and jumps to code for bar */
```

```
x = T4
```

1-D Representation of 3-D Array

3-dimensional array



1-D representation of 3-D array

$$\text{offset} = (((i * n_2) + j) * n_3) + k) * \text{ele_size}$$



Code Template for *Expressions and Assignments*

```
int a[10][20][35], b;
```

```
b = exp1;
```

```
code for evaluation of exp1 (result in T1)
```

```
b = T1
```

```
/* Assuming the array access to be, a[i][j][k] */
```

```
/* base address = addr(a), offset = (((i*n2)+j)*n3)+k)*ele_size */
```

```
a[exp2][exp3][exp4] = exp5;
```

```
10: code for exp2 (result in T2) | | 141: T8 = T7+T6
```

```
70: code for exp3 (result in T3) | | 142: T9 = T8*intsize
```

```
105: T4 = T2*20 | | 143: T10 = addr(a)
```

```
106: T5 = T4+T3 | | 144: code for exp5 (result in T11)
```

```
107: code for exp4 (result in T6) | | 186: T10[T9] = T11
```

```
140: T7 = T5*35
```

Short Circuit Evaluation for Boolean Expressions

- $(exp1 \ \&\& \ exp2)$: value = if $(\sim exp1)$ then FALSE else $exp2$
 - This implies that $exp2$ need not be evaluated if $exp1$ is FALSE
- $(exp1 \ || \ exp2)$: value = if $(exp1)$ then TRUE else $exp2$
 - This implies that $exp2$ need not be evaluated if $exp1$ is TRUE
- Since boolean expressions are used mostly in conditional and loop statements, it is possible to realize perform short circuit evaluation of expressions using control flow constructs
- In such a case, there are no explicit '||' and '&&' operators in the intermediate code (as earlier), but only jumps
- Much faster, since complete expression is not evaluated
- If unevaluated expressions have side effects, then program may have non-deterministic behaviour

Control-Flow Realization of Boolean Expressions

if ((a+b < c+d) || ((e==f) && (g > h-k))) A1; else A2; A3;

```
100:      T1 = a+b
101:      T2 = c+d
103:      if T1 < T2 goto L1
104:      goto L2
105:L2:    if e==f goto L3
106:      goto L4
107:L3:    T3 = h-k
108:      if g > T3 goto L5
109:      goto L6
110:L1:L5: code for A1
111:      goto L7
112:L4:L6: code for A2
113:L7:    code for A3
```