

Intermediate Code Generation - Part 3

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NPTEL Course on Principles of Compiler Design

Outline of the Lecture

- Introduction (covered in part 1)
- Different types of intermediate code (covered in part 1)
- Intermediate code generation for various constructs

Short Circuit Evaluation for Boolean Expressions

- $(exp1 \ \&\& \ exp2)$: value = if $(\sim exp1)$ then FALSE else $exp2$
 - This implies that $exp2$ need not be evaluated if $exp1$ is FALSE
- $(exp1 \ || \ exp2)$: value = if $(exp1)$ then TRUE else $exp2$
 - This implies that $exp2$ need not be evaluated if $exp1$ is TRUE
- Since boolean expressions are used mostly in conditional and loop statements, it is possible to realize perform short circuit evaluation of expressions using control flow constructs
- In such a case, there are no explicit '||' and '&&' operators in the intermediate code (as earlier), but only jumps
- Much faster, since complete expression is not evaluated
- If unevaluated expressions have side effects, then program may have non-deterministic behaviour

Control-Flow Realization of Boolean Expressions

if ((a+b < c+d) || ((e==f) && (g > h-k))) A1; else A2; A3;

```
100:      T1 = a+b
101:      T2 = c+d
103:      if T1 < T2 goto L1
104:      goto L2
105:L2:    if e==f goto L3
106:      goto L4
107:L3:    T3 = h-k
108:      if g > T3 goto L5
109:      goto L6
110:L1:L5: code for A1
111:      goto L7
112:L4:L6: code for A2
113:L7:    code for A3
```

Code Template for *Switch* Statement

```
switch (exp) {  
    case  $l_1$  :  $SL_1$   
    case  $l_{2_1}$  : case  $l_{2_2}$  :  $SL_2$   
    ...  
    case  $l_{n-1}$  :  $SL_{n-1}$   
    default:  $SL_n$   
}
```

This code template can be used for switch statements with 10-15 cases. Note that statement list SL_i must incorporate a 'break' statement, if necessary

```
code for exp (result in T)  
goto TEST  
 $L_1$ : code for  $SL_1$   
 $L_2$ : code for  $SL_2$   
...  
 $L_n$ : code for  $SL_n$   
goto NEXT  
TEST: if  $T == l_1$  goto  $L_1$   
      if  $T == l_{2_1}$  goto  $L_2$   
      if  $T == l_{2_2}$  goto  $L_2$   
      ...  
      if  $T == l_{n-1}$  goto  $L_{n-1}$   
      if default_yes goto  $L_n$   
NEXT:
```

Grammar for Switch Statement

The grammar for the 'switch' statement according to ANSI standard C is:

selection_statement \rightarrow **SWITCH** '(' expression ')' **statement**

However, a more intuitive form of the grammar is shown below

- $STMT \rightarrow SWITCH_HEAD \quad SWITCH_BODY$
- $SWITCH_HEAD \rightarrow switch (E) /* E must be int type */$
- $SWITCH_BODY \rightarrow \{ CASE_LIST \}$
- $CASE_LIST \rightarrow CASE_ST \mid CASE_LIST \quad CASE_ST$
- $CASE_ST \rightarrow CASE_LABELS \quad STMT_LIST ;$
- $CASE_LABELS \rightarrow \epsilon \mid CASE_LABELS \quad CASE_LABEL$
- $CASE_LABEL \rightarrow case \quad CONST_INTEGRAL_EXPR : \mid default :$
/* CONST_INTEGRAL_EXPR must be of int or char type */
- $STMT \rightarrow break /* also an option */$

The for-loop of C is very general

- *for (expression₁; expression₂; expression₃) statement*

This statement is equivalent to

expression₁;

while (expression₂) { statement expression₃ ; }

- All three expressions are optional and any one (or all) may be missing
- Code generation is non-trivial because the order of execution of *statement* and *expression₃* are reversed compared to their occurrence in the for-statement
- Difficulty is due to 1-pass bottom-up code generation
- Code generation during parse tree traversals mitigates this problem by generating code for *expression₃* before that of *statement*

Code Generation Template for *C For-Loop*

for (E_1 ; E_2 ; E_3) S

code for E_1

L1: code for E_2 (result in T)

goto L4

L2: code for E_3

goto L1

L3: code for S /* all jumps out of S goto L2 */

goto L2

L4: if T == 0 goto L5 /* if T is zero, jump to exit */

goto L3

L5: /* exit */

ALGOL For-Loop

- Let us also consider a more restricted form of the for-loop
 - $STMT \rightarrow \text{for } id = EXP_1 \text{ to } EXP_2 \text{ by } EXP_3 \text{ do } STMT_1$
where, EXP_1 , EXP_2 , and EXP_3 are all arithmetic expressions, indicating starting, ending and increment values of the iteration index
 - EXP_3 may have either positive or negative values
 - All three expressions are evaluated before the iterations begin and are stored. They are not evaluated again during the loop-run
 - All three expressions are mandatory (unlike in the C-for-loop)

Code Generation Template for ALGOL For-Loop

STMT \rightarrow for *id* = *EXP*₁ to *EXP*₂ by *EXP*₃ do *STMT*₁

Code for *EXP*₁ (result in T1)

Code for *EXP*₂ (result in T2)

Code for *EXP*₃ (result in T3)

goto L1

L0: Code for *STMT*₁

id = id + T3

goto L2

L1: id = T1

L2: if (T3 \leq 0) goto L3

if (id > T2) goto L4 /* positive increment */

goto L0

L3: if (id < T2) goto L4 /* negative increment */

goto L0

L4: