# EECE 7204-Assignment1

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# 1 Coding

#### 1.1 Insertion Sort

```
#include <iostream>
#include <time.h>
int *get_input(int limit){
   /**
    st Generating input - Worst case happens when the array is sorted in
    * reverse order. Sending numbers from 1000 to limit
    * decrementing by 1 for every element.
 int* input = new int[limit];
 int j = 0;
 for(int i=1000; i>1000-limit; i--){
    input[j] = i;
    j+=1;
 }
 return input;
int *insertion_sort(int input[], int limit){
    * Insertion sort algorithm
 int i, j, temp;
 for (i=0; i< limit; i++){</pre>
   for (j=i ; j>=0; j--){
     if(input[j] < input[j-1]){</pre>
```

```
temp = input[j];
        input[j] = input[j-1];
        input[j-1] = temp;
   }
 }
 return input;
int main(){
 int limit, i;
 std :: cout <<"Enter limit: ";</pre>
 std :: cin>> limit;
 int* input = get_input(limit);
  * Setting clocks before and after the sorting
 clock_t start = clock();
  int* sorted_array = insertion_sort(input, limit);
 clock_t end = clock();
 double cpu_time_used = ((double) (end - start)) / CLOCKS_PER_SEC;
 printf("Insertion Sort took %f seconds to finish \n", cpu_time_used);
}
1.2 Merge Sort
#include <iostream>
#include <time.h>
int *get_input(int limit){
   /**
    * Generating input - Worst case - Sending numbers from 1000 to limit
    * decrementing by 1 for every element.
 int* input = new int[limit];
 int j = 0;
 for(int i=1000; i>1000-limit; i--){
   input[j] = i;
```

```
j+=1;
 }
 return input;
void copy_merged_array_into_original(int original[], int merged[], int length, int left_low) {
   for (int i=0; i< length; ++i)</pre>
        original[left_low++] = merged[i];
}
void merge(int input[], int left_low, int left_high, int right_low, int right_high) {
    int length = right_high-left_low+1;
    int merged_array[length];
   int left = left_low;
   int right = right_low;
   auto left_array_exhausted = [&left, &left_high]() { return left > left_high;};
   auto right_array_exhausted = [&right, &right_high]() { return right > right_high;};
   for (int i = 0; i < length; ++i) {
        if (left_array_exhausted())
            merged_array[i] = input[right++];
        else if (right_array_exhausted())
            merged_array[i] = input[left++];
        else if (input[left] <= input[right])</pre>
            merged_array[i] = input[left++];
        else
            merged_array[i] = input[right++];
   }
    copy_merged_array_into_original(input, merged_array, length, left_low);
}
void merge_sort(int numbers[], int low, int high) {
    if (low >= high)
       return;
   else {
         * Recursive sorting and merging parts
```

```
*/
        int mid = (low + high)/2;
        merge_sort(numbers, low, mid);
        merge_sort(numbers, mid+1, high);
        merge(numbers, low, mid, mid+1, high);
    }
}
int main(){
    int limit;
    std :: cout << "Enter limit: ";</pre>
    std :: cin >> limit;
    int* input = get_input(limit);
    std :: cout << "Sorting using mergesort: ";</pre>
   * Setting clocks before and after the sorting
    clock_t start = clock();
    merge_sort(input, 0, limit-1);
    clock_t end = clock();
    double cpu_time_used = ((double) (end - start)) / CLOCKS_PER_SEC;
    printf("Merge Sort took %f seconds to finish \n", cpu_time_used);
```

# 2 Arrangement of Elements during Sorting

#### 2.1 Insertion Sort

}

## 2.2 Quicksort

Input: 10, 5, 7, 9, 8, 3 Iteration 1: 10

### 3 True or False

- $n + 3 \in \Omega(n)$  True, n+3 is lower bounded by  $\Omega(n)$
- n + 3  $\in \Omega(n^2)$  False, n+3 is not lower bounded by  $\Omega(n^2)$
- $n + 3 \in \Theta(n^2)$  False
- $2^{n+1} \in \mathcal{O}(n+1)$  False
- $2^{n+1} \in \Theta(2^n)$  True

## 4 Master Method

•  $T(n) = 8T(\frac{n}{2}) + n$ 

Here a = 8, b = 2 and f(n) = n

Considering Case I of the Master Method, i,e  $\mathcal{O}(n^{\log_b a - \varepsilon})$  and substituting for a and b, we have:

$$f(n) = n = \mathcal{O}(n^{\log_2 8 - \varepsilon}) = \mathcal{O}(n^{3 - \varepsilon})$$

For  $\varepsilon = 1$ , we have  $f(n) = \mathcal{O}(n^2)$ , which satisfies Case I

$$\therefore T(n) = \Theta(n^{\log_b a}) = \Theta(n^3)$$

•  $T(n) = 8T(\frac{n}{2}) + n^2$ 

Here a = 8, b = 2 and  $f(n) = n^2$ 

Considering Case II of the Master Method, i,e  $\Theta(n^{\log_b a})$  and substituting for a and b, we have:

$$f(n) = n^2 = \Theta(n^{\log_2 8}) = \Theta(n^3)$$
, which satisfies Case II

$$\therefore T(n) = \Theta(n^{\log_b a}.log^{k+1}.n) = \Theta(n^3.log n)$$

•  $T(n) = 8T(\frac{n}{2}) + n^3$ 

Here a = 8, b = 2 and  $f(n) = n^3$ 

Considering Case III of the Master Method, i,e  $\Theta(n^{\log_b a})$  and substituting for a and b, we have:

$$f(n) = n^3 = \Omega(n^{\log_2 8 + \varepsilon}) = \Omega(n^{3 + \varepsilon})$$
(1)

and,

$$8.f(\frac{n}{2}) \le (1 - \varepsilon').n^3 \tag{2}$$

$$\implies 8.\frac{n^3}{8} \le (1 - \varepsilon').n^3 \tag{3}$$

$$\implies n^3 \le (1 - \varepsilon').n^3 \tag{4}$$

(4) is true for many  $\varepsilon' s > 0$ 

From (1) and (4) Master Method Case III is satisfied.

$$\therefore \mathbf{T}(\mathbf{n}) = \Theta(\mathbf{f}(\mathbf{n})) = \Theta(n^3)$$

• 
$$T(n) = 8T(\frac{n}{2}) + n^4$$