

# NTRU Post-Quantum Encryption

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## 1 Example from Wikipedia

Let's try to re-create <https://en.wikipedia.org/wiki/NTRUEncrypt>. In this example the parameters  $(N, p, q)$  will have the values  $N = 11$   $p = 3$   $q = 2^5$ . The polynomials  $f$  and  $g$  are randomly chosen, so suppose they are represented by  $f = -1 + x + x^2 - x^4 + x^6 + x^9 - x^{10}$   
 $g = -1 + x^2 + x^3 + x^5 - x^8 - x^{10}$

### 1.1 `balancedmod(f,q)`

This function reduces every coefficient of a  $Z[x]$  polynomial  $f$  modulo  $q$  with additional balancing, so the result coefficients are integers in interval  $[-\frac{q}{2}, +\frac{q}{2}]$ . More specifically: for an odd  $q$   $[-\frac{q-1}{2}, +\frac{q-1}{2}]$ , for an even  $q$   $[-\frac{q}{2}, +\frac{q}{2} - 1]$ . It returns  $Z[x]$  reduced polynomial.

```
def balancedmod(f,q):  
    g = list(((f[i] + q//2) % q) - q//2 for i in range(N))  
    return Zx(g)
```

Example:

$$\text{balancedmod}(17 + 31x + 34x^2 + 33x^3 - 8x^4, 32) = -8x^4 + x^3 + 2x^2 - x - 15$$

### 1.2 `multiply(f,g)`

This function performs a multiplication operation specific for NTRU, which works like a traditional polynomial multiplication with additional reduction of the result by  $x^N - 1$  ( $x^n$  is replaced by 1,  $x^n - 1$  by  $x$ ,  $x^n - 2$  by  $x^2$ , ...) and returns  $Z[x]$  polynomial

```
def convolution(f,g):  
    return (f * g) % (x^N-1)
```

### 1.3 invertmodprime(f,p)

This function calculates an inversion of a polynomial modulo  $x^N - 1$  and then modulo  $p$  with assumption that  $p$  is prime. Also it returns a  $Z[x]$  polynomial  $h$  such as convolution of  $h \cdot f = 1 \pmod{p}$  and raises an exception if such  $Z[x]$  polynomial  $h$  doesn't exist.

```
def invertmodprime(f,p):
    Zq.<z> = PolynomialRing(Integers(p))
    ZQphi.<Z> = Zq.quotient(z^N-1)
    a = f % p
    a = a.subs(x=z)
    k = 0
    b = 1*z^0
    c = 0*z^0
    f = a
    g = z^N-1
    assert a.gcd(g) in {i for i in range(p)}

    while True:
        while list(f)[0] == 0:
            f = f / Z
            c = c * Z
            k += 1

        if find_degree(list(f)) == 0:
            b = 1/list(f)[0] * b
            ans = Z^(N-k) * b
            return Zx(ans.lift())

        if find_degree(list(f)) < find_degree(list(g)):
            f, g = g, f
            b, c = c, b

    u = list(f)[0] * (1/list(g)[0])
    f = f - u*g
    b = b - u*c
```

### 1.4 invertmodpowerof2(a, p)

This function calculates an inversion of a polynomial modulo  $x^N - 1$  and then modulo  $q$  with assumption that  $q$  is a power of 2. Returns a polynomial  $f_q \in Z[x]$  such as  $f \cdot f_q = 1 \pmod{q}$ . An exception is thrown if such polynomial  $Z[x]$  does not exist.

```
def invertmodpowerof2(a, p):
```

```

r = int(math.log(p, 2))
p = 2
q = p
b = invertmodprime(a, p)
while q < p^r:
    q = q^2
    b = b * (2 - a*b) % q % (x^N-1)

b = b % p^r % (x^N - 1)
return b

```

Example:

$$f = -x^{10} + x^9 + x^6 - x^4 + x^2 + x - 1, p = 3, N = 11, q = 2^5$$

$$f_q = \text{invertmodpowerof2}(f, q) = 30x^{10} + 18x^9 + 20x^8 + 22x^7 + 16x^6 + 15x^5 + 4x^4 + 16x^3 + 6x^2 + 9x + 5$$

Note that this is exactly the inverse mentioned in Wikipedia - NTRU.

## 1.5 generate\_keys(poly1 = None, poly2 = None)

In this section we generate a public and secret key. Function `generate_keys(poly1 = None, poly2 = None)` generates a public and private key pair, based on provided parameters. Returns  $Z[x]$  public key and a secret key as a tuple of  $Z[x]$   $f$  (private key) and  $Z[x] F_p$

```

def generate_keys(poly1 = None, poly2= None):
    if validate_params():
        while True:
            try:
                if poly1 is None or poly2 is None:
                    f = generate_polynomial(d+1, d)
                    g = generate_polynomial(d, d)
                else:
                    f = poly1
                    g = poly2
                f_q = invertmodpowerof2(f,q)
                f_p = invertmodprime(f,p)
                break
            except:
                pass
        public_key = balancedmod(p * convolution(f_q,g),q)
        secret_key = f,f_p
        return public_key,secret_key

    else:
        print("Provided params are not correct.")

```

Example: In this example the parameters (N, p, q) will have the values  $N = 11$ ,  $p = 3$  and  $q = 32$  and therefore the polynomials  $f$  and  $g$  are of degree at most 10. The system parameters (N, p, q) are known to everybody. The polynomials are randomly chosen, so suppose they are represented by  $f(x) = -1 + x + x^2 - x^4 + x^6 + x^9 - x^{10}$ ,  $g(x) = -1 + x^2 + x^3 + x^5 - x^8 - x^{10}$

$public\_key = get\_keys(f, g) = -16x^{10} - 13x^9 + 12x^8 - 13x^7 + 15x^6 - 8x^5 + 12x^4 - 12x^3 - 10x^2 - 7x + 8$

## 1.6 encrypt(message, public\_key, r = None)

This function performs encryption of a given message using a provided public key and returns  $Z[x]$  encrypted message.

```
def encrypt(message, public_key, r = None):
    if r is None:
        r = generate_polynomial(d, d-1)
    return balancedmod(convolution(public_key, r) + message, q)
```

Example. Let's choose message  $\mathbf{m} = -1 + x^3 - x^4 - x^8 + x^9 + x^{10}$  and random polynomial  $\mathbf{r} = -1 + x^2 + x^3 + x^4 - x^5 - x^7$

$\mathbf{e} = \text{encrypt}(\text{message}, \text{public\_key}) = 19x^{10} + 6x^9 + 25x^8 + 7x^7 + 30x^6 + 16x^5 + 14x^4 + 24x^3 + 26x^2 + 11x + 14$

## 1.7 decrypt(encrypted\_message, secret\_key)

Let's decrypt the message.  $\mathbf{a} = \mathbf{f} \cdot \mathbf{e} = \mathbf{pr} \cdot \mathbf{g} + \mathbf{f} \cdot \mathbf{m} \pmod{q}$ . The next step will be to calculate  $\mathbf{a}$  modulo p:  $\mathbf{b} = \mathbf{a} \pmod{p}$ .  $\mathbf{c} = \mathbf{f}_p \cdot \mathbf{b} = \mathbf{f}_p \cdot \mathbf{f} \cdot \mathbf{m} = \mathbf{m} \pmod{p}$ . This function performs decryption of a given ciphertext using an own private key. Returns  $Z[x]$  decrypted message

```
def decrypt(encrypted_message, secret_key):
    f, f_p = secret_key
    a = balancedmod(convolution(encrypted_message, f), q)
    return balancedmod(convolution(a, f_p), p)
```

Example.

$\mathbf{a} = \mathbf{f} \cdot \mathbf{e} = 25x^{10} + 29x^9 - 27x^8 + 7x^7 + 6x^6 + 7x^5 - 22x^4 - 11x^3 + 22x^2 - 7x + 3$

after central lift

$$\mathbf{b} = \mathbf{a} = x^{10} - x^9 + x^7 + x^5 - x^4 + x^3 + x^2 - x$$

$$\mathbf{c} = x^{10} + x^9 - x^8 - x^4 + x^3 - 1$$

We got the original message!

## References

- [1] Implementation by Elena Mashkina <https://github.com/elena-mashkina/ntru/blob/master/NTRU.sage>
- [2] Explanation <https://cr.yp.to/talks/2018.11.16/slides-djb-20181116-lattice-a4.pdf>