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January 22 2026

## Assignment 1 – Coding and Complexity

### **Problem 2:** Algorithm for Common Substring (pseudocode)

```
Algorithm LongestCommonSubstring (text1, text2) {
```

First convert text1 and text2 to character arrays A and B

Initialize: `longestLengthSub = ""` to store longest common substring

Next, iterate over every starting position, called  $i$ , in  $A$

for i from 0 to A.length - 1

for j from i to A.length - 1

`length = j - i + 1`

Next check of this substring is found in B

For q from 0 to B.length - length

**match = true**

for r from 0 to length -1

If  $A[i + r] \neq B[q + r]$

match = false

break

If match == true

If `length > longestLengthSub.length`

longest = substring of A from i to j

Break

Return longest

## **Problem 6: Algorithm Analysis**

Problem 1:

- Time Complexity is  $O(2^{m+n})$  where  $m$  is the length of the first string and  $n$  is the length of the second input string, because the algorithm branches off into two more calls when the base case is not met.
- Space Complexity is  $O(m+n)$ , where  $m$  and  $n$  are the lengths of the two input strings, because we are storing two input arrays.
- Big- $\Omega$  =  $\Omega(m+n)$  in the best case, where  $m$  and  $n$  are the lengths of the two input strings. When all characters match, we only make one recursive call per character.

Problem 2:

- Time Complexity is  $O(n^3)$  where  $n$  is the length of the first string and  $m$  is the length of the second string, because there are three nested loops that check every substring of the first against all spots in string  $m$ .
- Space Complexity is  $O(n+m)$ , where  $n$  is the length of the first string, and  $m$  is the length of the second string, because the solution iterates over all of the elements one time and stores both arrays.
- Big- $\Omega$  =  $\Omega(n^2)$  in the best case where  $n$  is the length of the first string, because a match is found immediately.

Problem 3:

- Time Complexity is  $O(n)$ , where  $n$  is the size of the array, because the loop iterates  $n-2$  times, doing constant-time operations each time.
- Space Complexity is  $O(n)$ , where  $n$  is the size of the array, because an array size of  $n$  stores the sequence of terms.
- Big- $\Omega$  =  $\Omega(n)$  where  $n$  is the size of the array, because in the best case, it still creates and stores all  $n$  elements, no matter what they are.

Problem 4:

- Time Complexity is  $O(n)$  where  $n$  is the size of the array because the NotFibonacci method takes  $O(n)$  and binary search takes  $O(\log n)$ , and  $n$  is higher order than  $\log n$ .
- Space Complexity is  $O(n)$ , where  $n$  is the size of the array, because the algorithm stores all  $n$  elements of the array.
- Big- $\Omega$  =  $\Omega(n)$  where  $n$  is the size of the array, because the algorithm creates all  $n$  elements before doing the search.

## Problem 5:

- Time Complexity is  $O(n)$  where  $n$  is the size of the array because each element is iterated through at least one time.
- Space Complexity is  $O(1)$ , where  $n$  is the size of the array, because there are a constant number of variables, and the array is modified in-place.
- Big- $\Omega$  =  $\Omega(n)$ , where  $n$  is the size of the array, because all elements in it must be looked at least once to see if it will remain.

## EC:

The screenshot shows a LeetCode submission page for the problem "removeElement". The submission was made by user "saraahamed" on Jan 22, 2026, at 19:57. It has passed 116 / 116 testcases. The runtime is 0 ms (Beats 100.00%) and the memory usage is 43.70 MB (Beats 15.01%). The code is written in Java and uses a two-pointer approach to remove elements from the array.

```
1 class Solution {
2     public int removeElement(int[] nums, int val) {
3         int k = 0;
4         for (int i = 0; i < nums.length; i++) {
5             if (nums[i] != val) {
6                 nums[k] = nums[i];
7                 k++;
8             }
9         }
10    return k;
11 }
```

The test case provided is [3,2,2,3] with value 3.