Artificial Intelligence

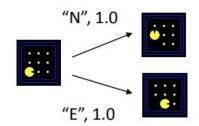
Lec 2: Uninformed Search

Pratik Mazumder

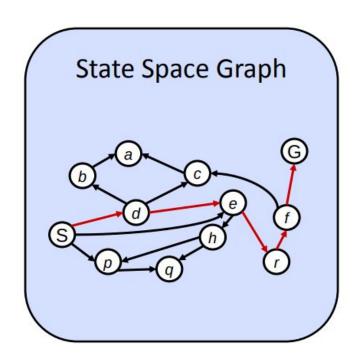
Search Problems

- Framework for solving problems
- A search problem consists of:
 - A state space: List of all possible states, i.e., all possible configurations of elements/conditions in the world or environment, e.g., passing one pen across the class, situations in a pacman game.
 - A successor function (with actions, costs):
 Action may also change the state.
 - A start state
 - A **goal test**: There can be multiple goal states
- A solution is a sequence of actions (a plan) which transforms the start state to a goal state



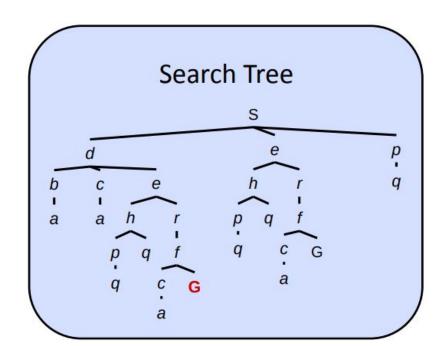


State Space Graphs vs. Search Trees

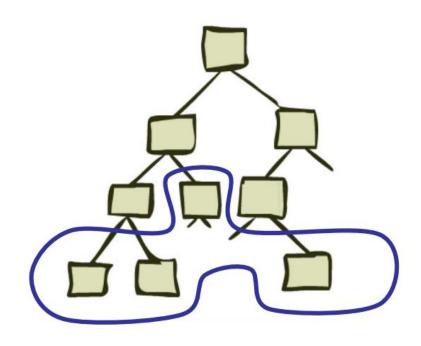


Each NODE in the search tree is an entire PATH in the state space graph.

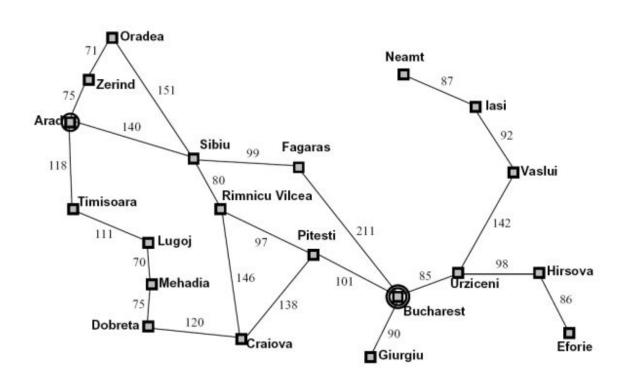
We construct both on demand – and we construct as little as possible.



Tree Search

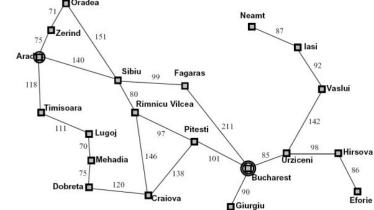


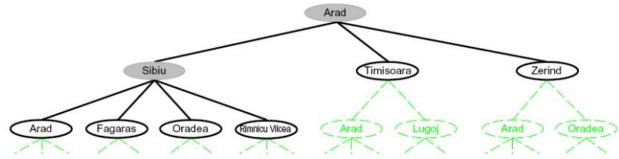
Searching with a Search Tree: Search Problem Example



Searching with a Search Tree: Search Problem Example

- Search (for goal) using a Search Tree:
 - Expand out potential plans (tree nodes)
 - Maintain a fringe of partial plans under consideration
 - Try to expand as few tree nodes as possible





General Tree Search

```
function TREE-SEARCH( problem, strategy) returns a solution, or failure initialize the search tree using the initial state of problem loop do

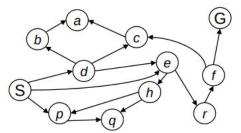
if there are no candidates for expansion then return failure choose a leaf node for expansion according to strategy

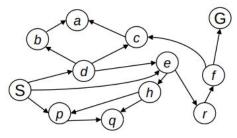
if the node contains a goal state then return the corresponding solution else expand the node and add the resulting nodes to the search tree end
```

Important ideas:

- Fringe (List of partial plans)
- Expansion
- Exploration strategy

Main question: which fringe nodes to explore?

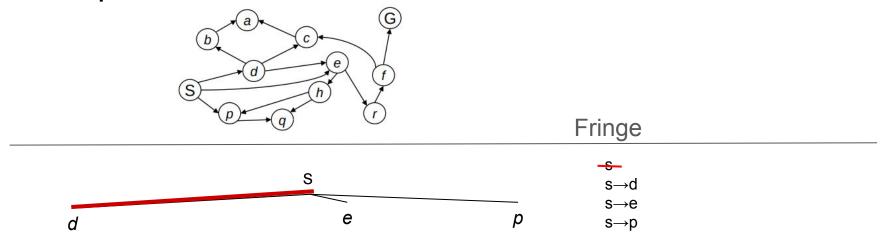


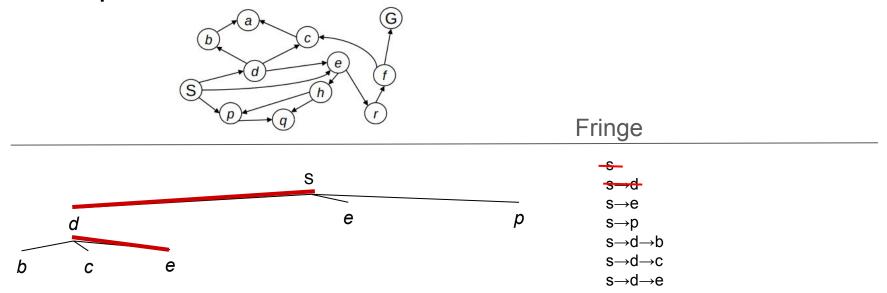


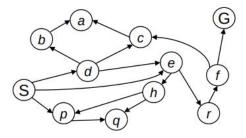
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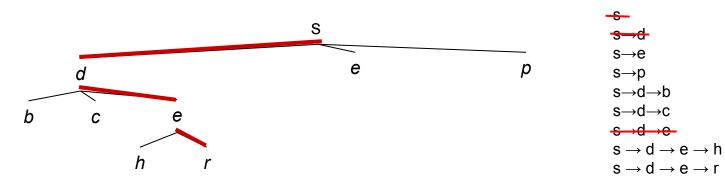
Fringe

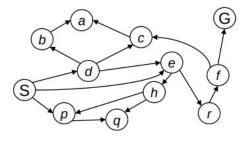
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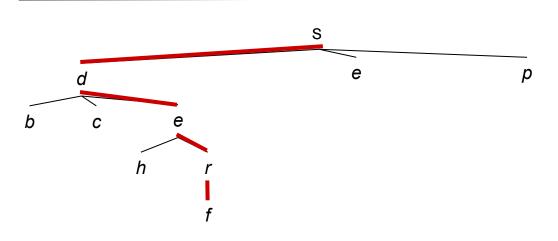


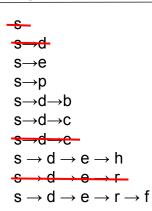


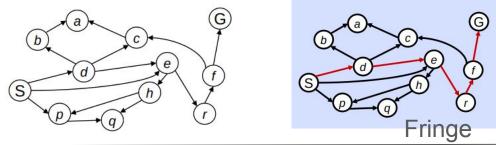


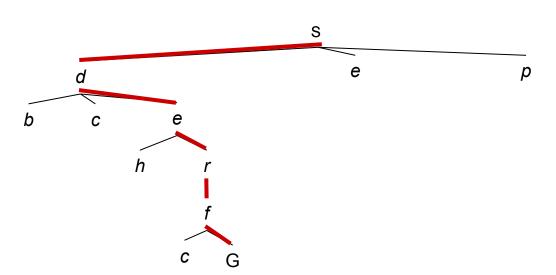


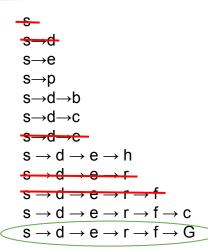


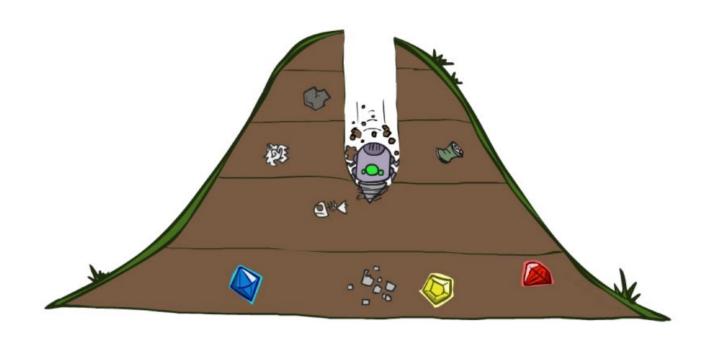


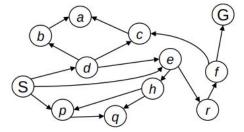




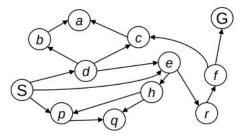








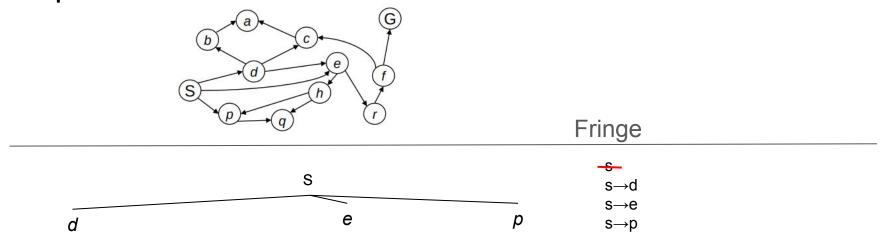
Strategy: expand a **deepest node first**

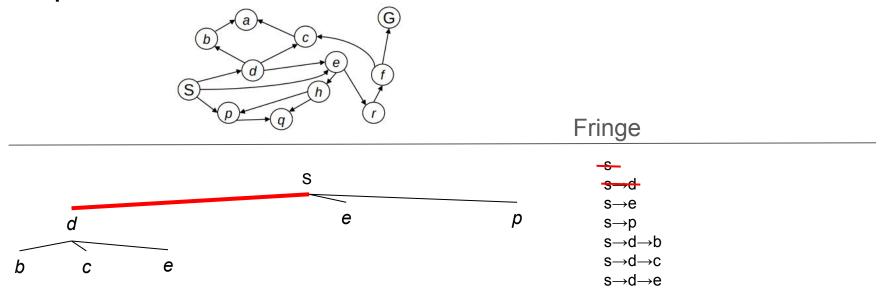


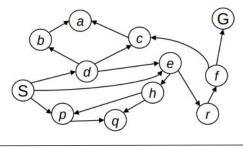
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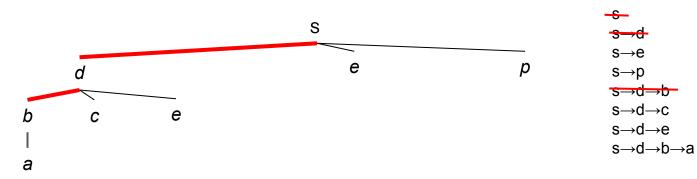
Fringe

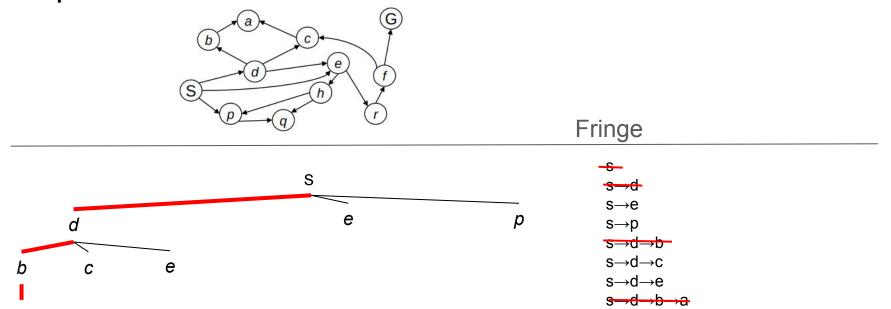
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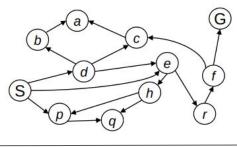


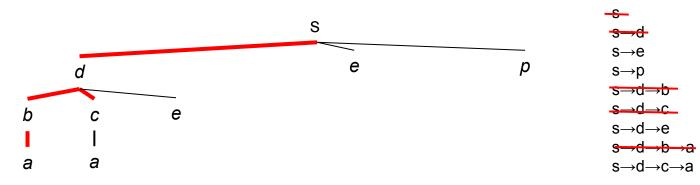


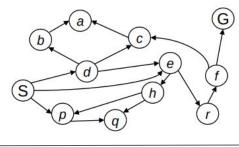


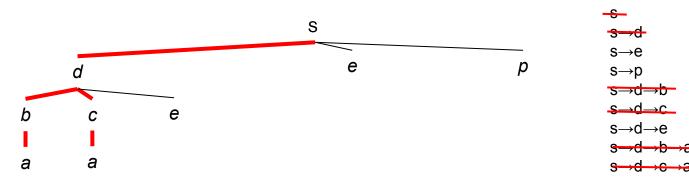


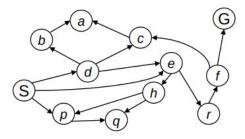


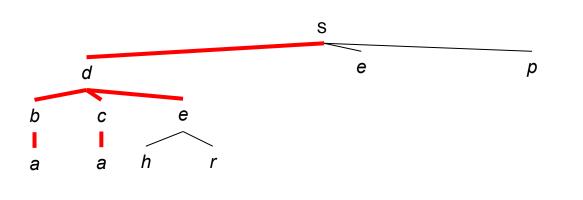


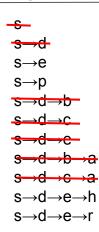


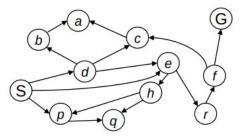


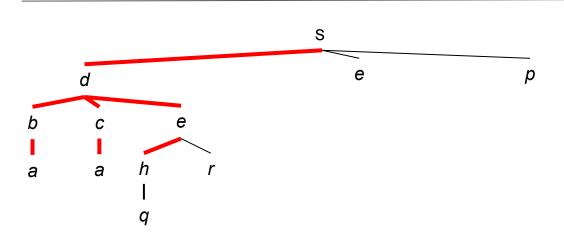


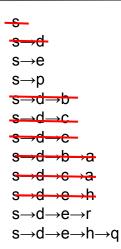


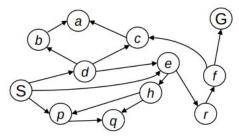


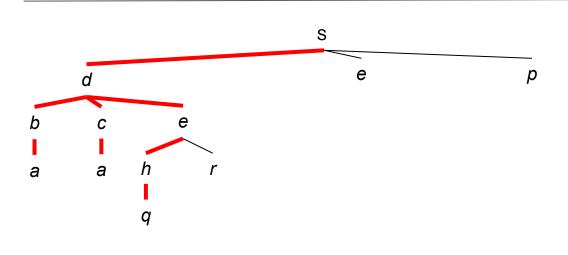


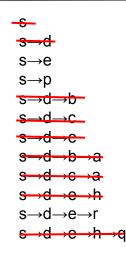


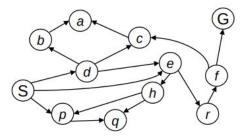


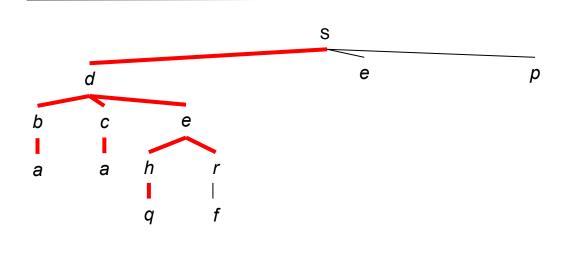


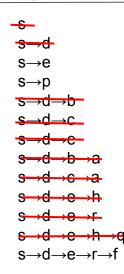


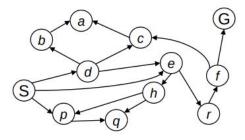


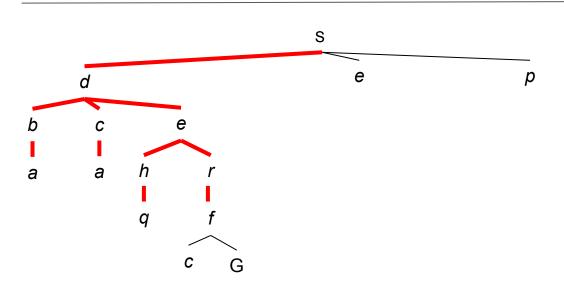


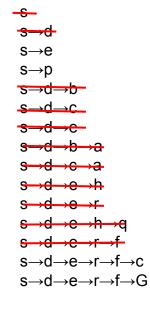


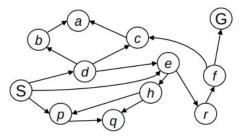


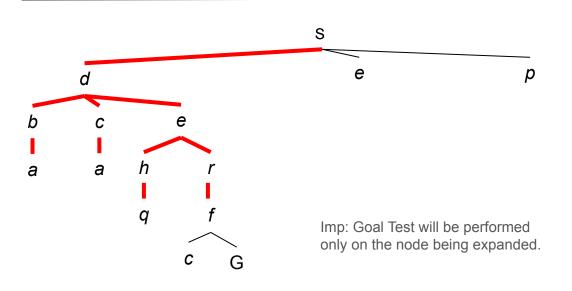


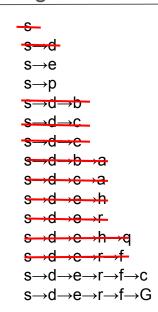


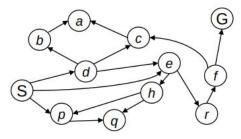


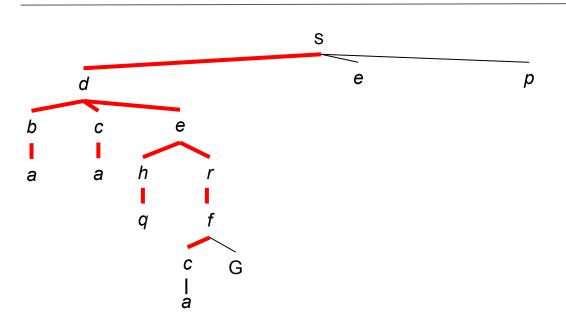


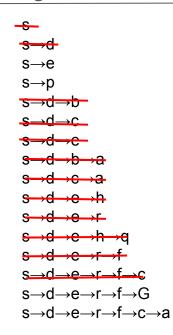


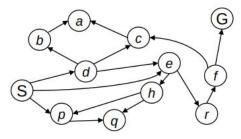


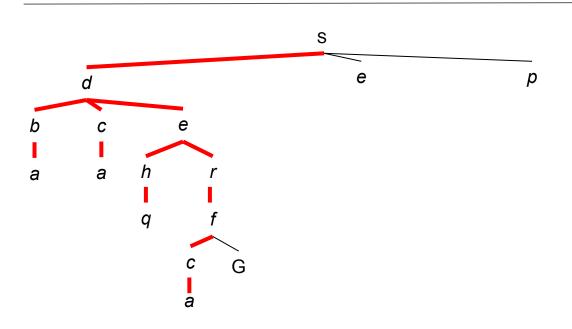


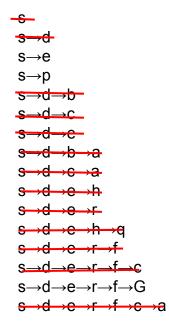


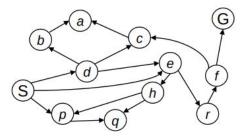


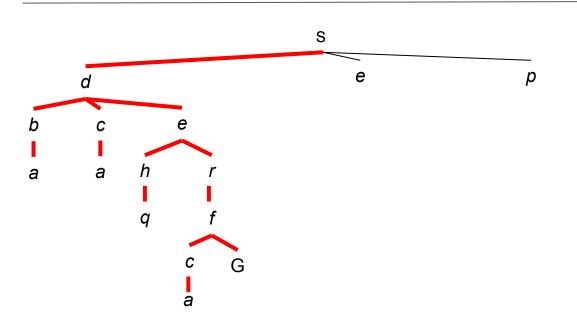


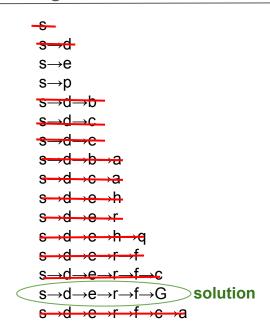


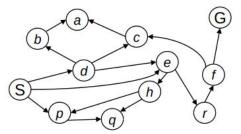








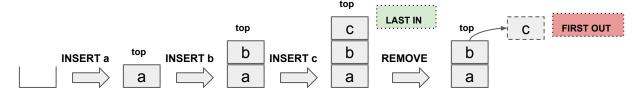




Strategy: expand a deepest node first

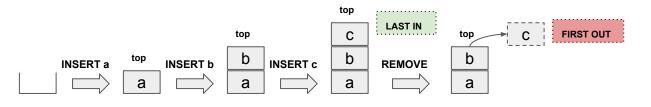
Implementation: List of partial plans/fringe is a LIFO stack

- Stack: Data structure or in simple terms a type of storage that follows the Last In First Out order.
 - LIFO last entry in the stack is the first to be removed from the stack



Aside: Stack Implementation

- Stack: Data structure or in simple terms a type of storage that follows the Last In First Out order.
 - LIFO last entry in the stack is the first to be removed from the stack



fringe = FringeUsingStack()

fringe.insert('a') fringe.insert('b') fringe.insert('c') fringe.insert('d')

print(fringe.remove()) # Output?????

```
class FringeUsingStack:
      def init (self):
             self.lt=[]
      def insert(self,i):
             self.lt.append(i)
      def remove(self):
             el = -999
             ln = len(self.lt)
             if(ln \ge 2):
                    el = self.lt[0]
                    self.lt = self.lt[1:]
             elif(ln==1):
                    el = self.lt[0]
                    self.lt = []
             else:
                    el = -999
             return el
```

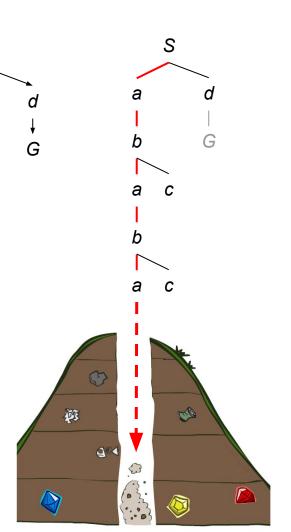
Depth First Search: Properties

Is it complete? i.e. Is it guaranteed to find a solution if one exists?

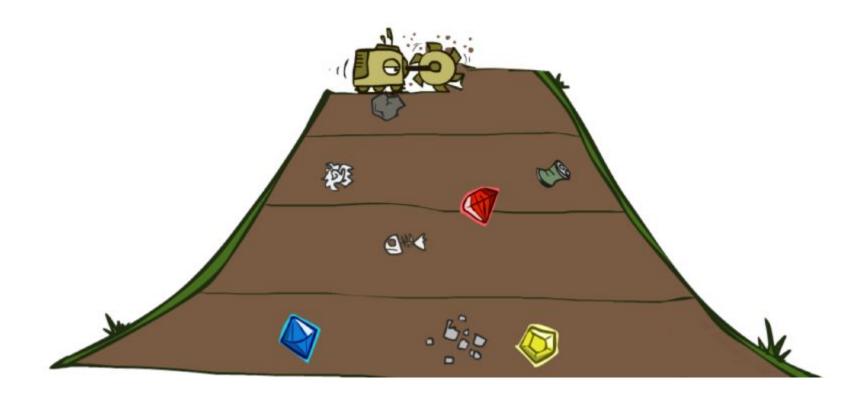
- No, not necessarily
 - If no cycles and finite no. of nodes then Yes

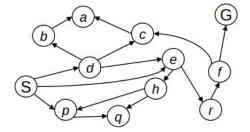
Is it optimal? i.e. Guaranteed to find the least cost path?

- No.
 - Algorithm makes no attempt to reduce number of actions/nodes in the solution

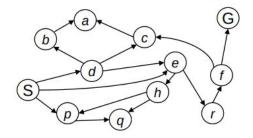


Breadth First Search





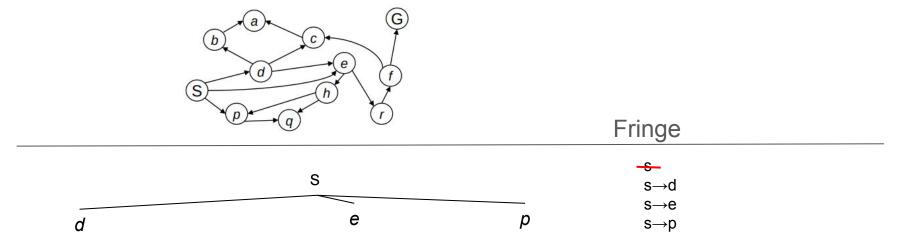
Strategy: expand a **shallowest node first**

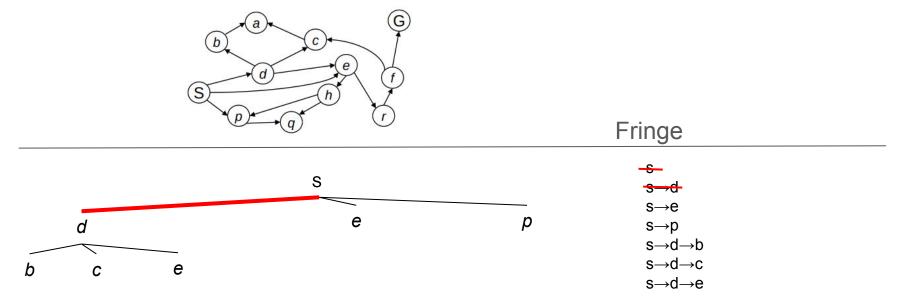


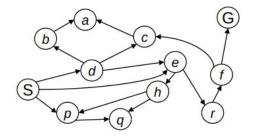
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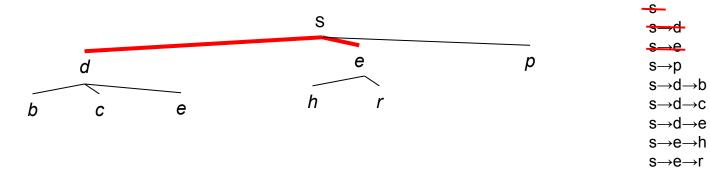
Fringe

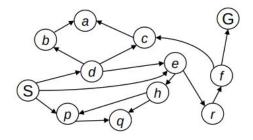
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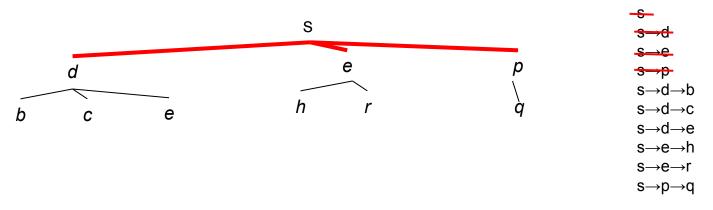


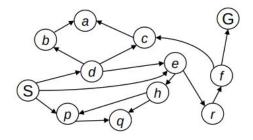


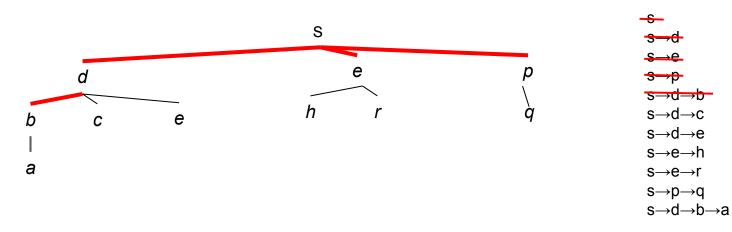


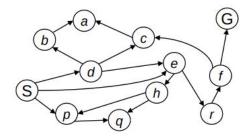


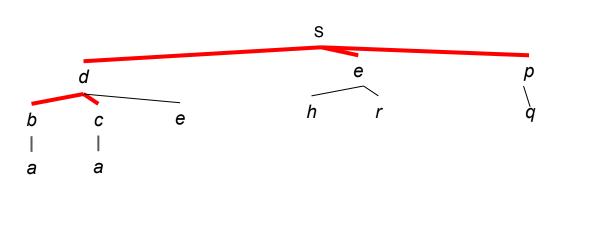


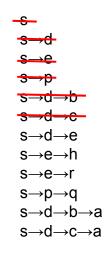


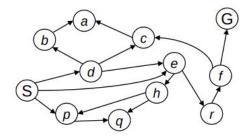




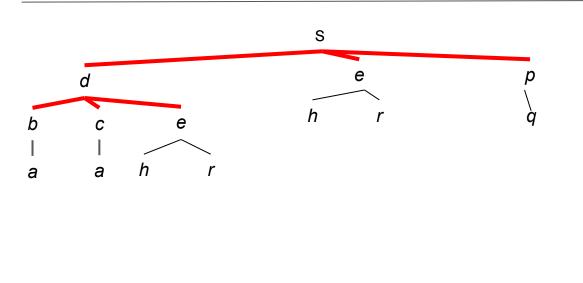




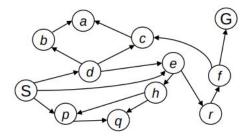


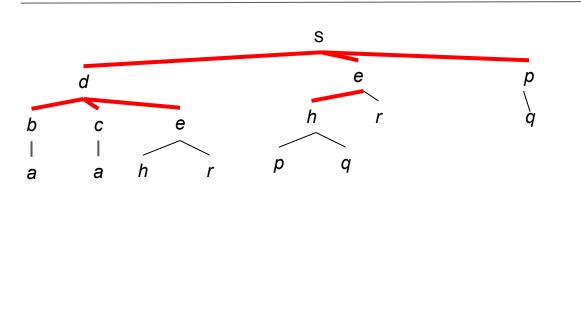


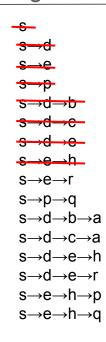
Fringe

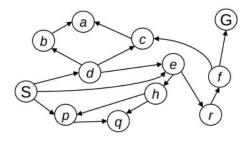


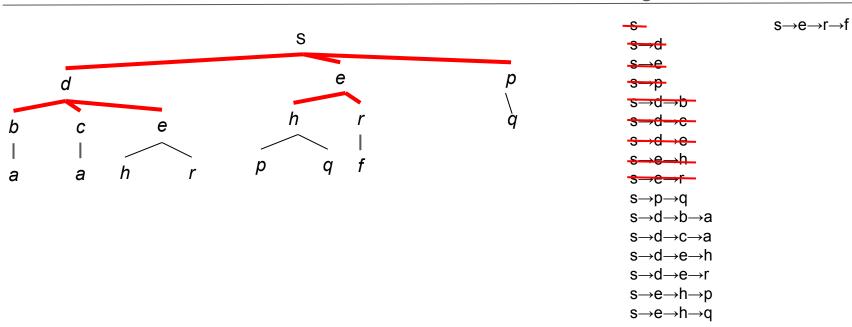
s→d s→e s→p s→d→b s→d→c s→d→c s→e→h s→e→r s→p→q s→d→b→a s→d→c→a s→d→c→a s→d→e→h s→d→e→r

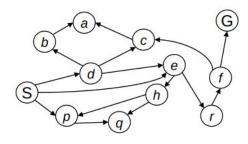




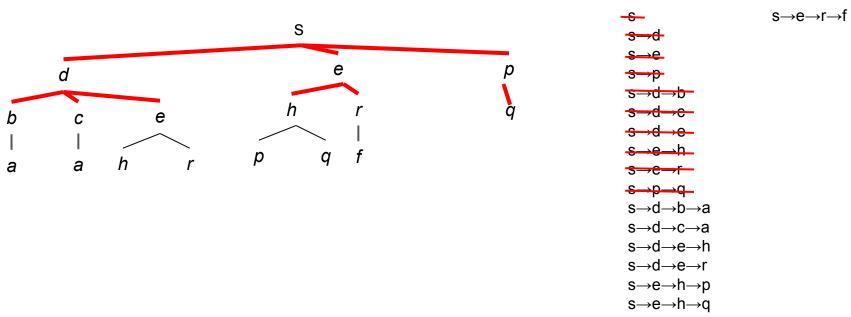


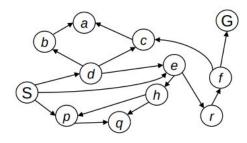


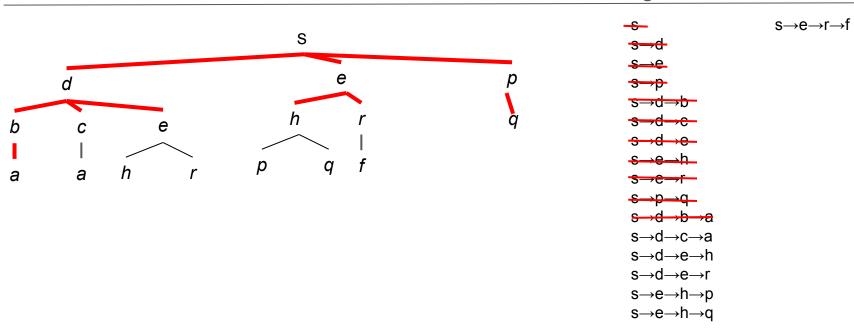


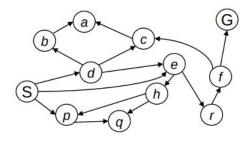


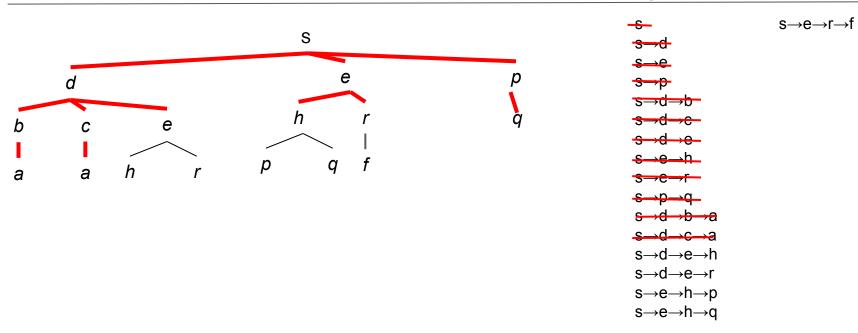


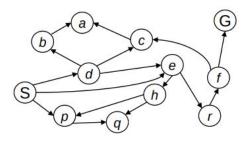








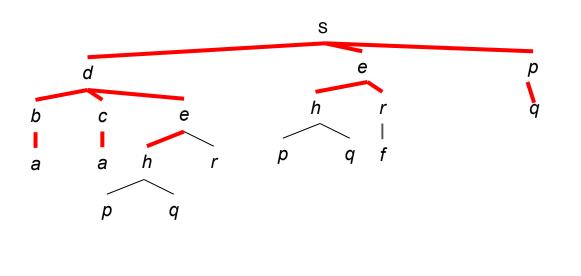




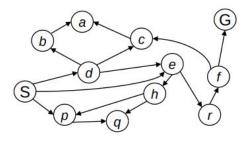
Fringe

s → d → c → a

 $s \rightarrow d \rightarrow e \rightarrow r$ $s \rightarrow e \rightarrow h \rightarrow p$ $s \rightarrow e \rightarrow h \rightarrow q$

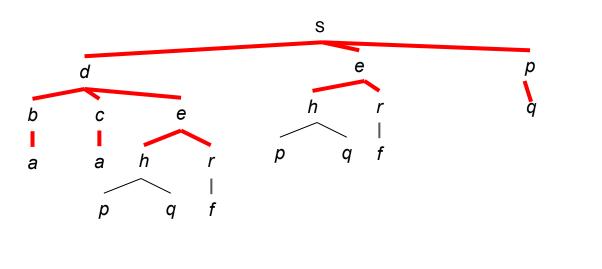


}_	$s \rightarrow e \rightarrow r \rightarrow f$
s →d	$s\rightarrow d\rightarrow e\rightarrow h\rightarrow p$
<u>e</u>	$s \rightarrow d \rightarrow e \rightarrow h \rightarrow e$
	

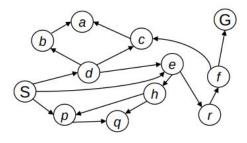


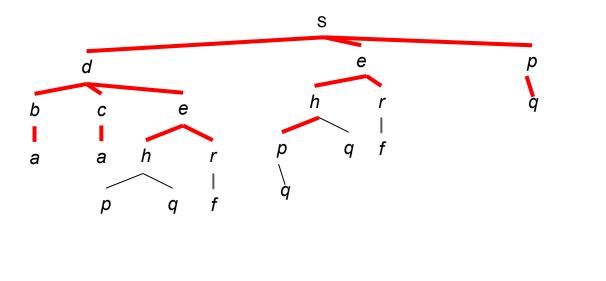
Fringe

 $s \rightarrow d \rightarrow e \rightarrow r$ $s \rightarrow e \rightarrow h \rightarrow p$ $s \rightarrow e \rightarrow h \rightarrow q$



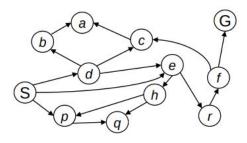
_	$s \rightarrow e \rightarrow r \rightarrow f$
->d	$s\rightarrow d\rightarrow e\rightarrow h\rightarrow p$
⊸e	$s \rightarrow d \rightarrow e \rightarrow h \rightarrow c$
→p	$s \rightarrow d \rightarrow e \rightarrow r \rightarrow f$
$\rightarrow d \rightarrow b$	

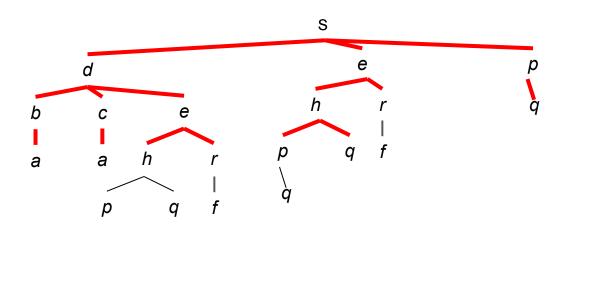




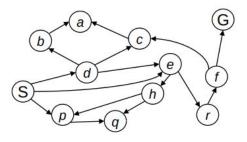
-S-
s→d
s →e
3 →p
s→d→b
s →d →c
s →d →e
s→e→ h_
s→e→r
s→p→q
s →d →b →a
s→d→c→a
s →d →e →h
sder
s→e→h→p
s→e→h→q

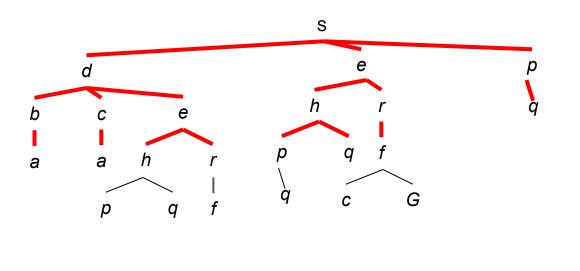
$$\begin{array}{c} s{\rightarrow}e{\rightarrow}r{\rightarrow}f\\ s{\rightarrow}d{\rightarrow}e{\rightarrow}h{\rightarrow}p\\ s{\rightarrow}d{\rightarrow}e{\rightarrow}h{\rightarrow}q\\ s{\rightarrow}d{\rightarrow}e{\rightarrow}r{\rightarrow}f\\ s{\rightarrow}e{\rightarrow}h{\rightarrow}p{\rightarrow}q\\ \end{array}$$

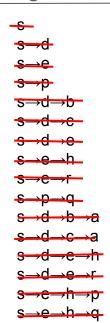


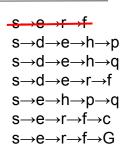


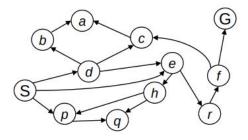
-	$s\rightarrow e\rightarrow r\rightarrow f$
->d	$s\rightarrow d\rightarrow e\rightarrow h\rightarrow$
<u>e</u>	$s\rightarrow d\rightarrow e\rightarrow h\rightarrow$
-→p	s→d→e→r→f
→d→b	$s\rightarrow e\rightarrow h\rightarrow p\rightarrow$
→d→c	
de-	
a .h	



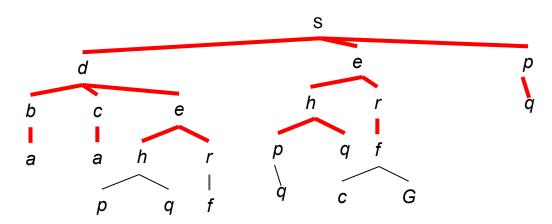




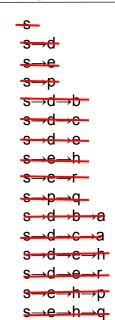


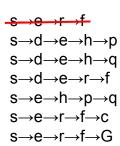


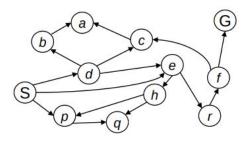
Fringe

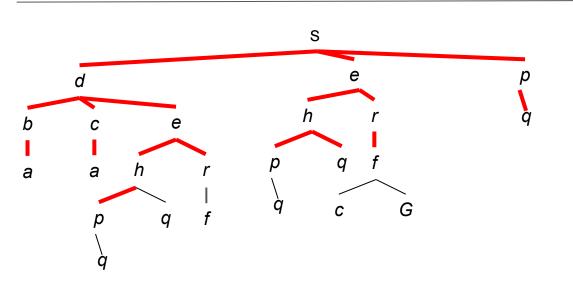


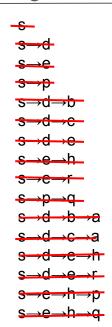
Imp: Goal Test will be performed only on the node being expanded.

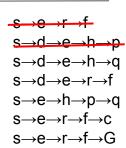


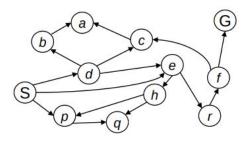


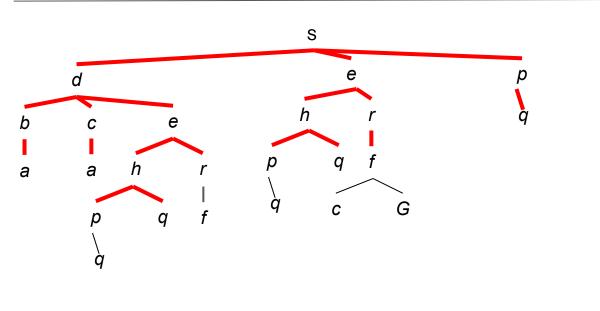


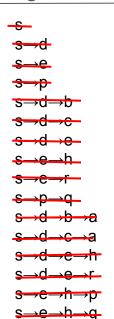


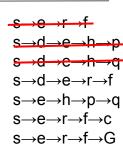


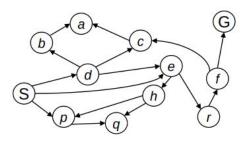


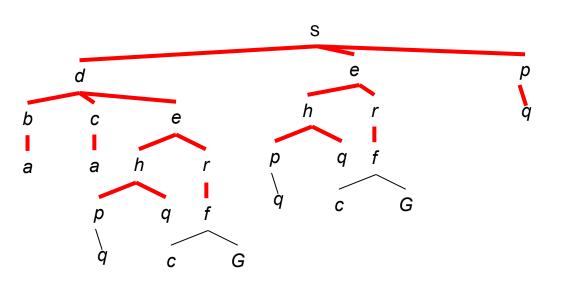


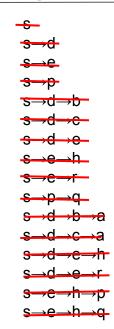


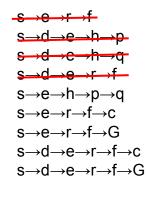


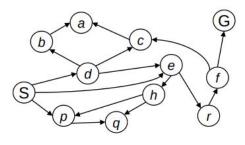


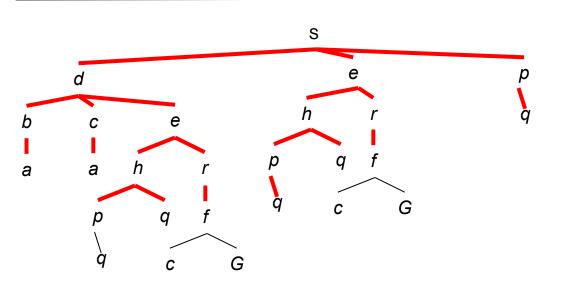




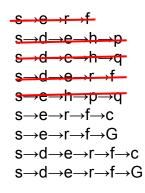


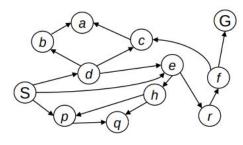


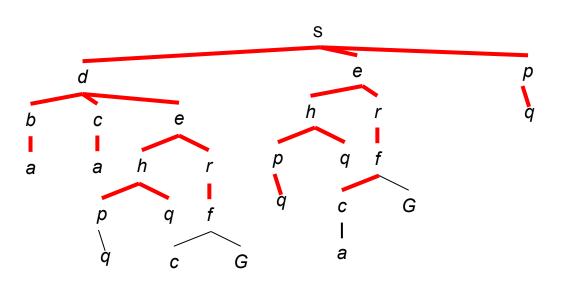


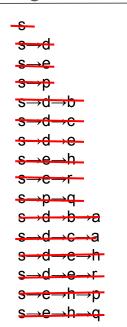


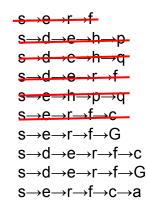


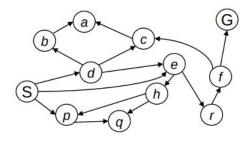


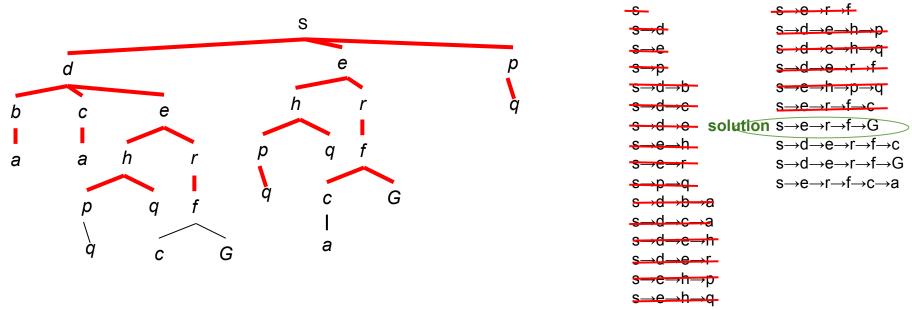


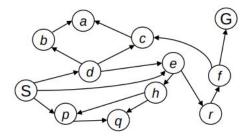


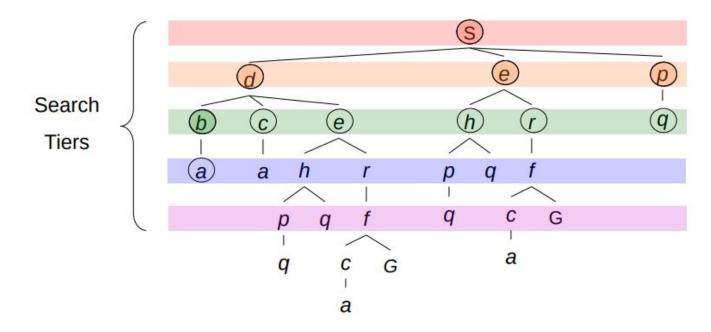






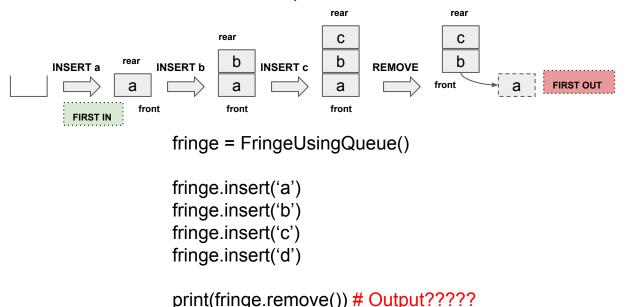






Aside: Queue Implementation

- Queue: Data structure or in simple terms a type of storage that follows the First In First Out order.
 - FIFO first entry in the queue is the first to be removed from the queue



```
class FringeUsingQueue:
      def init (self):
             self.lt=[]
      def insert(self,i):
             self.lt.append(i)
      def remove(self):
             el = -1
             ln = len(self.lt)
             if(ln \ge 2):
                   el = self.lt[0]
                    self.lt = self.lt[1:]
             elif(ln==1):
                   el = self.lt[0]
                    self.lt = []
             else:
                    el=-1
             return el
```

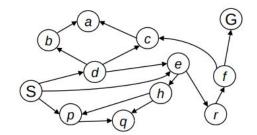
Breadth First Search: Properties

Is it complete? i.e. Is it guaranteed to find a solution if one exists?

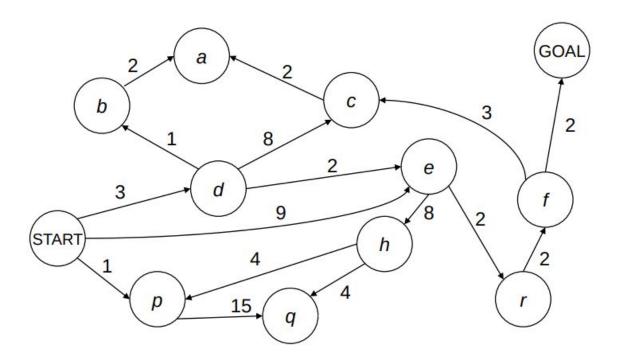
Yes. Will find the shallowest solution

Is it optimal? i.e. Guaranteed to find the least cost path?

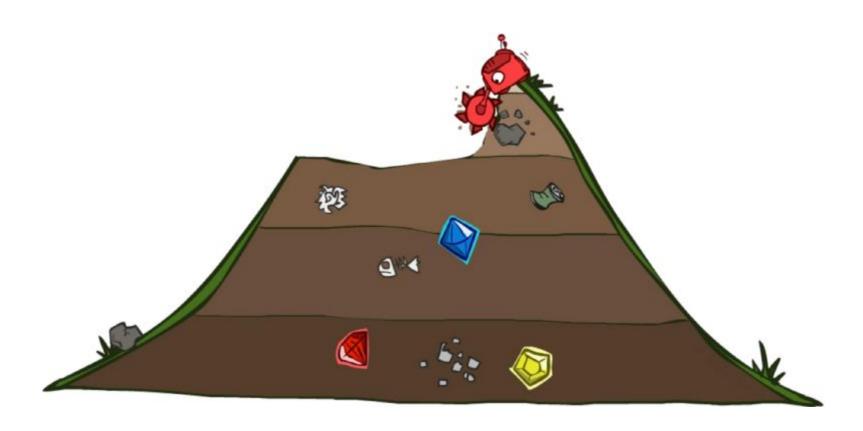
- Yes. Only if costs are all 1. Best in terms of number of actions
 - No attempt to search least cost action first

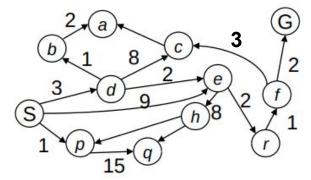


Cost Sensitive Search

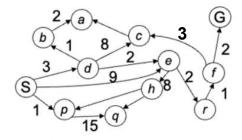


BFS finds the shortest path in terms of number of actions. It does not find the least-cost path. We will now cover a similar algorithm which does find the least-cost path.





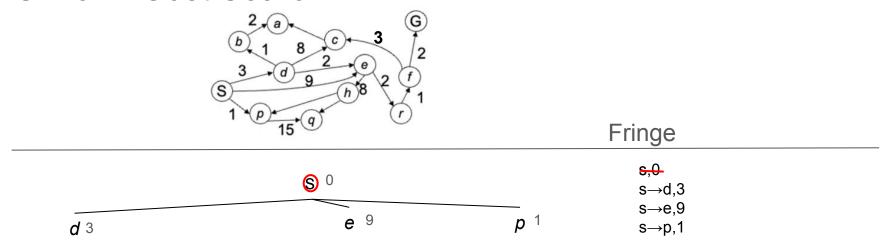
Strategy: expand a cheapest node first

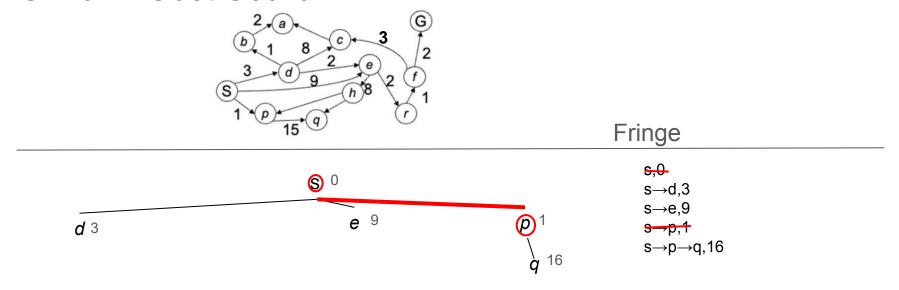


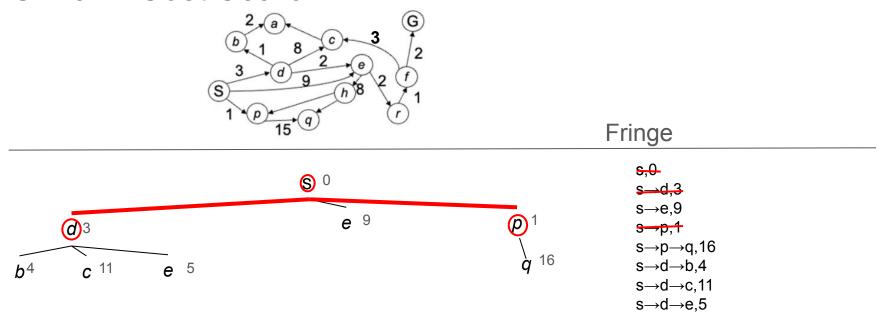
Fringe

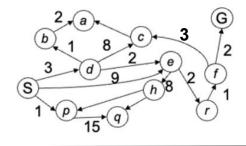
s,0

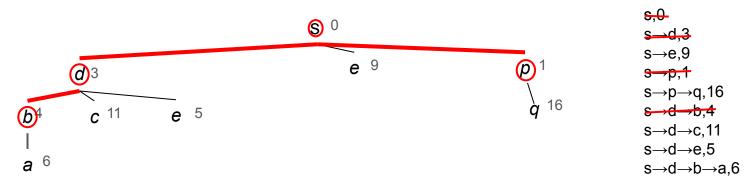
S

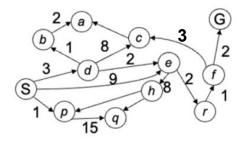




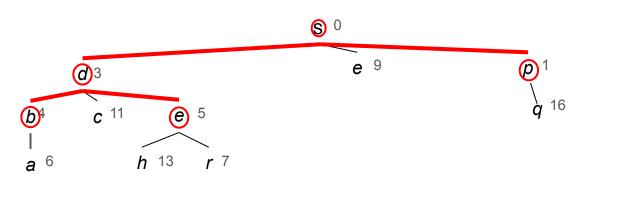












```
s,0

s\rightarrow d,3

s\rightarrow e,9

s\rightarrow p\rightarrow q,16

s\rightarrow d\rightarrow b,4

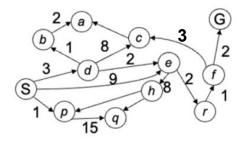
s\rightarrow d\rightarrow c,11

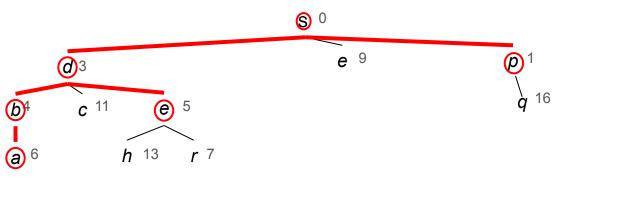
s\rightarrow d\rightarrow e,5

s\rightarrow d\rightarrow b\rightarrow a,6

s\rightarrow d\rightarrow e\rightarrow h,13

s\rightarrow d\rightarrow e\rightarrow r,7
```





```
s,0

s\rightarrow d,3

s\rightarrow e,9

s\rightarrow p,1

s\rightarrow p\rightarrow q,16

s\rightarrow d\rightarrow b,4

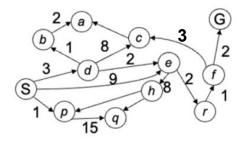
s\rightarrow d\rightarrow c,11

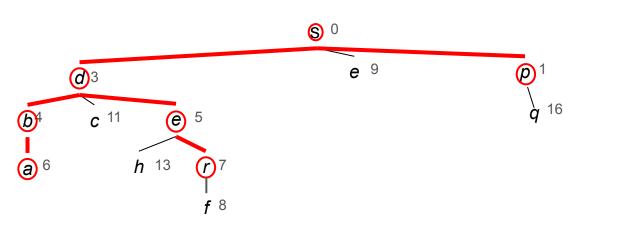
s\rightarrow d\rightarrow e,5

s\rightarrow d\rightarrow b\rightarrow a,6

s\rightarrow d\rightarrow e\rightarrow h,13

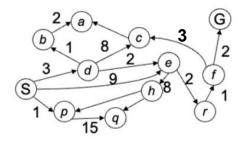
s\rightarrow d\rightarrow e\rightarrow r,7
```

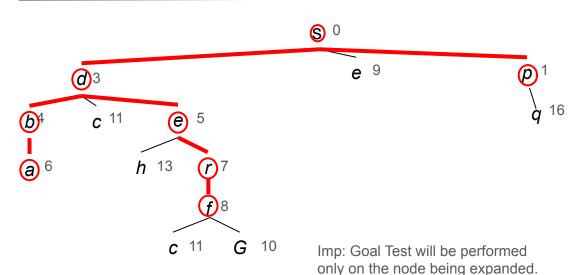




$$s,0$$

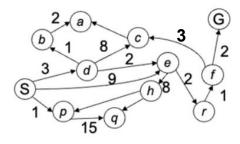
 $s\rightarrow d,3$
 $s\rightarrow e,9$
 $s\rightarrow p,1$
 $s\rightarrow p\rightarrow q,16$
 $s\rightarrow d\rightarrow b,4$
 $s\rightarrow d\rightarrow c,11$
 $s\rightarrow d\rightarrow c,5$
 $s\rightarrow d\rightarrow e\rightarrow h,13$
 $s\rightarrow d\rightarrow e\rightarrow r,7$
 $s\rightarrow d\rightarrow e\rightarrow r\rightarrow f,8$

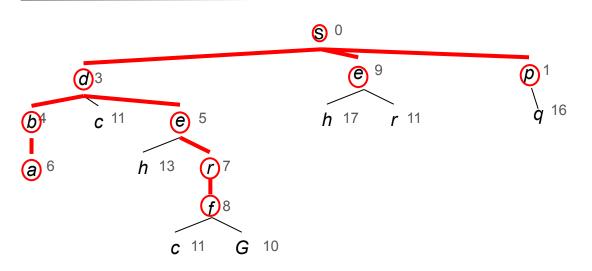




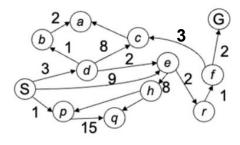
$$s,0$$

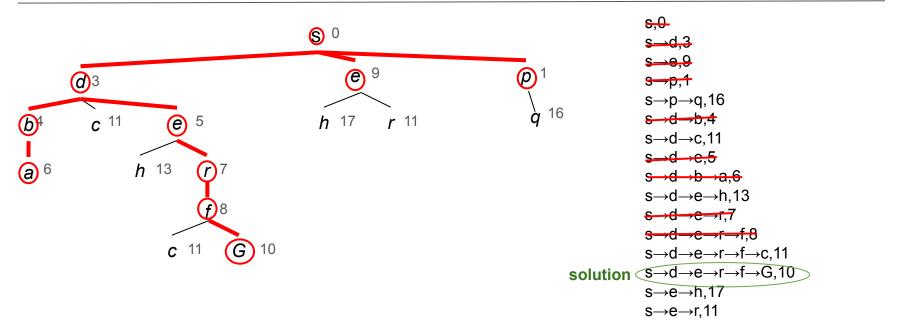
 $s\rightarrow d,3$
 $s\rightarrow e,9$
 $s\rightarrow p\rightarrow q,16$
 $s\rightarrow d\rightarrow b,4$
 $s\rightarrow d\rightarrow c,11$
 $s\rightarrow d\rightarrow e,5$
 $s\rightarrow d\rightarrow e\rightarrow h,13$
 $s\rightarrow d\rightarrow e\rightarrow r\rightarrow f,8$
 $s\rightarrow d\rightarrow e\rightarrow r\rightarrow f\rightarrow c,11$
 $s\rightarrow d\rightarrow e\rightarrow r\rightarrow f\rightarrow G,10$

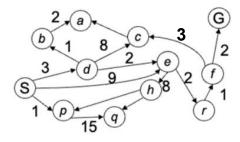


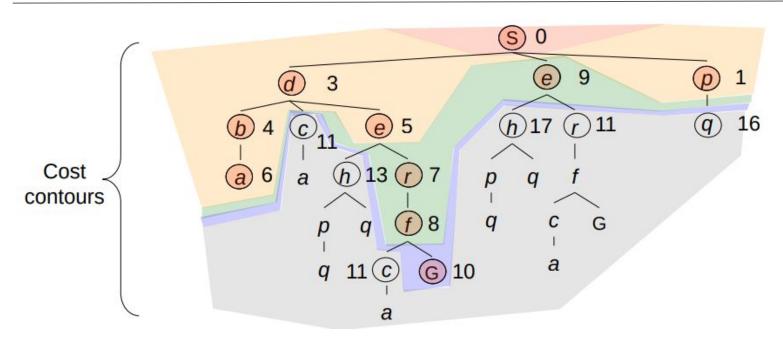


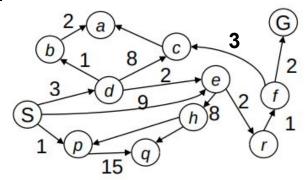
```
<del>s,0</del>
s<u>→d,3</u>
s→e,9
s→p,1
s \rightarrow p \rightarrow q,16
s \rightarrow d \rightarrow b, 4
s \rightarrow d \rightarrow c,11
s \rightarrow d \rightarrow e, 5
s \rightarrow d \rightarrow b \rightarrow a,6
s \rightarrow d \rightarrow e \rightarrow h,13
s \rightarrow d \rightarrow c \rightarrow r,7
s \rightarrow d \rightarrow c \rightarrow r \rightarrow f, \theta
s \rightarrow d \rightarrow e \rightarrow r \rightarrow f \rightarrow c, 11
s \rightarrow d \rightarrow e \rightarrow r \rightarrow f \rightarrow G,10
s\rightarrow e\rightarrow h,17
s\rightarrow e\rightarrow r,11
```











Strategy: expand a cheapest node first

Implementation: List of partial plans/fringe is a priority queue (priority: cumulative cost)

 Priority Queue: Data structure or in simple terms a type of storage from which you can get items in terms of priority irrespective of the order of insertion

Aside: Priority Queue Implementation

from queue import PriorityQueue

 Priority Queue: Data structure or in simple terms a type of storage from which you can get items in terms of priority irrespective of the order of insertion

```
class FringeUsingPriorityQueue:
      def init (self):
            self.lt=PriorityQueue()
      def insert(self,i,cost):
            self.lt.put((cost,i))
      def delete(self):
            if(self.lt.empty()):
                  el = -1
            else:
                   el = self.lt.get()
```

```
fringe = FringeUsingPriorityQueue()

fringe.insert(('a',10))
fringe.insert(('b',15))
fringe.insert(('c',5))
fringe.insert(('d',11))
```

For Background, you can look into Heap data structure and Priority Queue using Heap

return el

print(fringe.remove()) # Output?????

Uniform Cost Search: Properties

Is it complete? i.e. Is it guaranteed to find a solution if one exists?

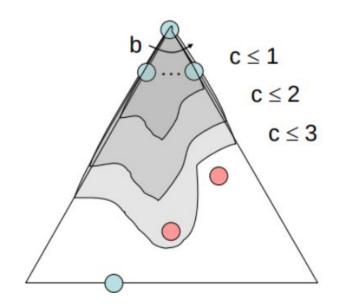
 Yes, assuming best solution has a finite cost and minimum arc cost is positive.

Is it optimal? i.e. Guaranteed to find the least cost path?

 Yes, since it is similar to BFS but traversing cost-wise instead of level-wise

The bad aspect:

- Explores options in every "direction"
- No information about goal location



Practice

Which chooses the shorter plan? BFS, DFS, UCS



Ques3

8	9	10(G)
4	11	7
3	5	6
2	1	S