



LECTURE 11 – DFS & TOPOLOGICAL SORT

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Course: Algorithms & Data Structures – Fall 2025

LECTURE 11



In current lecture, we'll continue to learn & implement algorithms upon graphs, and cover other traversing algorithm Depth-First Search (DFS), and Topological Sort.

TOPICS WE'LL COVER:

Depth-First Search

Use Cases of DFS

Topological Sort

Practice Problems

GOALS FOR THIS LECTURE:

- Understand and implement Depth-First Search (DFS) using recursion and stack.
- Understand the concept of Topological Sorting and when it applies.
- Implement Topological Sort using DFS on Directed Acyclic Graphs (DAGs).

DEPTH-FIRST SEARCH

Depth-First Search

Use Cases of DFS

Topological Sort

Practice Problems

Depth-First Search explores a graph **deeply**, going as far as possible along each branch before backtracking.

Key properties:

- Uses *stack behavior* (recursion or explicit stack);
- Visits nodes in *depth order*,

Algorithm:

1. Pick a starting vertex;
2. Visit it;
3. Recursively visit all unvisited neighbors.

Time complexity: $O(V + E)$.

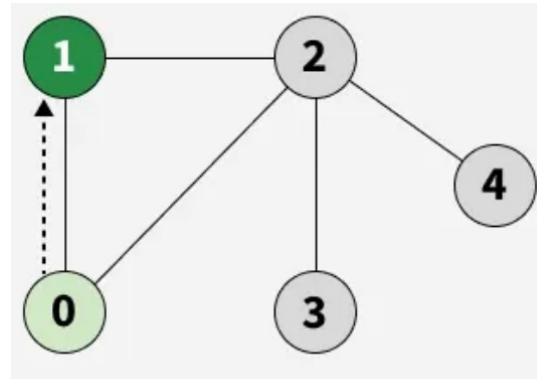


Figure 1 – Example graph [1]

visited[] =	0	1	2	3	4
	T	T	F	F	F
res[] =	0	1			

Figure 2 – Marker of visited marker

[1] <https://www.geeksforgeeks.org/dsa/depth-first-search-or-dfs-for-a-graph/>

DEPTH-FIRST SEARCH (CONT.)

We can implement DFS algorithm with a recursive way, and a stack way as shown in Figures 3, and 4, respectively.

Depth-First Search

Use Cases of DFS

Topological Sort

Practice Problems

```
void dfsRec(int v, vector<vector<int>>& adj, vector<bool>& visited) {
    if (visited[v]) return;

    visited[v] = true;
    cout << v << " ";
    for (int i=0; i < adj[v].size(); ++i) {
        int u = adj[v][i];
        if (!visited[u]) {
            dfsRec(u, adj, visited);
        }
    }
}

void dfs(int start, vector<vector<int>>& adj) {
    int n = adj.size();
    vector<bool> visited(n, false);
    dfsRec(start, adj, visited);
}
```

Figure 3 – DFS with recursion

```
void dfs(int start, vector<vector<int>>& adj) {
    int n = adj.size();
    vector<bool> visited(n, false);
    stack<int> st;

    st.push(start);

    while(!st.empty()) {
        int v = st.top();
        st.pop();

        if (!visited[v]) {
            visited[v] = true;
            cout << v << " ";
        }

        for (int i=adj[v].size()-1; i >= 0; --i) {
            int u = adj[v][i];
            if (!visited[u]) {
                st.push(u);
            }
        }
    }
    cout << endl;
}
```

Figure 4 – DFS with stack

In Figure 4, the algorithm seems like BFS, but for the place of queue uses stack.

USE CASES OF DFS

Depth-First Search

Use Cases of DFS

Topological Sort

Practice Problems

The DFS algorithm can be implemented for solving many complex algorithms, like listed below.

Use cases:

- Cycle detection;
- Topological sorting;
- Strongly connected components;
- Solving puzzles (backtracking).

Practice problem:

- [207. Course Schedule](#)

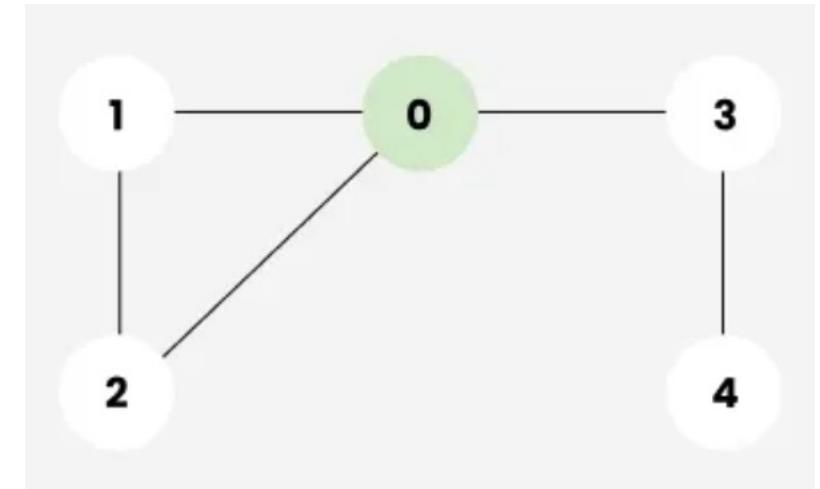


Figure 5 – Undirected graph with cycle [2]

[2] <https://www.geeksforgeeks.org/dsa/detect-cycle-undirected-graph/>

TOPOLOGICAL SORT

Depth-First Search

Use Cases of DFS

Topological Sort

Practice Problems

A **linear ordering of vertices** in a directed graph such that:

- For every directed edge $u \rightarrow v$, u appears before v in the ordering.
- ✓ Works only for Directed Acyclic Graphs (DAGs).

Graph representation for Figure 6: $\{\{1\}, \{2\}, \{\}, \{2, 4\}, \{\}\}$;

Applications:

- Ordering tasks with dependencies;
- Build systems (compilers);
- Course prerequisite problems;
- Scheduling.

Note: There may be several topological orderings for a graph.

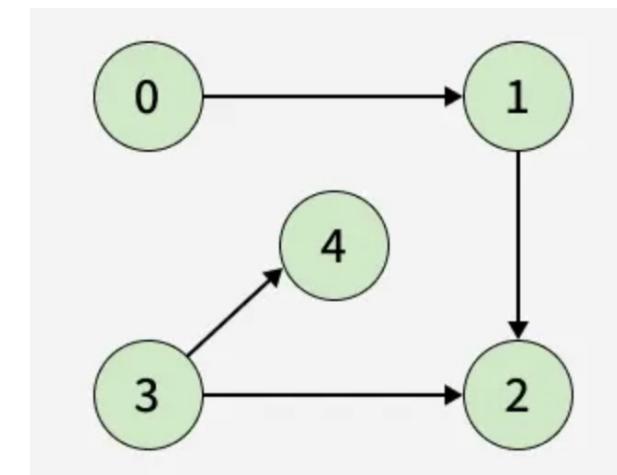


Figure 6 – Directed Acyclic Graph [3]

TOPOLOGICAL SORT (CONT.)

Depth-First Search

Use Cases of DFS

Topological Sort

Practice Problems

Idea behind Topological sort:

1. Perform DFS;
2. After exploring all neighbors of node u, push u into a stack;
3. Reverse stack → topological order.

```
void topoDFS(int v,
    vector<vector<int>>& adj,
    vector<bool>& visited,
    vector<int>& order) {

    visited[v] = true;

    for (int i=0; i < adj[v].size(); ++i) {
        int u = adj[v][i];
        if (!visited[u]) {
            topoDFS(u, adj, visited, order);
        }
    }
    order.push_back(v); // add after exploring children
}
```

Figure 7 – Topological Sort (1st function)

```
void topologicalSort(vector<vector<int>>& adj) {
    int n = adj.size();

    vector<int> order;
    vector<bool> visited(n, false);

    for (int i=0; i < n; ++i) {
        if (!visited[i]) {
            topoDFS(i, adj, visited, order);
        }
    }
    reverse(order.begin(), order.end());

    for (int i=0; i < order.size(); ++i) {
        cout << order[i] << " ";
    }
    cout << endl;
}
```

Figure 8 – Topological Sort (main function)

Time Complexity: $O(V+E)$

PRACTICE PROBLEMS

Depth-First Search

Use Cases of DFS

Topological Sort

Practice Problems

We have two tasks to solve:

- [112. Path Sum](#);
- [802. Find Eventual Safe States](#).

Q & A