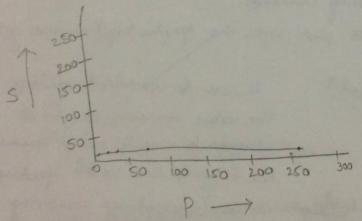
Problem 5.5: - when considering the example 5.1, the standard speedup for adding 'n' numbers using 'n' processing elements is  $S = \frac{O(n)}{T_0}$  Since  $T_p = O(\log n)$  :  $S = O(\frac{n}{\log n})$ 

Speedup S is defined on the ratio of secial runtime of the best ceruntial algorithm for solving a problem to the timetaken by the parallel algorithm to solve the same problem on it processing elements. The processing elements and by the parallel algorithm are assumed to be identical to the one parallel algorithm are assumed to be identical to the one and by the securitial algorithm.

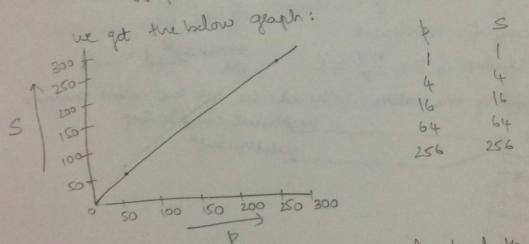
Below is the graph obtained to standard gredup:



The value of scaled speedup errains constant though the number of processors is invested. Though there are many processors but while doing summation in figure 5.2 in text processors but while doing summation in figure 5.2 in text processors but while doing summation in figure 5.2 in text processors but while doing summation in figure 5.2 in text processors but while doing summation in figure 5.2 in text processors but while doing summation in figure 5.2 in text processors by the summing 16 doments there are a total of 4 sommunion steps to be considered, the speedup does not depend on the num steps to be considered, the speedup does not depend on the num of procurous, but its only dependent on ruember of procuring elements Guiven scaled speedup = +W w: base problem size the a single procuring when plotted graph for p from 1 to 256, below is the Gudup where Tp = (1/2-1)+11/69/ curve stotained: 200 The above value get a linear graph. In this can, scold gudup is the spedup that can At calculated when the problem is it is increased and we runned from yours. comparision, lits plot both the glaphs together, we get In care of standard gudup, the value remains constant where in can of scaled speedup, gre value of the oreall gudup and increased by investing 100 the number or processors. 50 100 150 200 250 300 For Standard Epidup, the spendup does not depend on the number of procurous and its value remains constant though the number of processors is increased whereas the value of scaled speedup can se increared by increasing the number of processors. problem 6.6: In this case the modern size is scaled up according to bookfriency duration of plogby where p is trumber it proving elements. To = (np-1) + 11 logb

we also have (so efficient Scaled speedup = \$\frac{1}{12}(\rho\nu)\frac\

of promes (b)



From the above graph, we can understand that the squally with the number of processors. Spendup increases linearly with the number of processors. Using (societaient turntion, we can make the scaled squally as the number of processors squally to increase linearly as the number of processors increases.

problem 5.8: - Though the number of processing elements are increased, if the total workload is not increased then the speedup does not increase himself with the number of processing elements and when the efficiency is calculated which is the settle of guiday to the number of processor. The efficiency decreases apparatus from the problems 5.5 and 5.7; it can be obscured apparatusly. From the problems 5.5 and 5.7; it can be obscured that the scaled speedup absents increase linearly but it has that the scaled speedup absents however him the himser graph. So, Saled speedup absents when the problem of devening efficiency through the number of problem of devening efficiency through the number of problem is devening an invested. The scaled speedup came is linear only if the isosefficiency densition of a parallel septem is linear.