

Q. Describe an $O(n)$ - time algorithm that, given a set S of n distinct numbers and a positive integer $k \leq n$, determine the k numbers in S that are closest to the median of S .

To compute the k numbers near the median of a set S of n distinct numbers, it is possible to combine methods for finding medians and methods for partitioning into an $O(n)$ -time algorithm.

1. **Find Median:** Using the median of medians algorithm, find the median of the set S in $O(n)$ time. The median is the element at position $n/2$ (or $(n-1)/2$, if n is odd) when the set is sorted.
2. **Find the Absolute Differences:** Find the difference between each element of the set S and its median, in absolute value. Now, you will have a set D of absolute differences.
3. **Find the k -th Smallest Absolute Difference:** The quickselect algorithm is used to find the k -th smallest absolute difference from D by doing this in linear time of $O(n)$.
4. **Partition the Set:** Having found the k -th smallest absolute difference, partition the set S in such a way that the k elements with the least absolute differences can be selected. This can easily be done using a simple linear scan in $O(n)$ time.
5. **Returning Result:** That will give us the k elements closest to the median, the ones that have the smallest absolute differences to the median.

Example

```
int S[] = {12, 3, 17, 13, 26, 9, 5, 8, 15, 10, 7};
int n = sizeof(S) / sizeof(S[0]);
int k = 4; // Find 4 numbers closest to the median

findKClosestToMedian(S, n, k);
```

Output

```
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Median: 10
The 4 numbers closest to the median are:
12 9 8 10

...Program finished with exit code 0
Press ENTER to exit console.
```