Chapter 3 outline

- 3.1 Transport-layer services
- 3.2 Multiplexing and demultiplexing
- 3.3 Connectionless transport: UDP
- 3.4 Principles of reliable data transfer
- 3.5 Connection-oriented transport: TCP
  - segment structure
  - reliable data transfer
  - flow control
  - connection management
- 3.6 Principles of congestion control
- 3.7 TCP congestion control

TCP
unreliable network
⇩
reliable network

흐름제어

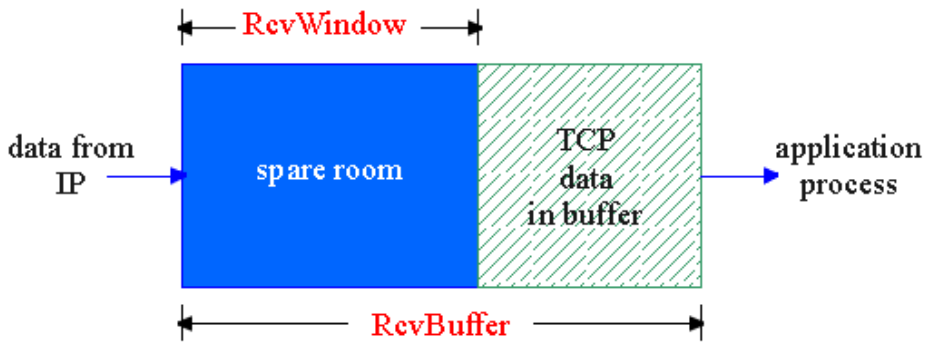
TCP Flow Control

sender와 receiver의 data 처리 속도 차이를 제어
rcv가 너무 많은 pkt을 보낼 경우 조절

flow control
sender won't overflow receiver's buffer by transmitting too much, too fast

rcv가 sender에게 자신의 상태를 feedback

- receive side of TCP connection has a receive buffer:



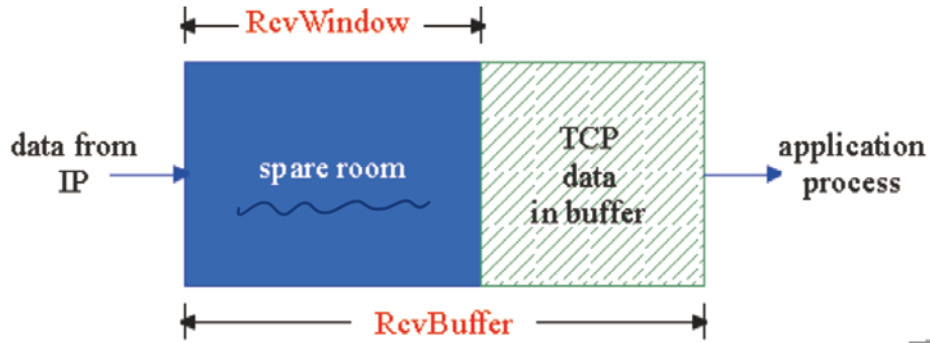
- app process may be slow at reading from buffer

- speed-matching service: matching the send rate to the receiving app's drain rate

데이터 처리 속도 : 수신 > 송신 ok
수신 < 송신 Not ok
⇩
데이터 손실

∴ 송신측의 데이터 전송량을 수신측에 따라 조절

TCP Flow control: how it works



(Suppose TCP receiver discards out-of-order segments)

- spare room in buffer

= **RcvWindow**

= **RcvBuffer-[LastByteRcvd - LastByteRead]**

- Rcvr advertises spare room by including value of **RcvWindow** in segments

- Sender limits unACKed data to **RcvWindow**
  - guarantees receive buffer doesn't overflow

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TCP Connection Management

Recall: TCP sender, receiver establish "connection" before exchanging data segments

- initialize TCP variables:
  - seq. #s
  - buffers, flow control info (e.g. **RcvWindow**)
- *client:* connection initiator

```
Socket clientSocket = new
Socket("hostname","port
number");
```

- *server:* contacted by client

```
Socket connectionSocket =
welcomeSocket.accept();
```

Three way handshake:

Step 1: client host sends TCP SYN segment to server
- specifies initial seq #
- no data

Step 2: server host receives SYN, replies with SYNACK segment
- server allocates buffers
- specifies server initial seq. #

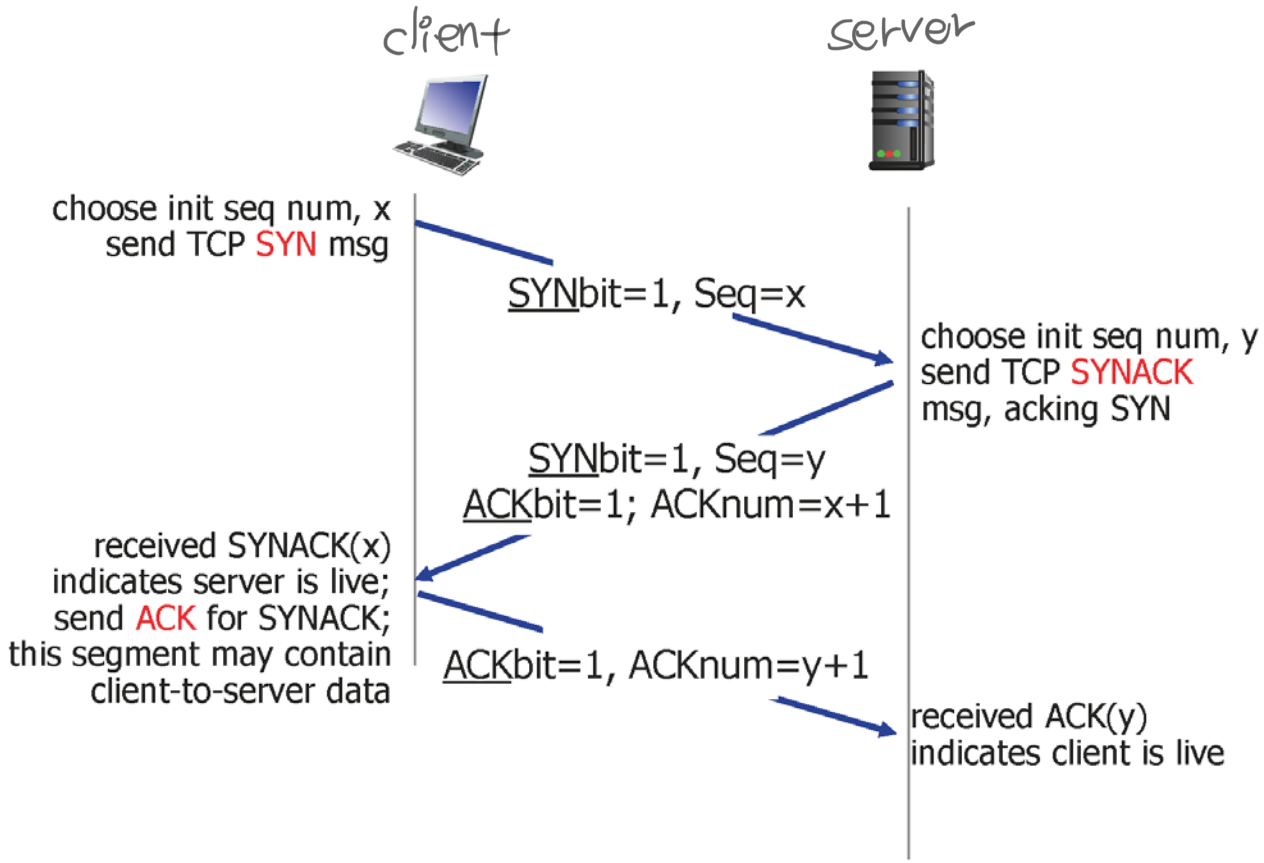
Step 3: client receives SYNACK, replies with ACK segment, which may contain data

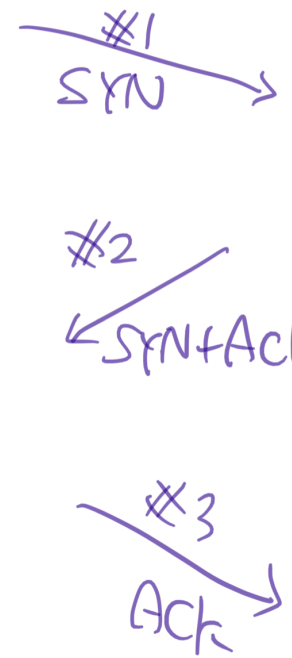
TCP 3-way handshake

양쪽 모두 데이터를 송수신 준비가 되었음을 보장,
서로 버퍼 현황을 알 수 있게 한다.
상대편 초기 순차 일련
번호를 알 수 있게한다.

A
client는
B
server에
접속요청 packet을보냄.

A클라이언트는
SYN을 보내고
SYN/ACK응답
기다리는 → B server → Bserver
SYN-SENT상태. SYN RECEIVED상태. ESTABLISHED.





Closing TCP Connection

4.

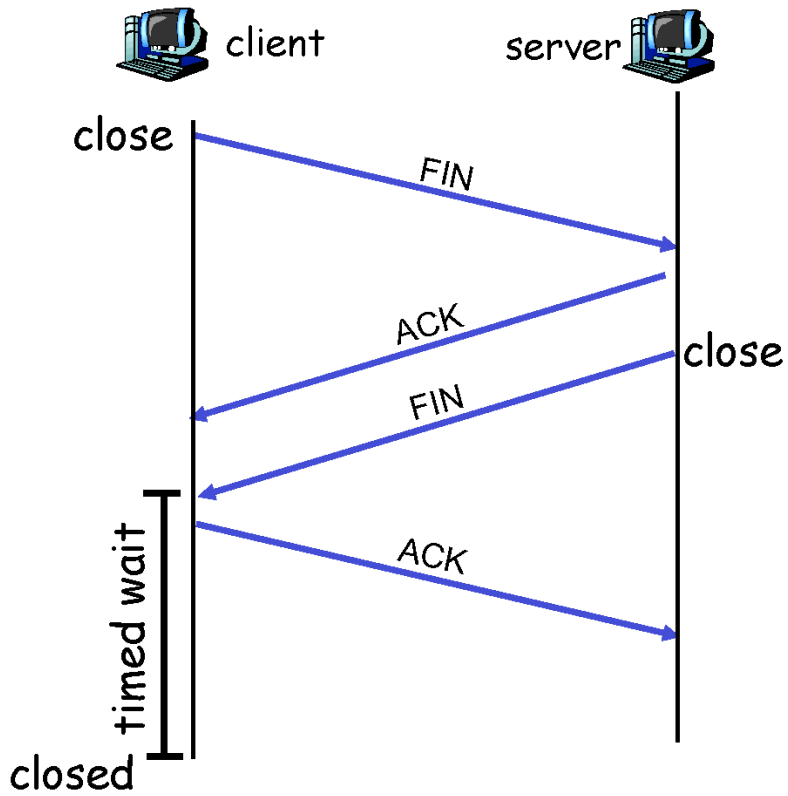
Closing a connection:

client closes socket:

```
clientSocket.close();
```

Step 1: client end system sends TCP FIN control segment to server

Step 2: server receives FIN, replies with ACK. Closes connection, sends FIN.

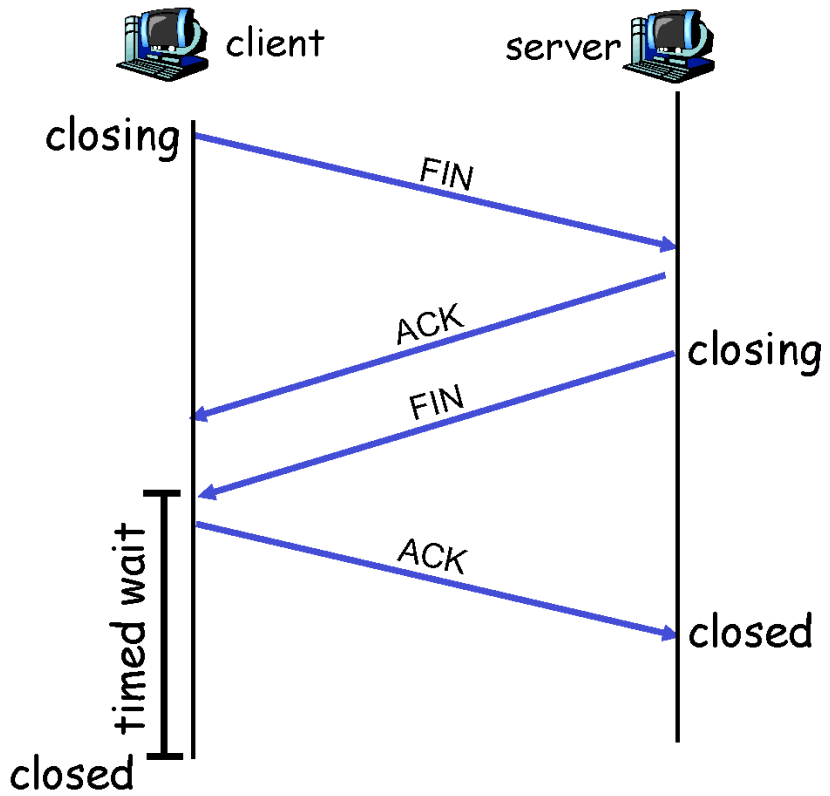


TCP Connection Management (cont.)

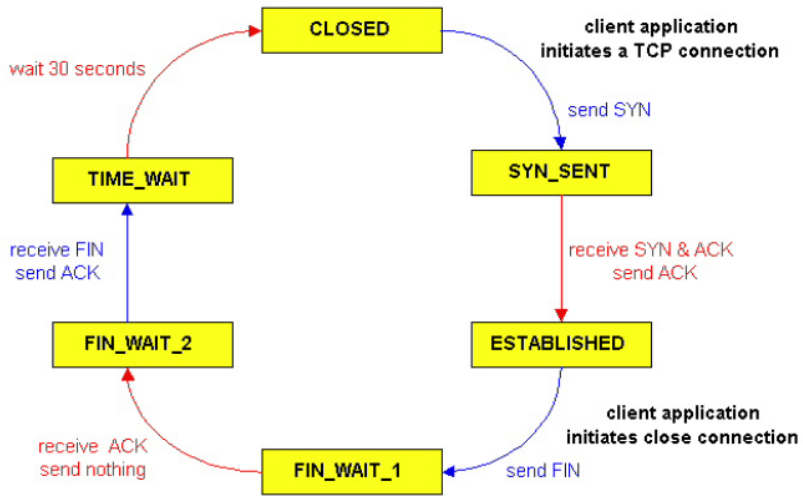
Step 3: client receives FIN, replies with ACK.

- Enters "timed wait" - will respond with ACK to received FINs

Step 4: server, receives ACK. Connection closed.

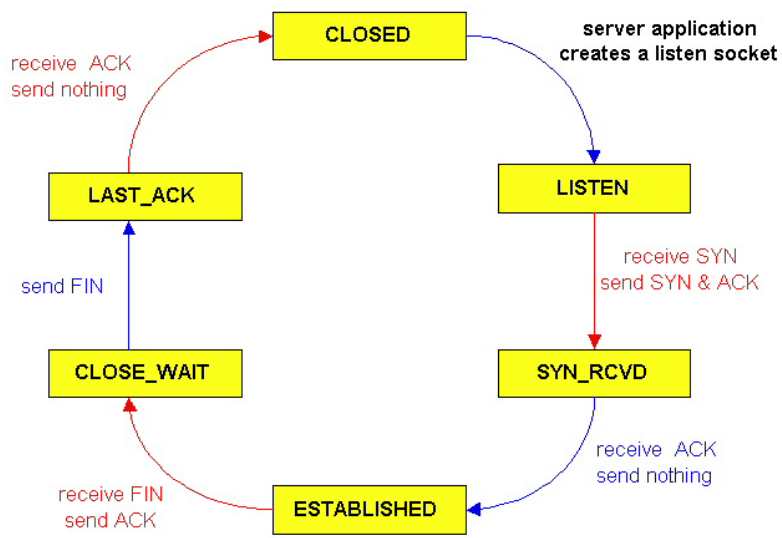


TCP Connection Management (cont)



TCP client lifecycle

TCP server lifecycle



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Approaches towards congestion control

Two broad approaches towards congestion control:

End-end congestion control:

- no explicit feedback from network
- congestion inferred from end-system observed loss, delay
- approach taken by TCP

Network-assisted congestion control:

- routers provide feedback to end systems
  - single bit indicating congestion (SNA, DECbit, TCP/IP ECN, ATM)
  - explicit rate sender should send at