



# LECTURE 10 – GRAPHS & BFS

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**Course:** Algorithms & Data Structures – Fall 2025

# LECTURE 10

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In this lecture, we'll cover the basic notations on graphs: matrix representation, adjacency & edge list. Then, we'll look up how to traverse the given graph with a help of Breadth-First Search (BFS) algorithm.

## TOPICS WE'LL COVER:

Introduction to Graphs

Matrix representation

Adjacency & Edge list

Breadth-First Search

Practice Problems

## GOALS FOR THIS LECTURE:

- Understand how to represent graphs with matrices and lists.
- Implement traversing upon the graph using BFS algorithm.
- Solving some tasks using BFS algorithm.

# INTRODUCTION TO GRAPHS

## Introduction to Graphs

## Matrix representation

## Adjacency & Edge list

## Breadth-First Search

## Practice Problems

A **graph** is a collection of *vertices* (nodes) and *edges* (connections) between them.

$$G = (V, E) \quad (1)$$

where,  $V$  – set of vertices, and  $E$  – set of edges (pair of vertices).

### Types:

- Directed / Undirected;
- Weighted / Unweighted.

### Example:

Vertices = {1, 2, 3, 4, 5, 6};

Edges = {(1, 2), (1, 3), (2, 4), (2, 6), (3, 4), (3, 5)}

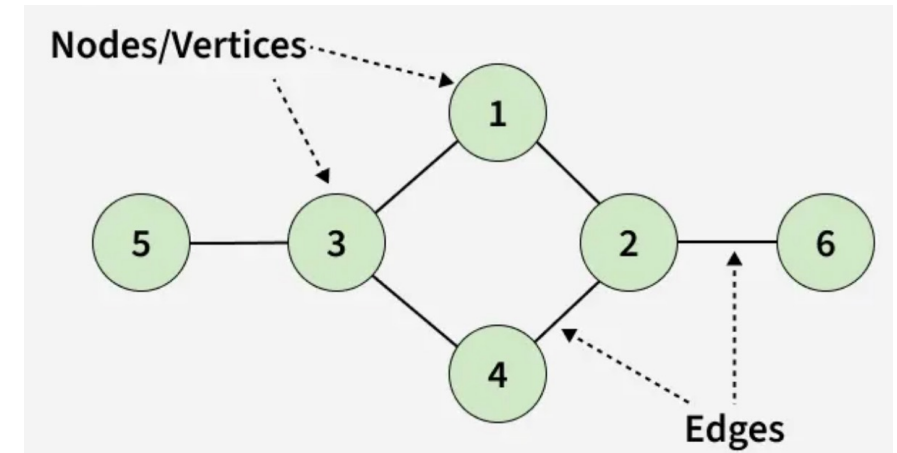


Figure 1 – Example of unweighted, undirected graph[1]

[1] <https://www.geeksforgeeks.org/dsa/introduction-to-graphs-data-structure-and-algorithm-tutorials/>

# MATRIX REPRESENTATION

Introduction to  
Graphs

Matrix  
representation

Adjacency & Edge  
list

Breadth-First  
Search

Practice Problems

One of the simple representation of the graph is to use **2D matrix** of size  $V \times V$ , where each cell

$$(i, j) = \begin{cases} 1, & \text{if there is an edge from vertex } i \text{ to } j \\ 0, & \text{otherwise} \end{cases} \quad (2)$$

Example:

$$AdjM = \begin{bmatrix} 0 & 1 & 0 & 0 & 1 \\ 1 & 0 & 1 & 1 & 1 \\ 0 & 1 & 0 & 1 & 1 \\ 0 & 1 & 1 & 0 & 1 \\ 1 & 1 & 1 & 1 & 0 \end{bmatrix} \quad (3)$$

Pros:

- Easy to check if edge exists  $\rightarrow O(1)$ ;
- Simple implementation.

Cons:

- Requires  $O(V^2)$  space.

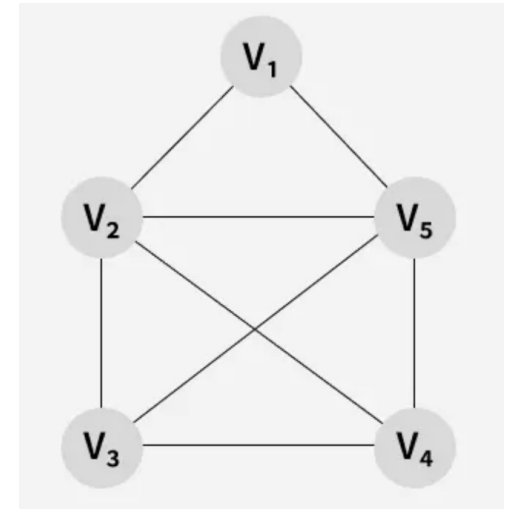


Figure 2 – Example graph [1]

# ADJACENCY & EDGE LIST

Introduction to  
Graphs

Matrix  
representation

Adjacency & Edge  
list

Breadth-First  
Search

Practice Problems

In **adjacency list** each vertex stores a list of all the vertices it's connected to.

**Example:**

0: [1, 2],  
1: [0, 2, 3],  
2: [0, 1, 4],  
3: [1, 4],  
4: [2, 3]

*Pros:*

- Space-efficient for sparse graphs ( $O(V + E)$ )
- Easier for traversal (like BFS/DFS)

*Cons:*

- Slower edge existence check  $\rightarrow O(\deg(V))$

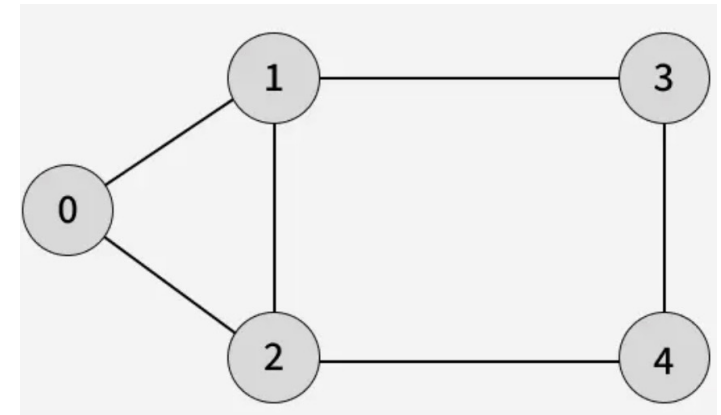


Figure 3 – Example undirected graph [2]

[2] <https://www.geeksforgeeks.org/dsa/breadth-first-search-or-bfs-for-a-graph/>

# ADJACENCY & EDGE LIST (CONT.)

Introduction to  
Graphs

Matrix  
representation

Adjacency & Edge  
list

Breadth-First  
Search

Practice Problems

The **edge list** store all edges as pairs  $(u, v)$ .

**Example:**

Edges =  $[(0, 1), (0, 2),$   
 $(1, 2), (1, 3),$   
 $(3, 1), (3, 4)]$

*Use Cases:*

- Easy for algorithms like Kruskal's MST
- Compact representation

*Pros: Simple & compact.*

*Cons: Not efficient for neighbor lookups.*

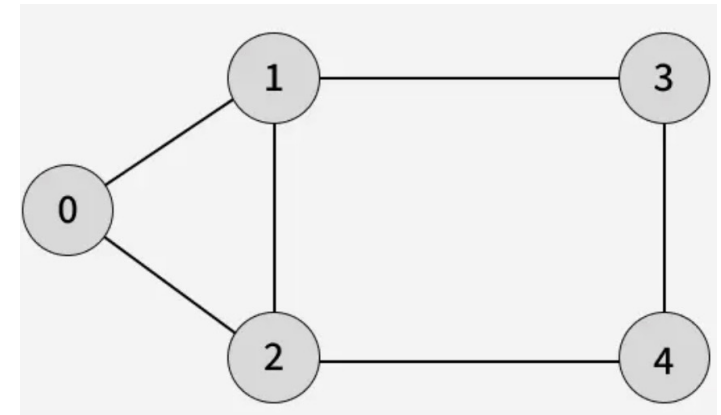


Figure 3 – Example undirected graph [2]

# BREADTH-FIRST SEARCH

Introduction to  
Graphs

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list

Breadth-First  
Search

Practice Problems

The Breadth-First Search (BFS) algorithm uses a **queue** to traverse a graph **level by level**, and great for finding shortest paths in **unweighted graphs**.

**Steps of algorithm:**

1. Starts from initial vertex;
2. Visit its all neighbors;
3. Move to next level neighbors;
4. Repeat the 2<sup>nd</sup> & 3<sup>rd</sup> steps until all reachable nodes are visited.

Time Complexity:  $O(V+E)$

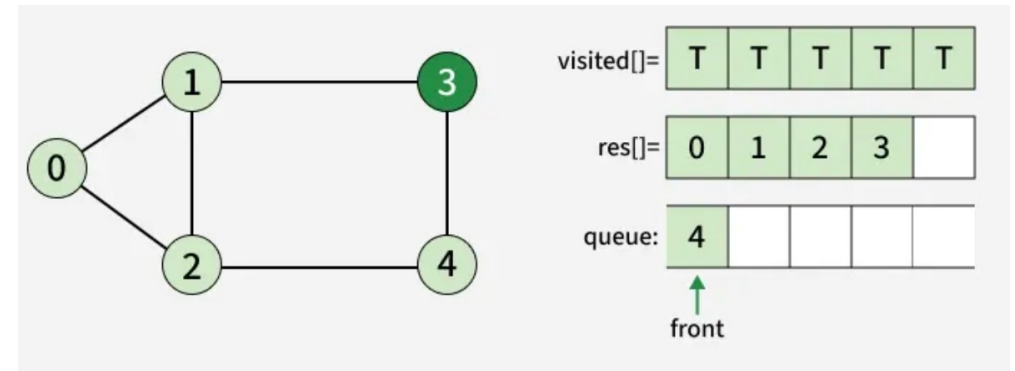


Figure 4 – Algorithm representation of BFS

# BREADTH-FIRST SEARCH (CONT.)

Given the graph as shown in Figure 5, and we should traverse upon that graph.

Introduction to  
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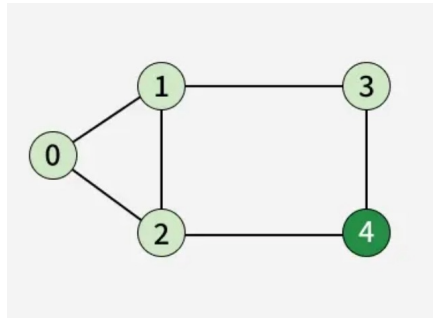


Figure 5 – Given graph [2].



```
vector<vector<int>> adj = {
    {1, 2},
    {0, 2, 3},
    {0, 1, 4},
    {1, 3},
    {2, 3}
};
```

Figure 6 – Graph representation.

```
void bfs(int start, vector<vector<int>> adj) {
    vector<bool> visited(adj.size(), false);
    queue<int> q;

    visited[start] = true;
    q.push(start);

    while(!q.empty()) {
        int vertex = q.front();
        q.pop();
        cout << vertex << " ";

        for (int i = 0; i < adj[vertex].size(); ++i) {
            int neighbor = adj[vertex][i];
            if (!visited[neighbor]) {
                visited[neighbor] = true;
                q.push(neighbor);
            }
        }
    }
}
```

Figure 7 – BFS algorithm.

```
(base) assylzhan@AIs-MacBook-Air L10 % ./a.out
0 1 2 3 4 %
```

Figure 8 – output of the BFS algorithm.

[2] <https://www.geeksforgeeks.org/dsa/breadth-first-search-or-bfs-for-a-graph/>



# PRACTICE PROBLEMS

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Introduction to  
Graphs

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Practice Problems

We have two tasks to solve:

- 104. Maximum Depth of Binary Tree;
- 463. Island Perimeter;
- [547. Number of Provinces.](#)

Q & A