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B.Sc. CS 2nd Sem

Fundamentals of Data Structures and Algorithms (USA23202J)

Lab Manual

Lab 1: Recursion

Title: Understanding Recursion through Factorial Calculation

Aim: To implement and understand the concept of recursion by calculating the factorial of a given number.

Procedure:

- 1. Define a recursive function, factorial (n), that takes an integer n as input.
- 2. Establish the base case: if n is 0 or 1, the factorial is 1.
- 3. Define the recursive step: if n is greater than 1, the factorial is n multiplied by the factorial of n-1.
- 4. Call the function with a sample input and print the result.

```
# factorial.py
def factorial(n):
   Calculates the factorial of a non-negative integer using recursion.
       n: A non-negative integer.
   Returns:
      The factorial of n.
    # Base case: Factorial of 0 or 1 is 1
    if n == 0 or n == 1:
       return 1
    # Recursive step: n * factorial(n-1)
       return n * factorial(n - 1)
# Main part of the program to test the function
if name == " main ":
    # Get input from the user
       num = int(input("Enter a non-negative integer to calculate its
factorial: "))
       if num < 0:
```

```
print("Factorial is not defined for negative numbers.")
else:
    result = factorial(num)
    print(f"The factorial of {num} is {result}")
except ValueError:
    print("Invalid input. Please enter an integer.")
```

Input:

Enter a non-negative integer to calculate its factorial: 5

Expected Output:

The factorial of 5 is 120

Lab 2: Arrays

Title: Array Operations: Sum and Average of Elements

Aim: To implement basic array operations, specifically calculating the sum and average of elements in an array.

Procedure:

- 1. Initialize an array (list in Python) with a set of integer values.
- 2. Iterate through the array to sum all its elements.
- 3. Calculate the average by dividing the sum by the number of elements.
- 4. Print the sum and average.

```
# array operations.py
def calculate_array_stats(arr):
   Calculates the sum and average of elements in a given array.
   Args:
       arr: A list of numbers.
       A tuple containing the sum and average (sum, average).
       Returns (0, 0) if the array is empty.
   if not arr:
       return 0, 0 # Handle empty array case
   total sum = 0
    for element in arr:
       total sum += element
   num elements = len(arr)
   average = total sum / num elements
   return total sum, average
# Main part of the program to test the function
if __name__ == "__main_ ":
    # Example array
   my array = [10, 20, 30, 40, 50]
   print(f"Original Array: {my array}")
    # Calculate sum and average
    array sum, array average = calculate array stats(my array)
    # Print the results
   print(f"Sum of array elements: {array sum}")
   print(f"Average of array elements: {array average}")
    # Another example with different values
    another array = [2, 4, 6, 8]
   print(f"\nOriginal Array: {another array}")
   array sum_2, array_average_2 = calculate_array_stats(another_array)
   print(f"Sum of array elements: {array sum 2}")
   print(f"Average of array elements: {array average 2}")
```

Input: (No direct user input for this program, array is hardcoded)

Expected Output:

Original Array: [10, 20, 30, 40, 50] Sum of array elements: 150 Average of array elements: 30.0

Original Array: [2, 4, 6, 8] Sum of array elements: 20 Average of array elements: 5.0

Lab 3: Implementation of LinkedList

Title: Singly Linked List Implementation: Insertion and Traversal

Aim: To implement a singly linked list with functionalities for inserting new nodes at the end and traversing the list to print its elements.

Procedure:

- 1. Define a Node class with data and next (pointer to the next node) attributes.
- 2. Define a LinkedList class with a head attribute, initially None.
- 3. Implement an append method to add a new node to the end of the list.
- 4. Implement a display method to traverse the list from the head and print the data of each node.

```
# linked list.py
class Node:
    ** ** **
   Represents a node in a singly linked list.
   Each node stores data and a reference to the next node.
    def init (self, data):
        self.data = data
        self.next = None # Pointer to the next node
class LinkedList:
    Represents a singly linked list.
   Manages the head of the list and provides methods for operations.
        init (self):
    def
        self.head = None # Initially, the list is empty
    def append(self, data):
        Adds a new node with the given data to the end of the linked list.
        Args:
           data: The data to be stored in the new node.
        new node = Node (data)
        if self.head is None:
            self.head = new node # If list is empty, new node becomes the
head
           return
        last node = self.head
        while last_node.next: # Traverse to the last node
            last node = last node.next
        last node.next = new node # Link the new node to the end
    def display(self):
        Traverses the linked list from the head and prints the data of each
node.
        current node = self.head
        if current node is None:
            print("Linked List is empty.")
```

```
return
       elements = []
       while current node:
            elements.append(str(current_node.data))
            current_node = current_node.next
       print("Linked List: " + " -> ".join(elements))
# Main part of the program to test the linked list implementation
if __name__ == "__main__":
   my list = LinkedList()
   print("Appending elements to the linked list:")
   my list.append(10)
   my list.append(20)
   my list.append(30)
   my list.display() # Expected: 10 -> 20 -> 30
   my list.append(40)
   print("\nAfter appending 40:")
   my list.display() # Expected: 10 -> 20 -> 30 -> 40
   empty list = LinkedList()
   print("\nTrying to display an empty list:")
    empty list.display() # Expected: Linked List is empty.
```

Input: (No direct user input for this program, operations are hardcoded)

```
Appending elements to the linked list:
Linked List: 10 -> 20 -> 30

After appending 40:
Linked List: 10 -> 20 -> 30 -> 40

Trying to display an empty list:
Linked List is empty.
```

Lab 4: Implementation of Stack and its Applications

Title: Stack Implementation using List and Parenthesis Checker Application

Aim: To implement a stack data structure using a Python list and demonstrate its application in checking for balanced parentheses in an expression.

Procedure:

- 1. Implement a Stack class with push, pop, peek, is_empty, and size methods. Use a Python list internally.
- 2. Implement a function is balanced parentheses that takes a string expression as input.
- 3. Use the stack to process the expression:
 - o If an opening parenthesis (, {, [is encountered, push it onto the stack.
 - o If a closing parenthesis), },] is encountered, pop from the stack. If the stack is empty or the popped element doesn't match the opening counterpart, the parentheses are unbalanced.
- 4. After processing the entire expression, if the stack is empty, the parentheses are balanced; otherwise, they are not.

```
# stack parenthesis checker.py
class Stack:
   A simple Stack implementation using a Python list.
        __init__(self):
        self.items = []
    def push(self, item):
        """Adds an item to the top of the stack."""
        self.items.append(item)
    def pop(self):
        """Removes and returns the item from the top of the stack.
       Raises IndexError if the stack is empty.
        if self.is empty():
            raise IndexError("pop from empty stack")
        return self.items.pop()
    def peek(self):
        """Returns the item at the top of the stack without removing it.
       Raises IndexError if the stack is empty.
        if self.is empty():
            raise IndexError("peek from empty stack")
        return self.items[-1]
    def is empty(self):
        """Checks if the stack is empty."""
        return len(self.items) == 0
    def size(self):
        """Returns the number of items in the stack."""
        return len(self.items)
def is balanced parentheses (expression):
    11 11 TT
```

```
Checks if the parentheses in an expression are balanced.
    Args:
       expression: A string containing parentheses.
    Returns:
       True if parentheses are balanced, False otherwise.
    s = Stack()
   mapping = {")": "(", "}": "{", "]": "["}
    open brackets = set(mapping.values())
    close brackets = set(mapping.keys())
    for char in expression:
        if char in open brackets:
            s.push(char)
        elif char in close brackets:
            if s.is empty():
                return False # Closing bracket without a corresponding
opening bracket
            top element = s.pop()
            if mapping[char] != top element:
                return False # Mismatched brackets
    return s.is empty() # True if stack is empty (all brackets matched)
# Main part of the program to test the stack and parenthesis checker
if __name__ == "__main__":
    # Test Stack operations
   my stack = Stack()
   print("Stack operations:")
   print(f"Is stack empty? {my_stack.is_empty()}") # True
   my stack.push(10)
   my stack.push(20)
   print(f"Stack size: {my stack.size()}") # 2
   print(f"Top element: {my_stack.peek()}") # 20
   print(f"Popped element: {my_stack.pop()}") # 20
   print(f"Stack size after pop: {my_stack.size()}") # 1
   print(f"Is stack empty? {my_stack.is_empty()}") # False
   my_stack.pop()
   print(f"Is stack empty after all pops? {my stack.is empty()}") # True
    # Test Parenthesis Checker
    print("\nParenthesis Checker:")
    expressions = [
                        # Balanced
        "([{}])",
        "({[()]})",
                      # Balanced
        "(()",
                        # Unbalanced (missing closing)
        "())",
                        # Unbalanced (extra closing)
        "[{}]",
                        # Balanced
        "{[}]",
                       # Unbalanced (mismatched)
                        # Balanced (empty string)
        "abc(def[ghi])", # Balanced (non-bracket chars ignored)
    1
    for expr in expressions:
        status = "Balanced" if is balanced parentheses(expr) else
"Unbalanced"
        print(f"Expression: '{expr}' is {status}")
```

Input: (No direct user input for this program, expressions are hardcoded)

Expected Output:

Stack operations:

Is stack empty? True

Stack size: 2 Top element: 20 Popped element: 20 Stack size after pop: 1 Is stack empty? False

Is stack empty after all pops? True

Parenthesis Checker:

Expression: '([{}])' is Balanced Expression: '({[()]})' is Balanced Expression: '({[()]})' is Balanced
Expression: '(()' is Unbalanced
Expression: '())' is Unbalanced
Expression: '[{}]' is Balanced
Expression: '{[}]' is Unbalanced
Expression: 'is Balanced
Expression: 'abc(def[ghi])' is Balanced

Lab 5: Queue implementation using array and pointers

Title: Queue Implementation using List (Array) and Front/Rear Pointers

Aim: To implement a queue data structure using a Python list (simulating an array) and manage front and rear pointers for enqueue and dequeue operations.

Procedure:

- 1. Implement a Queue class.
- 2. Initialize an empty list (items) and front and rear pointers (integers, e.g., -1 or 0).
- 3. Implement enqueue (add to rear): Add an element to the end of the list and update rear.
- 4. Implement dequeue (remove from front): Remove an element from the beginning of the list and update front. Handle underflow (empty queue).
- 5. Implement is empty, size, and peek (view front element).

```
# queue array pointers.py
class Queue:
   A Queue implementation using a Python list, simulating array behavior
   with conceptual front and rear pointers.
   def init (self):
        self.items = [] # The underlying list to store queue elements
        # For a simple list-based queue, Python's list methods handle
        # the 'pointer' logic internally for efficiency.
        # Conceptually, front is always at index 0, rear is at len(items) -
1.
    def enqueue (self, item):
        Adds an item to the rear of the queue.
        self.items.append(item)
        print(f"Enqueued: {item}")
    def dequeue (self):
       Removes and returns the item from the front of the queue.
       Raises IndexError if the queue is empty.
        if self.is empty():
           raise IndexError("dequeue from empty queue")
        dequeued item = self.items.pop(0) # Remove from the beginning (front)
        print(f"Dequeued: {dequeued item}")
       return dequeued item
    def peek(self):
        Returns the item at the front of the queue without removing it.
        Raises IndexError if the queue is empty.
        if self.is empty():
           raise IndexError("peek from empty queue")
        return self.items[0]
    def is empty(self):
        Checks if the queue is empty.
```

```
return len(self.items) == 0
   def size(self):
       Returns the number of items in the queue.
       return len(self.items)
# Main part of the program to test the queue implementation
if name == " main ":
   \overline{my} queue = \overline{Queue} ()
   print("Queue operations:")
   print(f"Is queue empty? {my queue.is empty()}") # True
   print(f"Queue size: {my queue.size()}") # 0
   my queue.enqueue("Task A")
   my queue.enqueue("Task B")
   my queue.enqueue("Task C")
   print(f"Queue size: {my queue.size()}") # 3
   print(f"Front element (peek): {my_queue.peek()}") # Task A
   my queue.dequeue() # Task A
   print(f"Queue size after dequeue: {my queue.size()}") # 2
   print(f"Front element (peek): {my_queue.peek()}") # Task B
   my queue.enqueue("Task D")
   my queue.display() # Custom display function to show current queue state
   my_queue.dequeue() # Task B
   my queue.dequeue() # Task C
   my queue.dequeue() # Task D
   print(f"Is queue empty? {my_queue.is_empty()}") # True
   try:
       my_queue.dequeue() # This should raise an error
   except IndexError as e:
       print(f"Error: {e}") # Expected: Error: dequeue from empty queue
```

Input: (No direct user input for this program, operations are hardcoded)

```
Queue operations:
Is queue empty? True
Queue size: 0
Enqueued: Task A
Enqueued: Task B
Enqueued: Task C
Queue size: 3
Front element (peek): Task A
Dequeued: Task A
Queue size after dequeue: 2
Front element (peek): Task B
Enqueued: Task D
Queue: ['Task B', 'Task C', 'Task D']
Dequeued: Task B
Dequeued: Task C
Dequeued: Task D
Is queue empty? True
Error: dequeue from empty queue
```

Note: I've added a simple <code>display</code> method to the Queue class for better visualization in the output, although it wasn't explicitly requested in the procedure. This helps in understanding the queue's state.

Lab 6: Implementation of binary tree using Arrays

Title: Binary Tree Implementation using Arrays (List Representation)

Aim: To implement a binary tree using a Python list (array-based representation) and demonstrate insertion and traversal.

Procedure:

- 1. Represent the binary tree using a list where:
 - The root is at index 0.
 - o For a node at index i, its left child is at 2*i + 1.
 - o Its right child is at 2*i + 2.
- 2. Initialize the list with None or a placeholder for empty nodes.
- 3. Implement an insert function to add nodes at specific positions (or implicitly by their parent's index).
- 4. Implement a print tree function to display the array representation.

```
# binary tree array.py
class BinaryTreeArray:
    Implements a binary tree using an array (Python list) representation.
    - Root is at index 0.
    - Left child of node at index i is at 2*i + 1.
    - Right child of node at index i is at 2*i + 2.
       __init__(self, max_size):
    def
        Initializes the tree with a fixed maximum size.
        None indicates an empty node.
        self.tree = [None] * max size
        self.max size = max size
       self.size = 0 # Current number of nodes in the tree
    def insert(self, data, parent_index=0, position='root'):
        Inserts data into the tree.
       For simplicity, this example allows insertion at a specific index,
       or as a child of a given parent index (left/right).
       A more robust implementation would manage insertion based on tree
properties (e.g., BST).
        if self.size >= self.max size:
            print("Tree is full. Cannot insert more elements.")
            return False
        if position == 'root':
            if self.tree[0] is None:
               self.tree[0] = data
               self.size += 1
               print(f"Inserted {data} at root (index 0).")
               return True
            else:
                print(f"Root already exists. Cannot insert {data} as root.")
                return False
        elif position == 'left child':
            left child index = 2 * parent index + 1
```

```
if left child index < self.max size and
self.tree[left child index] is None:
                self.tree[left child index] = data
                self.size += 1
               print(f"Inserted {data} as left child of index {parent index}
(at index {left child index}).")
               return True
            else:
               print(f"Cannot insert {data} as left child of {parent index}.
Position occupied or out of bounds.")
                return False
        elif position == 'right_child':
            right child index = 2 * parent index + 2
            if right child index < self.max size and
self.tree[right_child_index] is None:
                self.tree[right child index] = data
                self.size += 1
                print(f"Inserted {data} as right child of index
{parent index} (at index {right child index}).")
                return True
            else:
                print(f"Cannot insert {data} as right child of
{parent index}. Position occupied or out of bounds.")
                return False
        else:
            print("Invalid position specified for insertion.")
            return False
    def print tree array(self):
        Prints the array representation of the binary tree.
        print("Binary Tree (Array Representation):")
       print(self.tree)
       print(f"Current number of nodes: {self.size}")
# Main part of the program to test the binary tree array implementation
if __name__ == "__main__":
    # Create a tree with a maximum size of 10 nodes
   tree = BinaryTreeArray(10)
    # Insert root
   tree.insert('A', position='root') # Root
    # Insert children of root (A at index 0)
    tree.insert('B', parent index=0, position='left child') # Left child of
A (index 1)
    tree.insert('C', parent index=0, position='right child') # Right child of
A (index 2)
    # Insert children of B (B at index 1)
    tree.insert('D', parent index=1, position='left child') # Left child of
B (index 3)
    tree.insert('E', parent index=1, position='right child') # Right child of
B (index 4)
    # Insert children of C (C at index 2)
    tree.insert('F', parent index=2, position='left child') # Left child of
C (index 5)
    tree.print tree array()
   print("\nAttempting to insert where position is occupied or out of
   tree.insert('X', parent index=0, position='left child') # Try to insert
over B
```

```
tree.insert('Y', parent_index=5, position='right_child') # Parent index
5, right child index 12 (out of bounds)
    tree.insert('Z', parent_index=4, position='left_child') # Left child of E
(index 9)
    tree.print_tree_array()
```

Input: (No direct user input for this program, operations are hardcoded)

```
Inserted A at root (index 0).
Inserted B as left child of index 0 (at index 1).
Inserted C as right child of index 0 (at index 2).
Inserted D as left child of index 1 (at index 3).
Inserted E as right child of index 1 (at index 4).
Inserted F as left child of index 2 (at index 5).
Binary Tree (Array Representation):
['A', 'B', 'C', 'D', 'E', 'F', None, None, None, None]
Current number of nodes: 6
Attempting to insert where position is occupied or out of bounds:
Cannot insert {\tt X} as left child of {\tt O}. Position occupied or out of bounds.
Cannot insert Y as right child of 5. Position occupied or out of bounds.
Inserted Z as left child of index 4 (at index 9).
Binary Tree (Array Representation):
['A', 'B', 'C', 'D', 'E', 'F', None, None, None, 'Z']
Current number of nodes: 7
```

Lab 7: Implement all the three type of Tree Traversals

Title: Binary Tree Traversals: Inorder, Preorder, and Postorder

Aim: To implement a binary tree using a node-based approach and demonstrate the three standard tree traversal algorithms: Inorder, Preorder, and Postorder.

Procedure:

- 1. Define a TreeNode class with data, left (left child), and right (right child) attributes.
- 2. Implement a BinaryTree class with a root attribute.
- 3. Implement methods for inserting nodes to build a sample tree (e.g., a simple insertion for demonstration, or a more complex BST insertion).
- 4. Implement three recursive functions for traversal:

```
o inorder_traversal(node): Left -> Root -> Right
o preorder_traversal(node): Root -> Left -> Right
o postorder_traversal(node): Left -> Right -> Root
```

```
# tree traversals.py
class TreeNode:
    11 11 11
    Represents a node in a binary tree.
    def __init__(self, data):
       \overline{\text{self.data}} = \text{data}
       self.left = None
        self.right = None
class BinaryTree:
    A simple Binary Tree class for demonstrating traversals.
    def init (self):
        self.root = None
    def insert(self, data):
        Inserts a new node into the binary tree.
        This is a simple insertion for demonstration purposes,
        not a balanced or BST-specific insertion.
        It adds nodes level by level (like a complete binary tree for
simplicity).
        For a more robust solution, a queue would be used for level-order
insertion.
        Here, we'll manually build a small tree for traversal examples.
        # For this lab, we'll manually build a sample tree for clarity
        # rather than a generic insert method that balances or orders.
    def build_sample_tree(self):
        Manually builds a sample binary tree for traversal demonstration.
        Tree structure:
             1
             / \
            2 3
           / \
```

```
,, ,, ,,
        self.root = TreeNode(1)
        self.root.left = TreeNode(2)
        self.root.right = TreeNode(3)
        self.root.left.left = TreeNode(4)
        self.root.left.right = TreeNode(5)
        print("Sample tree built.")
    def inorder traversal(self, node):
        Performs an Inorder traversal (Left -> Root -> Right).
        if node:
            self.inorder traversal(node.left)
            print(node.data, end=" ")
            self.inorder traversal(node.right)
    def preorder traversal(self, node):
        Performs a Preorder traversal (Root -> Left -> Right).
        if node:
            print(node.data, end=" ")
            self.preorder_traversal(node.left)
            self.preorder traversal(node.right)
    def postorder traversal(self, node):
        Performs a Postorder traversal (Left -> Right -> Root).
        11 11 11
        if node:
            self.postorder_traversal(node.left)
            self.postorder traversal(node.right)
            print(node.data, end=" ")
# Main part of the program to test tree traversals
if __name__ == "__main__":
    \overline{\text{tree}} = \overline{\text{BinaryTree}}()
   tree.build sample tree()
    print("\nInorder Traversal:")
    tree.inorder traversal(tree.root) # Expected: 4 2 5 1 3
   print()
   print("\nPreorder Traversal:")
    tree.preorder traversal(tree.root) # Expected: 1 2 4 5 3
   print()
   print("\nPostorder Traversal:")
    tree.postorder traversal(tree.root) # Expected: 4 5 2 3 1
   print()
```

Input: (No direct user input for this program, tree is hardcoded)

```
Sample tree built.

Inorder Traversal:
4 2 5 1 3

Preorder Traversal:
1 2 4 5 3
```

Postorder Traversal: 4 5 2 3 1

Lab 8: Implementation of BST HeapDataStructure

Title: Binary Search Tree (BST) Implementation: Insertion and Search

Aim: To implement a Binary Search Tree (BST) with functionalities for inserting new nodes while maintaining BST properties and searching for a specific value.

Procedure:

- 1. Define a Node class for the BST, similar to a binary tree node, but with the understanding that left children are smaller and right children are larger.
- 2. Implement a BST class with a root attribute.
- 3. Implement an insert method:
 - o If the tree is empty, the new node becomes the root.
 - o Otherwise, traverse the tree to find the correct position: go left if the new data is smaller, go right if it's larger.
- 4. Implement a search method:
 - Start from the root.
 - o If the current node's data matches the search value, return the node.
 - o If the search value is smaller, go left; otherwise, go right.
 - o If a None child is reached, the value is not in the tree.
- 5. Implement an inorder traversal for verification.

```
# bst implementation.py
class Node:
   Represents a node in a Binary Search Tree (BST).
    def init (self, key):
       self.key = key
        self.left = None
        self.right = None
class BST:
    11 11 11
    Implements a Binary Search Tree.
    def init (self):
        self.root = None
    def insert(self, key):
        Inserts a new key into the BST.
        new node = Node(key)
        if self.root is None:
            self.root = new node
            print(f"Inserted {key} as root.")
            return
        current = self.root
        while True:
            if key < current.key:
                if current.left is None:
                    current.left = new node
                    print(f"Inserted {key} as left child of {current.key}.")
                    return
```

```
current = current.left
            elif key > current.key:
                if current.right is None:
                    current.right = new node
                    print(f"Inserted {key} as right child of {current.key}.")
                current = current.right
            else:
                print(f"Key {key} already exists in the BST. Not inserted.")
                return
    def search(self, key):
        Searches for a key in the BST.
        Returns:
          The Node object if found, None otherwise.
        current = self.root
        while current:
            if key == current.key:
                return current
            elif key < current.key:</pre>
               current = current.left
            else: # key > current.key
               current = current.right
        return None # Key not found
    def inorder traversal(self, node):
        Performs an Inorder traversal of the BST (Left -> Root -> Right).
        This prints elements in sorted order.
        11 11 11
        if node:
            self.inorder_traversal(node.left)
            print(node.key, end=" ")
            self.inorder_traversal(node.right)
# Main part of the program to test the BST implementation
if __name__ == "__main__":
   bst = BST()
    # Insert elements
   bst.insert(50)
   bst.insert(30)
   bst.insert(70)
   bst.insert(20)
   bst.insert(40)
   bst.insert(60)
   bst.insert(80)
   bst.insert(30) # Attempt to insert duplicate
   print("\nInorder Traversal of BST (should be sorted):")
   bst.inorder traversal(bst.root)
   print()
    # Search for elements
   print("\nSearching for elements:")
    search keys = [40, 90, 50, 10]
    for key in search keys:
       node = bst.search(kev)
        if node:
           print(f"Key {key} found in BST.")
        else:
            print(f"Key {key} not found in BST.")
```

Input: (No direct user input for this program, operations are hardcoded)

```
Inserted 50 as root.
Inserted 30 as left child of 50.
Inserted 70 as right child of 50.
Inserted 20 as left child of 30.
Inserted 40 as right child of 30.
Inserted 60 as left child of 70.
Inserted 80 as right child of 70.
Key 30 already exists in the BST. Not inserted.

Inorder Traversal of BST (should be sorted):
20 30 40 50 60 70 80

Searching for elements:
Key 40 found in BST.
Key 90 not found in BST.
Key 50 found in BST.
Key 10 not found in BST.
```

Lab 9: Implementation of Min and Max Heap

Title: Min-Heap and Max-Heap Implementation using List

Aim: To implement both Min-Heap and Max-Heap data structures using a Python list (array-based representation) and demonstrate their basic operations (insertion and extraction).

Procedure:

- 1. **Heap Property:** Understand that in a Max-Heap, parent nodes are always greater than or equal to their children, and in a Min-Heap, parent nodes are always less than or equal to their children.
- 2. Array Representation: For a node at index i:

```
    Left child: 2*i + 1
    Right child: 2*i + 2
    Parent: (i - 1) // 2
```

- 3. Implement MinHeap class:
 - o insert (key): Add to the end, then heapify_up (bubble up) to maintain minheap property.
 - o extract_min(): Remove root, replace with last element, then heapify_down (bubble down) to maintain min-heap property.
 - o _heapify_up(index): Helper to move element up.
 - o heapify down(index): Helper to move element down.
- 4. Implement MaxHeap class: (Similar to MinHeap, but comparison logic is reversed for heapify up and heapify down).

```
# min_max_heap.py
class MinHeap:
    11 11 11
    Implements a Min-Heap data structure.
    def __init__(self):
        self.heap = []
    def _parent(self, i):
        return (i - 1) // 2
    def left child(self, i):
        return 2 * i + 1
    def right child(self, i):
        return 2 * i + 2
    def swap(self, i, j):
        self.heap[i], self.heap[j] = self.heap[j], self.heap[i]
    def _heapify_up(self, i):
        Moves the element at index i up the heap to maintain the min-heap
property.
        while i > 0 and self.heap[self. parent(i)] > self.heap[i]:
            self. swap(i, self. parent(i))
            i = self. parent(i)
    def heapify down(self, i):
```

```
Moves the element at index i down the heap to maintain the min-heap
property.
        min index = i
        left = self. left child(i)
        right = self. right child(i)
        n = len(self.heap)
        if left < n and self.heap[left] < self.heap[min index]:</pre>
            min index = left
        if right < n and self.heap[right] < self.heap[min index]:</pre>
            min index = right
        if min index != i:
            self. swap(i, min index)
            self. heapify down (min index)
    def insert(self, key):
        Inserts a new key into the min-heap.
        self.heap.append(key)
        self. heapify up(len(self.heap) - 1)
        print(f"Inserted {key} into Min-Heap. Current Heap: {self.heap}")
    def extract min(self):
        Removes and returns the minimum element from the min-heap (root).
        if not self.heap:
            raise IndexError("Heap is empty")
        if len(self.heap) == 1:
            min val = self.heap.pop()
            print(f"Extracted Min: {min_val}. Current Heap: {self.heap}")
            return min_val
        min val = self.heap[0]
        self.heap[0] = self.heap.pop() # Move last element to root
        self. heapify down(0)
        print(f"Extracted Min: {min val}. Current Heap: {self.heap}")
        return min val
    def peek min(self):
        """Returns the minimum element without removing it."""
        if not self.heap:
            raise IndexError("Heap is empty")
        return self.heap[0]
    def is empty(self):
        return len(self.heap) == 0
    def size(self):
       return len(self.heap)
class MaxHeap:
    Implements a Max-Heap data structure.
    def init (self):
        self.heap = []
    def parent(self, i):
        return (i - 1) // 2
```

```
def left child(self, i):
        return 2 * i + 1
    def right child(self, i):
        return 2 * i + 2
    def swap(self, i, j):
        self.heap[i], self.heap[j] = self.heap[j], self.heap[i]
    def _heapify_up(self, i):
        Moves the element at index i up the heap to maintain the max-heap
property.
        while i > 0 and self.heap[self. parent(i)] < self.heap[i]:</pre>
            self. swap(i, self. parent(i))
            i = self. parent(i)
    def _heapify_down(self, i):
        Moves the element at index i down the heap to maintain the max-heap
property.
        \max index = i
        left = self._left_child(i)
        right = self. right child(i)
        n = len(self.heap)
        if left < n and self.heap[left] > self.heap[max index]:
            max index = left
        if right < n and self.heap[right] > self.heap[max index]:
            max index = right
        if max_index != i:
            self._swap(i, max_index)
            self._heapify_down(max_index)
    def insert(self, key):
        11 11 11
        Inserts a new key into the max-heap.
        self.heap.append(key)
        self._heapify_up(len(self.heap) - 1)
        print(f"Inserted {key} into Max-Heap. Current Heap: {self.heap}")
    def extract max(self):
        Removes and returns the maximum element from the max-heap (root).
        if not self.heap:
            raise IndexError("Heap is empty")
        if len(self.heap) == 1:
            max val = self.heap.pop()
            print(f"Extracted Max: {max val}. Current Heap: {self.heap}")
            return max val
        max val = self.heap[0]
        self.heap[0] = self.heap.pop() # Move last element to root
        self. heapify down(0)
        print(f"Extracted Max: {max val}. Current Heap: {self.heap}")
        return max val
    def peek max(self):
        """Returns the maximum element without removing it."""
        if not self.heap:
```

```
raise IndexError("Heap is empty")
        return self.heap[0]
   def is empty(self):
       return len(self.heap) == 0
   def size(self):
       return len(self.heap)
# Main part of the program to test Min-Heap and Max-Heap implementations
if name == " main ":
   print("--- Min-Heap Operations ---")
   min heap = MinHeap()
   min heap.insert(3)
   min heap.insert(1)
   min heap.insert(4)
   min heap.insert(2)
   min heap.insert(5)
   print(f"Min-Heap size: {min heap.size()}")
   print(f"Min element (peek): {min heap.peek min()}")
   min heap.extract min() # Should extract 1
   min heap.extract min() # Should extract 2
   print(f"Min-Heap size: {min heap.size()}")
   print(f"Is Min-Heap empty? {min heap.is empty()}")
   min heap.extract min() # Should extract 3
   min_heap.extract_min() # Should extract 4
   min heap.extract min() # Should extract 5
   print(f"Is Min-Heap empty? {min heap.is empty()}")
   try:
       min heap.extract min()
   except IndexError as e:
       print(f"Error: {e}")
   print("\n--- Max-Heap Operations ---")
   max heap = MaxHeap()
   max heap.insert(3)
   max heap.insert(1)
   max heap.insert(4)
   max heap.insert(2)
   max heap.insert(5)
   print(f"Max-Heap size: {max heap.size()}")
   print(f"Max element (peek): {max heap.peek max()}")
   max heap.extract max() # Should extract 5
   max heap.extract max() # Should extract 4
   print(f"Max-Heap size: {max heap.size()}")
   print(f"Is Max-Heap empty? {max heap.is empty()}")
   max heap.extract max() # Should extract 3
   max heap.extract max() # Should extract 2
   max heap.extract max() # Should extract 1
   print(f"Is Max-Heap empty? {max heap.is empty()}")
   try:
       max heap.extract max()
   except IndexError as e:
       print(f"Error: {e}")
```

Input: (No direct user input for this program, operations are hardcoded)

```
--- Min-Heap Operations ---
```

```
Inserted 3 into Min-Heap. Current Heap: [3]
Inserted 1 into Min-Heap. Current Heap: [1, 3]
Inserted 4 into Min-Heap. Current Heap: [1, 3, 4]
Inserted 2 into Min-Heap. Current Heap: [1, 2, 4, 3]
Inserted 5 into Min-Heap. Current Heap: [1, 2, 4, 3, 5]
Min-Heap size: 5
Min element (peek): 1
Extracted Min: 1. Current Heap: [2, 3, 4, 5]
Extracted Min: 2. Current Heap: [3, 5, 4]
Min-Heap size: 3
Is Min-Heap empty? False
Extracted Min: 3. Current Heap: [4, 5]
Extracted Min: 4. Current Heap: [5]
Extracted Min: 5. Current Heap: []
Is Min-Heap empty? True
Error: Heap is empty
--- Max-Heap Operations ---
Inserted 3 into Max-Heap. Current Heap: [3]
Inserted 1 into Max-Heap. Current Heap: [3, 1]
Inserted 4 into Max-Heap. Current Heap: [4, 1, 3]
Inserted 2 into Max-Heap. Current Heap: [4, 2, 3, 1]
Inserted 5 into Max-Heap. Current Heap: [5, 4, 3, 1, 2]
Max-Heap size: 5
Max element (peek): 5
Extracted Max: 5. Current Heap: [4, 2, 3, 1]
Extracted Max: 4. Current Heap: [3, 2, 1]
Max-Heap size: 3
Is Max-Heap empty? False
Extracted Max: 3. Current Heap: [2, 1]
Extracted Max: 2. Current Heap: [1]
Extracted Max: 1. Current Heap: []
Is Max-Heap empty? True
Error: Heap is empty
```

Lab 10: Implementation of Bubble and Insertion Sort

Title: Sorting Algorithms: Bubble Sort and Insertion Sort

Aim: To implement and compare the Bubble Sort and Insertion Sort algorithms for arranging elements in ascending order.

Procedure:

1. Bubble Sort:

- Iterate through the array from the first element to the second to last.
- o In each pass, compare adjacent elements and swap them if they are in the wrong order (larger element bubbles up).
- o Repeat passes until no swaps are needed in a full pass.

2. Insertion Sort:

- o Start from the second element.
- o For each element, compare it with elements to its left.
- Shift elements greater than the current element one position to the right to make space.
- o Insert the current element into its correct sorted position.
- 3. Test both algorithms with sample arrays and print the sorted results.

```
# bubble insertion sort.py
def bubble sort(arr):
   Sorts an array using the Bubble Sort algorithm.
   Time Complexity: O(n^2) in worst and average case.
   Space Complexity: O(1).
   n = len(arr)
    # Traverse through all array elements
    for i in range (n - 1):
        # Last i elements are already in place
        swapped = False
        for j in range(n - 1 - i):
            # Traverse the array from 0 to n-i-1
            # Swap if the element found is greater than the next element
            if arr[j] > arr[j + 1]:
                arr[j], arr[j + 1] = arr[j + 1], arr[j]
                swapped = True
        # If no two elements were swapped by inner loop, then break
        if not swapped:
           break
    return arr
def insertion sort(arr):
    Sorts an array using the Insertion Sort algorithm.
    Time Complexity: O(n^2) in worst and average case, O(n) in best case.
    Space Complexity: O(1).
    # Traverse through 1 to len(arr)
    for i in range(1, len(arr)):
        key = arr[i]
        \# Move elements of arr[0..i-1], that are greater than key,
        # to one position ahead of their current position
        j = i - 1
        while j \ge 0 and key < arr[j]:
```

```
arr[j + 1] = arr[j]
            j -= 1
        arr[j + 1] = key
    return arr
# Main part of the program to test sorting algorithms
if name == " main ":
    # Test Bubble Sort
   arr bubble = [64, 34, 25, 12, 22, 11, 90]
   print(f"Original array for Bubble Sort: {arr bubble}")
   bubble sort(arr bubble)
   print(f"Sorted array (Bubble Sort): {arr bubble}") # Expected: [11, 12,
22, 25, 34, 64, 90]
   arr bubble 2 = [5, 1, 4, 2, 8]
    print(f"Original array for Bubble Sort: {arr bubble 2}")
   bubble sort(arr bubble 2)
   print(f"Sorted array (Bubble Sort): {arr_bubble 2}") # Expected: [1, 2,
4, 5, 8]
   print("-" * 30)
    # Test Insertion Sort
   arr insertion = [12, 11, 13, 5, 6]
    print(f"Original array for Insertion Sort: {arr insertion}")
   insertion sort(arr insertion)
   print(f"Sorted array (Insertion Sort): {arr insertion}") # Expected: [5,
6, 11, 12, 13]
   arr insertion_2 = [7, 8, 9, 1, 2, 3]
    print(f"Original array for Insertion Sort: {arr insertion 2}")
   insertion_sort(arr_insertion_2)
   print(f"Sorted array (Insertion Sort): {arr insertion 2}") # Expected:
[1, 2, 3, 7, 8, 9]
```

Input: (No direct user input for this program, arrays are hardcoded)

Lab 11: Implementation of Quicksort and Mergesort

Title: Advanced Sorting Algorithms: Quicksort and Mergesort

Aim: To implement and understand the principles of Quicksort and Mergesort, two efficient comparison-based sorting algorithms.

Procedure:

1. Ouicksort:

- Choose a pivot element from the array.
- o Partition the array into two sub-arrays: elements smaller than the pivot and elements greater than the pivot.
- Recursively apply Quicksort to the two sub-arrays.

2. Mergesort:

- o Divide the unsorted list into n sublists, each containing one element (a list of one element is considered sorted).
- Repeatedly merge sublists to produce new sorted sublists until there is only one sublist remaining. This will be the sorted list.
- 3. Test both algorithms with sample arrays and print the sorted results.

```
# quick merge sort.py
def quick sort(arr):
    Sorts an array using the Quicksort algorithm.
    Time Complexity: O(n \log n) on average, O(n^2) in worst case.
    Space Complexity: O(log n) for recursion stack on average.
    if len(arr) <= 1:
       return arr
    else:
        pivot = arr[len(arr) // 2] # Choose middle element as pivot
        left = [x for x in arr if x < pivot]</pre>
       middle = [x for x in arr if x == pivot]
       right = [x \text{ for } x \text{ in arr if } x > pivot]
        return quick sort(left) + middle + quick sort(right)
def merge sort(arr):
    Sorts an array using the Mergesort algorithm.
    Time Complexity: O(n log n) in all cases.
    Space Complexity: O(n) for temporary arrays.
    if len(arr) <= 1:
        return arr
   mid = len(arr) // 2
    left half = arr[:mid]
    right half = arr[mid:]
    left half = merge sort(left half)
    right half = merge sort(right half)
    return merge(left half, right half)
def merge(left, right):
    Merges two sorted lists into a single sorted list.
```

```
Helper function for merge sort.
    result = []
    i = j = 0
    while i < len(left) and j < len(right):
        if left[i] < right[j]:</pre>
            result.append(left[i])
            i += 1
        else:
            result.append(right[j])
            j += 1
    # Append remaining elements
    result.extend(left[i:])
    result.extend(right[j:])
    return result
# Main part of the program to test sorting algorithms
if __name__ == "__main__":
    # Test Quicksort
   arr quick = [10, 7, 8, 9, 1, 5]
    print(f"Original array for Quicksort: {arr quick}")
    sorted quick = quick sort(arr quick)
   print(f"Sorted array (Quicksort): {sorted quick}") # Expected: [1, 5, 7,
8, 9, 10]
    arr quick 2 = [3, 1, 4, 1, 5, 9, 2, 6]
    print(f"Original array for Quicksort: {arr quick 2}")
    sorted quick 2 = quick sort(arr quick 2)
   print(f"Sorted array (Quicksort): {sorted quick 2}") # Expected: [1, 1,
2, 3, 4, 5, 6, 9]
   print("-" * 30)
    # Test Mergesort
    arr_merge = [38, 27, 43, 3, 9, 82, 10]
    print(f"Original array for Mergesort: {arr_merge}")
    sorted_merge = merge_sort(arr_merge)
    print(f"Sorted array (Mergesort): {sorted merge}") # Expected: [3, 9, 10,
27, 38, 43, 82]
    arr merge 2 = [6, 5, 12, 10, 9, 1]
    print(f"Original array for Mergesort: {arr merge 2}")
    sorted merge 2 = merge sort(arr merge 2)
   print(f"Sorted array (Mergesort): {sorted merge 2}") # Expected: [1, 5,
6, 9, 10, 12]
```

Input: (No direct user input for this program, arrays are hardcoded)

Lab 12: Linear search and Binary search

Title: Searching Algorithms: Linear Search and Binary Search

Aim: To implement and compare Linear Search and Binary Search algorithms for finding a specific element in a list.

Procedure:

1. Linear Search:

- o Iterate through each element of the list from the beginning.
- o Compare each element with the target value.
- o If a match is found, return its index. If the end of the list is reached without a match, return an indicator (e.g., -1).

2. Binary Search:

- o **Pre-requisite:** The list must be sorted.
- o Find the middle element of the list.
- o If the middle element is the target, return its index.
- o If the target is smaller, search in the left half.
- o If the target is larger, search in the right half.
- o Repeat until the element is found or the sub-array becomes empty.
- 3. Test both algorithms with sample lists and print the results.

```
# linear binary search.py
def linear search(arr, target):
    11 11 11
    Performs a linear search to find the target element in the array.
   Time Complexity: O(n) in worst and average case.
    Space Complexity: O(1).
    for i in range(len(arr)):
        if arr[i] == target:
           return i # Return the index if found
    return -1 # Return -1 if not found
def binary search(arr, target):
    Performs a binary search to find the target element in a sorted array.
   Time Complexity: O(log n).
    Space Complexity: O(1) (iterative) or O(log n) (recursive).
    left, right = 0, len(arr) - 1
    while left <= right:
       mid = (left + right) // 2
        if arr[mid] == target:
            return mid # Return the index if found
        elif arr[mid] < target:</pre>
            left = mid + 1 # Search in the right half
        else: # arr[mid] > target
           right = mid - 1 # Search in the left half
    return -1 # Return -1 if not found
# Main part of the program to test search algorithms
if __name_ == " main ":
    # Test Linear Search
   my_list_linear = [5, 1, 9, 2, 7, 3, 8, 4, 6]
```

```
target1 = 7
target2 = 10
print(f"List for Linear Search: {my list linear}")
index1 = linear search(my list linear, target1)
if index1 != -1:
   print(f"Linear Search: {target1} found at index {index1}")
else:
    print(f"Linear Search: {target1} not found")
index2 = linear search(my list linear, target2)
if index2 !=-1:
    print(f"Linear Search: {target2} found at index {index2}")
else:
   print(f"Linear Search: {target2} not found")
print("-" * 30)
# Test Binary Search (requires sorted list)
my list binary = [1, 2, 3, 4, 5, 6, 7, 8, 9, 10]
target3 = 4
target4 = 11
print(f"List for Binary Search (must be sorted): {my list binary}")
index3 = binary search(my list binary, target3)
if index3 != -1:
    print(f"Binary Search: {target3} found at index {index3}")
else:
    print(f"Binary Search: {target3} not found")
index4 = binary_search(my_list_binary, target4)
if index4 !=-1:
    print(f"Binary Search: {target4} found at index {index4}")
else:
   print(f"Binary Search: {target4} not found")
```

Input: (No direct user input for this program, lists and targets are hardcoded)

Lab 13: Implementation of Graph using Array

Title: Graph Representation using Adjacency Matrix (Array)

Aim: To implement a graph data structure using an adjacency matrix (2D array/list of lists) to represent connections between vertices.

Procedure:

- 1. Represent the graph as a 2D list (matrix) where matrix[i][j] = 1 if there is an edge from vertex i to vertex j, and 0 otherwise.
- 2. Initialize the matrix with all zeros.
- 3. Implement a add_edge method to add a connection between two vertices by setting the corresponding matrix entry to 1. For an undirected graph, set both matrix[u][v] and matrix[v][u] to 1.
- 4. Implement a print graph method to display the adjacency matrix.

```
# graph adjacency matrix.py
class Graph:
    ** ** **
    Implements a graph using an adjacency matrix representation.
    Assumes a fixed number of vertices.
    def __init__(self, num_vertices):
        self.num vertices = num vertices
        # Initialize the adjacency matrix with all zeros
        self.adj matrix = [[0] * num_vertices for _ in range(num_vertices)]
        print(f"Graph initialized with {num vertices} vertices.")
    def add edge(self, u, v, weight=1, directed=False):
        Adds an edge between vertex u and vertex v.
            u: Starting vertex (0-indexed).
            v: Ending vertex (0-indexed).
            weight: Weight of the edge (default 1 for unweighted).
            directed: If True, adds a directed edge (u -> v).
                      If False, adds an undirected edge (u \leftarrow v).
        ,, ,, ,,
        if 0 \le u \le self.num vertices and 0 \le v \le self.num vertices:
            self.adj matrix[u][v] = weight
            if not directed:
                self.adj matrix[v][u] = weight # For undirected graph
            print(f"Added edge: {u} {'->' if directed else '<->'} {v}
(Weight: {weight})")
            print(f"Error: Vertices {u}, {v} are out of bounds.")
    def print graph(self):
        Prints the adjacency matrix representation of the graph.
        print("\nAdjacency Matrix:")
        for row in self.adj matrix:
            print(row)
# Main part of the program to test graph implementation
if __name__ == "__main__":
```

```
# Create a graph with 5 vertices
g = Graph(5)
# Add undirected edges
g.add edge(0, 1) \# 0 <-> 1
g.add_edge(0, 4) # 0 <-> 4
g.add edge(1, 2) \# 1 <-> 2
g.add_edge(1, 3) # 1 <-> 3
g.add edge(1, 4) \# 1 <-> 4
g.add edge(2, 3) \# 2 < -> 3
g.add edge(3, 4) \# 3 <-> 4
g.print_graph()
print("\n--- Directed Graph Example ---")
q directed = Graph(3)
g directed.add edge(0, 1, directed=True) # 0 -> 1
g_directed.add_edge(1, 2, directed=True) # 1 -> 2
g directed.add edge(2, 0, directed=True) # 2 -> 0
g directed.add edge(0, 2, weight=5, directed=True) # 0 -> 2 with weight 5
g_directed.print_graph()
```

Input: (No direct user input for this program, graph structure is hardcoded)

```
Graph initialized with 5 vertices.
Added edge: 0 <-> 1 (Weight: 1)
Added edge: 0 <-> 4 (Weight: 1)
Added edge: 1 <-> 2 (Weight: 1)
Added edge: 1 <-> 3 (Weight: 1)
Added edge: 1 <-> 4 (Weight: 1)
Added edge: 2 <-> 3 (Weight: 1)
Added edge: 3 <-> 4 (Weight: 1)
Adjacency Matrix:
[0, 1, 0, 0, 1]
[1, 0, 1, 1, 1]
[0, 1, 0, 1, 0]
[0, 1, 1, 0, 1]
[1, 1, 0, 1, 0]
--- Directed Graph Example ---
Graph initialized with 3 vertices.
Added edge: 0 -> 1 (Weight: 1)
Added edge: 1 -> 2 (Weight: 1)
Added edge: 2 -> 0 (Weight: 1)
Added edge: 0 -> 2 (Weight: 5)
Adjacency Matrix:
[0, 1, 5]
[0, 0, 1]
[1, 0, 0]
```

Lab 14: Implementation of shortest path algorithm

Title: Shortest Path Algorithm: Dijkstra's Algorithm

Aim: To implement Dijkstra's algorithm to find the shortest path from a single source vertex to all other vertices in a weighted, directed graph.

Procedure:

- 1. Represent the graph using an adjacency matrix or adjacency list (adjacency matrix is used here for consistency with Lab 13).
- 2. Initialize distances to infinity for all vertices except the source (0 for source).
- 3. Initialize visited set to keep track of processed vertices.
- 4. Use a min-priority queue (or simply iterate to find minimum distance) to select the unvisited vertex with the smallest distance.
- 5. For the selected vertex, update the distances of its neighbors if a shorter path is found through the current vertex.
- 6. Repeat until all vertices are visited or all reachable vertices have minimum distances calculated.

```
# dijkstra shortest path.py
import sys
class Graph:
   Represents a graph using an adjacency matrix for Dijkstra's algorithm.
    def
        __init__(self, vertices):
        self.V = vertices
        self.graph = [[0 for column in range(vertices)]
                     for row in range(vertices)]
    def add edge(self, u, v, weight):
        """Adds a directed edge with a given weight."""
        self.graph[u][v] = weight
    def print solution(self, dist, src):
        """Prints the calculated shortest distances from the source."""
       print(f"\nShortest distances from source vertex {src}:")
        for node in range (self.V):
            if dist[node] == sys.maxsize:
               print(f"Vertex {node}: INF (unreachable)")
               print(f"Vertex {node}: {dist[node]}")
    def dijkstra(self, src):
        Implements Dijkstra's shortest path algorithm for a single source.
           src: The source vertex (0-indexed).
        dist = [sys.maxsize] * self.V # Initialize distances to infinity
                                    # Distance from source to itself is 0
       dist[src] = 0
       spt set = [False] * self.V # spt set[i] is True if vertex i is
included in shortest path tree
        for count in range(self.V):
```

```
# Pick the minimum distance vertex from the set of vertices not
yet processed.
            # u is always equal to src in first iteration
            min dist = sys.maxsize
            min index = -1
            for v in range(self.V):
                if dist[v] < min_dist and not spt_set[v]:</pre>
                    min dist = dist[v]
                    min index = v
            # If no reachable unvisited vertex is found, break
            if min index == -1:
                break
            u = \min index
            spt set[u] = True # Mark the picked vertex as processed
            # Update dist value of the adjacent vertices of the picked vertex
            for v in range(self.V):
                # Only update if:
                # 1. v is not in spt set
                \# 2. There is an edge from u to v
                # 3. Total weight of path from src to v through u is smaller
than current dist[v]
                if (not spt set[v] and self.graph[u][v] > 0 and
                        dist[u] != sys.maxsize and
                        dist[u] + self.graph[u][v] < dist[v]):</pre>
                    dist[v] = dist[u] + self.graph[u][v]
        self.print solution(dist, src)
# Main part of the program to test Dijkstra's algorithm
if __name__ == "__main__":
    # Create a graph with 6 vertices
    # Example graph from GeeksforGeeks
    # (0) --10--> (1) --10--> (2)
    # | \
               1
                  15
    # 6 \
                              4
    # | \
                  - 1
    # (3) --4--> (4) --5--> (5)
   g = Graph(6)
   g.add_edge(0, 1, 10)
   g.add_edge(0, 3, 6)
   g.add_edge(0, 4, 15) # Example: direct edge to 4
   g.add edge(1, 2, 10)
   g.add edge(1, 4, 15)
   g.add edge(2, 5, 4)
   g.add edge(3, 4, 4)
   g.add edge(4, 5, 5)
    source vertex = 0
   print(f"Running Dijkstra's algorithm from source vertex {source vertex}")
   g.dijkstra(source vertex)
   print("\n--- Another Example ---")
   g2 = Graph(4)
   g2.add edge(0, 1, 1)
   g2.add edge(0, 2, 4)
   g2.add edge(1, 2, 2)
   g2.add edge(1, 3, 5)
   q2.add edge(2, 3, 1)
   g2.dijkstra(0)
```

Input: (No direct user input for this program, graph and source are hardcoded)

```
Running Dijkstra's algorithm from source vertex 0

Shortest distances from source vertex 0:

Vertex 0: 0

Vertex 1: 10

Vertex 2: 20

Vertex 3: 6

Vertex 4: 10

Vertex 5: 15

--- Another Example ---

Shortest distances from source vertex 0:

Vertex 0: 0

Vertex 1: 1

Vertex 2: 3

Vertex 3: 4
```

Lab 15: Implementation of minimum spanning tree

Title: Minimum Spanning Tree: Prim's Algorithm

Aim: To implement Prim's algorithm to find the Minimum Spanning Tree (MST) of a connected, undirected, weighted graph.

Procedure:

- 1. Represent the graph using an adjacency matrix.
- 2. Initialize key values (minimum edge weight to connect to MST) to infinity for all vertices, and 0 for the starting vertex.
- 3. Initialize mst set to keep track of vertices already included in MST.
- 4. Initialize parent array to store the MST structure.
- 5. Repeat v times (where v is the number of vertices):
 - o Select the vertex u not yet in mst set that has the minimum key value.
 - o Add u to mst set.
 - o For each neighbor v of u: if v is not in mst_set and the weight of edge (u, v) is less than key[v], update key[v] and set parent[v] = u.
- 6. Print the edges of the MST and its total weight.

```
# prims_mst.py
import sys
class Graph:
   Represents a graph for Prim's algorithm using an adjacency matrix.
   def __init__ (self, vertices):
    self.V = vertices
        self.graph = [[0 for column in range(vertices)]
                      for row in range(vertices)]
    def add edge(self, u, v, weight):
        """Adds an undirected edge with a given weight."""
        self.graph[u][v] = weight
        self.graph[v][u] = weight # For undirected graph
    def print mst(self, parent):
        """Prints the constructed MST edges and total weight."""
        print("\nEdges in Minimum Spanning Tree:")
        total weight = 0
        for i in range(1, self.V):
            print(f"{parent[i]} - {i} Weight: {self.graph[i][parent[i]]}")
            total weight += self.graph[i][parent[i]]
        print(f"\nTotal weight of MST: {total_weight}")
    def prim mst(self):
        Implements Prim's algorithm to find the Minimum Spanning Tree.
        key = [sys.maxsize] * self.V # Key values used to pick minimum weight
edge in cut
       parent = [-1] * self.V # Array to store constructed MST
       mst set = [False] * self.V # To represent set of vertices not yet
included in MST
```

```
parent[0] = -1 # First node is always root of MST
        for count in range(self.V):
            # Pick the minimum key vertex from the set of vertices
            # not yet included in MST
           min_key = sys.maxsize
           min index = -1
            for v in range(self.V):
                if key[v] < min key and not mst set[v]:
                   \min \text{ key} = \ker[v]
                   min index = v
            # If no reachable unvisited vertex is found, break
            if min index == -1:
               break
            u = \min index
           mst set[u] = True # Add the picked vertex to the MST set
            # Update key value and parent index of the adjacent vertices of
the picked vertex.
            # Consider only those vertices which are not yet in MST set.
            for v in range(self.V):
                # graph[u][v] is non-zero only for adjacent vertices of u
                # mst set[v] is false for vertices not yet included in MST
                # Update the key only if graph[u][v] is smaller than key[v]
                if self.graph[u][v] > 0 and not mst set[v] and key[v] >
self.graph[u][v]:
                   key[v] = self.graph[u][v]
                   parent[v] = u
       self.print mst(parent)
# Main part of the program to test Prim's algorithm
if __name__ == "__main_ ":
    # Create a graph with 5 vertices
    # Example graph:
    # (0) --2-- (1) --3-- (2)
   # (3) --9-- (4) -----
   g = Graph(5)
   g.add edge(0, 1, 2)
   g.add_edge(0, 3, 6)
   g.add edge(1, 2, 3)
   g.add edge(1, 3, 8) # Not part of typical MST for this example
   g.add edge(1, 4, 5)
   g.add edge(2, 4, 7)
   g.add edge(3, 4, 9)
   print("Running Prim's algorithm for MST...")
   g.prim mst()
   print("\n--- Another Example ---")
   q2 = Graph(4)
   g2.add edge(0, 1, 10)
   g2.add edge(0, 2, 6)
   g2.add edge(0, 3, 5)
   g2.add edge(1, 3, 15)
   g2.add edge(2, 3, 4)
   g2.prim mst()
```

key[0] = 0 # Make key 0 so that this vertex is picked first

Input: (No direct user input for this program, graph structure is hardcoded)