

ASSIGNMENT-4

1.

$$\mathbb{Z} = \{\dots, -3, -2, -1, 0, 1, 2, 3, \dots\}$$

a.

$+$ \rightarrow addition

\times \rightarrow multiplication.

a. $\forall a, b \in \mathbb{Z} \quad ((a+b = b+a) \wedge (a \times b = b \times a))$

b. $\forall a, b \in \mathbb{Z} \quad ((a+b \in \mathbb{Z}) \wedge (a \times b \in \mathbb{Z}))$

c. $\forall a, b, c \in \mathbb{Z} \quad (a \times (b+c) = (a \times b) + (a \times c))$

d. $\forall a, b, c \in \mathbb{Z} \quad (((a+b)+c = a+(b+c)) \wedge ((a \times b) \times c = a \times (b \times c)))$

e. $\forall a \in \mathbb{Z} \quad ((a+0 = a) \wedge (a \times 1 = a))$

2.

$$\alpha = \forall x (P(x) \vee Q(x))$$

$$\beta = \forall x P(x) \vee \forall x Q(x)$$

Let's take

$$P(x) = x$$

$$Q(x) = \neg x$$

and let's assume ~~α is true~~
 $P(x) = \text{true}$
 $\Rightarrow Q(x) = \text{false}$
 for α .

$$\Rightarrow \alpha = \text{True}$$

But since

$\forall x P(x)$ and $\forall x Q(x)$ are independent statements

and $\forall x P(x)$ can be false and $\forall x Q(x)$ can be false

In such a case $\beta = \text{False}$

$\therefore \alpha$ does not entail β .

3.

(a) Owns (Han Solo, Millennium Falcon)

(b) Unhappy (Princess Leia) c) loves (Princess Leia, Han Solo)

(d) $\forall x ((\text{Owns}(x, \text{Millennium Falcon}) \vee \text{Unhappy}(x)) \Rightarrow \text{visits}(x, \text{Obi-wan Kenobi}))$

(e) $\forall x (\text{visits}(x, \text{Obi-wan Kenobi}) \Rightarrow \text{wise}(x))$

(f) $\forall x ((\text{Owns}(x, \text{Millennium Falcon}) \wedge \text{visits}(x, \text{Obi-wan Kenobi})) \Rightarrow \text{teaches to use lightsaber}(\text{Obi-wan Kenobi}, x))$

(g) $\forall x (((\text{Unhappy}(x) \vee \text{Owns}(x, \text{Millennium Falcon})) \wedge \text{teaches to use lightsaber}(\text{Obi-wan Kenobi}, x)) \Rightarrow \text{joins}(x, \text{Rebel Alliance}))$

(h) $\forall x, y ((\text{Unhappy}(x) \wedge \text{loves}(x, y)) \Rightarrow \text{declares}_{\text{love}}(x, y))$

(i) $\forall x, y ((\text{teaches to use lightsaber}(\text{Obi-wan Kenobi}, x) \wedge \text{declares love}(y, x) \wedge \text{wise}(y)) \Rightarrow \text{friend}(x, \text{Chewbacca}))$

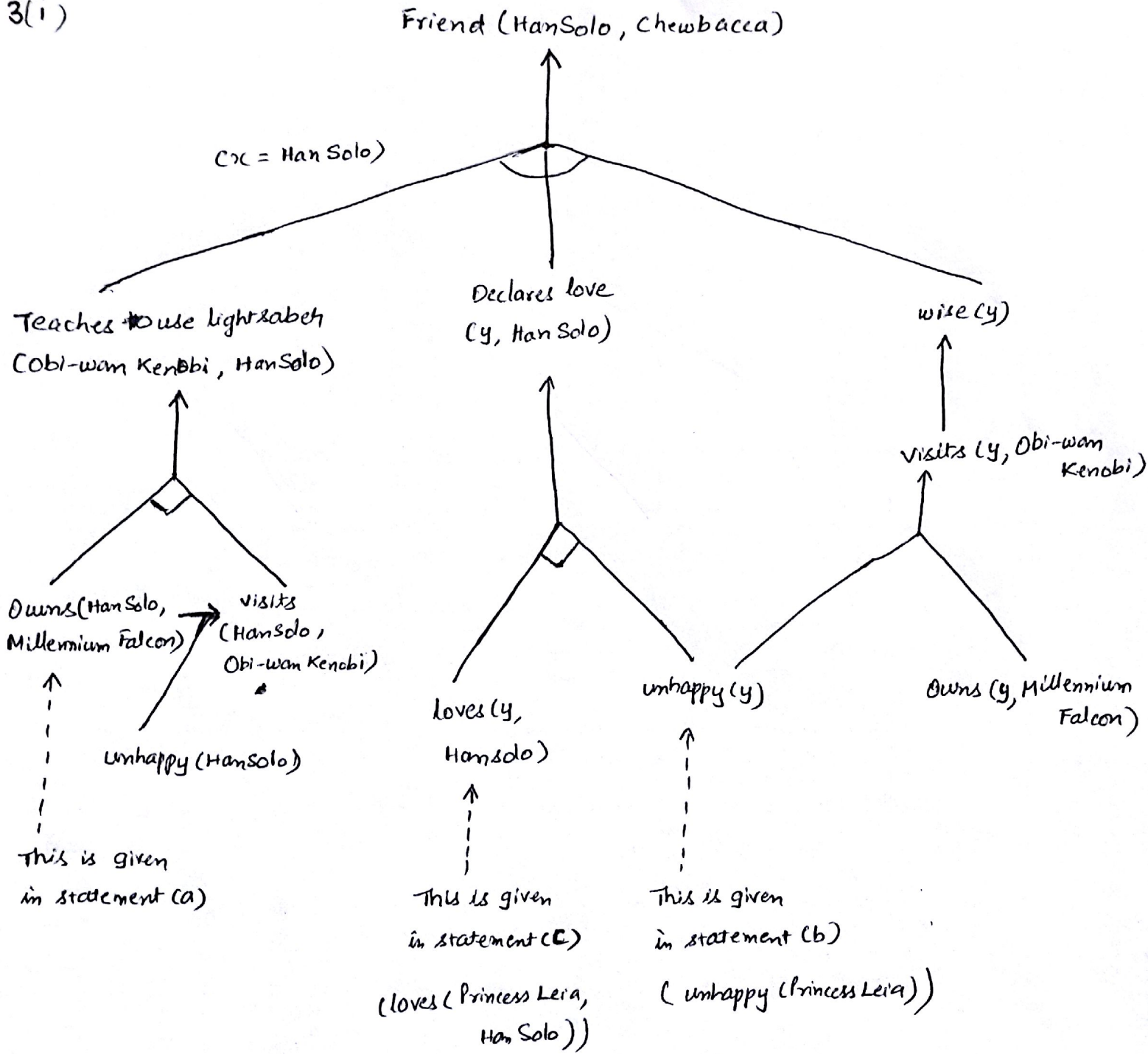
(i) Backward Chaining:



Please Turn Over

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3(i)

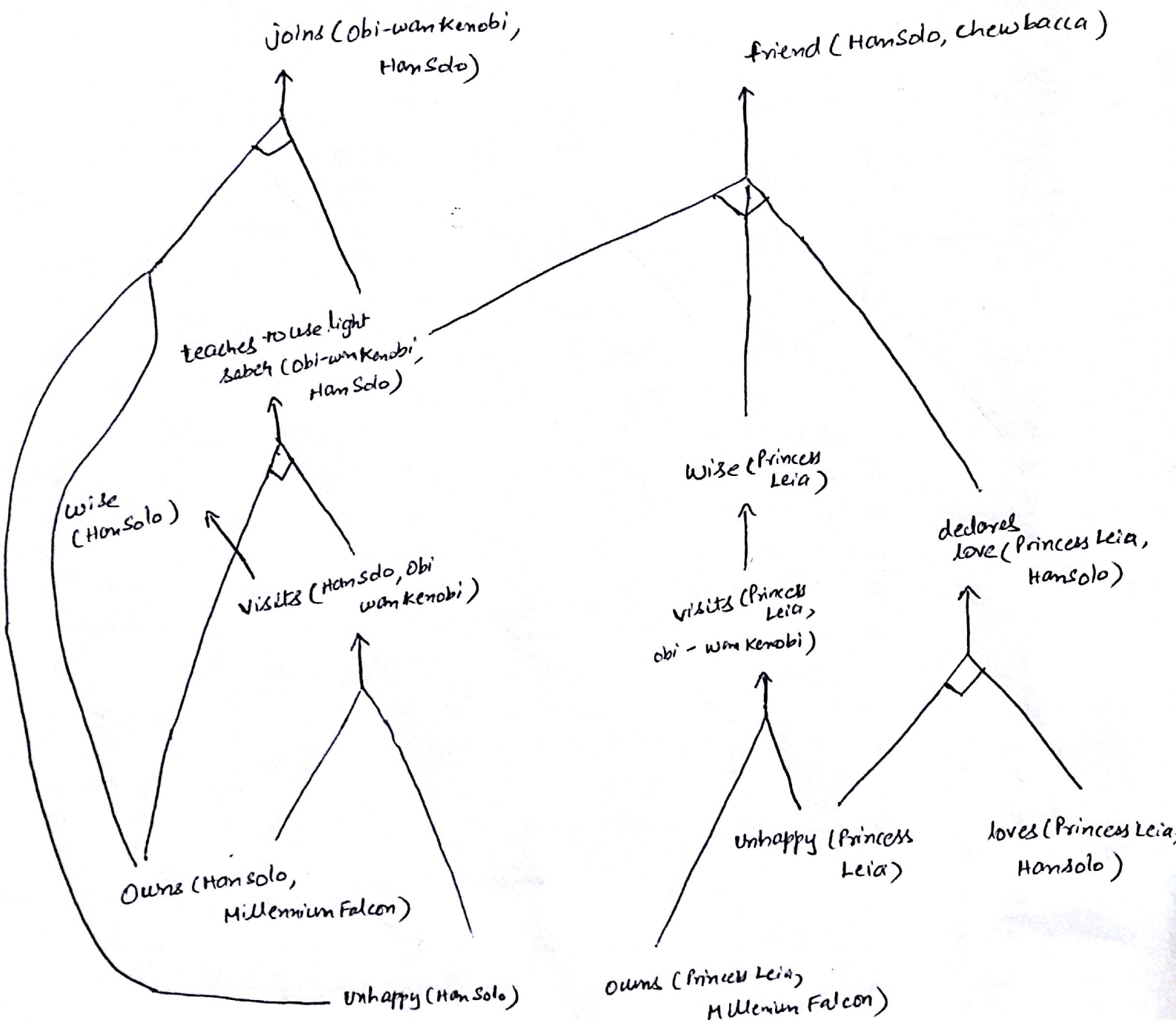


$\Rightarrow (y = \text{Princess Leia})$

unhappy (HanSolo) & Owns (y, Millennium Falcon) are not given and cannot be evaluated.

3(ii) Forward Chaining:

- unhappy (Han Solo) and owns (Princess Leia, Millennium Falcon) are not given and have only taken for completing the diagram.

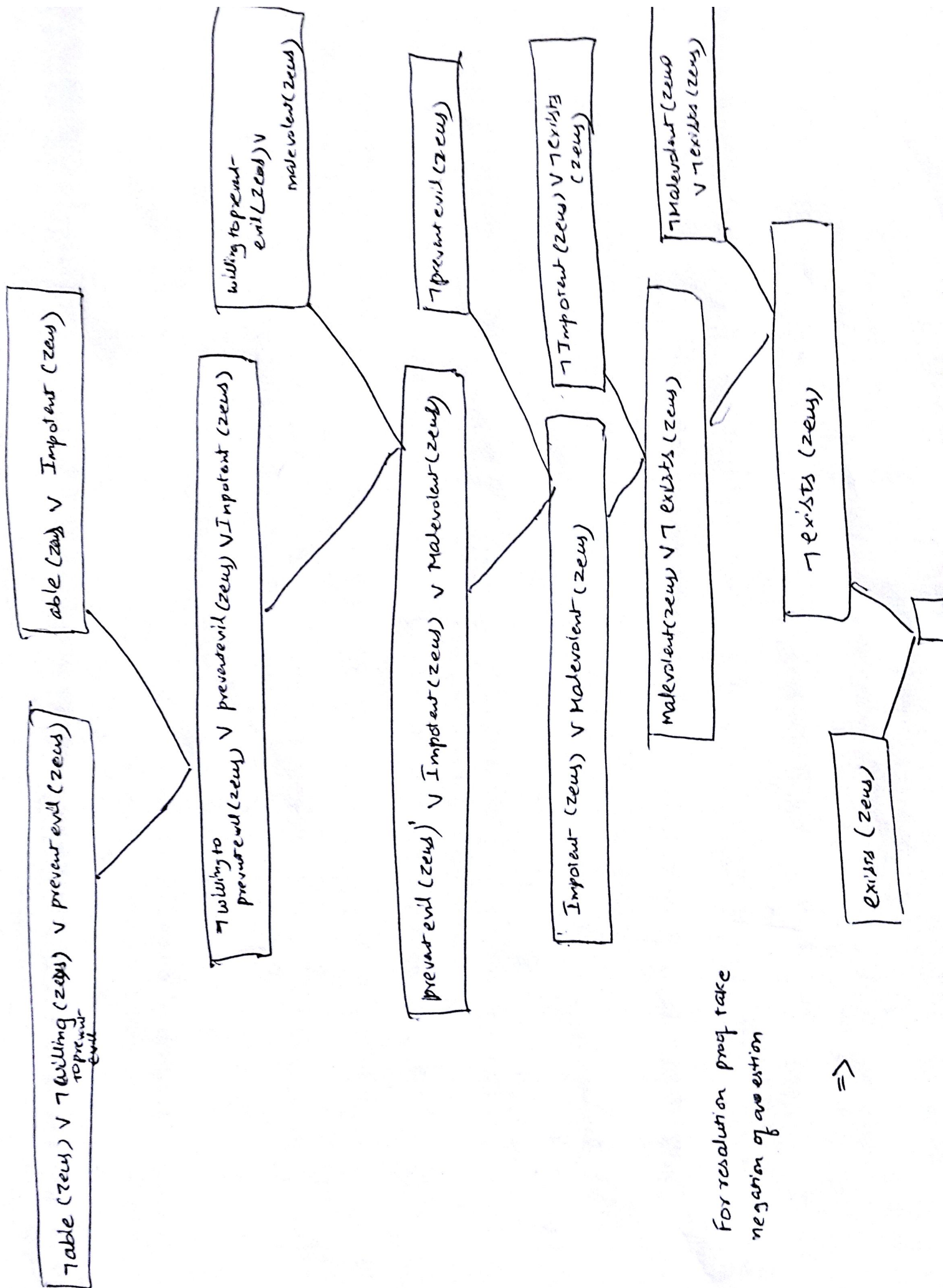


4:

- a) $\text{able}(\text{zeus}) \wedge \text{willing to prevent evil}(\text{zeus}) \Rightarrow \text{prevent evil}(\text{zeus})$
- b) $\neg \text{able}(\text{zeus}) \Rightarrow \text{Impotent}(\text{zeus})$
- c) $\neg \text{willing to prevent evil}(\text{zeus}) \Rightarrow \text{Malevolent}(\text{zeus})$
- d) $\neg \text{prevent evil}(\text{zeus})$
- e) $\text{exists}(\text{zeus}) \Rightarrow \neg \text{Impotent}(\text{zeus}) \wedge \neg \text{Malevolent}(\text{zeus})$

1) CNF

- a) $\neg (\text{able}(\text{zeus}) \wedge \text{willing to prevent evil}(\text{zeus})) \vee \neg \text{prevent evil}(\text{zeus})$
 $\Rightarrow \neg \text{able}(\text{zeus}) \vee \neg \text{willing to prevent evil}(\text{zeus}) \vee \neg \text{prevent evil}(\text{zeus})$
- b) $\neg (\neg \text{able}(\text{zeus})) \vee \text{Impotent}(\text{zeus})$
 $\Rightarrow \text{able}(\text{zeus}) \vee \text{Impotent}(\text{zeus})$
- c) $\neg (\neg \text{willing to prevent evil}(\text{zeus})) \vee \text{malevolent}(\text{zeus})$
 $\Rightarrow \text{willing to prevent evil}(\text{zeus}) \vee \text{malevolent}(\text{zeus})$
- d) $\neg \text{prevent evil}(\text{zeus})$
- e) $\neg \text{exists}(\text{zeus}) \vee (\neg \text{Impotent}(\text{zeus}) \wedge \neg \text{malevolent}(\text{zeus}))$
 $\Rightarrow (\neg \text{exists}(\text{zeus}) \vee \neg \text{Impotent}(\text{zeus})) \wedge$
 $(\neg \text{exists}(\text{zeus}) \vee \neg \text{malevolent}(\text{zeus}))$



For resolution proof take
negation of assertion

=>

5.

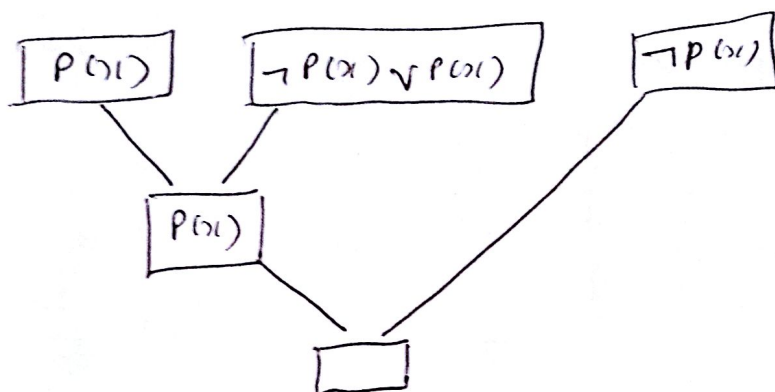
a) $\forall x (P(x) \Rightarrow P(x))$

CNF

$$\forall x (\neg P(x) \vee P(x))$$

$$\forall x \neg P(x) \vee \forall x P(x)$$

$$\neg P(x) \vee P(x) \quad (\text{Removing universal quantifier})$$



Empty set proves the statement

b) $(\neg \exists x P(x)) \Rightarrow (\forall x \neg P(x))$

Negation
↓
conclusion \Rightarrow

$$\neg (\forall x (\neg P(x)))$$

$$\exists x P(x)$$

$$P(f(x))$$

Knowledge Base :

$$(\neg \exists x P(x))$$

$$\forall x \neg P(x)$$

$$\neg P(x)$$

(removing universal quantifiers)

CNF

$$\neg \exists x P(x) \Rightarrow \forall x \neg P(x)$$

$$\neg (\neg \exists x P(x)) \vee \forall x \neg P(x)$$

$$\exists x \neg P(x) \vee \forall x \neg P(x)$$

$$\neg P(f(x)) \vee \neg P(x)$$

