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Department of Computer Applications

Circular - 2023-24

B.Sc. CS 4th Sem

OPERATING SYSTEMS (USA23402J)- Lab Manual

Lab 1: Comparison between various Operating Systems

Title

Comparison between various Operating Systems

Aim

To understand and compare the key features, architectures, and functionalities of different operating systems (e.g., Windows, macOS, Linux, Android, iOS).

Procedure

- 1. **Research:** Gather information on the history, design philosophy, kernel types (monolithic, microkernel, hybrid), file systems, process management, memory management, security features, and user interfaces of at least three different operating systems.
- 2. **Categorization:** Classify the operating systems based on their target devices (desktop, mobile, server) and primary use cases.
- 3. **Comparison Table:** Create a comparative table highlighting the similarities and differences across various parameters.
- 4. **Analysis:** Analyze the strengths and weaknesses of each operating system in different scenarios.
- 5. **Conclusion:** Summarize the findings and discuss which operating system might be best suited for specific applications or user needs.

Source Code

(N/A - This lab is primarily theoretical and research-based, not involving direct coding.)

Input

 $(N/A - This \ lab \ does \ not \ require \ specific \ program \ input.)$

Expected Output

A comprehensive report or presentation detailing the comparison, including a comparative table and analysis of the chosen operating systems.

Lab 2: Booting process in GNU/Linux OS

Title

Booting process in GNU/Linux OS

Aim

To understand the sequence of events and components involved in the booting process of a GNU/Linux operating system.

Procedure

- 1. **BIOS/UEFI Initialization:** Describe the role of BIOS/UEFI in the initial boot sequence, including POST (Power-On Self-Test) and loading the boot loader.
- 2. **Boot Loader (GRUB/LILO):** Explain how boot loaders like GRUB or LILO take control, locate the kernel, and load it into memory.
- 3. **Kernel Initialization:** Detail the steps involved in kernel initialization, including hardware detection, module loading, and setting up the root filesystem.
- 4. **Init/Systemd:** Explain the role of the init process (or systemd in modern Linux distributions) as the first user-space process, responsible for starting other services and bringing the system to a usable state.
- 5. **Runlevels/Targets:** Discuss the concept of runlevels (SysVinit) or targets (systemd) and how they define the system's operational state.
- 6. **Observation:** (Optional) If possible, observe the boot messages during a Linux system startup to identify key stages.

Source Code

(N/A - This lab is observational and conceptual, not involving direct coding.)

Input

(N/A - This lab does not require specific program input.)

Expected Output

A detailed explanation of the GNU/Linux booting process, potentially including a flowchart or diagram illustrating the sequence.

Lab 3: Multi-thread Programming

Title

Multi-thread Programming

Aim

To implement a program demonstrating multi-threading concepts, including thread creation, synchronization, and communication.

Procedure

- 1. **Choose a Language:** Select a programming language that supports multi-threading (e.g., C, C++, Java, Python).
- 2. **Problem Definition:** Define a simple problem that can benefit from multi-threading (e.g., parallel computation, concurrent task execution).
- 3. **Thread Creation:** Implement code to create multiple threads.
- 4. **Task Assignment:** Assign different parts of the problem to each thread.
- 5. **Synchronization (if needed):** If threads share resources, implement synchronization mechanisms (e.g., mutexes, semaphores, locks) to prevent race conditions.
- 6. **Thread Join:** Ensure the main thread waits for all child threads to complete their execution.
- 7. **Compile and Run:** Compile and execute the program.

```
// Example: C program for simple multi-threading with pthreads
#include <stdio.h>
#include <pthread.h>
#include <unistd.h> // For sleep
// Function to be executed by a thread
void *myThreadFunction(void *threadid) {
    long tid;
   tid = (long)threadid;
   printf("Hello from thread #%ld!\n", tid);
    sleep(1); // Simulate some work
   printf("Thread #%ld exiting.\n", tid);
   pthread exit(NULL);
int main() {
   pthread t threads[3]; // Array to hold thread IDs
   int rc;
   long t;
   printf("Main: Creating threads...\n");
    for (t = 0; t < 3; t++) {
        printf("Main: Creating thread %ld\n", t);
        rc = pthread create(&threads[t], NULL, myThreadFunction, (void *)t);
           printf("ERROR; return code from pthread create() is %d\n", rc);
            return 1;
        }
    }
    printf("Main: All threads created. Waiting for them to finish...\n");
```

```
// Wait for all threads to complete
for (t = 0; t < 3; t++) {
    pthread_join(threads[t], NULL);
}

printf("Main: All threads finished. Exiting.\n");
pthread_exit(NULL); // Exit the main thread gracefully
}</pre>
```

(N/A - This example program does not require user input.)

Expected Output

```
Main: Creating threads...
Main: Creating thread 0
Main: Creating thread 1
Main: Creating thread 2
Main: All threads created. Waiting for them to finish...
Hello from thread #0!
Hello from thread #1!
Hello from thread #2!
Thread #0 exiting.
Thread #1 exiting.
Thread #2 exiting.
Main: All threads finished. Exiting.
```

(Note: The order of "Hello from thread" and "Thread exiting" messages might vary due to thread scheduling.)

Lab 4: Simulation of FCFS CPU scheduling algorithm

Title

Simulation of FCFS (First-Come, First-Served) CPU scheduling algorithm

Aim

To simulate the First-Come, First-Served (FCFS) CPU scheduling algorithm and calculate performance metrics like average waiting time and average turnaround time.

Procedure

- 1. **Data Structure:** Define a structure or class to represent a process, including attributes like process ID, arrival time, and burst time.
- 2. **Input:** Get the number of processes and their arrival and burst times from the user.
- 3. **Sorting:** (Implicit in FCFS) Processes are scheduled in the order of their arrival. If arrival times are the same, order by process ID.
- 4. Calculation:
 - Completion Time: For each process, calculate its completion time (start time + burst time).
 - o **Turnaround Time:** Calculate turnaround time (completion time arrival time).
 - o Waiting Time: Calculate waiting time (turnaround time burst time).
- 5. **Performance Metrics:** Calculate the average waiting time and average turnaround time for all processes.
- 6. **Output:** Display the results in a clear tabular format.

```
// Example: C program for FCFS CPU scheduling simulation
#include <stdio.h>
struct Process {
   int pid; // Process ID
   int arrival time;
   int burst time;
   int start time;
   int completion time;
   int turnaround time;
   int waiting time;
};
void calculate times(struct Process p[], int n) {
   int current time = 0;
   int total waiting time = 0;
   int total turnaround time = 0;
printf("\nPID\tArrival\tBurst\tStart\tCompletion\tTurnaround\tWaiting\n");
  printf("-----
----\n");
   for (int i = 0; i < n; i++) {
       // Start time is the maximum of current time and process's
arrival time
       p[i].start time = (current time > p[i].arrival time) ? current time :
p[i].arrival time;
       p[i].completion_time = p[i].start_time + p[i].burst_time;
       p[i].turnaround time = p[i].completion time - p[i].arrival time;
```

```
p[i].waiting time = p[i].turnaround time - p[i].burst time;
        // Update current time for the next process
        current_time = p[i].completion_time;
        total waiting time += p[i].waiting time;
        total turnaround time += p[i].turnaround time;
        printf("%d\t%d\t%d\t%d\t\t%d\t\t%d\t\t
               p[i].pid, p[i].arrival time, p[i].burst time, p[i].start time,
               p[i].completion_time, p[i].turnaround_time,
p[i].waiting time);
    }
    printf("\nAverage Waiting Time: %.2f\n", (float)total waiting time / n);
    printf("Average Turnaround Time: %.2f\n", (float)total_turnaround_time /
n);
}
int main() {
    int n;
    printf("Enter the number of processes: ");
    scanf("%d", &n);
    struct Process processes[n];
    printf("Enter Arrival Time and Burst Time for each process:\n");
    for (int i = 0; i < n; i++) {
        processes[i].pid = i + 1;
        printf("Process %d (Arrival Time Burst Time): ", i + 1);
        scanf("%d %d", &processes[i].arrival time, &processes[i].burst time);
    \ensuremath{//} FCFS assumes processes are already sorted by arrival time.
    // If not, you would need to sort them here.
    // For simplicity, we assume input is given in arrival order.
    calculate_times(processes, n);
    return 0;
}
Input
Enter the number of processes: 3
Enter Arrival Time and Burst Time for each process:
Process 1 (Arrival Time Burst Time): 0 5
Process 2 (Arrival Time Burst Time): 1 3
Process 3 (Arrival Time Burst Time): 2 8
```

Expected Output

PID	Arrival	Burst	Start	Completion	Turnaround	Waiting
1 2	0	5	0	5	5 7	0 4
3	2	8	8	16	14	6

Average Waiting Time: 3.33 Average Turnaround Time: 8.67

Lab 5: Priority CPU scheduling algorithm

Title

Priority CPU scheduling algorithm

Aim

To simulate the Priority CPU scheduling algorithm (non-preemptive) and calculate performance metrics like average waiting time and average turnaround time.

Procedure

- 1. **Data Structure:** Define a structure or class for a process, including process ID, arrival time, burst time, and priority. (Lower priority number usually means higher priority).
- 2. **Input:** Get the number of processes and their arrival time, burst time, and priority from the user.
- 3. **Sorting/Selection:** At each time step, select the process with the highest priority among those that have arrived and are not yet completed.
- 4. Calculation:
 - o Completion Time: For each process, calculate its completion time.
 - o **Turnaround Time:** Calculate turnaround time (completion time arrival time).
 - o Waiting Time: Calculate waiting time (turnaround time burst time).
- 5. **Performance Metrics:** Calculate the average waiting time and average turnaround time.
- 6. **Output:** Display the results in a clear tabular format.

Source Code

nq n");

```
// Example: C program for Non-Preemptive Priority CPU scheduling simulation
#include <stdio.h>
#include <limits.h> // For INT MAX
struct Process {
   int pid;
   int arrival time;
   int burst time;
   int priority;
   int start time;
   int completion time;
   int turnaround time;
   int waiting time;
   int is completed; // Flag to track if process is completed
};
void find avg time(struct Process p[], int n) {
   int current time = 0;
   int completed processes = 0;
   int total_waiting_time = 0;
   int total turnaround time = 0;
    // Initialize all processes as not completed
    for (int i = 0; i < n; i++) {
       p[i].is_completed = 0;
printf("\nPID\tArrival\tBurst\tPriority\tStart\tCompletion\tTurnaround\tWaiti
```

```
printf("-----
    while (completed_processes < n) {</pre>
        int best priority = INT MAX;
        int selected process index = -1;
        // Find the process with the highest priority that has arrived and is
not completed
        for (int i = 0; i < n; i++) {
            if (p[i].arrival time <= current time && p[i].is completed == 0)
                if (p[i].priority < best priority) {</pre>
                    best priority = p[i].priority;
                    selected_process_index = i;
                // If priorities are same, FCFS among them (based on arrival
time)
                else if (p[i].priority == best priority) {
                    if (p[i].arrival time <
p[selected_process_index].arrival_time) {
                        selected process index = i;
            }
        }
        if (selected process index == -1) {
            // No process is ready, increment time
            current time++;
        } else {
            int i = selected process index;
            p[i].start time = current time;
            p[i].completion time = p[i].start time + p[i].burst time;
            p[i].turnaround_time = p[i].completion_time - p[i].arrival_time;
            p[i].waiting_time = p[i].turnaround_time - p[i].burst_time;
            p[i].is\_completed = 1;
            current time = p[i].completion time;
            completed processes++;
            total_waiting_time += p[i].waiting_time;
            total_turnaround_time += p[i].turnaround_time;
            printf("%d\t%d\t%d\t\t%d\t\t%d\t\t%d\t\t%d\t\t
                   p[i].pid, p[i].arrival time, p[i].burst time,
p[i].priority,
                   p[i].start time, p[i].completion time,
p[i].turnaround time, p[i].waiting time);
    printf("\nAverage Waiting Time: %.2f\n", (float)total waiting time / n);
   printf("Average Turnaround Time: %.2f\n", (float)total turnaround time /
n);
int main() {
   printf("Enter the number of processes: ");
   scanf("%d", &n);
    struct Process processes[n];
```

```
printf("Enter Arrival Time, Burst Time, and Priority for each
process:\n");
    for (int i = 0; i < n; i++) {
       processes[i].pid = i + 1;
       printf("Process %d (Arrival Time Burst Time Priority): ", i + 1);
       scanf("%d %d %d", &processes[i].arrival_time,
&processes[i].burst_time, &processes[i].priority);
    find_avg_time(processes, n);
   return 0;
}
Input
Enter the number of processes: 4
Enter Arrival Time, Burst Time, and Priority for each process:
Process 1 (Arrival Time Burst Time Priority): 0 6 2
Process 2 (Arrival Time Burst Time Priority): 1 8 1
Process 3 (Arrival Time Burst Time Priority): 2 7 3
Process 4 (Arrival Time Burst Time Priority): 3 3 0
```

Expected Output

PID Waiting		Burst	Priority	Start	Completion	Turnaround
4	3	3	0	3	6	3
2	1	8	1	6	14	13
5 1	0	6	2	14	20	20
14 3	2	7	3	20	27	25
18						

Average Waiting Time: 9.25
Average Turnaround Time: 15.25

Lab 6: Simulation of Round Robin CPU scheduling algorithm

Title

Simulation of Round Robin CPU scheduling algorithm

Aim

To simulate the Round Robin (RR) CPU scheduling algorithm and calculate performance metrics like average waiting time and average turnaround time.

Procedure

- 1. **Data Structure:** Define a structure or class for a process, including process ID, arrival time, burst time, and remaining burst time.
- 2. **Input:** Get the number of processes, their arrival times, burst times, and the time quantum from the user.
- 3. **Ready Queue:** Implement a queue to manage processes ready for execution.
- 4. Scheduling Logic:
 - o Maintain a current_time variable.
 - o At each time step, add processes that have arrived to the ready queue.
 - o If the ready queue is not empty, dequeue a process.
 - Execute the process for the time quantum or until its burst time is completed, whichever is shorter.
 - o If the process completes, record its completion time, turnaround time, and waiting time.
 - o If the process doesn't complete, enqueue it back to the ready queue.
 - o Increment current time.
- 5. **Performance Metrics:** Calculate the average waiting time and average turnaround time.
- 6. **Output:** Display the results in a clear tabular format.

```
// Example: C program for Round Robin CPU scheduling simulation
#include <stdio.h>
#include <stdbool.h> // For boolean type
struct Process {
   int pid;
   int arrival time;
   int burst time;
   int remaining burst time;
   int start time; // First time process gets CPU
   int completion time;
    int turnaround time;
   int waiting time;
} ;
void find avg time(struct Process p[], int n, int quantum) {
    int current time = 0;
    int completed processes = 0;
    int total waiting time = 0;
    int total turnaround time = 0;
    // Array to keep track of whether a process has entered the ready queue
   bool in_ready_queue[n];
    for (int i = 0; i < n; i++) {
       p[i].remaining burst time = p[i].burst time;
```

```
p[i].start time = -1; // Initialize start time
       in ready queue[i] = false;
    // Simple queue implementation (using an array for simplicity)
    int ready queue[n];
    int front = 0, rear = -1; // Queue pointers
printf("\nPID\tArrival\tBurst\tRemaining\tStart\tCompletion\tTurnaround\tWait
ing\n");
   printf("-----
----\n");
    while (completed processes < n) {</pre>
       // Add newly arrived processes to the ready queue
        for (int i = 0; i < n; i++) {
           if (p[i].arrival time <= current time && !in ready queue[i] &&
p[i].remaining burst time > 0) {
               rear = (rear + 1) % n; // Circular queue
               ready queue[rear] = i;
               in ready queue[i] = true;
           }
        }
       if (front > rear && completed processes < n) { // If queue is empty
and not all processes completed
           // No process is ready, increment time to the next arrival
           int next arrival = -1;
           for (int i = 0; i < n; i++) {
               if(p[i].remaining_burst_time > 0) {
                   if(next_arrival == -1 || p[i].arrival_time <</pre>
next arrival) {
                       next arrival = p[i].arrival time;
                   }
               }
           if(next_arrival != -1 && current_time < next_arrival) {</pre>
               current time = next arrival;
           } else {
               current time++; // Increment if no next arrival or already
past it
           continue; // Go to next iteration to check for new arrivals
       }
       int current process index = ready queue[front];
        front = (front + 1) % n; // Dequeue
        if (p[current process index].start time == -1) {
           p[current process index].start time = current time;
       int execution time = (p[current process index].remaining burst time <
quantum) ?
                            p[current process index].remaining burst time :
quantum;
       current time += execution time;
       p[current process index].remaining burst time -= execution time;
       // Add newly arrived processes during this quantum to the ready queue
        for (int i = 0; i < n; i++) {
           if (p[i].arrival time <= current time && !in ready queue[i] &&
p[i].remaining burst time > 0) {
               rear = (rear + 1) % n;
```

```
ready queue[rear] = i;
                in ready queue[i] = true;
            }
        }
        if (p[current_process_index].remaining_burst_time == 0) {
            p[current_process_index].completion_time = current_time;
            p[current_process_index].turnaround_time =
p[current process index].completion time -
p[current process index].arrival time;
            p[current_process_index].waiting_time =
p[current process index].turnaround time -
p[current process index].burst time;
            completed processes++;
            total waiting time += p[current process index].waiting time;
            total turnaround time +=
p[current process index].turnaround time;
            printf("%d\t%d\t%d\t\t%d\t\t%d\t\t%d\t\t%d\n",
                   p[current_process_index].pid,
p[current_process_index].arrival_time, p[current_process_index].burst_time,
                   p[current_process_index].remaining_burst_time,
p[current_process_index].start_time,
                   p[current_process_index].completion_time,
p[current process index].turnaround time,
                   p[current_process_index].waiting time);
        } else {
            // Process not completed, add back to ready queue
            rear = (rear + 1) % n;
            ready queue[rear] = current process index;
        }
    }
    printf("\nAverage Waiting Time: %.2f\n", (float)total_waiting_time / n);
    printf("Average Turnaround Time: %.2f\n", (float)total_turnaround_time /
n);
int main() {
    int n:
    printf("Enter the number of processes: ");
    scanf("%d", &n);
    struct Process processes[n];
    printf("Enter Arrival Time and Burst Time for each process:\n");
    for (int i = 0; i < n; i++) {
       processes[i].pid = i + 1;
        printf("Process %d (Arrival Time Burst Time): ", i + 1);
        scanf("%d %d", &processes[i].arrival time, &processes[i].burst time);
    }
    int quantum;
    printf("Enter the time quantum: ");
    scanf("%d", &quantum);
    find avg time(processes, n, quantum);
    return 0;
```

Enter the number of processes: 3
Enter Arrival Time and Burst Time for each process:

Process 1 (Arrival Time Burst Time): 0 10 Process 2 (Arrival Time Burst Time): 1 4 Process 3 (Arrival Time Burst Time): 2 6 Enter the time quantum: 3

Expected Output

PID Waiting		Burst	Remaining	Start	Completion	Turnaround
2	1	4	0	1	7	6
2						
3	2	6	0	7	15	13
1	0	10	0	Ο	20	20
10	O	10	O	O	20	20

Average Waiting Time: 6.33 Average Turnaround Time: 13.00

Lab 7: Write a procedure for timer interrupt handler

Title

Procedure for Timer Interrupt Handler

Aim

To understand the role and implementation aspects of a timer interrupt handler in an operating system context.

Procedure

- 1. **Concept of Timer Interrupts:** Explain what a timer interrupt is, its purpose (e.g., time slicing for CPU scheduling, system clock updates), and how it's generated by hardware.
- 2. Interrupt Vector Table (IVT)/Interrupt Descriptor Table (IDT): Describe how the operating system sets up the IVT/IDT to map interrupt numbers to specific interrupt service routines (ISRs).
- 3. **Timer ISR Flow:** Outline the typical steps involved when a timer interrupt occurs:
 - o CPU saves context of the currently running process.
 - o CPU jumps to the timer interrupt handler's address.
 - Handler's Responsibilities:
 - Acknowledge the interrupt to the Programmable Interrupt Controller (PIC).
 - Update system time/tick count.
 - Decrement process time slice (if used for scheduling).
 - If time slice expires, trigger the scheduler.
 - Perform any other periodic tasks.
 - CPU restores context of the next process (or the interrupted process if not switched).
 - o CPU returns from interrupt.
- 4. **Pseudocode/Conceptual Code:** Provide pseudocode or a conceptual C-like structure for a timer interrupt handler.
- 5. **Kernel Context:** Discuss how the handler operates in kernel mode and its implications for system stability.

```
// Pseudocode for a conceptual Timer Interrupt Handler
// Global variables (part of OS kernel data)
unsigned long system ticks = 0;
int current process timeslice = 0;
const int TIME QUANTUM = 10; // Example time quantum for scheduling
// Function prototypes (assuming they exist in the kernel)
void save cpu context(void *context ptr);
void restore cpu_context(void *context ptr);
void acknowledge pic interrupt(int irq number);
void schedule next process(); // Function to invoke the scheduler
void *get_current_process_context();
void *get_next_process_context();
// The Timer Interrupt Service Routine (ISR)
// This function is invoked by the CPU when a timer interrupt (e.g., IRQ0)
occurs.
void timer_interrupt_handler() {
```

```
// 1. Save the context of the currently executing process/thread
       (Registers, program counter, stack pointer, etc.)
   //
         This is often done by hardware or assembly stub before calling this
C function.
   // For conceptual understanding, we assume a function does this.
    // save cpu context(get current process context()); // Conceptual
    // 2. Acknowledge the interrupt to the Programmable Interrupt Controller
(PIC)
          This tells the PIC that the interrupt has been handled and it can
   //
send more.
    acknowledge pic interrupt(0); // Assuming timer uses IRQ0
    // 3. Update system time/tick count
    system ticks++;
    // 4. Handle CPU scheduling (time slicing)
    current_process_timeslice--;
    if (current process timeslice <= 0) {
        // Time slice for the current process has expired
        \ensuremath{//} Trigger the scheduler to pick the next process
        schedule_next_process();
        // Reset time slice for the new process
        current process timeslice = TIME QUANTUM;
    }
    // 5. Perform other periodic tasks (e.g., updating sleep timers, flushing
buffers)
       ... (other kernel tasks) ...
   //
    // 6. Restore the context of the next process to be executed
    // (If scheduler switched processes, restore the new one; otherwise,
restore the original)
   // This is also often done by hardware or assembly stub after this C
function returns.
    // restore cpu context(get next process context()); // Conceptual
// Example of how the timer might be initialized (conceptual)
void initialize timer() {
    // Set up the timer hardware to generate interrupts periodically
    // (e.g., configure Programmable Interval Timer - PIT)
    // Register the timer interrupt handler with the Interrupt Descriptor
Table (IDT)
    // so that the CPU knows where to jump when IRQO occurs.
    // current process timeslice = TIME QUANTUM; // Initialize for first
process
```

(N/A - This lab is theoretical and conceptual, not requiring direct input.)

Expected Output

A detailed written procedure and conceptual pseudocode outlining the steps and responsibilities of a timer interrupt handler within an operating system.

Lab 8: Classical inter process communication problem (Producer consumer)

Title

Classical Inter-Process Communication Problem: Producer-Consumer

Aim

To implement a solution to the Producer-Consumer problem using inter-process communication (IPC) mechanisms like semaphores and shared memory (or mutexes and condition variables for threads).

Procedure

- 1. **Problem Definition:** Understand the Producer-Consumer problem: a producer generates data and puts it into a buffer, and a consumer takes data from the buffer. The buffer has a finite size.
- 2. **IPC Mechanisms:** Choose appropriate IPC mechanisms:
 - o **For Processes:** Shared memory for the buffer, and semaphores (e.g., empty, full, mutex) for synchronization.
 - For Threads: A shared buffer, and mutexes and condition variables for synchronization.
- 3. **Buffer Implementation:** Create a shared buffer (e.g., a circular array).
- 4. Producer Logic:
 - o Generate an item.
 - o Wait if the buffer is full (using empty semaphore or condition variable).
 - o Acquire mutex (or mutex semaphore).
 - o Add item to buffer.
 - o Release mutex.
 - o Signal that the buffer is not empty (using full semaphore or condition variable).
- 5. Consumer Logic:
 - o Wait if the buffer is empty (using full semaphore or condition variable).
 - o Acquire mutex (or mutex semaphore).
 - o Remove item from buffer.
 - o Release mutex.
 - o Signal that the buffer is not full (using empty semaphore or condition variable).
- 6. **Compile and Run:** Compile the producer and consumer programs (or a single program with producer/consumer threads) and observe their synchronized execution.

```
// Example: C program for Producer-Consumer problem using pthreads, mutexes,
and condition variables

#include <stdio.h>
#include <stdlib.h>
#include <pthread.h>
#include <unistd.h> // For sleep

#define BUFFER_SIZE 5
int buffer[BUFFER_SIZE];
int in = 0; // Index for producer
int out = 0; // Index for consumer
```

```
int count = 0; // Number of items in buffer
pthread mutex t mutex; // Mutex for critical section
pthread cond t full; // Condition variable for consumer (buffer is full)
pthread cond t empty; // Condition variable for producer (buffer is empty)
void *producer(void *param) {
    int item;
    for (int i = 0; i < 10; i++) { // Produce 10 items
        item = rand() % 100; // Generate a random item
        pthread mutex lock(&mutex); // Acquire mutex
        while (count == BUFFER SIZE) { // If buffer is full, wait
            printf("Producer: Buffer is full. Waiting...\n");
            pthread cond wait (&empty, &mutex);
        buffer[in] = item;
        in = (in + 1) % BUFFER SIZE;
        count++;
        printf("Producer: Produced item %d. Buffer count: %d\n", item,
count);
        pthread cond signal (&full); // Signal consumer that buffer is not
empty
        pthread mutex unlock(&mutex); // Release mutex
        sleep(rand() % 2); // Simulate production time
    pthread exit(0);
}
void *consumer(void *param) {
    int item;
    for (int i = 0; i < 10; i++) { // Consume 10 items
        pthread_mutex_lock(&mutex); // Acquire mutex
        while (count == 0) \{ // \text{ If buffer is empty, wait } \}
            printf("Consumer: Buffer is empty. Waiting...\n");
            pthread cond wait(&full, &mutex);
        }
        item = buffer[out];
        out = (out + 1) % BUFFER SIZE;
        count--;
        printf("Consumer: Consumed item %d. Buffer count: %d\n", item,
count);
        pthread cond signal (&empty); // Signal producer that buffer is not
full
        pthread mutex unlock(&mutex); // Release mutex
        sleep(rand() % 3); // Simulate consumption time
    pthread exit(0);
}
int main() {
    pthread t prod tid, cons tid;
    pthread mutex init (&mutex, NULL);
    pthread cond init(&full, NULL);
    pthread cond init(&empty, NULL);
    pthread create(&prod tid, NULL, producer, NULL);
    pthread create(&cons tid, NULL, consumer, NULL);
```

```
pthread_join(prod_tid, NULL);
pthread_join(cons_tid, NULL);

pthread_mutex_destroy(&mutex);
pthread_cond_destroy(&full);
pthread_cond_destroy(&empty);

printf("\nProducer-Consumer simulation finished.\n");
return 0;
}
```

(N/A - This program does not require user input.)

Expected Output

(Output will vary slightly due to thread scheduling and random delays, but will follow this pattern:)

```
Producer: Produced item 84. Buffer count: 1
Consumer: Consumed item 84. Buffer count: 0
Producer: Produced item 75. Buffer count: 1
Producer: Produced item 28. Buffer count: 2
Consumer: Consumed item 75. Buffer count: 1
Producer: Produced item 7. Buffer count: 2
Consumer: Consumed item 28. Buffer count: 1
Producer: Produced item 90. Buffer count: 2
Producer: Produced item 16. Buffer count: 3
Consumer: Consumed item 7. Buffer count: 2
Producer: Produced item 12. Buffer count: 3
Producer: Produced item 34. Buffer count: 4
Producer: Produced item 56. Buffer count: 5
Producer: Buffer is full. Waiting...
Consumer: Consumed item 90. Buffer count: 4
Producer: Produced item 93. Buffer count: 5
Producer: Buffer is full. Waiting...
Consumer: Consumed item 16. Buffer count: 4
Consumer: Consumed item 12. Buffer count: 3
Consumer: Consumed item 34. Buffer count: 2
Consumer: Consumed item 56. Buffer count: 1
Consumer: Consumed item 93. Buffer count: 0
```

Producer-Consumer simulation finished.

Lab 9: Write a procedure to make message passing in inter process communication

Title

Procedure for Message Passing in Inter-Process Communication

Aim

To understand and implement a basic message passing mechanism for inter-process communication (IPC) using system calls like msgget, msgsnd, and msgrcv (in Unix-like systems).

Procedure

- 1. **Concept of Message Queues:** Explain message queues as a linked list of messages stored within the kernel, allowing processes to send and receive messages.
- 2. **System Calls:** Describe the purpose of key system calls:
 - o msgget (): Creates a new message queue or gets an ID of an existing one.
 - o msgsnd(): Sends a message to a message queue.
 - o msgrcv(): Receives a message from a message queue.
 - o msgctl(): Performs control operations on a message queue (e.g., delete).
- 3. **Message Structure:** Define a structure for the messages to be sent, including a mtype (message type) and mtext (message text/data).
- 4. Sender Process Logic:
 - o Get/create message queue ID using msgget ().
 - o Prepare the message structure.
 - o Send the message using msgsnd().
 - o (Optional) Delete the message queue using msgctl() after communication is done.
- 5. Receiver Process Logic:
 - o Get message queue ID using msgget ().
 - o Receive the message using msgrcv(), specifying the desired message type.
 - o Process the received message.
 - o (Optional) Delete the message queue using msgctl() after communication is done.
- 6. **Compile and Run:** Compile the sender and receiver programs separately and run them concurrently to observe message exchange.

```
// Example: C program for Message Passing (Sender) using System V IPC
#include <stdio.h>
#include <stdlib.h>
#include <string.h>
#include <sys/ipc.h>
#include <sys/msg.h>
#include <sys/types.h>

// Define a structure for the message
struct message_buffer {
   long mtype; // Message type (must be > 0)
   char mtext[100]; // Message data
};
```

```
int main() {
    key t key;
    int msgid;
   struct message buffer message;
    // 1. Generate a unique key for the message queue
    // ftok(pathname, proj_id) generates a System V IPC key
    key = ftok("progfile", 65); // "progfile" can be any existing file
    // 2. Create or get the message queue ID
    // msgget(key, msgflg) returns the message queue identifier
    // IPC CREAT: Create a new queue if it doesn't exist
    // 0666: Read/write permissions for owner, group, others
   msgid = msgget(key, 0666 | IPC_CREAT);
    if (msgid == -1) {
        perror("msgget failed");
        exit(EXIT_FAILURE);
    }
   printf("Message Queue ID: %d\n", msgid);
    // 3. Prepare the message to be sent
   message.mtype = 1; // Message type 1
    strcpy(message.mtext, "Hello from Sender!");
    // 4. Send the message to the message queue
    // msgsnd(msqid, msgp, msgsz, msgflg)
    // msqid: Message queue ID
    // msgp: Pointer to the message buffer
    // msgsz: Size of the message text (mtext)
    // msgflg: Flags (e.g., IPC_NOWAIT)
    if (msgsnd(msgid, &message, sizeof(message.mtext), 0) == -1) {
       perror("msgsnd failed");
       exit(EXIT FAILURE);
    }
   printf("Message sent: %s\n", message.mtext);
    // Optional: Delete the message queue after sending
    // msgctl(msqid, cmd, buf)
    // cmd: IPC RMID to remove the queue
    // if (msgctl(msgid, IPC RMID, NULL) == -1) {
           perror("msgctl (IPC RMID) failed");
    //
    //
           exit(EXIT FAILURE);
    // }
    // printf("Message queue removed.\n");
   return 0;
}
// Example: C program for Message Passing (Receiver) using System V IPC
#include <stdio.h>
#include <stdlib.h>
#include <string.h>
#include <sys/ipc.h>
#include <sys/msq.h>
#include <sys/types.h>
// Define a structure for the message (must be same as sender)
struct message buffer {
   long mtype;
   char mtext[100];
} ;
```

```
int main() {
    key t key;
    int msgid;
    struct message buffer message;
    // 1. Generate the same unique key as the sender
    key = ftok("progfile", 65);
    // 2. Get the message queue ID (it must already exist)
   msgid = msgget(key, 0666); // No IPC CREAT here, just get existing
    if (msgid == -1) {
       perror("msgget failed (queue might not exist or permissions are
wrong)");
       exit(EXIT FAILURE);
    }
    printf("Message Queue ID: %d\n", msgid);
    // 3. Receive the message from the message queue
    // msgrcv(msqid, msgp, msgsz, msgtyp, msgflg)
    // msgtyp: 0 for first message, >0 for specific type, <0 for first
message <= abs(msgtyp)</pre>
   if (msgrcv(msgid, &message, sizeof(message.mtext), 1, 0) == -1) { //
Receive message of type 1
        perror("msgrcv failed");
       exit(EXIT FAILURE);
    }
    printf("Message received: %s\n", message.mtext);
    // Optional: Delete the message queue after receiving
    // if (msgctl(msgid, IPC RMID, NULL) == -1) {
          perror("msgctl (IPC RMID) failed");
    //
           exit(EXIT FAILURE);
    // }
    // printf("Message queue removed.\n");
   return 0;
}
```

(Run the sender program first, then the receiver program. No direct input is given to the programs themselves.)

Expected Output

For Sender Program:

```
Message Queue ID: <some_integer_id>
Message sent: Hello from Sender!
```

For Receiver Program:

```
Message Queue ID: <same_integer_id>
Message received: Hello from Sender!
```

Lab 10: Program to implement Bankers Algorithm

Title

Program to implement Banker's Algorithm

Aim

To implement the Banker's Algorithm for deadlock avoidance and determine if a system is in a safe state.

Procedure

- 1. **Algorithm Concept:** Understand the Banker's Algorithm, which checks for a safe state by simulating the allocation of resources to processes.
- 2. **Data Structures:** Define arrays/matrices to store:
 - o Available: Number of available instances of each resource type.
 - o Max: Maximum demand of each resource type for each process.
 - Allocation: Number of resources of each type currently allocated to each process.
 - o Need: Remaining resources needed by each process (Need = Max Allocation).
- 3. **Input:** Get the number of processes, number of resource types, and the initial values for Available, Max, and Allocation matrices.
- 4. Safety Algorithm Implementation:
 - o Initialize Work = Available and Finish[i] = false for all processes i.
 - o Find a process i such that Finish[i] == false and Need[i] <= Work.
 - o If such a process is found:
 - Work = Work + Allocation[i]
 - Finish[i] = true
 - Repeat step 2.
 - o If no such process is found, check if Finish[i] == true for all i. If yes, the system is in a safe state; otherwise, it's in an unsafe state.
- 5. **Output:** Display whether the system is in a safe state and, if so, print a safe sequence of processes.

```
need[i][j] = max[i][j] - allocation[i][j];
        }
    }
}
bool is_safe() {
    int work[MAX RESOURCES];
    bool finish[MAX_PROCESSES];
    int safe sequence[MAX PROCESSES];
    int count = 0;
    // Initialize work and finish
    for (int j = 0; j < n resources; j++) {
        work[j] = available[j];
    for (int i = 0; i < n processes; i++) {
        finish[i] = false;
    int iter count = 0; // To prevent infinite loops in case of no safe
sequence
    // Find a process that can be executed
    while (count < n processes && iter_count < n processes * n processes) {
// Added iter count for safety
        bool found process = false;
        for (int p = 0; p < n_processes; p++) {</pre>
            if (!finish[p]) {
                bool can_execute = true;
                for (int r = 0; r < n resources; r++) {
                    if (need[p][r] > work[r]) {
                        can execute = false;
                        break;
                    }
                }
                if (can_execute) {
                    // Execute process p
                    for (int r = 0; r < n_resources; r++) {
                        work[r] += allocation[p][r];
                    finish[p] = true;
                    safe sequence[count++] = p;
                    found process = true;
                    break; // Found one, restart search for next
                }
            }
        }
        if (!found process) {
            // No process found that can be executed
            break;
        iter count++;
    }
    if (count == n processes) {
        printf("\nSystem is in a SAFE STATE.\n");
        printf("Safe Sequence: < ");</pre>
        for (int i = 0; i < n processes; i++) {
            printf("P%d ", safe sequence[i]);
        printf(">\n");
        return true;
    } else {
        printf("\nSystem is in an UNSAFE STATE.\n");
```

```
return false;
    }
}
int main() {
    printf("Enter the number of processes: ");
    scanf("%d", &n_processes);
    printf("Enter the number of resource types: ");
    scanf("%d", &n resources);
    printf("\nEnter Available resources:\n");
    for (int j = 0; j < n resources; j++) {
        printf("Resource %d: ", j);
        scanf("%d", &available[j]);
    }
    printf("\nEnter Max matrix:\n");
    for (int i = 0; i < n_processes; i++) {
        printf("Process %d: ", i);
        for (int j = 0; j < n_resources; j++) {</pre>
            scanf("%d", &max[i][j]);
    }
    printf("\nEnter Allocation matrix:\n");
    for (int i = 0; i < n processes; i++) {
        printf("Process %d: ", i);
        for (int j = 0; j < n_resources; j++) {
            scanf("%d", &allocation[i][j]);
    }
    calculate need();
    printf("\nNeed Matrix:\n");
    for (int i = 0; i < n_processes; i++) {</pre>
        printf("P%d: ", i);
        for (int j = 0; j < n_resources; j++) {
            printf("%d ", need[i][j]);
        printf("\n");
    }
    is safe();
    return 0;
}
Input
Enter the number of processes: 5
Enter the number of resource types: 3
Enter Available resources:
Resource 0: 3
Resource 1: 3
Resource 2: 2
Enter Max matrix:
Process 0: 7 5 3
Process 1: 3 2 2
Process 2: 9 0 2
Process 3: 2 2 2
Process 4: 4 3 3
Enter Allocation matrix:
```

```
Process 0: 0 1 0
Process 1: 2 0 0
Process 2: 3 0 2
Process 3: 2 1 1
Process 4: 0 0 2
```

Expected Output

```
Need Matrix:
P0: 7 4 3
P1: 1 2 2
P2: 6 0 0
P3: 0 1 1
P4: 4 3 1

System is in a SAFE STATE.
```

Safe Sequence: < P1 P3 P4 P0 P2 >

Lab 11: Program to implement memory allocation with pages

Title

Program to implement Memory Allocation with Pages (Paging Simulation)

Aim

To simulate a basic memory allocation scheme using paging, demonstrating how logical addresses are translated to physical addresses.

Procedure

- 1. **Paging Concept:** Understand paging: dividing logical memory into fixed-size blocks called pages and physical memory into fixed-size blocks called frames.
- 2. **Page Table:** Implement a page table data structure, which maps page numbers to frame numbers.
- 3. **Input:**
 - o Page size.
 - o Number of pages in logical address space.
 - o Mapping of logical pages to physical frames (the page table entries).
 - A logical address to translate.
- 4. Address Translation Logic:
 - o Given a logical address, calculate the page number and offset within the page.
 - Use the page table to find the corresponding frame number for the calculated page number.
 - o Calculate the physical address: Physical Address = (Frame Number * Page Size) + Offset.
- 5. **Output:** Display the page number, offset, frame number, and the resulting physical address. Handle cases where the logical address is invalid (e.g., page not found, address out of bounds).

```
// Example: C program for Paging Simulation
#include <stdio.h>
#include <stdlib.h>
#define MAX PAGES 10
#define MAX FRAMES 10
// Page Table Entry: maps logical page to physical frame
struct PageTableEntry {
   int page number;
   int frame number;
   int valid bit; // 1 if valid, 0 if invalid/not in memory
};
int main() {
   int page size;
   int num logical pages;
   int num physical frames;
   struct PageTableEntry page table[MAX PAGES];
   printf("Enter page size (e.g., 4096 bytes): ");
    scanf("%d", &page size);
```

```
printf("Enter number of logical pages: ");
    scanf("%d", &num logical pages);
    printf("Enter number of physical frames: ");
    scanf("%d", &num physical frames);
    // Initialize page table (assume no pages are mapped initially)
    for (int i = 0; i < num logical pages; i++) {</pre>
        page table[i].page number = i;
        page table[i].frame number = -1; // -1 indicates not mapped
       page table[i].valid bit = 0;
    printf("\nEnter page to frame mappings (Logical Page Number Physical
Frame Number, -1 -1 to stop):\n");
    int lp, pf;
    while (1) {
        printf("Mapping (LP PF): ");
        scanf("%d %d", &lp, &pf);
        if (lp == -1 \&\& pf == -1) {
            break;
        if (lp >= 0 && lp < num logical pages && pf >= 0 && pf <
num physical frames) {
            page table[lp].frame number = pf;
            page_table[lp].valid_bit = 1;
        } else {
            printf("Invalid page or frame number. Please re-enter.\n");
        }
    }
    printf("\n--- Page Table ---\n");
    printf("Page No.\tFrame No.\tValid Bit\n");
    for (int i = 0; i < num logical pages; i++) {</pre>
        printf("%d\t\t%d\n", page table[i].page number,
page_table[i].frame_number, page_table[i].valid_bit);
    int logical address;
    printf("\nEnter a logical address to translate (-1 to exit): ");
    while (scanf("%d", &logical_address) == 1 && logical_address != -1) {
        if (logical address < 0 || logical address >= (num logical pages *
page size)) {
           printf("Error: Logical address out of bounds.\n");
        } else {
            int page number = logical address / page size;
            int offset = logical address % page size;
            if (page number >= num logical pages ||
page table[page number].valid bit == 0) {
                printf("Error: Page %d is not in memory or invalid.\n",
page number);
            } else {
                int frame number = page table[page number].frame number;
                int physical address = (frame number * page size) + offset;
                printf("Logical Address: %d\n", logical address);
                printf(" Page Number: %d\n", page number);
                printf(" Offset: %d\n", offset);
                printf(" Frame Number: %d\n", frame number);
                printf("Physical Address: %d\n", physical address);
        printf("\nEnter a logical address to translate (-1 to exit): ");
    }
```

```
}
Input
Enter page size (e.g., 4096 bytes): 100
Enter number of logical pages: 5
Enter number of physical frames: 3
Enter page to frame mappings (Logical Page Number Physical Frame Number, -1 -
1 to stop):
Mapping (LP PF): 0 2
Mapping (LP PF): 1 0
Mapping (LP PF): 3 1
Mapping (LP PF): -1 -1
Enter a logical address to translate (-1 to exit): 50
Enter a logical address to translate (-1 to exit): 120
Enter a logical address to translate (-1 to exit): 380
Enter a logical address to translate (-1 to exit): 450
Enter a logical address to translate (-1 to exit): 250
Enter a logical address to translate (-1 to exit): -1
Expected Output
--- Page Table ---
Page No.
               Frame No.
                               Valid Bit
                2
               0
1
                                1
2
                -1
                                0
3
                1
                                1
4
                -1
                                0
Enter a logical address to translate (-1 to exit): 50
Logical Address: 50
  Page Number: 0
  Offset: 50
  Frame Number: 2
Physical Address: 250
Enter a logical address to translate (-1 to exit): 120
Logical Address: 120
  Page Number: 1
  Offset: 20
  Frame Number: 0
Physical Address: 20
Enter a logical address to translate (-1 to exit): 380
Logical Address: 380
  Page Number: 3
  Offset: 80
  Frame Number: 1
Physical Address: 180
Enter a logical address to translate (-1 to exit): 450
Error: Page 4 is not in memory or invalid.
Enter a logical address to translate (-1 to exit): 250
Error: Page 2 is not in memory or invalid.
Enter a logical address to translate (-1 to exit): -1
```

return 0;

Lab 12: Simulation of FIFO page replacement algorithm

Title

Simulation of FIFO (First-In, First-Out) Page Replacement Algorithm

Aim

To simulate the First-In, First-Out (FIFO) page replacement algorithm and calculate the number of page faults for a given page reference string and a fixed number of frames.

Procedure

1. **FIFO Concept:** Understand FIFO: when a page fault occurs and the memory is full, the oldest page (the one that has been in memory the longest) is replaced.

2. Data Structures:

- o An array or list to represent the main memory frames.
- A queue to keep track of the order in which pages entered the frames (for FIFO replacement).

3. Input:

- o Number of available frames in physical memory.
- o A page reference string (sequence of page numbers).

4. Simulation Logic:

- o Iterate through the page reference string.
- o For each page:
 - Check if the page is already in a frame (page hit). If yes, continue.
 - If not (page fault):
 - Increment page fault count.
 - If frames are available, add the page to an empty frame and enqueue it.
 - If frames are full, dequeue the oldest page from the queue, replace it in the corresponding frame with the new page, and enqueue the new page.
- 5. **Output:** Display the state of frames at each step, indicating page hits/faults, and the total number of page faults.

```
// Example: C program for FIFO Page Replacement Algorithm Simulation
#include <stdio.h>
#include <stdbool.h>
#include <string.h> // For memset
#define MAX PAGES 20
#define MAX FRAMES 10
int main() {
   int page reference string[MAX PAGES];
   int num pages;
   int num frames;
   int frames[MAX FRAMES]; // Represents physical memory frames
   int page faults = 0;
   int front = 0; // For FIFO queue logic (index of oldest page in frames)
   int rear = -1; // For FIFO queue logic (index where new page would be
added)
   int current frame count = 0; // How many frames are currently occupied
```

```
printf("Enter the number of frames: ");
   scanf("%d", &num frames);
   printf("Enter the number of pages in the reference string: ");
   scanf("%d", &num pages);
   printf("Enter the page reference string (space separated): ");
   for (int i = 0; i < num pages; <math>i++) {
       scanf("%d", &page reference string[i]);
   // Initialize frames with -1 (empty)
   memset(frames, -1, sizeof(frames));
   printf("\nPage Reference String\tFrames\t\t\tHit/Fault\n");
   printf("-----
\n");
   for (int i = 0; i < num pages; <math>i++) {
        int current page = page reference string[i];
       bool is hit = false;
       // Check if page is already in frames (page hit)
        for (int j = 0; j < num\_frames; j++) {
            if (frames[j] == current page) {
               is hit = true;
               break;
           }
       printf("%d\t\t", current page);
       if (is hit) {
           printf(" (Hit)\n");
        } else { // Page Fault
           page_faults++;
           printf(" (Fault)\n");
           if (current frame count < num frames) {</pre>
               // Frames are not full, add to an empty frame
               frames[rear + 1] = current_page;
               rear = (rear + 1); // Update rear for next insertion
               current frame count++;
            } else {
               // Frames are full, replace the oldest page (FIFO)
               frames[front] = current page; // Replace the page at 'front'
               front = (front + 1) % num frames; // Move front to next
oldest
               rear = (rear + 1) % num frames; // Move rear to the newly
inserted position
           }
       }
       // Print current state of frames
       printf("\t\t\t");
       for (int j = 0; j < num frames; <math>j++) {
           if (frames[j] == -1) {
               printf(" - ");
           } else {
               printf(" %d ", frames[j]);
       printf("\n");
   }
```

```
printf("\nTotal Page Faults: %d\n", page_faults);
return 0;
}
```

```
Enter the number of frames: 3   
Enter the number of pages in the reference string: 12   
Enter the page reference string (space separated): 0 1 2 3 0 1 4 0 1 2 3 4
```

Expected Output

Page Reference String	Frames	Hit/Fault
0	(Fault)	
1	0 (Fault)	
2	0 1 - (Fault) 0 1 2	
3	(Fault) 3 1 2	
0	(Fault) 3 0 2	
1	(Fault)	
4	3 0 1 (Fault)	
0	4 0 1 (Hit) 4 0 1	
1	(Hit)	
2	4 0 1 (Fault)	
3	4 2 1 (Fault)	
4	3 2 1 (Fault) 3 2 4	

Total Page Faults: 10

Lab 13: multiple partition (dynamic Memory allocation method)

Title

Multiple Partition (Dynamic Memory Allocation Method) Simulation

Aim

To simulate dynamic memory allocation using multiple partitions, specifically focusing on the First Fit, Best Fit, and Worst Fit algorithms.

Procedure

- 1. **Dynamic Allocation Concept:** Understand how dynamic partitioning allows memory to be divided into variable-sized partitions as needed by processes.
- 2. Algorithms:
 - o **First Fit:** Allocate the first hole that is big enough.
 - o **Best Fit:** Allocate the smallest hole that is big enough.
 - Worst Fit: Allocate the largest hole that is big enough.
- 3. **Data Structures:** Represent memory as a list of holes (free blocks) and allocated blocks. Each block should have a starting address and size.
- 4. Input:
 - o Initial memory partitions (free blocks).
 - o Sequence of process requests (process ID, size).
 - o Choice of allocation algorithm (First Fit, Best Fit, Worst Fit).
- 5. Allocation Logic (for each algorithm):
 - When a process requests memory:
 - Search for a suitable hole according to the chosen algorithm.
 - If a hole is found:
 - Allocate the requested size from the hole.
 - If the hole is larger than needed, split it, leaving a smaller free hole.
 - Mark the allocated block.
 - If no suitable hole is found, the process cannot be allocated.
- 6. **Deallocation Logic:** When a process finishes, its memory block becomes a free hole. Implement merging of adjacent free holes to prevent fragmentation.
- 7. **Output:** Display the memory state (allocated and free blocks) after each allocation/deallocation, and indicate if a process could not be allocated.

```
// Example: C program for Dynamic Memory Allocation Simulation (First Fit)
#include <stdio.h>
#include <stdib.h>
#include <stdbool.h>

#define MAX_BLOCKS 10

// Structure to represent a memory block (either free or allocated)
struct MemoryBlock {
   int id; // -1 for free block, process ID for allocated block
   int start_address;
   int size;
};
```

```
// Global array to simulate memory blocks
struct MemoryBlock memory[MAX BLOCKS];
int num memory blocks = 0; // Current number of blocks in memory array
// Function to display current memory state
void display memory() {
   printf("\n--- Current Memory State ---\n");
   printf("ID\tStart\tSize\tStatus\n");
   printf("----\n");
   for (int i = 0; i < num_memory_blocks; i++) {
       printf("%d\t%d\t%d\t%s\n",
              memory[i].id,
              memory[i].start address,
              memory[i].size,
               (memory[i].id == -1) ? "FREE" : "ALLOCATED (P)");
   printf("----\n");
}
// Function to merge adjacent free blocks (to reduce fragmentation)
void merge free blocks() {
    // A more robust merge would sort by start address first.
    // For simplicity, this assumes blocks are generally ordered.
    for (int i = 0; i < num_memory_blocks - 1; i++) {</pre>
        if (memory[i].id == -1 \&\& memory[i+1].id == -1) {
            // Merge memory[i+1] into memory[i]
           memory[i].size += memory[i+1].size;
            // Shift subsequent blocks to fill the gap
            for (int j = i + 1; j < num memory blocks - 1; <math>j++) {
               memory[j] = memory[j+1];
           num memory blocks--;
            i--; // Recheck the current block as it might merge again
        }
    }
}
// First Fit Allocation
void first_fit(int process_id, int request_size) {
    int allocated = -1;
    for (int i = 0; i < num memory blocks; i++) {</pre>
        if (memory[i].id == -1 && memory[i].size >= request size) {
            // Found a suitable free block
            if (memory[i].size == request size) {
               // Exact fit
               memory[i].id = process id;
            } else {
               // Split the block
                // Create a new allocated block
                for (int j = num\ memory\ blocks;\ j > i + 1;\ j--) {
                   memory[j] = memory[j-1];
               memory[i+1].id = -1; // New free block
               memory[i+1].start address = memory[i].start address +
request size;
               memory[i+1].size = memory[i].size - request size;
               // Update the original block to be allocated
               memory[i].id = process id;
               memory[i].size = request size;
               num memory blocks++;
            allocated = i;
           printf("Process P%d allocated %d units using First Fit.\n",
process id, request size);
           break;
```

```
}
    if (allocated == -1) {
        printf("Process P%d (size %d) could not be allocated (No suitable
block found).\n", process id, request size);
}
// Deallocation
void deallocate(int process id) {
    bool found = false;
    for (int i = 0; i < num memory blocks; i++) {</pre>
        if (memory[i].id == process_id) {
   memory[i].id = -1; // Mark as free
            printf("Process P%d deallocated.\n", process id);
            found = true;
            merge_free_blocks(); // Attempt to merge adjacent free blocks
            break;
        }
    if (!found) {
        printf("Process P%d not found for deallocation.\n", process id);
}
int main() {
    int total_memory_size;
    printf("Enter total memory size: ");
    scanf("%d", &total memory size);
    // Initially, all memory is one large free block
    memory[0].id = -1;
    memory[0].start address = 0;
    memory[0].size = total_memory_size;
    num memory blocks = 1;
    display_memory();
    int choice;
    int pid, size;
    while (1) {
        printf("\n1. Allocate (First Fit)\n");
        printf("2. Deallocate\n");
        printf("3. Exit\n");
        printf("Enter your choice: ");
        scanf("%d", &choice);
        switch (choice) {
            case 1:
                printf("Enter Process ID and size to allocate: ");
                scanf("%d %d", &pid, &size);
                first fit(pid, size);
                display memory();
                break;
            case 2:
                printf("Enter Process ID to deallocate: ");
                scanf("%d", &pid);
                deallocate (pid);
                display memory();
                break;
            case 3:
                printf("Exiting.\n");
                exit(0);
            default:
```

```
printf("Invalid choice. Please try again.\n");
       }
   return 0;
Input
Enter total memory size: 1000

    Allocate (First Fit)

2. Deallocate
3. Exit
Enter your choice: 1
Enter Process ID and size to allocate: 101 200
1. Allocate (First Fit)
2. Deallocate
3. Exit
Enter your choice: 1
Enter Process ID and size to allocate: 102 150

    Allocate (First Fit)

2. Deallocate
3. Exit
Enter your choice: 1
Enter Process ID and size to allocate: 103 300
1. Allocate (First Fit)
2. Deallocate
3. Exit
Enter your choice: 2
Enter Process ID to deallocate: 102
1. Allocate (First Fit)
2. Deallocate
3. Exit
Enter your choice: 1
Enter Process ID and size to allocate: 104 100
1. Allocate (First Fit)
2. Deallocate
3. Exit
Enter your choice: 3
Expected Output
Enter total memory size: 1000
--- Current Memory State ---
ID Start Size Status
    0 1000 FREE
1. Allocate (First Fit)
2. Deallocate
3. Exit
Enter your choice: 1
Enter Process ID and size to allocate: 101 200
Process P101 allocated 200 units using First Fit.
--- Current Memory State ---
ID Start Size Status
```

101	0	200	ALLOCATED	(P)
-1	200	800	FREE	

- Allocate (First Fit)
- 2. Deallocate
- 3. Exit

Enter your choice: 1

Enter Process ID and size to allocate: 102 150 Process P102 allocated 150 units using First Fit.

--- Current Memory State ---

ID	Start	Size	Status	_
101 102 -1	0 200 350	200 150 650	ALLOCATED (FALLOCATED (FFREE	•

- Allocate (First Fit)
- 2. Deallocate
- 3. Exit

Enter your choice: 1

Enter Process ID and size to allocate: 103 300 Process P103 allocated 300 units using First Fit.

--- Current Memory State ---

ID	Start	Size	Status	
101 102 103 -1	0 200 350 650	200 150 300 350	ALLOCATED ALLOCATED ALLOCATED FREE	(P)

- 1. Allocate (First Fit)
- 2. Deallocate
- 3. Exit

Enter your choice: 2

Enter Process ID to deallocate: 102

Process P102 deallocated.

--- Current Memory State ---

ID	Start	Size	Status	
101 -1 103 -1	0 200 350 650	200 150 300 350	ALLOCATED FREE ALLOCATED FREE	. ,

- 1. Allocate (First Fit)
- 2. Deallocate
- 3. Exit

Enter your choice: 1

Enter Process ID and size to allocate: 104 100 Process P104 allocated 100 units using First Fit.

--- Current Memory State ---

ID	Start	Size	Status
101	0	200	ALLOCATED (P) ALLOCATED (P) FREE ALLOCATED (P)
104	200	100	
-1	300	50	
103	350	300	

-1 650 350 FREE

1. Allocate (First Fit)

- 2. Deallocate
- 3. Exit

Enter your choice: 3

Exiting.

Lab 14: Simulation of FIFO page replacement algorithm Paging

Title

Simulation of FIFO (First-In, First-Out) Page Replacement Algorithm with Paging Context

Aim

This lab is a duplicate of Lab 12. It aims to simulate the First-In, First-Out (FIFO) page replacement algorithm and calculate the number of page faults for a given page reference string and a fixed number of frames, reinforcing the concepts from Lab 12 within the broader context of paging.

Procedure

(Refer to Lab 12 Procedure)

Source Code

(Refer to Lab 12 Source Code)

Input

(Refer to Lab 12 Input)

Expected Output

(Refer to Lab 12 Expected Output)

Lab 15: Simulation of optimal page replacement algorithm

Title

Simulation of Optimal Page Replacement Algorithm

Aim

To simulate the Optimal (OPT) page replacement algorithm and calculate the number of page faults for a given page reference string and a fixed number of frames.

Procedure

- 1. **Optimal Concept:** Understand Optimal: when a page fault occurs and the memory is full, the page that will not be used for the longest period of time in the future is replaced. This is a theoretical algorithm as it requires future knowledge.
- 2. Data Structures:
 - o An array or list to represent the main memory frames.
- 3. Input:
 - o Number of available frames in physical memory.
 - o A page reference string (sequence of page numbers).
- 4. Simulation Logic:

- o Iterate through the page reference string.
- o For each page:
 - Check if the page is already in a frame (page hit). If yes, continue.
 - If not (page fault):
 - Increment page fault count.
 - If frames are available, add the page to an empty frame.
 - If frames are full, determine which page in the current frames will not be used for the longest time in the *future* of the reference string. Replace that page with the new page.
- 5. **Output:** Display the state of frames at each step, indicating page hits/faults, and the total number of page faults.

```
// Example: C program for Optimal Page Replacement Algorithm Simulation
#include <stdio.h>
#include <stdbool.h>
#include <limits.h> // For INT MAX
#define MAX PAGES 20
#define MAX FRAMES 10
// Function to find the optimal page to replace
int find_optimal_page(int page_reference_string[], int num_pages, int
frames[], int num_frames, int current_index) {
    int farthest index = current index;
    int page to \overline{replace} = -1;
    for (int i = 0; i < num frames; <math>i++) {
        int found future use = -1; // Index of next use in future
        // Find when frames[i] will be used next
        for (j = current index + 1; j < num pages; j++) {</pre>
            if (frames[i] == page reference string[j]) {
                found future use = j;
                break;
            }
        }
        if (found future use == -1) {
            // This page will not be used again, so it's the best to replace
            return i;
            // This page will be used, compare its next use with others
            if (found future use > farthest index) {
                farthest index = found future use;
                page to replace = i;
            }
        }
    }
    return page_to_replace;
}
int main() {
    int page reference string[MAX PAGES];
    int num pages;
    int num frames;
    int frames[MAX FRAMES]; // Represents physical memory frames
    int page faults = 0;
    int current frame count = 0; // How many frames are currently occupied
    printf("Enter the number of frames: ");
```

```
scanf("%d", &num frames);
   printf("Enter the number of pages in the reference string: ");
   scanf("%d", &num pages);
   printf("Enter the page reference string (space separated): ");
   for (int i = 0; i < num pages; i++) {</pre>
        scanf("%d", &page reference string[i]);
    }
    // Initialize frames with -1 (empty)
    for (int i = 0; i < num frames; <math>i++) {
        frames[i] = -1;
    printf("\nPage Reference String\tFrames\t\t\tHit/Fault\n");
   printf("-----
\n");
    for (int i = 0; i < num pages; i++) {</pre>
        int current page = page reference string[i];
        bool is hit = false;
        // Check if page is already in frames (page hit)
        for (int j = 0; j < num_frames; j++) {
   if (frames[j] == current_page) {</pre>
                is hit = true;
                break;
            }
        printf("%d\t\t", current page);
        if (is hit) {
            printf("(Hit)\n");
        } else { // Page Fault
            page_faults++;
            printf(" (Fault)\n");
            if (current frame count < num frames) {</pre>
                // Frames are not full, add to an empty frame
                frames[current frame count] = current page;
                current frame count++;
            } else {
                // Frames are full, replace the optimal page
                int replace index = find optimal page(page reference string,
num pages, frames, num frames, i);
                frames[replace index] = current page;
        }
        // Print current state of frames
        printf("\t\t\t");
        for (int j = 0; j < num frames; <math>j++) {
            if (frames[j] == -1) {
               printf(" - ");
            } else {
                printf(" %d ", frames[j]);
       printf("\n");
   printf("\nTotal Page Faults: %d\n", page faults);
   return 0;
```

}

Input

```
Enter the number of frames: 3   
Enter the number of pages in the reference string: 12   
Enter the page reference string (space separated): 0 1 2 3 0 1 4 0 1 2 3 4
```

Expected Output

Page Reference String	Frames	Hit/Fault
0	(Fault)	
	0	
1	(Fault)	
_	0 1 -	
2	(Fault)	
	0 1 2	
3	(Fault)	
	0 1 3	
0	(Hit)	
	0 1 3	
1	(Hit)	
	0 1 3	
4	(Fault)	
	0 4 3	
0	(Hit)	
	0 4 3	
1	(Fault)	
	0 4 1	
2	(Fault)	
	2 4 1	
3	(Fault)	
	2 3 1	
4	(Hit)	
	2 3 4	

Total Page Faults: 8