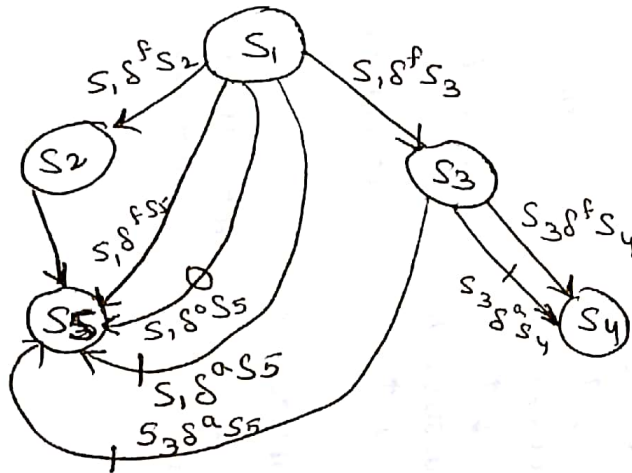


# HOMEWORK-4

SHIVANI SABHLOK  
500237896

①	Process	Read	Write	Dependency
	S1	a	a	
	S2	a, b	b	$S_1 \delta^f S_2$
	S3	a, d	c	$S_1 \delta^f S_3$
	S4	c	d	$S_3 \delta^f S_4, S_3 \delta^a S_4$
	S5	a	a	$S_1 \delta^f S_5, S_1 \delta^a S_5, S_1 \delta^a S_5, S_2 \delta^a S_5, S_3 \delta^a S_5$

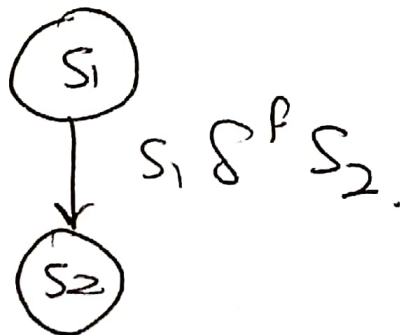


② for (i = 0; i < max; i++)

for (j = 0; j < max; j++)

S1: A[i][j] = A[i+1][j-1]

S2: B[i][j] = A[i][j]



	Dist Vector	Direction Vector	
$S_1$	$(-1, 1)$	$(-, +)$	$\Rightarrow - \Rightarrow \text{Anti-dependency.}$
$S_2$	NA	NA	

- if we reorder  $S_1$ , sign of the direction vector would change  $\Rightarrow$  no reordering possible.
- for  $S_2$ , since there are no dependencies, 2! re-orderings are allowed. There will be no memory conflicts.

Thus, loops cannot be reordered because it will change the sign or direction vector for  $S_1$ .

3. for ( $i=0; i < \text{max}; i++$ )  
     for ( $j=0; j < \text{max}; j++$ )  
         for ( $k=0; k < \text{max}; k++$ )  
             for ( $m=0; m < \text{max}; m++$ )  
                 BODY ( $i, j, k, m$ )

Solu

$$(a) \underline{A(i)(j)(k)(m)} = B(i-1)(j-1)(k-1)(m-1)$$

since the two are independent memory locations, all combinations of loop ordering

are allowed

$\Rightarrow$   $4!$  reordering combinations are allowed

(b)  $A(i)(j)(k)(m) = A(i-1)(j)(k)(m)$

Direction Vector:  $(+, 0, 0, 0)$

Dist. vector:  $(1, 0, 0, 0)$

Since direction vector will always be  $+$ ve,  
all combinations of reordering are allowed

$\Rightarrow 4! (= 24)$  possible reorderings are possible

(c)  $A(i)(j)(k)(M+1) = A(i)(j)(k)(M) + B(i)$

Direction vector for A:  $(0, 0, 0, +)$

Distance vector for A:  $(0, 0, 0, 1)$

B has no effect here

Thus, since the sign of the direction vector won't change,  $4! (= 24)$  reorderings are possible.