

(iii)

+, * distribute

$$a * (b + c) = (a * b) + (a * c)$$

6-Feb-25

Recall: Group: $(G, *)$

Necessary

(i) closure: $\forall a, b \in G, a * b \in G$

(ii) Associativity: $\forall a, b, c \in G,$
 $a * (b * c) = (a * b) * c$

(iii) Identity: $\exists e \in G$ s.t.
 $\forall a \in G$
 $e * a = a * e = a$

(iv) Inverse: $\forall a \in G$
 $\exists b \in G$
s.t. $a * b = e$

Additional: (v) Commutativity

$$\forall a, b \in G, a * b = b * a$$

when it is satisfied for a group G ,
 G is called an Abelian group

Consider:

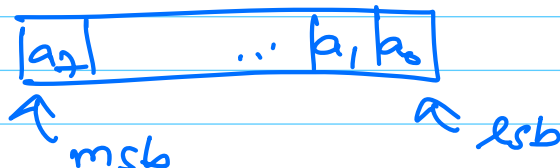
$$G = \{ P(x) \text{ of degree } \leq 7 \}$$

where coefficients
 $\in GF(2)$

$$G = \{ (a_0 + a_1x + a_2x^2 + \dots + a_7x^7) \mid a_0, a_1, \dots, a_7 \in \{0, 1\} \}$$

Representation:

1 byte



\neq

$$a(x) \neq b(x)$$

$$= a(x) + b(x) \pmod{2}$$

for each coeff

$$P_1(x) = x + 1$$

$$P_2(x) = x^2 + x$$

\Rightarrow

$$P_1(x) + P_2(x) = x^2 + 1$$

Notice that this creates a group.

\rightarrow Abelian group

What if we wanted multiplication of polynomials?

(terms mod 2)

$$P_1(x) \cdot P_2(x) \pmod{\text{degree 8}}$$

poly.

$P(x)$

$$(x^8 + \dots)$$

Consider

$$\mathbb{Z}_p$$

$$\mathbb{Z}_n = \{0, 1, 2, \dots, (n-1)\}$$

$*$ = multip. mod n

Is it a group?

if n is prime then already proven

what if n is composite?

irreducible polynomial

$$\mathbb{Z}_6 - \{0\} = \{1, 2, 3, 4, 5\} \pmod{6}$$

$$2 \times 3 = 0 \rightarrow \text{Not closed}$$

$$\mathbb{Z}_n^* = \{ x : 1 \leq x < n \text{ s.t. } \gcd(x, n) = 1 \}$$

$$\mathbb{Z}_6^* = \{ 1, 5 \} \pmod{6}$$

$$\mathbb{Z}_{10}^* = \{ 1, 3, 7, 9 \}$$

$$\mathbb{Z}_{12}^* = \{ 1, 5, 7, 11 \}$$

$$\begin{aligned} 7^{-1} &= 3 \\ 3^{-1} &= 7 \end{aligned}$$

AES polynomial:

$$= \begin{array}{c} \begin{array}{ccc} 11 & 11 & B \\ \swarrow & \downarrow & \\ 0001 & 0001 & 1011 \end{array} \\ \begin{array}{ccc} 1 & 1 & B \\ \downarrow & \downarrow & \downarrow \end{array} \\ (x^8 + x^4 + x^3 + x + 1) \end{array}$$

Coeffs are mod (2) \equiv Coeff are in GF(2)
 \downarrow
Galois Field

Field: $(G, +, \cdot)$

- (i) $(G, +)$ is an Abelian group with identity as 0
- (ii) $(G - \{0\}, \cdot)$ is an Abelian group

(iii)

Distributive operation

$$a \cdot (b+c) = (a \cdot b) + (a \cdot c)$$

$$\forall a, b, c \in G$$

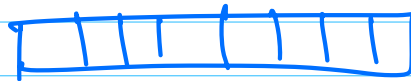
GF(2)

$\{0,1\}$ + mod 2
 \times mod 2

\oplus AND

+	0	1
0	0	1
1	1	0

\times	0	1
0	0	0
1	0	1

Rijndael finite field / AES F.F.Contains = 2^8 elements

← 8 bits → 0 or 1

 $+$ \equiv usual addition of polynomials
 with coeff $\in GF(2)$

① Every finite field of same size is isomorphic to each other

② The no. of elements in a finite field are p^n for some prime p & some int n .

Most common finite fields in Crypto are

$$GF(2^n)$$

$$\text{AES finite field} = GF(2^8)$$

Polynomial of degree ≤ 7 with coeff $\in GF(2)$
 mod 11B

S-box in AES.

8 bit \rightarrow 8 bit

$$\boxed{x} \rightarrow \text{S box} \rightarrow \left[\underset{8 \times 8}{A} \right] \bar{x}^{-1} + \left[\underset{b}{\quad} \right]$$

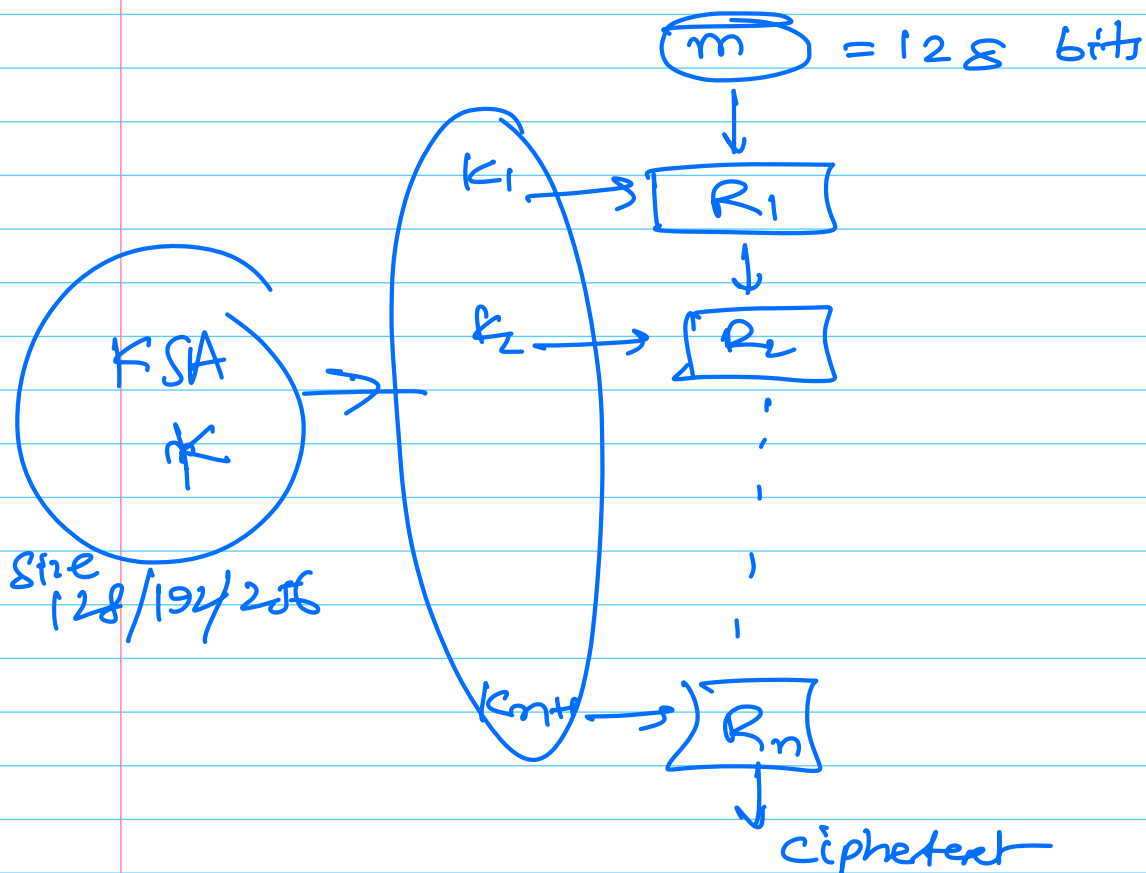
Where we take $\bar{0}^{-1}$ as

AES : Iterated block cipher

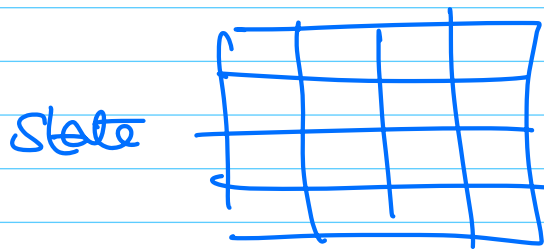
AES-128 : 10 rounds

AES-192 : 12 rounds

AES-256 : 14 rounds



Each round has 4 operations, which are implemented on a state array of 128 bits

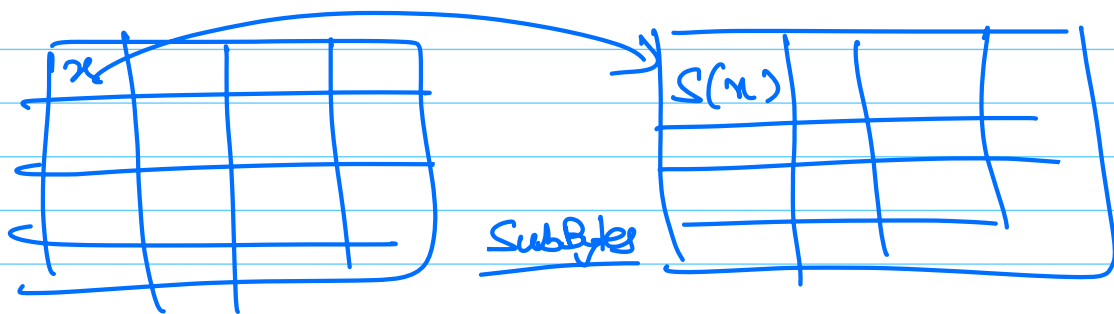


4x4 matrix of bytes

$$16 \times 8 = 128 \text{ bits}$$

Round function: Round i

1. Sub Bytes (S)
2. Shift Rows (S)
3. Mix Columns (S)
4. Add Round Keys (S, K_i)



$$S(x) = A \vec{x} + b$$

↓ ↓
some const.

Precomputed S Box size:

2⁸ inputs \times 1 byte each = 256 bytes