





Decision Analytics

Lecture 23: Wrap-up

Wrap-up: module overview

Introduction & Tools

Modelling

Constraint programming

Linear programming

Tools

- NumPy
 - Matrix operations
 - Linear algebra



- OR Tools
 - Constraint programming
 - Linear programming



Mathematical Optimisation

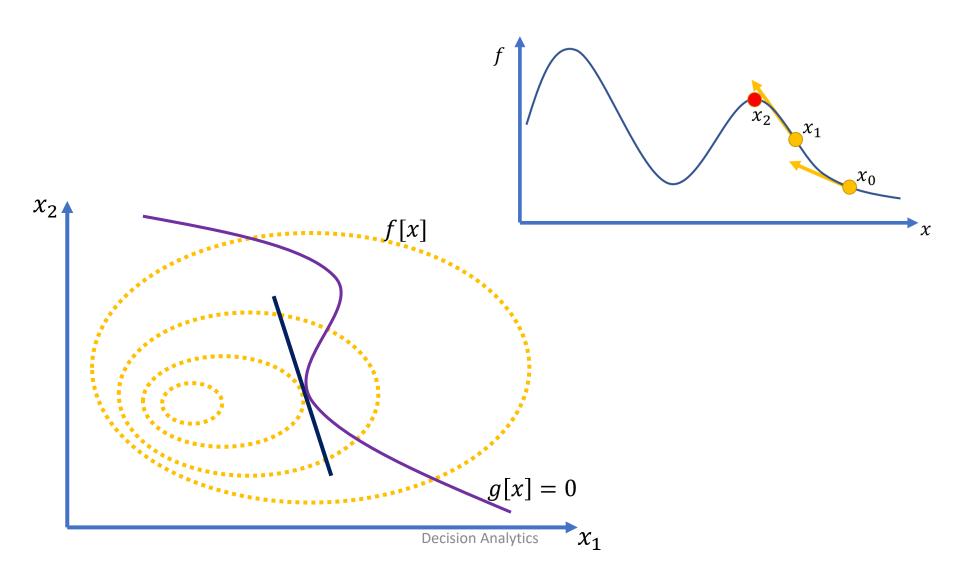
Given a domain D and an objective function

$$f:D\to\mathbb{R}$$

find $\hat{x} \in D$ that minimises f, i.e. $\forall x \in D : f(\hat{x}) \leq f(x)$

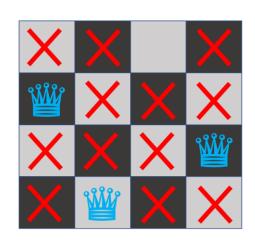
(for maximisation problems replace f with -f)

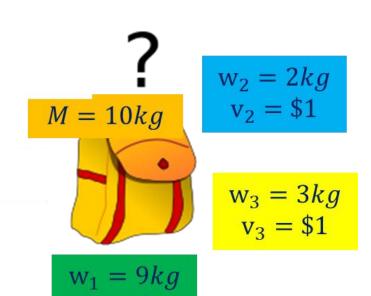
Gradient-based search & Lagrange multipliers



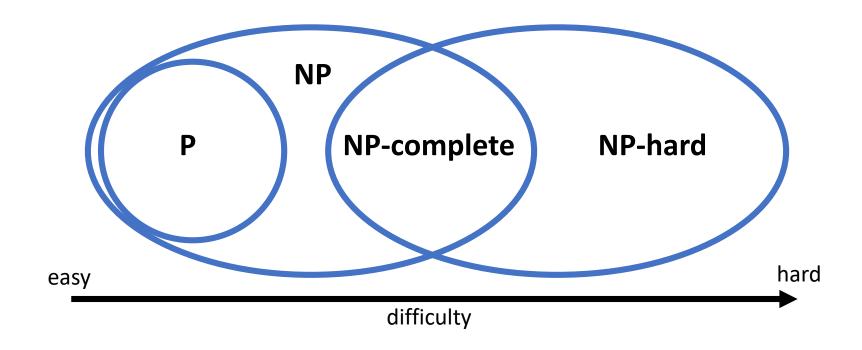
Constraint Satisfaction Problems

$$\begin{split} X = &< x_1, \dots, x_n > \\ D = &< D_1, \dots, D_n > \\ x_i \in D_i \\ C = &< C_1, \dots, C_t > \\ C_i = &< R_{S_i}, S_i > \\ S_i \subset X, R_{S_i} \subset D_{S_{i_1}} \times \dots \times D_{S_{i_{|S_i|}}} \end{split}$$





Complexity classes



Boolean algebra and binary domains

$$(x_1 \lor \neg x_2 \lor x_3) \land (\neg y_1 \lor y_2) \land (z_1 \lor z_2 \lor z_3 \lor z_4)$$

```
model.AddBoolOr([x1,x2.Not(),x3])
model.AddBoolOr([y1.Not(),y2])
model.AddBoolOr([z1,z2,z3,z4])
```

x	у	$x \wedge y$	$x \vee y$	$\neg x$
False	False	False	False	True
True	False	False	True	False
False	True	False	True	
True	True	True	True	

Planning problems and SAT



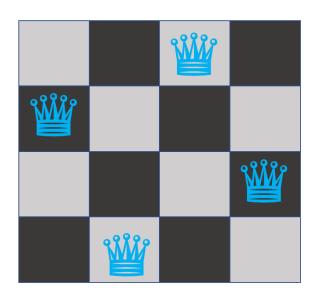


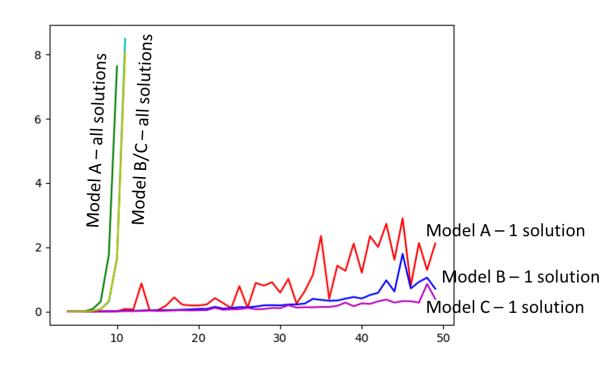


First-order logic and SAT

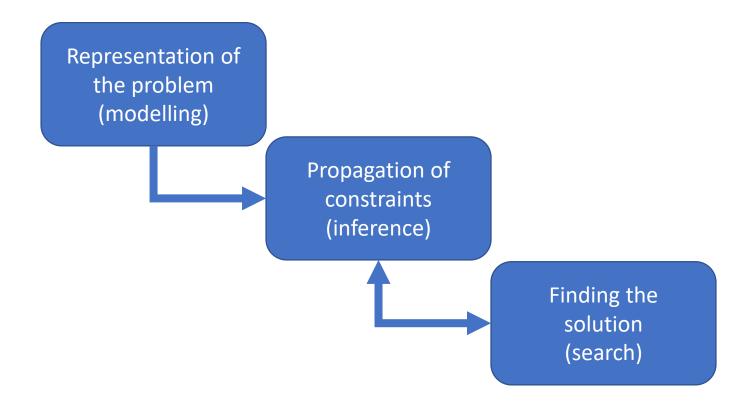


Modelling beyond SAT





Constraint programming

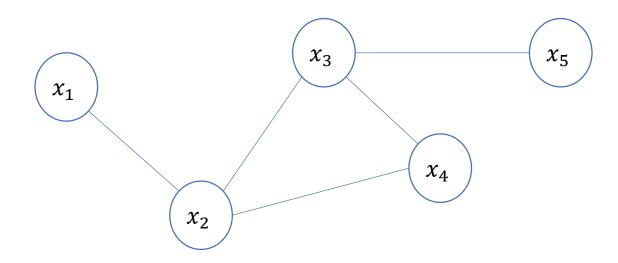


Constraint networks

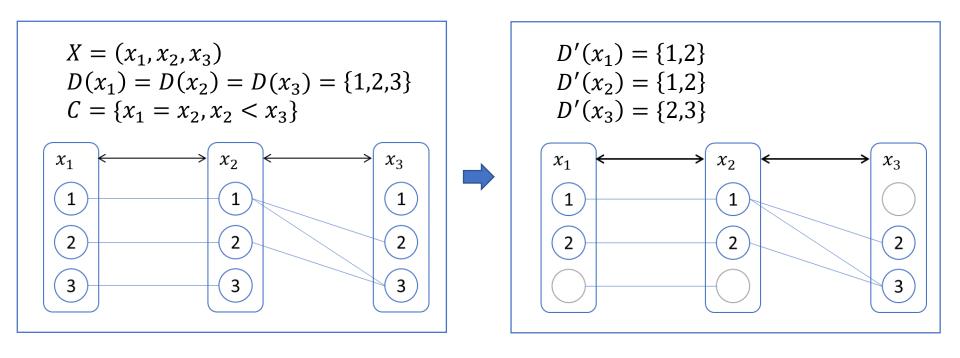
$$X = (x_1, x_2, x_3, x_4, x_5)$$

$$D(x_i) = \{1, 2, 3, 4, 5\}$$

$$C = \{x_1 < x_2, x_2 = x_3, x_3 \ge x_4, x_3 \ge x_5, x_2 \le x_4\}$$



Arc consistency

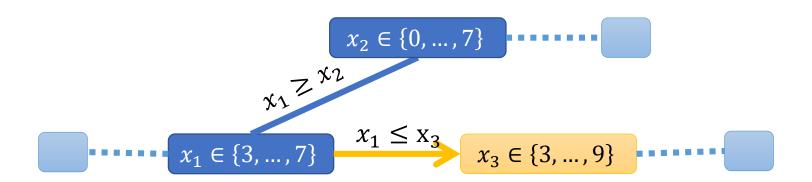


AC3 Algorithm

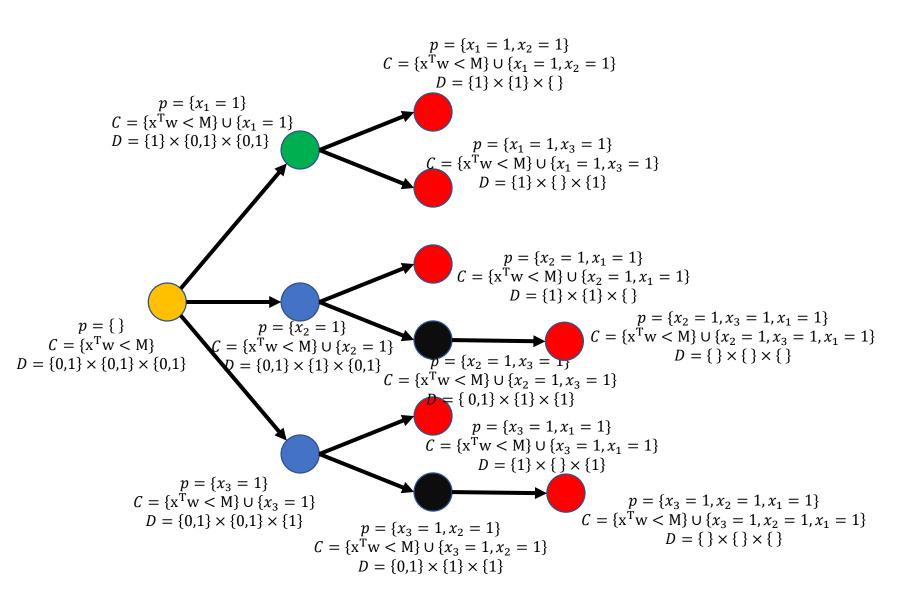
```
function Revise3(in x_i: variable; c: constraint): Boolean;
   begin
        CHANGE \leftarrow false;
        foreach v_i \in D(x_i) do
 2
             if \exists \tau \in c \cap \pi_{X(c)}(D) with \tau[x_i] = v_i then
 3
                  remove v_i from D(x_i);
 4
                  CHANGE ← true;
 5
        return CHANGE;
   end
function AC3/GAC3(in X: set): Boolean;
   begin
        /* initalisation */:
    Q \leftarrow \{(x_i, c) \mid c \in C, x_i \in X(c)\};
        /* propagation */;
        while Q \neq \emptyset do
 8
             select and remove (x_i, c) from Q;
 9
             if Revise(x_i, c) then
10
                  if D(x_i) = \emptyset then return false;
11
                  else Q \leftarrow Q \cup \{(x_i, c') \mid c' \in C \land c' \neq c \land x_i, x_i \in X(c') \land j \neq i\};
12
13
        return true;
    end
```

Propagator iteration

$$D''(x_i) = \pi_{\{x_i\}}(c_j \cap \pi_{X(c_j)}(D'))$$



Posting constraints



Branching strategies

```
CHOOSE_FIRST
CHOOSE_LOWEST_MIN
CHOOSE_HIGHEST_MAX
CHOOSE_MIN_DOMAIN_SIZE
CHOOSE_MAX_DOMAIN_SIZE
```

```
SELECT_MIN_VALUE

SELECT_MAX_VALUE

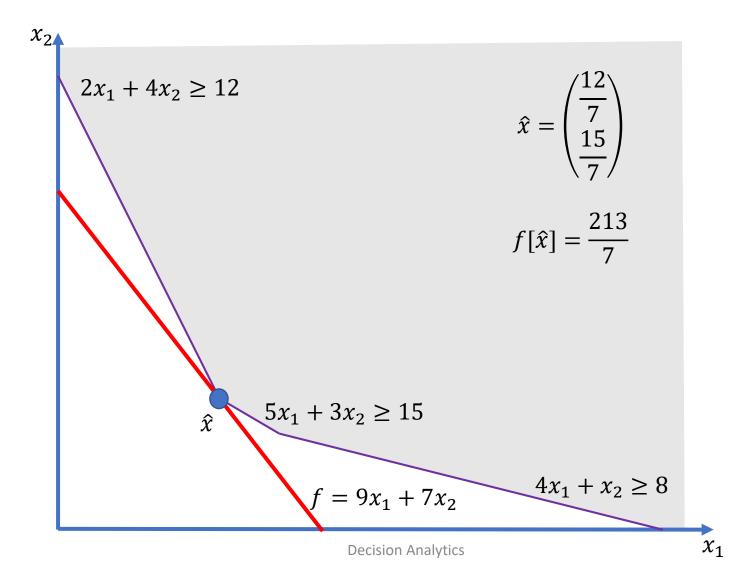
SELECT_LOWER_HALF

SELECT_UPPER_HALF
```

Maintaining Arc Consistency (MAC)

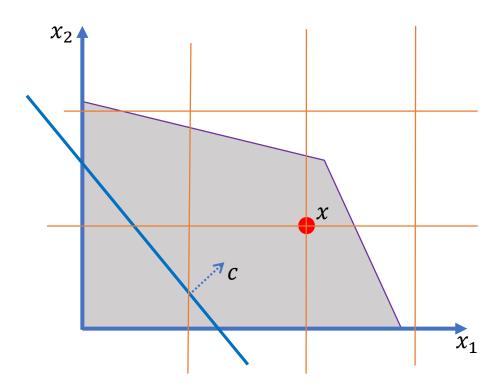
```
Search (p, (X, D, C)):
         (X,D',C) = make arc consistent (X,D,C)
         if \exists D'(x_i): |D'(x_i)| = 0
                   (X,D',C) is a no-good
                   return
         else if \forall D'(x_i): |D'(x_i)| = 1
                   (X,D',C) is a solution
                   return
         else
                   choose x_i so that |D'(x_i)| > 1
                   choose v for b^1 = \{x_i \le v\}, b^2 = \{x_i > v\}
                   Search (p \cup \{b^1\}, (X, D', C \cup \{b^1\}))
                   Search (p \cup \{b^2\}, (X, D', C \cup \{b^2\}))
```

Linear programming



The simplex algorithm

Integer Linear Programming



Thank you for your attention!