

## **Database Management System 1 CPS 542 – M3**

# Database Management System of a Blood Bank

**Term Project Final Report** 

By: Sriram Sandilya Kambhampati Teja Molaganuru

Under the Guidance of:

DR.CEMIL KIRBAS

CPS 542
DEPARTMENT OF COMPUTER SCIENCE

## Database Management System of a Blood Bank

#### **Introduction:**

In this innovative concept, a sophisticated database management system plays a pivotal role in enhancing the efficiency and accuracy of managing blood supply levels, donor registrations, recipient records, and the seamless distribution of life-saving blood products. This comprehensive database comprises eight distinct entities and seven intricate relationships, forming a robust foundation for streamlining the entire process. The overarching objective of this system is to provide invaluable support in meticulously tracking and organizing every facet of blood donation, storage, and distribution. Anticipated outcomes include a substantial improvement in donor registration procedures, a finely tuned inventory management system, and a real-time tracking mechanism for blood stock levels. Ultimately, this visionary system is poised to expedite the critical process of delivering safe and timely blood to patients in need, thereby significantly enhancing healthcare outcomes.

The Proposed database consists of the following entities and relationships:

#### **Entities:**

**Donor:** A module used to register new donors and record their personal details and medical background. Information on donor eligibility is also included.

Donor (Name, Age, Gender, Phone Number, date of last donation, Blood type, **Don id**)

**Patient:** A module for registering patients, including their names, contact information, necessary blood type, and quantity of blood.

Patient (Name, Age, Gender, Blood type, Phone Number, Pat id)

**Blood Bank:** A module for registering all blood banks, together with their current blood supply, blood needs, blood types, and blood volume.

BBank (Bank name, Phone Number, **BBank id**)

*Inventory:* A module for keeping track of the various blood kinds and quantities that are available for donation.

Inventory (Blood Type, expiry date, Quantity, Inven Id)

*Facility:* A module that is superset for two additional entities to conduct transactions with blood banks and donors is referred to as a facility.

Facility (Name, Phone number, Operating hours, Fac id)

Hospital: A subset of facility entity that conducts hospital business.

Hospital (request date, blood type, request Id, quantity)

**Blood Camp:** A subdivision of a facilities entity that does blood camp business.

BCamp (SDate, EDate)

**Employee:** A module for managing tasks in the blood bank, facility, and inventory.

Employee (Name, gender, age, phone number, **Emp id**)

#### **Relations:**

- 1. Donor donates to BC to Blood camp.
- 2. Donor donates to BB to Blood bank.
- 3. Patients admit to Hospital.
- 4. Facility requests Blood bank.
- 5. Blood bank maintains Inventory.
- 6. Blood bank manages Employee.
- 7. Employee works at Facility.
- 8. Employee handles Inventory.

#### **Transactions:**

- 1. Get a list of all blood banks along with the details of employees working there and the facilities they are linked to, including blood banks without any employees.
- 2. Display all blood type requests from hospitals, including those that have not yet been fulfilled.
- 3. List all donors and their donation dates at blood camps, including those who have not yet donated.
- 4. Update contact information for a donor in the Donor table.
- 5. Find the total quantity of each blood type available in the inventory.
- 6. Get the names of donors who are older than the average age of all donors.
- 7. List each blood bank along with the count of donations it has received.
- 8. Increase the age of all employees working at a specific facility by 1 year.
- 9. Find the names of all donors who have donated at facilities where a blood camp was held.
- 10. Retrieve each donor's most recent donation date.

#### **ER Diagram:**

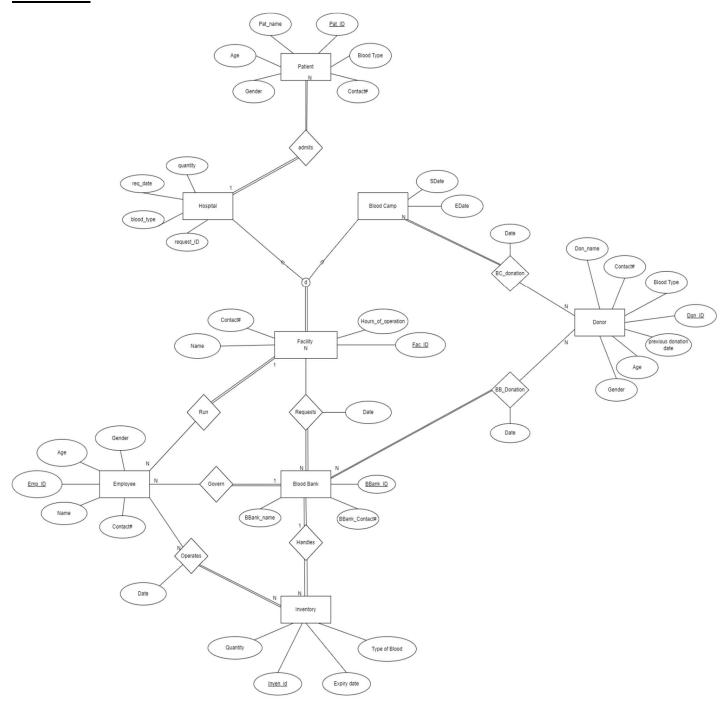
#### **Description:**

An Entity-Relationship (ER) diagram in Database Management Systems (DBMS) is a visual representation used to model and describe the structure of a database. It consists of entities, which are objects or concepts in the real world, and the relationships between them. Entities are represented as rectangles, and relationships are depicted as lines connecting these rectangles. Each entity is defined by attributes that describe its properties. ER diagrams help in understanding the database's schema, illustrating how different entities are related, and serving as a blueprint for designing the database schema. They are a vital tool for database designers and developers to ensure data integrity, efficient data retrieval, and a clear understanding of the database's structure.

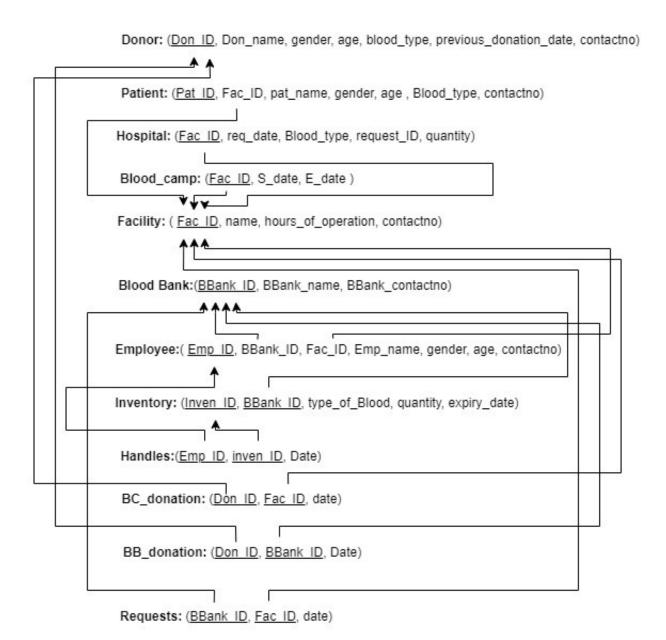
In our database management system for a blood bank, we have the following entities:

- Donor (Name, Age, Gender, Contact Number, date of last donation, Blood type, Don id)
- ➤ Patient (Name, Age, Gender, Blood type, Contact Number, Pat id)
- ➤ Blood Bank (Bank name, Contact Number, **BBank id**)
- ➤ Inventory (Blood Type, expiry date, Quantity, **Inven Id**)
- Facility (Name, Contact number, Operating hours, Fac id)
- ➤ Hospital (request date, blood type, request Id, quantity)
- ➤ Blood Camp (SDate, EDate)
- Employee (Name, gender, age, Contact number, Emp id)

### ER Model:



## Relation Schema of the Database:



## Normalization of the Relational Schema:

Normalization is done for minimizing the data redundancy in the database model of the Database Management System of a Blood Bank.

For normalization you need to check if the system is in 1NF, 2NF, 3NF and BCNF.

- ➤ 1NF: A relation is in 1NF if it contains an atomic value. For each multi-valued attribute, create a new table, in which you place the key to the original table and the mutli-valued attribute. Keep the original table, with its key.
- ➤ **2NF:** A relation will be 2NF if it is in 1NF, and all non-key attributes are fully functional and dependent on the primary key.
- ➤ **3NF:** As it should be in 2NF and every non-prime attribute of the relation is non-transitively dependent on every key in the relation also whenever a non-trivial functional dependency X—A exists, then either X is asuper key or A is a member of some candidate key.
- $\triangleright$  BCNF (Boyce codd's normal form): Stronger than 3NF, it should be in 3NF and whenever a non-trivial functional dependency  $X \rightarrow A$  exists, X should be a super key.

#### Donor Table:

**Donor:**(<u>Don\_ID</u>,Don\_name,gender,age,blood\_type,Previous\_donation\_date, contactno)

The above relation's attributes equal the letters {ABCDEFGH} in the following order R

(A|B|C|D|E|F|G)

The List of Functional Dependency: A -->BCDEFG

#### First Normal Form (1NF):

- The above relational model violates 1NF as attribute contactno is a multi-valued attribute. In order to make it atomicwe split the contactno to contactno 1 and contactno 2. Donor:(Don\_ID,Don\_name,gender,age,Blood\_Type,contactno 1,contactno 2, Previous\_donation\_date, contactno)
- The above relational model is in **1NF**. This is because every relation listed above has no multivalued attributes(i.e., atomic) and each relation has a primary key as well as a correct domain type.

#### **Second Normal Form (2NF):**

• The above relational model is in 2NF because it has a single candidate key, and all non-key attributes depend on the primary key.

#### Third Normal Form (3NF):

• Now the Above model is in 3NF as it is in 2NF and every non-prime attribute of the relation is non-transitivelydependent on every key in the relation also whenever a non-trivial functional dependency X→A exists, then either X is a super key or A is a member of some candidate key.

#### **BCNF** (Boyce Codd's normal form):

#### Patient Table:

**Patient:**(Pat\_Id, Fac\_ID, Pat\_name, gender, age, Blood\_Type, contactno)
The above relation's attributes equal the letters {ABCDEFG} in the following order R (A|B|C|D|E|F|G)
The List of Functional Dependency: A -->BCDEFG

#### First Normal Form (1NF):

- The above relational model violates 1NF as attribute contactno is a multi-valued attribute. Inorder to make it atomicwe split the contactno to contactno1 and contactno2.
- Patient:(<u>Pat\_Id</u>, Fac\_ID,Pat\_Name, gender, age, Blood Type, contactno1,contactno2)
- The above relational model is in 1NF. This is because every relation listed above has no multivalued attributes(i.e., atomic) and each relation has a primary key as well as a correct domain type.

#### **Second Normal Form (2NF):**

• The above relational model is in 2NF because it has a single candidate key, and all non-key attributes depend on the primary key.

#### Third Normal Form (3NF):

• Now the Above model is in 3NF as it is in 2NF and every non-prime attribute of the relation is non-transitively dependent on every key in the relation also whenever a non-trivial functional dependency X→A exists, then either X is a super key or A is a member of some candidate key.

#### **BCNF** (Boyce Codd's normal form):

• The above relational model is in BCNF because there are no overlapping candidate keys or functional dependencies.

#### Hospital Table:

**Hospital:**(Fac\_ID, req\_date, Blood\_type, request\_ID, quantity) The above relation's attributes equal the letters  $\{ABCDE\}$  in the following order R (A|B|C|D|E). The List of Functional Dependency: A -->BCDE

#### First Normal Form (1NF):

• The above relational model is in 1NF. This is because every relation listed above has no multivalued attributes(i.e., atomic) and each relation has a primary key as well as a correct domain type.

#### **Second Normal Form (2NF):**

• The above relational model is in 2NF because it has a single candidate key, and all non-key attributes depend on the primary key.

#### Third Normal Form (3NF):

• Now the Above model is in 3NF as it is in 2NF and every non-prime attribute of the relation is non-transitivelydependent on every key in the relation also whenever a non-trivial functional dependency X→A exists, then either X is a super key or A is a member of some candidate key.

#### **BCNF** (Boyce Codd's normal form):

#### **Blood Camp Table:**

#### **Blood\_Camp:**(Fac\_ID, S\_date, E\_date)

The above relation's attributes equal the letters  $\{ABCD\}$  in the following order R (A|B|C). The List of Functional Dependency: A -->BC

#### First Normal Form (1NF):

• The above relational model is in 1NF. This is because every relation listed above has no multivalued attributes(i.e., atomic) and each relation has a primary key as well as a correct domain type.

#### **Second Normal Form (2NF):**

• The above relational model is in 2NF because it has a single candidate key, and all non-key attributes depend on the primary key.

#### Third Normal Form (3NF):

• Now the Above model is in 3NF as it is in 2NF and every non-prime attribute of the relation is non-transitively dependent on every key in the relation also whenever a non-trivial functional dependency X→A exists, then either X is a super key or A is a member of some candidate key.

#### **BCNF** (Boyce Codd's normal form):

• The above relational model is in BCNF because there are no overlapping candidate keys or functional dependencies.

#### Facility Table:

Facility: (Fac ID , name, hours of operation, contactno)

The above relation's attributes equal the letters  $\{ABCD\}$  in the following order R (A|B|C|D).

The List of Functional Dependency: A -->BCD

#### First Normal Form (1NF):

- The above relational model violates 1NF as attribute contactno is a multi-valued attribute. Inorder to make it atomicwe split the contactno to contactno 1 and contactno 2
- Facility:(Fac\_ID ,name, hours\_of\_operation, contactno1, contactno2)
- The above relational model is in 1NF. This is because every relation listed above has no multivalued attributes(i.e., atomic) and each relation has a primary key as well as a correct domain type.

#### **Second Normal Form (2NF):**

• The above relational model is in 2NF because it has a single candidate key, and all non-key attributes depend on the primary key.

#### Third Normal Form (3NF):

• Now the Above model is in 3NF as it is in 2NF and every non-prime attribute of the relation is non-transitivelydependent on every key in the relation also whenever a non-trivial functional dependency X→A exists, then either X is a super key or A is a member of some candidate key.

#### **BCNF** (Bovce Codd's normal form):

#### Blood Bank Table:

**Blood Bank:** (BBank\_ID, BBank\_name, BBank\_contactno, Fac\_ID) The above relation's attributes equal the letters  $\{ABCD\}$  in the following order R (A|B|C|D). The List of Functional Dependency: A -->BCD

#### First Normal Form (1NF):

- The above relational model violates 1NF as attribute BB\_contactno is a multi-valued attribute. Inorder to make it atomic we split the BB\_contactno to BB\_contactno 1 and BB\_contactno 2.
- Blood Bank: (BBank ID, BBank name, BBank contactno1, BBank contactno2, Fac ID)
- The above relational model is in 1NF. This is because every relation listed above has no multivalued attributes(i.e., atomic) and each relation has a primary key as well as a correct domain type.

#### **Second Normal Form (2NF):**

• The above relational model is in 2NF because it has a single candidate key, and all non-key attributes dependen the primary key.

#### Third Normal Form (3NF):

• Now the Above model is in 3NF as it is in 2NF and every non-prime attribute of the relation is non-transitively dependent on every key in the relation also whenever a non-trivial functional dependency X→A exists, then either X is a super key or A is a member of some candidate key.

#### **BCNF** (Boyce Codd's normal form):

• The above relational model is in BCNF because there are no overlapping candidate keys or functional dependencies.

#### Employee Table:

**Employee:**(Emp\_ID, BBank\_ID, Fac\_ID, Emp\_name, gender, age, contactno)
The above relation's attributes equal the letters {ABCDEFG} in the following order R (A|B|C|D|E|F|G)
The List of Functional Dependency: A -->BCDEFG

#### First Normal Form (1NF):

- The above relational model violates 1NF as attribute contactno is a multi-valued attribute. Inorder to make it atomicwe split the contactno to contactno 1 and contactno 2.
- Employee:(Emp ID, BBank ID, Fac ID, Emp name, gender, age, contactno1, contactno2)
- The above relational model is in 1NF. This is because every relation listed above has no multivalued attributes(i.e., atomic) and each relation has a primary key as well as a correct domain type.

#### **Second Normal Form (2NF):**

• The above relational model is in 2NF because it has a single candidate key, and all non-key attributes depend on the primary key.

#### Third Normal Form (3NF):

• Now the Above model is in 3NF as it is in 2NF and every non-prime attribute of the relation is non-transitively dependent on every key in the relation also whenever a non-trivial functional dependency X→A exists, then either X is a super key or A is a member of some candidate key.

#### **BCNF** (Boyce Codd's normal form):

#### *Inventory Table:*

**Inventory:**(Inven\_ID, BBank\_ID, type\_of\_Blood, quantity, expiry\_date)
The above relation's attributes equal the letters {ABCD} in the following order R (A|B|C|D|E)
The List of Functional Dependency: A -->BCDE

#### First Normal Form (1NF):

• The above relational model is in 1NF. This is because every relation listed above has no multivalued attributes(i.e., atomic) and each relation has a primary key as well as a correct domain type.

#### **Second Normal Form (2NF):**

• The above relational model is in 2NF because it has a single candidate key, and all non-key attributes depend on the primary key.

#### Third Normal Form (3NF):

• Now the Above model is in 3NF as it is in 2NF and every non-prime attribute of the relation is non-transitively dependent on every key in the relation also whenever a non-trivial functional dependency X→A exists, then either X is a super key or A is a member of some candidate key.

#### **BCNF** (Boyce Codd's normal form):

• The above relational model is in BCNF because there are no overlapping candidate keys or functional dependencies.

#### Handles Table:

#### Handles: (Emp ID, invent ID, Date)

The above relation's attributes equal the letters {ABC} in the following order R (AB|C) The List of FunctionalDependency: A B-->C

#### First Normal Form (1NF):

• The above relational model is in 1NF. This is because every relation listed above has no multivalued attributes(i.e., atomic) and each relation has a primary key as well as a correct domain type.

#### **Second Normal Form (2NF):**

• The above relational model is in 2NF because it has a single candidate key, and all non-key attributes depend on the primary key.

#### Third Normal Form (3NF):

• Now the Above model is in 3NF as it is in 2NF and every non-prime attribute of the relation is non-transitivelydependent on every key in the relation also whenever a non-trivial functional dependency X→A exists, then either X is a super key or A is a member of some candidate key.

#### **BCNF** (Boyce Codd's normal form):

#### **BC** donation Table:

#### **BC** donation:(Don ID, Fac ID, Date)

The above relation's attributes equal the letters {ABC} in the following order R (AB|C) The List of Functional Dependency: A B-->C

#### First Normal Form (1NF):

• The above relational model is in 1NF. This is because every relation listed above has no multivalued attributes(i.e., atomic) and each relation has a primary key as well as a correct domain type.

#### **Second Normal Form (2NF):**

• The above relational model is in 2NF because it has a single candidate key, and all non-key attributes depend on the primary key.

#### Third Normal Form (3NF):

• Now the Above model is in 3NF as it is in 2NF and every non-prime attribute of the relation is non-transitively dependent on every key in the relation also whenever a non-trivial functional dependency X→A exists, then either X is a super key or A is a member of some candidate key.

#### **BCNF** (Boyce Codd's normal form):

• The above relational model is in BCNF because there are no overlapping candidate keys or functional dependencies.

#### <u>BB\_donation Table:</u>

#### **BB\_donation:** (Don\_ID, BBank\_ID, Date)

The above relation's attributes equal the letters {ABC} in the following order R (AB|C) The List of Functional Dependency: A B-->C

#### First Normal Form (1NF):

• The above relational model is in 1NF. This is because every relation listed above has no multivalued attributes(i.e., atomic) and each relation has a primary key as well as a correct domain type.

#### **Second Normal Form (2NF):**

• The above relational model is in 2NF because it has a single candidate key, and all non-key attributes depend on the primary key.

#### Third Normal Form (3NF):

• Now the Above model is in 3NF as it is in 2NF and every non-prime attribute of the relation is non-transitively dependent on every key in the relation also whenever a non-trivial functional dependency X→A exists, then either X is a super key or A is a member of some candidate key.

#### **BCNF** (Boyce Codd's normal form):

#### Requests Table:

#### Requests: (BBank ID, Fac ID, Date)

The above relation's attributes equal the letters {ABC} in the following order R (AB|C) The List of Functional Dependency: A B-->C

#### First Normal Form (1NF):

• The above relational model is in 1NF. This is because every relation listed above has no multivalued attributes(i.e., atomic) and each relation has a primary key as well as a correct domain type.

#### **Second Normal Form (2NF):**

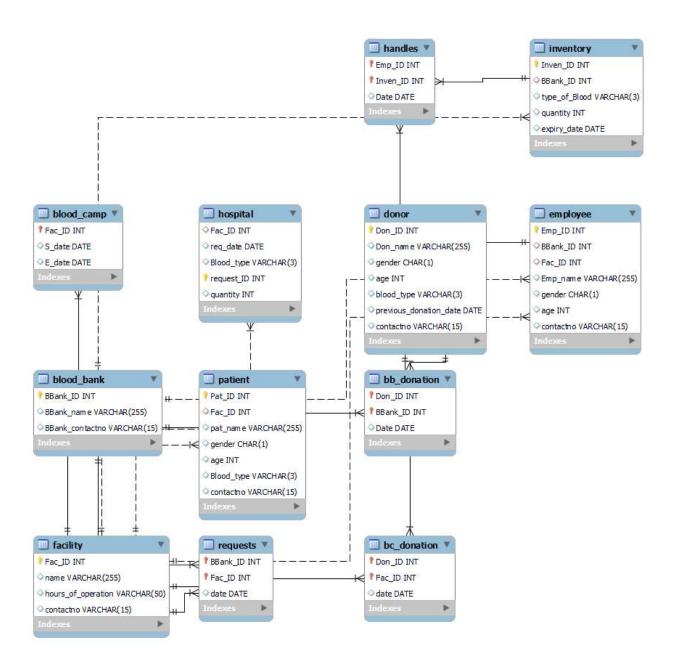
• The above relational model is in 2NF because it has a single candidate key, and all non-key attributes depend on the primary key.

#### Third Normal Form (3NF):

• Now the Above model is in 3NF as it is in 2NF and every non-prime attribute of the relation is non-transitivelydependent on every key in the relation also whenever a non-trivial functional dependency X→A exists, then either X is a super key or A is a member of some candidate key.

#### **BCNF** (Boyce Codd's normal form):

## **Logical Model:**



## **Physical Model:**

-- MySQL Workbench Forward Engineering SET @OLD UNIQUE CHECKS=@@UNIQUE CHECKS, UNIQUE CHECKS=0; SET @OLD FOREIGN KEY CHECKS=@@FOREIGN KEY CHECKS, FOREIGN\_KEY\_CHECKS=0; SET @OLD SQL MODE=@@SQL MODE, SQL MODE='ONLY FULL GROUP BY,STRICT TRANS TABLES,NO ZER O IN DATE, NO ZERO DATE, ERROR FOR DIVISION BY ZERO, NO ENG INE SUBSTITUTION'; -- Schema new blood bank .. ------- Schema new blood bank CREATE SCHEMA IF NOT EXISTS 'new blood bank' DEFAULT CHARACTER SET utf8; -- Schema blood bank

```
-- Schema blood bank
CREATE SCHEMA IF NOT EXISTS 'blood bank' DEFAULT CHARACTER
SET utf8mb4 COLLATE utf8mb4 0900 ai ci;
USE 'new blood bank';
-- Table 'new blood bank'. 'facility'
CREATE TABLE IF NOT EXISTS 'new blood bank'. 'facility' (
 'Fac ID' INT NOT NULL AUTO INCREMENT,
 'name' VARCHAR(45) NULL,
 'hours of operation' VARCHAR(45) NULL,
 'contactno1' VARCHAR(45) NULL,
 `contactno2` VARCHAR(45) NULL,
 PRIMARY KEY ('Fac ID'))
ENGINE = InnoDB;
-- Table 'new blood bank'. 'Hospital'
CREATE TABLE IF NOT EXISTS 'new blood bank'. 'Hospital' (
 'req date' VARCHAR(45) NULL,
 'Blood type' VARCHAR(45) NULL,
 'request ID' INT NULL,
```

```
'quantity' VARCHAR(45) NULL,
 'facility Fac ID' INT NOT NULL,
 PRIMARY KEY ('facility Fac ID'),
 INDEX 'fk Hospital facility1 idx' ('facility Fac ID' ASC) VISIBLE,
 CONSTRAINT 'fk Hospital facility1'
 FOREIGN KEY ('facility Fac ID')
  REFERENCES 'new blood bank'. 'facility' ('Fac ID')
  ON DELETE NO ACTION
  ON UPDATE NO ACTION)
ENGINE = InnoDB;
-- Table 'new blood bank'.'patient'
CREATE TABLE IF NOT EXISTS 'new blood bank'.'patient' (
 'pat ID' INT NOT NULL AUTO INCREMENT,
 'Pat Name' VARCHAR(45) NULL,
 'Gender' VARCHAR(45) NULL,
 'Age' INT NULL,
 'Blood Type' VARCHAR(45) NULL,
 'contactno1' VARCHAR(45) NULL,
 'contactno2' VARCHAR(45) NULL,
 'Fac ID' INT NOT NULL,
 'Hospital facility Fac ID' INT NOT NULL,
```

```
PRIMARY KEY ('pat ID'),
 INDEX 'fk patient Hospital1 idx' ('Hospital facility Fac ID' ASC) VISIBLE,
 CONSTRAINT 'fk patient Hospital1'
 FOREIGN KEY ('Hospital facility Fac ID')
  REFERENCES 'new blood bank'. 'Hospital' ('facility Fac ID')
  ON DELETE NO ACTION
  ON UPDATE NO ACTION)
ENGINE = InnoDB;
-- Table 'new blood bank'.'Donor'
CREATE TABLE IF NOT EXISTS 'new blood bank'.'Donor' (
 'Don ID' INT NOT NULL AUTO INCREMENT,
 'Don name' VARCHAR(45) NULL,
 'gender' VARCHAR(45) NULL,
 'age' INT NULL,
 'blood type' VARCHAR(45) NULL,
 'contactno1' VARCHAR(45) NULL,
 'contactno2' VARCHAR(45) NULL,
 'Previous donation date' VARCHAR(45) NULL,
 PRIMARY KEY ('Don ID'))
ENGINE = InnoDB;
```

```
-- Table 'new blood bank'. 'Facility'
CREATE TABLE IF NOT EXISTS 'new blood bank'. 'Facility' (
 'Fac ID' INT NOT NULL AUTO INCREMENT,
 'Name' VARCHAR(45) NULL,
 'Hours of operation' VARCHAR(45) NULL,
 'Contactno1' VARCHAR(45) NULL,
 'Contactno2' VARCHAR(45) NULL,
PRIMARY KEY ('Fac ID'))
ENGINE = InnoDB;
-- Table 'new blood bank'. 'Blood Bank'
CREATE TABLE IF NOT EXISTS 'new blood bank'. 'Blood Bank' (
 'BBank ID' INT NOT NULL AUTO INCREMENT,
 'BBank name' VARCHAR(45) NULL,
 'BBank contactno1' VARCHAR(45) NULL,
 'BBank contactno2' VARCHAR(45) NULL,
 PRIMARY KEY ('BBank ID'))
ENGINE = InnoDB;
```

```
-- Table 'new blood bank'. 'Employee'
CREATE TABLE IF NOT EXISTS 'new blood bank'. 'Employee' (
 'Emp ID' INT NOT NULL AUTO INCREMENT,
 'Emp name' VARCHAR(45) NULL,
 'gender' VARCHAR(45) NULL,
 'age' INT NULL,
 'contactno1' VARCHAR(45) NULL,
 `contactno2` VARCHAR(45) NULL,
 'facility Fac ID' INT NOT NULL,
 'Blood Bank BBank ID' INT NOT NULL,
 PRIMARY KEY ('Emp ID'),
INDEX 'fk Employee facility1 idx' ('facility Fac ID' ASC) VISIBLE,
INDEX 'fk Employee Blood Bank1 idx' ('Blood Bank BBank ID' ASC)
VISIBLE,
CONSTRAINT 'fk Employee facility1'
 FOREIGN KEY ('facility Fac ID')
  REFERENCES 'new blood bank'. 'facility' ('Fac ID')
  ON DELETE NO ACTION
  ON UPDATE NO ACTION,
CONSTRAINT 'fk_Employee Blood Bank1'
 FOREIGN KEY ('Blood Bank BBank ID')
  REFERENCES 'new blood bank'. 'Blood Bank' ('BBank ID')
  ON DELETE NO ACTION
```

```
ENGINE = InnoDB;
-- Table 'new blood bank'.'Inventory'
CREATE TABLE IF NOT EXISTS 'new blood bank'. 'Inventory' (
 'Inven ID' INT NOT NULL AUTO INCREMENT,
 'type of Blood' VARCHAR(45) NULL,
 'quantity' INT NULL,
 'expiry date' DATE NULL,
 'Blood Bank BBank ID' INT NOT NULL,
PRIMARY KEY ('Inven ID'),
INDEX 'fk Inventory Blood Bank1 idx' ('Blood Bank BBank ID' ASC)
VISIBLE,
CONSTRAINT 'fk Inventory Blood Bank1'
 FOREIGN KEY ('Blood Bank BBank ID')
 REFERENCES 'new blood bank'. 'Blood Bank' ('BBank ID')
  ON DELETE NO ACTION
  ON UPDATE NO ACTION)
ENGINE = InnoDB;
```

-- Table 'new blood bank'.'Donor Blood Bank'

ON UPDATE NO ACTION)

```
CREATE TABLE IF NOT EXISTS 'new blood bank'.'Donor Blood Bank' (
 'Donor ID' INT NOT NULL,
 'BBank ID' INT NOT NULL,
 'Date' DATE NULL,
PRIMARY KEY ('Donor ID', 'BBank ID'),
INDEX 'fk Donor has Blood Bank Blood Bank1 idx' ('BBank ID' ASC)
VISIBLE,
 INDEX 'fk Donor has Blood Bank Donor1 idx' ('Donor ID' ASC) VISIBLE,
 CONSTRAINT 'fk Donor has Blood Bank Donor1'
 FOREIGN KEY ('Donor ID')
 REFERENCES 'new blood bank'. 'Donor' ('Don ID')
  ON DELETE NO ACTION
  ON UPDATE NO ACTION,
CONSTRAINT 'fk Donor has Blood Bank Blood Bank1'
 FOREIGN KEY ('BBank ID')
 REFERENCES 'new blood bank'. 'Blood Bank' ('BBank ID')
  ON DELETE NO ACTION
  ON UPDATE NO ACTION)
ENGINE = InnoDB;
-- Table 'new blood bank'.'Requests'
CREATE TABLE IF NOT EXISTS 'new blood bank'. 'Requests' (
```

```
'facility ID' INT NOT NULL,
 'Blood Bank BBank ID' INT NOT NULL,
 'Date' DATE NULL,
 PRIMARY KEY ('facility ID', 'Blood Bank BBank ID'),
 INDEX 'fk facility has Blood Bank Blood Bank1 idx'
('Blood Bank BBank ID' ASC) VISIBLE,
 INDEX 'fk facility has Blood Bank facility1 idx' ('facility ID' ASC)
VISIBLE,
 CONSTRAINT 'fk facility has Blood Bank facility1'
  FOREIGN KEY ('facility ID')
  REFERENCES 'new blood bank'. 'facility' ('Fac ID')
  ON DELETE NO ACTION
  ON UPDATE NO ACTION,
 CONSTRAINT 'fk facility has Blood Bank Blood Bank1'
  FOREIGN KEY ('Blood Bank BBank ID')
  REFERENCES 'new blood bank'. 'Blood Bank' ('BBank ID')
  ON DELETE NO ACTION
  ON UPDATE NO ACTION)
ENGINE = InnoDB;
-- Table 'new blood bank'. 'Blood camp'
CREATE TABLE IF NOT EXISTS 'new blood bank'. 'Blood camp' (
 'S date' VARCHAR(45) NULL,
```

```
'E date' VARCHAR(45) NULL,
 `facility Fac ID` INT NOT NULL,
 PRIMARY KEY ('facility Fac ID'),
 INDEX 'fk Blood camp facility1 idx' ('facility Fac ID' ASC) VISIBLE,
 CONSTRAINT 'fk Blood camp facility1'
 FOREIGN KEY ('facility Fac ID')
  REFERENCES 'new blood bank'. 'facility' ('Fac ID')
  ON DELETE NO ACTION
  ON UPDATE NO ACTION)
ENGINE = InnoDB;
-- Table 'new blood bank'.'BC donation'
CREATE TABLE IF NOT EXISTS 'new blood bank'. 'BC donation' (
 'Date' DATE NULL,
 'Donor Don ID' INT NOT NULL,
 'Blood camp facility Fac ID' INT NOT NULL,
 PRIMARY KEY ('Donor Don ID', 'Blood camp facility Fac ID'),
 INDEX 'fk BC donation Donor1 idx' ('Donor Don ID' ASC) VISIBLE,
 INDEX 'fk BC donation Blood camp1 idx' ('Blood camp facility Fac ID'
ASC) VISIBLE,
 CONSTRAINT 'fk BC donation Donor1'
  FOREIGN KEY ('Donor Don ID')
  REFERENCES 'new blood bank'.'Donor' ('Don ID')
```

```
ON DELETE NO ACTION
  ON UPDATE NO ACTION,
 CONSTRAINT 'fk BC donation Blood camp1'
 FOREIGN KEY ('Blood camp facility Fac ID')
  REFERENCES 'new blood bank'. 'Blood camp' ('facility Fac ID')
  ON DELETE NO ACTION
  ON UPDATE NO ACTION)
ENGINE = InnoDB;
-- Table 'new blood bank'. 'Handles'
CREATE TABLE IF NOT EXISTS 'new_blood_bank'. 'Handles' (
 'Date' DATE NULL,
 'Employee Emp ID' INT NOT NULL,
 'Inventory Inven ID' INT NOT NULL,
PRIMARY KEY ('Employee Emp ID', 'Inventory Inven ID'),
INDEX 'fk Handles Employee1 idx' ('Employee Emp ID' ASC) VISIBLE,
INDEX 'fk Handles Inventory1 idx' ('Inventory Inven ID' ASC) VISIBLE,
CONSTRAINT 'fk Handles Employee1'
 FOREIGN KEY ('Employee Emp ID')
 REFERENCES 'new blood bank'.'Employee' ('Emp ID')
  ON DELETE NO ACTION
  ON UPDATE NO ACTION,
```

```
CONSTRAINT 'fk Handles Inventory1'
 FOREIGN KEY ('Inventory Inven ID')
 REFERENCES 'new blood bank'.'Inventory' ('Inven ID')
  ON DELETE NO ACTION
  ON UPDATE NO ACTION)
ENGINE = InnoDB;
USE 'blood bank';
-- Table 'blood bank'.'donor'
CREATE TABLE IF NOT EXISTS 'blood bank'.'donor' (
 'Don ID' INT NOT NULL,
 'Don name' VARCHAR(255) NULL DEFAULT NULL,
 'gender' CHAR(1) NULL DEFAULT NULL,
 'age' INT NULL DEFAULT NULL,
 'blood type' VARCHAR(3) NULL DEFAULT NULL,
 'previous donation date' DATE NULL DEFAULT NULL,
 'contactno' VARCHAR(15) NULL DEFAULT NULL,
PRIMARY KEY ('Don ID'))
ENGINE = InnoDB
DEFAULT CHARACTER SET = utf8mb4
COLLATE = utf8mb4 0900 ai ci;
```

```
-- Table 'blood bank'. 'blood bank'
CREATE TABLE IF NOT EXISTS 'blood bank'.'blood bank' (
 'BBank ID' INT NOT NULL,
 'BBank name' VARCHAR(255) NULL DEFAULT NULL,
 'BBank contactno' VARCHAR(15) NULL DEFAULT NULL,
 PRIMARY KEY ('BBank ID'))
ENGINE = InnoDB
DEFAULT CHARACTER SET = utf8mb4
COLLATE = utf8mb4 0900 ai ci;
-- Table 'blood bank'.'bb donation'
CREATE TABLE IF NOT EXISTS 'blood bank'.'bb donation' (
 'Don ID' INT NOT NULL,
 'BBank ID' INT NOT NULL,
 'Date' DATE NULL DEFAULT NULL,
PRIMARY KEY ('Don ID', 'BBank ID'),
INDEX 'BBank ID' ('BBank ID' ASC) VISIBLE,
CONSTRAINT 'bb donation ibfk 1'
 FOREIGN KEY ('Don ID')
```

```
REFERENCES 'blood bank'.'donor' ('Don ID'),
CONSTRAINT 'bb_donation_ibfk_2'
 FOREIGN KEY ('BBank ID')
 REFERENCES 'blood bank'.'blood bank' ('BBank ID'))
ENGINE = InnoDB
DEFAULT CHARACTER SET = utf8mb4
COLLATE = utf8mb4 0900 ai ci;
-- Table 'blood bank'. 'facility'
CREATE TABLE IF NOT EXISTS 'blood bank'. 'facility' (
 'Fac ID' INT NOT NULL,
 'name' VARCHAR(255) NULL DEFAULT NULL,
 'hours of operation' VARCHAR(50) NULL DEFAULT NULL,
 'contactno' VARCHAR(15) NULL DEFAULT NULL,
PRIMARY KEY ('Fac ID'))
ENGINE = InnoDB
DEFAULT CHARACTER SET = utf8mb4
COLLATE = utf8mb4 0900 ai ci;
-- Table 'blood bank'.'bc donation'
```

```
CREATE TABLE IF NOT EXISTS 'blood bank'.'bc donation' (
 'Don ID' INT NOT NULL,
 'Fac ID' INT NOT NULL,
 'date' DATE NULL DEFAULT NULL,
PRIMARY KEY ('Don_ID', 'Fac_ID'),
INDEX 'Fac ID' ('Fac ID' ASC) VISIBLE,
 CONSTRAINT 'bc donation ibfk 1'
 FOREIGN KEY ('Don ID')
 REFERENCES 'blood bank'.'donor' ('Don ID'),
CONSTRAINT 'bc donation ibfk 2'
 FOREIGN KEY ('Fac ID')
 REFERENCES 'blood bank'.'facility' ('Fac ID'))
ENGINE = InnoDB
DEFAULT CHARACTER SET = utf8mb4
COLLATE = utf8mb4 0900 ai ci;
-- Table 'blood bank'. 'blood camp'
CREATE TABLE IF NOT EXISTS 'blood bank'.'blood camp' (
 'Fac ID' INT NOT NULL,
 'S date' DATE NULL DEFAULT NULL,
 'E date' DATE NULL DEFAULT NULL,
```

```
PRIMARY KEY ('Fac ID'),
 CONSTRAINT 'blood camp ibfk 1'
 FOREIGN KEY ('Fac ID')
  REFERENCES 'blood bank'.'facility' ('Fac ID'))
ENGINE = InnoDB
DEFAULT CHARACTER SET = utf8mb4
COLLATE = utf8mb4 0900 ai ci;
-- Table 'blood bank'.'employee'
CREATE TABLE IF NOT EXISTS 'blood bank'.'employee' (
 'Emp ID' INT NOT NULL,
 'BBank ID' INT NULL DEFAULT NULL,
 'Fac ID' INT NULL DEFAULT NULL,
 'Emp name' VARCHAR(255) NULL DEFAULT NULL,
 'gender' CHAR(1) NULL DEFAULT NULL,
 'age' INT NULL DEFAULT NULL,
 'contactno' VARCHAR(15) NULL DEFAULT NULL,
PRIMARY KEY ('Emp ID'),
 INDEX 'BBank ID' ('BBank ID' ASC) VISIBLE,
INDEX 'Fac ID' ('Fac ID' ASC) VISIBLE,
CONSTRAINT 'employee ibfk 1'
 FOREIGN KEY ('BBank ID')
```

```
REFERENCES 'blood bank'.'blood bank' ('BBank ID'),
CONSTRAINT 'employee ibfk 2'
 FOREIGN KEY ('Fac ID')
 REFERENCES 'blood bank'.'facility' ('Fac ID'))
ENGINE = InnoDB
DEFAULT CHARACTER SET = utf8mb4
COLLATE = utf8mb4 0900 ai ci;
-- Table 'blood bank'.'inventory'
CREATE TABLE IF NOT EXISTS 'blood bank'. 'inventory' (
 'Inven ID' INT NOT NULL,
 'BBank ID' INT NULL DEFAULT NULL,
 'type of Blood' VARCHAR(3) NULL DEFAULT NULL,
 'quantity' INT NULL DEFAULT NULL,
 `expiry_date` DATE NULL DEFAULT NULL,
PRIMARY KEY ('Inven ID'),
INDEX 'BBank ID' ('BBank ID' ASC) VISIBLE,
CONSTRAINT 'inventory ibfk 1'
 FOREIGN KEY ('BBank ID')
 REFERENCES 'blood bank'. 'blood bank' ('BBank ID'))
ENGINE = InnoDB
DEFAULT CHARACTER SET = utf8mb4
```

```
COLLATE = utf8mb4_0900_ai_ci;
```

```
-- Table 'blood bank'.'handles'
CREATE TABLE IF NOT EXISTS 'blood bank'. 'handles' (
 'Emp ID' INT NOT NULL,
 'Inven ID' INT NOT NULL,
 'Date' DATE NULL DEFAULT NULL,
 PRIMARY KEY ('Emp ID', 'Inven ID'),
 INDEX 'Inven ID' ('Inven ID' ASC) VISIBLE,
 CONSTRAINT 'handles ibfk 1'
 FOREIGN KEY ('Emp_ID')
 REFERENCES 'blood bank'.'employee' ('Emp ID'),
 CONSTRAINT 'handles ibfk 2'
 FOREIGN KEY ('Inven ID')
  REFERENCES 'blood bank'.'inventory' ('Inven ID'))
ENGINE = InnoDB
DEFAULT CHARACTER SET = utf8mb4
COLLATE = utf8mb4 0900 ai ci;
-- Table 'blood bank'. 'hospital'
```

```
CREATE TABLE IF NOT EXISTS 'blood bank'.'hospital' (
 'Fac ID' INT NULL DEFAULT NULL,
 'req date' DATE NULL DEFAULT NULL,
 'Blood type' VARCHAR(3) NULL DEFAULT NULL,
 'request ID' INT NOT NULL,
 'quantity' INT NULL DEFAULT NULL,
PRIMARY KEY ('request ID'),
INDEX 'Fac ID' ('Fac ID' ASC) VISIBLE,
 CONSTRAINT 'hospital ibfk 1'
 FOREIGN KEY ('Fac ID')
 REFERENCES 'blood bank'.'facility' ('Fac ID'))
ENGINE = InnoDB
DEFAULT CHARACTER SET = utf8mb4
COLLATE = utf8mb4 0900 ai ci;
-- Table 'blood bank'. 'patient'
CREATE TABLE IF NOT EXISTS 'blood bank'.'patient' (
 'Pat ID' INT NOT NULL,
 'Fac ID' INT NULL DEFAULT NULL,
 'pat name' VARCHAR(255) NULL DEFAULT NULL,
 'gender' CHAR(1) NULL DEFAULT NULL,
```

```
'age' INT NULL DEFAULT NULL,
 'Blood type' VARCHAR(3) NULL DEFAULT NULL,
 'contactno' VARCHAR(15) NULL DEFAULT NULL,
 PRIMARY KEY ('Pat ID'),
INDEX 'Fac ID' ('Fac ID' ASC) VISIBLE,
CONSTRAINT `patient_ibfk_1`
 FOREIGN KEY ('Fac ID')
 REFERENCES 'blood bank'.'facility' ('Fac ID'))
ENGINE = InnoDB
DEFAULT CHARACTER SET = utf8mb4
COLLATE = utf8mb4 0900 ai ci;
-- Table 'blood bank'.'requests'
CREATE TABLE IF NOT EXISTS 'blood bank'.'requests' (
 'BBank ID' INT NOT NULL,
 'Fac ID' INT NOT NULL,
 'date' DATE NULL DEFAULT NULL,
PRIMARY KEY ('BBank ID', 'Fac ID'),
 INDEX 'Fac ID' ('Fac ID' ASC) VISIBLE,
CONSTRAINT 'requests ibfk 1'
 FOREIGN KEY ('BBank ID')
  REFERENCES 'blood bank'.'blood bank' ('BBank ID'),
```

```
CONSTRAINT `requests_ibfk_2`

FOREIGN KEY (`Fac_ID`)

REFERENCES `blood_bank`.`facility` (`Fac_ID`))

ENGINE = InnoDB

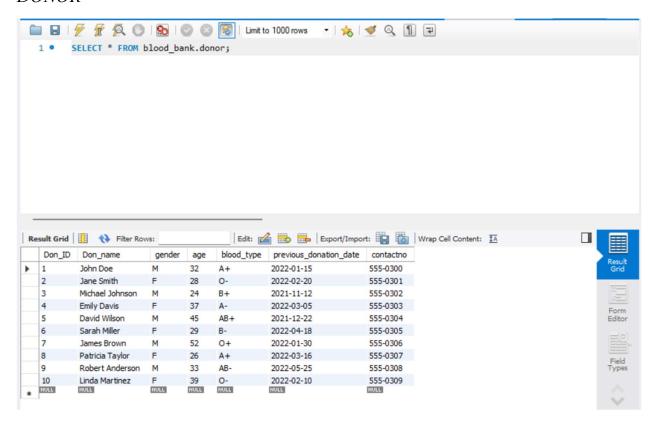
DEFAULT CHARACTER SET = utf8mb4

COLLATE = utf8mb4_0900_ai_ci;
```

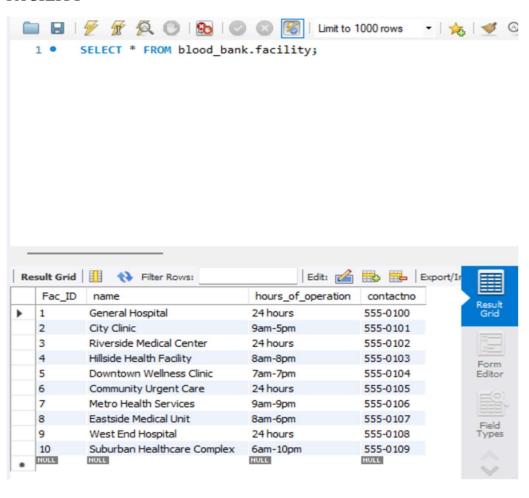
SET SQL\_MODE=@OLD\_SQL\_MODE;
SET FOREIGN\_KEY\_CHECKS=@OLD\_FOREIGN\_KEY\_CHECKS;
SET UNIQUE\_CHECKS=@OLD\_UNIQUE\_CHECKS;

## TABLE VIEWS:

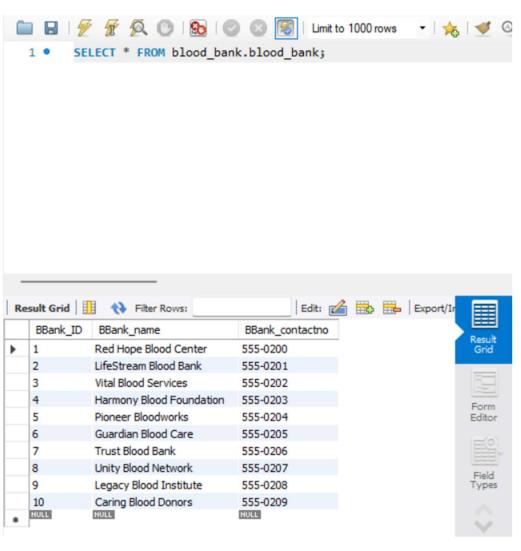
#### **DONOR**



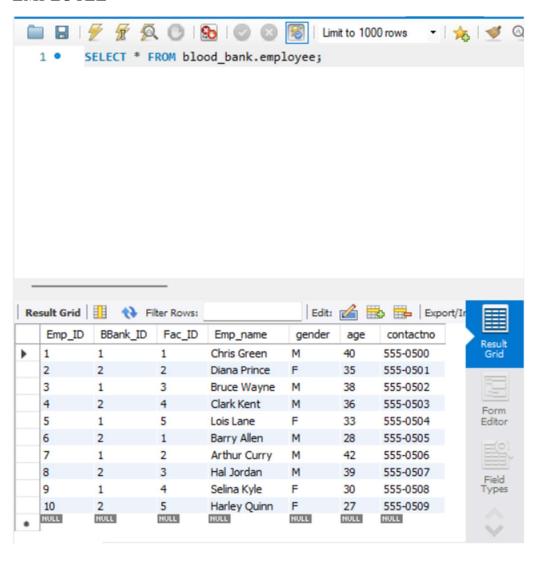
#### **FACILITY**



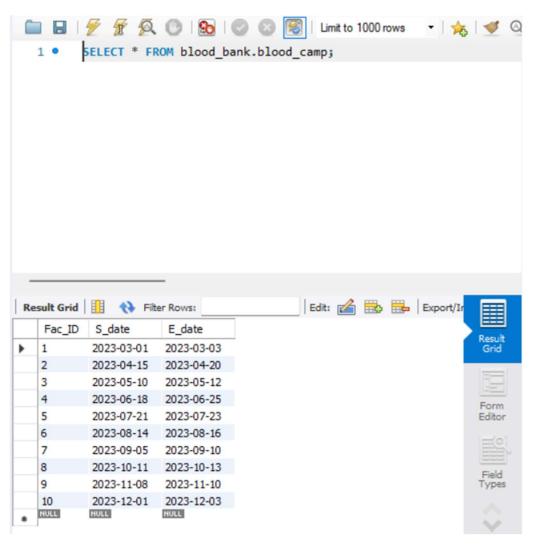
### BLOOD BANK



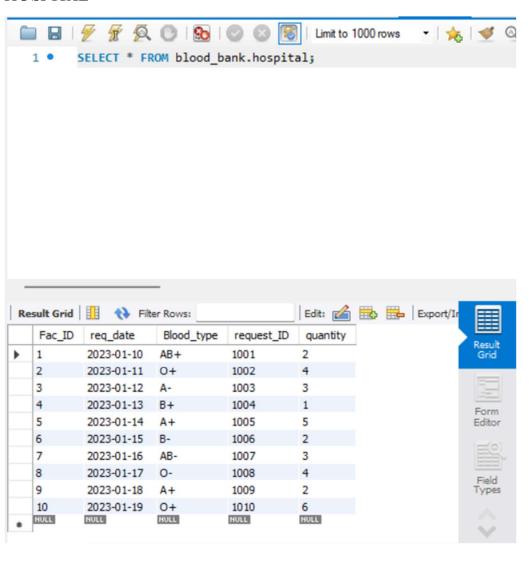
#### **EMPLOYEE**



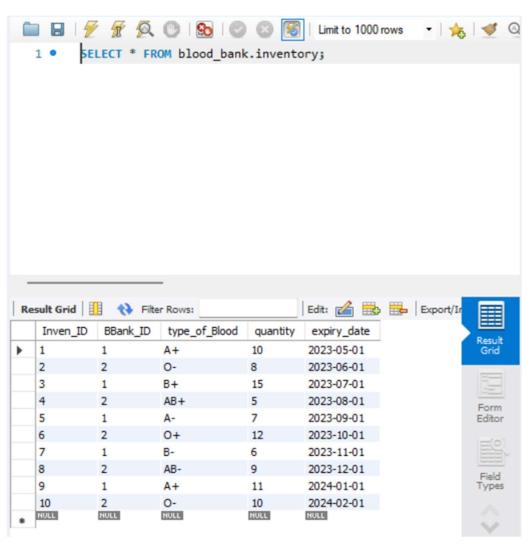
#### **BLOOD CAMP**



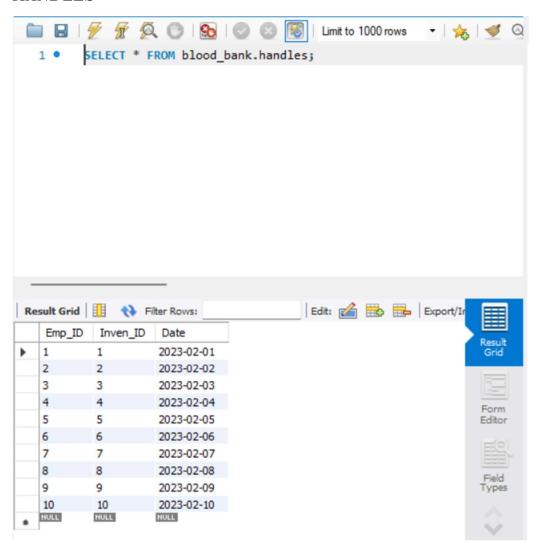
#### **HOSPITAL**



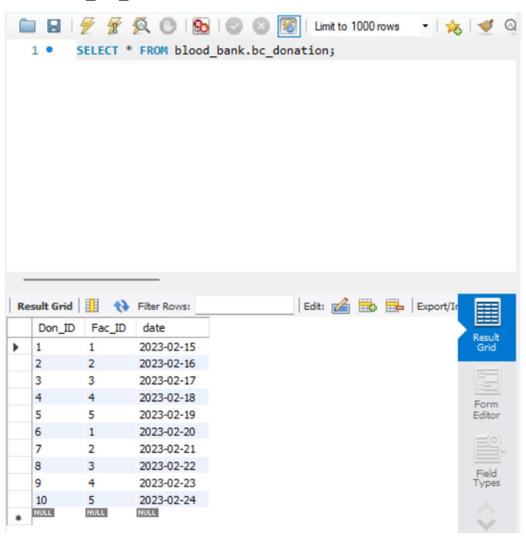
#### *INVENTORY*



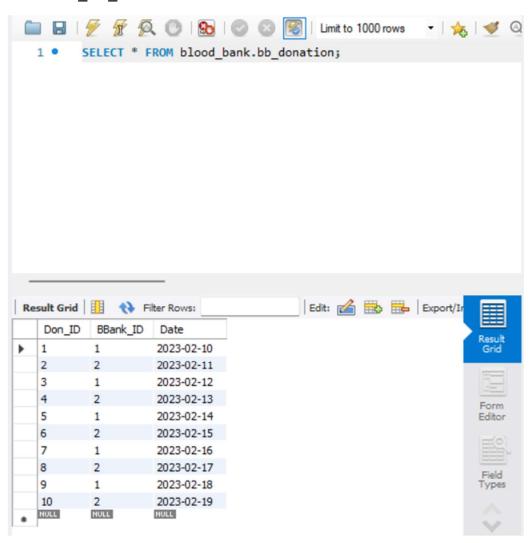
#### **HANDLES**



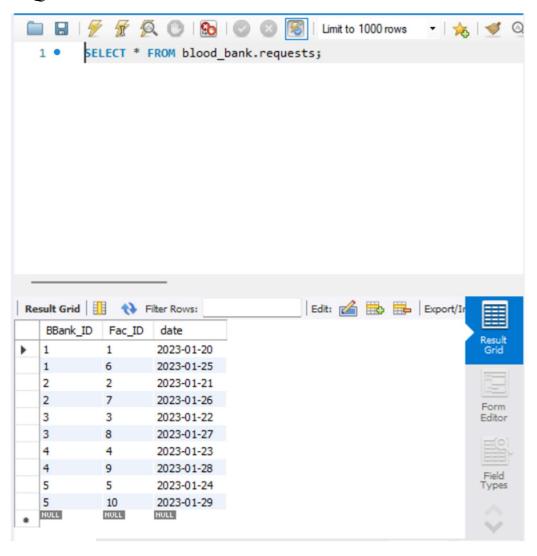
## DONATES TO BC



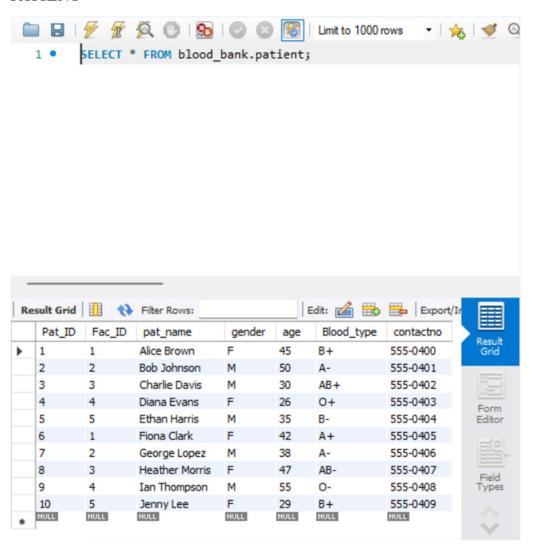
### DONATES\_TO\_BB



## **REQUESTS**



#### **PATIENT**



# **Transactions:**

1) Get a list of all blood banks along with the details of employees working there and the facilities they are linked to, including blood banks without any employees.

#### Query:

```
SELECT

BB.BBank_name,

E.Emp_name,

F.name AS 'Facility_Name'

FROM

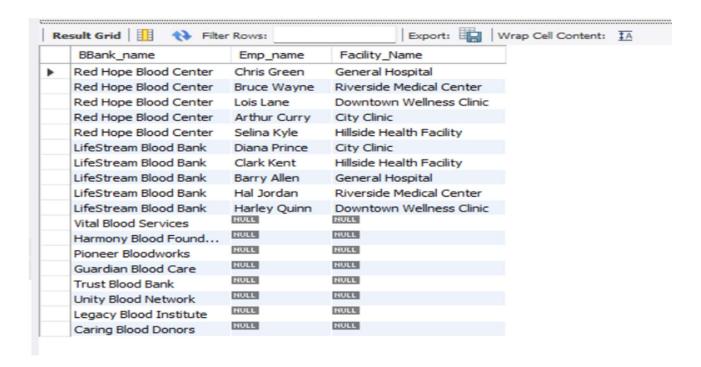
Blood_Bank BB

LEFT JOIN

Employee E ON BB.BBank_ID = E.BBank_ID

LEFT JOIN

Facility F ON E.Fac_ID = F.Fac_ID;
```



2) Display all blood type requests from hospitals, including those that have not yet been fulfilled.

#### Query:

```
SELECT

H.Blood_type,

H.req_date,

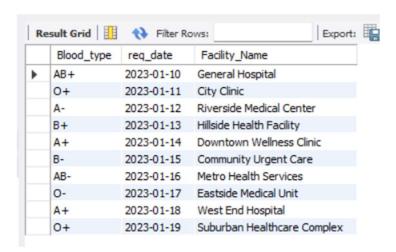
F.name AS 'Facility_Name'

FROM

Hospital H

RIGHT JOIN

Facility F ON H.Fac_ID = F.Fac_ID;
```



3) List all donors and their donation dates at blood camps, including those who have not yet donated.

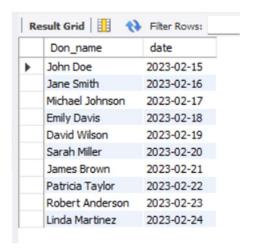
# Query:

```
D.Don_name,
BC.date
FROM
```

Donor D

**LEFT JOIN** 

BC\_donation BC ON D.Don\_ID = BC.Don\_ID;



4) Update contact information for a donor in the Donor table:

### Query:

**UPDATE** Donor

SET contactno = '555-0333'

WHERE Don ID = 1;

9 18:46:12 UPDATE Donor SET contactno = '555-0333' WHERE Don\_ID = 1

1 row(s) affected Rows matched: 1 Changed: 1 Warnings: 0

5) Find the total quantity of each blood type available in the inventory.

# Query:

SELECT type\_of\_Blood, SUM(quantity) AS Total\_Quantity

**FROM Inventory** 

GROUP BY type\_of\_Blood;

	type_of_Blood	Total_Quantity
•	A+	21
	0-	18
	B+	15
	AB+	5
	A-	7
	0+	12
	B-	6
	AB-	9

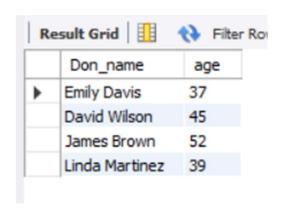
6) Get the names of donors who are older than the average age of all donors.

## Query:

SELECT Don\_name, age

**FROM Donor** 

WHERE age > (SELECT AVG(age) FROM Donor);



7) List each blood bank along with the count of donations it has received.

# Query:

SELECT BB.BBank\_name, COUNT(BD.BBank\_ID) AS Donation\_Count FROM Blood\_Bank BB

LEFT JOIN BB\_donation BD ON BB.BBank\_ID = BD.BBank\_ID
GROUP BY BB.BBank\_name;

	BBank_name	Donation_Count
١	Red Hope Blood Center	5
	LifeStream Blood Bank	5
	Vital Blood Services	0
	Harmony Blood Foundation	0
	Pioneer Bloodworks	0
	Guardian Blood Care	0
	Trust Blood Bank	0
	Unity Blood Network	
	Legacy Blood Institute	0
	Caring Blood Donors	0

8) Increase the age of all employees working at a specific facility by 1 year.

# Query:

**UPDATE** Employee

SET age = age + 1

WHERE Fac\_ID IN (SELECT Fac\_ID FROM Facility WHERE name = 'City Clinic');

2 18 18:58:29 UPDATE Employee SET age = age + 1 WHERE Fac\_ID IN (SELECT Fac\_ID FROM Facility WHERE name ... 2 row(s) affected Rows matched: 2 Changed: 2 Warnings: 0

9) Find the names of all donors who have donated at facilities where a blood camp was held.

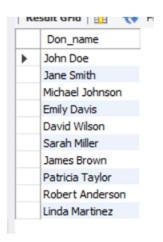
### Query:

SELECT DISTINCT D.Don\_name

FROM Donor D

JOIN BC\_donation BCD ON D.Don\_ID = BCD.Don\_ID

JOIN Blood\_camp BC ON BCD.Fac\_ID = BC.Fac\_ID;



10) Retrieve each donor's most recent donation date.

# Query:

SELECT D.Don\_ID, D.Don\_name, MAX(BD.Date) AS Last\_Donation\_Date
FROM Donor D

LEFT JOIN BB\_donation BD ON D.Don\_ID = BD.Don\_ID

GROUP BY D.Don\_ID, D.Don\_name;

	Don_ID	Don_name	Last_Donation_Date
•	1	John Doe	2023-02-10
	2	Jane Smith	2023-02-11
	3	Michael Johnson	2023-02-12
	4	Emily Davis	2023-02-13
	5	David Wilson	2023-02-14
	6	Sarah Miller	2023-02-15
	7	James Brown	2023-02-16
	8	Patricia Taylor	2023-02-17
	9	Robert Anderson	2023-02-18
	10	Linda Martinez	2023-02-19

# Conclusion:

In conclusion, this project offers a comprehensive exploration of the intricacies involved in developing a database, with a specific focus on creating a blood bank database—a relevant and practical application in the modern world. Throughout the development process, the project team encountered various challenges, including the complexities of designing an effective ER diagram and establishing relationships between different entities. One particular challenge involved managing subsets within the facility entity, which demanded a nuanced understanding of specialization and generalization concepts. Consequently, the final blood bank database stands as a robust and efficient solution that meets all the project's initial criteria.

In summary, this project provides a valuable and informative overview of the multifaceted nature of database development. From the challenges encountered to the effective solutions devised, it offers a wealth of insights and lessons applicable to a wide array of real-world scenarios. Given the increasing importance of data management and analysis in today's world, the skills and techniques demonstrated in this project are likely to be in high demand among students and professionals alike.