



Introduction

CSC/ECE 570 Computer Networks
Fall 2024



Agenda

1. History
2. Internet Structure
3. Packet vs. circuit switching
4. Internet Structure Continued

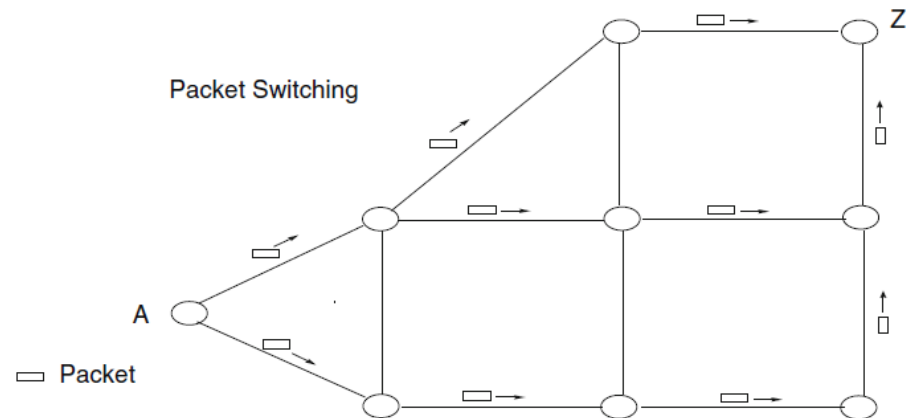
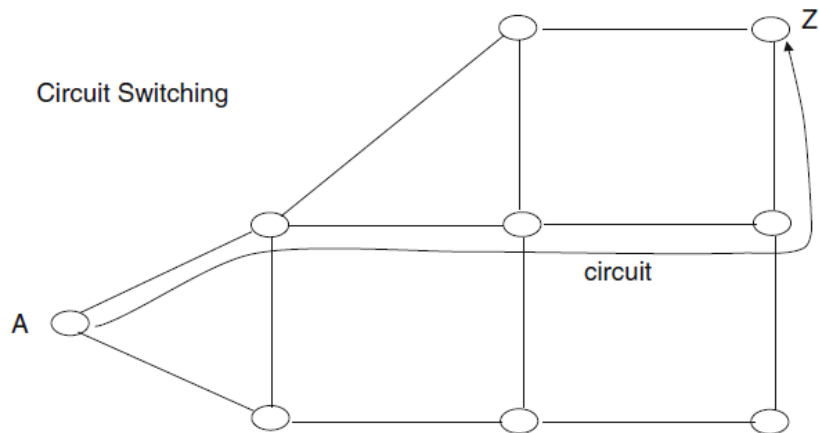


History

- Cerf and Kahn are considered the **Internet founding fathers**
 - Vinton Cerf -> Born in 1943, DARPA program manager, funded and chairman of ICANN, president of ACM
 - Bob Kahn -> Born in 1938, Ph.D. from City University of NY, worked at AT&T, and was a professor at MIT
 - Established internetworking principles:
 1. *minimalistic, automatic way* - no internal changes required to interconnect networks
 2. *best effort* service model
 3. *stateless routers*
- define **Internet architecture**

Architectures for Networking

- Switched networks
 - Circuit switched vs. Packet switched





1960's: Early Packet-Switching Principles

- 1961: Kleinrock - queueing theory shows effectiveness of packet-switching
- 1964: Baran - packet-switching in military nets
- 1967: ARPAnet conceived by Advanced Research Projects Agency
- 1969: first ARPAnet node operational
- 1970: ALOHAnet wireless network in Hawaii



1970's: Internetworking, proprietary nets

- 1972:
 - ARPAnet public demo
 - NCP (Network Control Protocol) first host-host protocol
 - first e-mail program
 - ARPAnet has 15 nodes
- 1974: Cerf and Kahn - architecture for interconnecting networks
- 1976: Ethernet at Xerox PARC
- Late 70's: proprietary architectures: DECnet, SNA, XNA
- late 70's: switching fixed length packets (ATM precursor)
- 1979: ARPAnet has 200 nodes



1980's: new protocols, a proliferation of networks

- 1983: deployment of TCP/IP
- 1982: SMTP e-mail protocol defined
- 1983: DNS defined for name-to-IP-address translation
- 1985: FTP protocol defined
- 1988: TCP congestion control
- new national networks: Csnet, BITnet, NSFnet, Minitel
- 100,000 hosts connected to confederation of networks



1990, 2000's: Commercialization, the Web, new apps

- 1991: NSF lifts restrictions on commercial use of NSFnet (decommissioned, 1995)
- early 1990s: Web
 - hypertext [Bush 1945, Nelson 1960' s]
 - HTML, HTTP: Berners-Lee
 - 1994: Mosaic, later Netscape
 - late 1990' s: commercialization of the Web
- late 1990' s – 2000' s:
 - more killer apps: instant messaging, P2P file sharing
 - network security to forefront
 - est. 50 million hosts, 100 million+ users
 - backbone links running at Gbps



2005 - present

- ~750 million hosts
- Smartphones and tablets
- Aggressive deployment of broadband access
- Increasing ubiquity of high-speed wireless access
- Emergence of online social networks:
 - Facebook: soon one billion users
- Service providers (Google, Microsoft) create their own networks
 - Bypass Internet, providing “instantaneous” access to search, email, etc.
- E-commerce, universities, enterprises running their services in “cloud” (eg, Amazon EC2)



IOSC, IETF and RFCs

- **ISOC:** the Internet Society (founded by Cerf, Kahn and Lyman Chapin), main support body of IETF
- **IETF:** Internet Engineering Task Force, major standardization body of the Internet
- **RFCs (Request For Comments):** formal documentation of Internet standards submitted to/approved by IETF; all open access at https://www.rfc-editor.org/search/rfc_search.php
 - RFC 768 (Aug. 1980): UDP
 - RFC 791 (Sep. 1981): IP
 - RFC 793 (Sep. 1981): TCP
 - RFC 1034/1035 (Nov. 1987): DNS
 - ...
 - RFC 8888 (Jan. 2021): RTP Control Protocol Congestion Control



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Network devices

- Endhosts (computers and clients)
- Routers, switches, and hubs to connect endhosts
- Links of different media: wired/wireless
- A **network** consists of some number of endhosts interconnected through routers/switches/hubs
- The **Internet** interconnects networks



Types of Networks

- Not all networks are created equal!

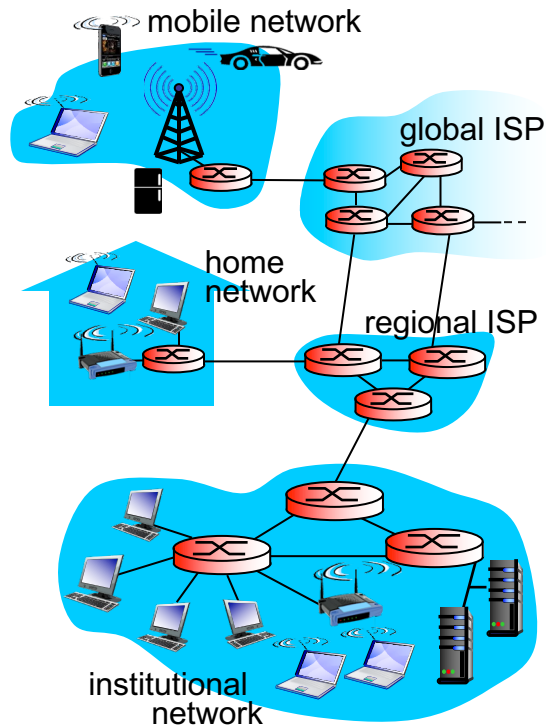
1. **Edge Networks:** Carry own traffic

Examples: Customer networks such as home networks and NCSU network

2. **Core Networks:** Carry transit traffic; i.e. traffic for their customers (in addition to own traffic)

Examples: ISP networks such as Sprint, AT&T, etc.

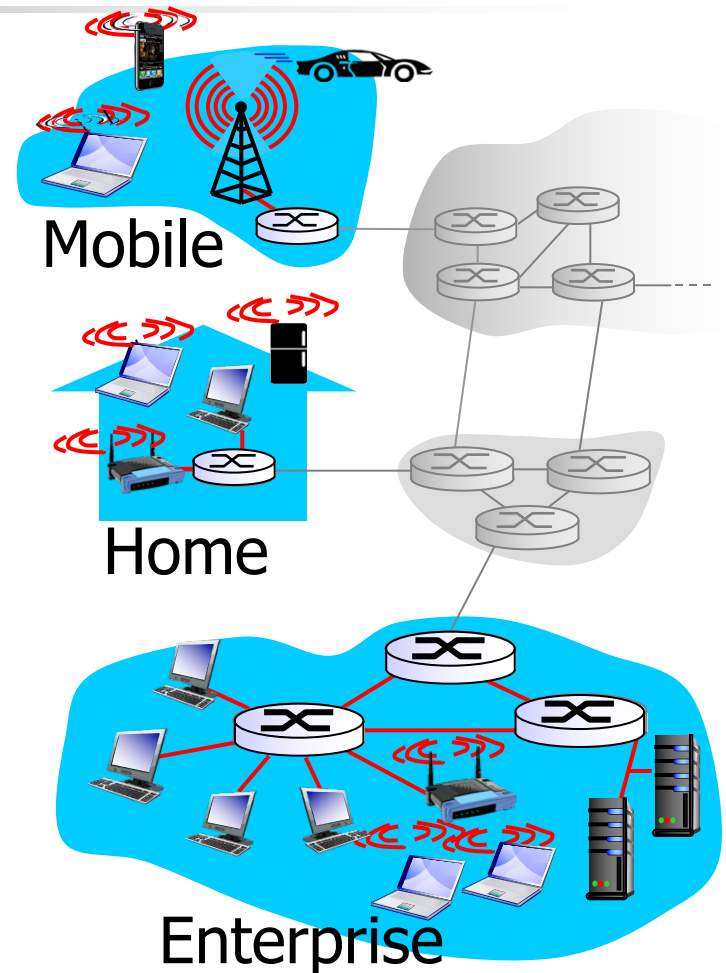
A closer look at network structure:



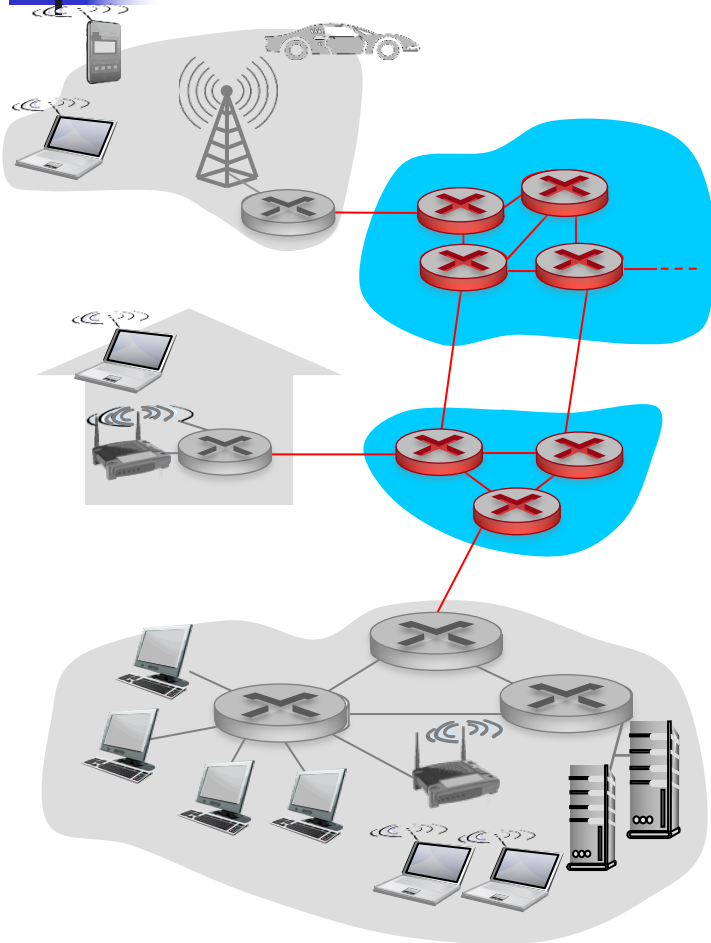
- *network edge:*
 - end systems/hosts, applications
 - clients & servers
- *access networks, physical media:* wired, wireless communication links
- *network core:*
 - interconnected routers
 - network of networks

Edge/Access Networks

- Connecting end-devices to edge routers
 - home access nets
 - enterprise access networks (school, company)
 - mobile access networks
- Heterogeneity, usually centralized
 - Media, hardware/devices, protocols
 - Bandwidth, delay, characteristics, ...



Core Networks



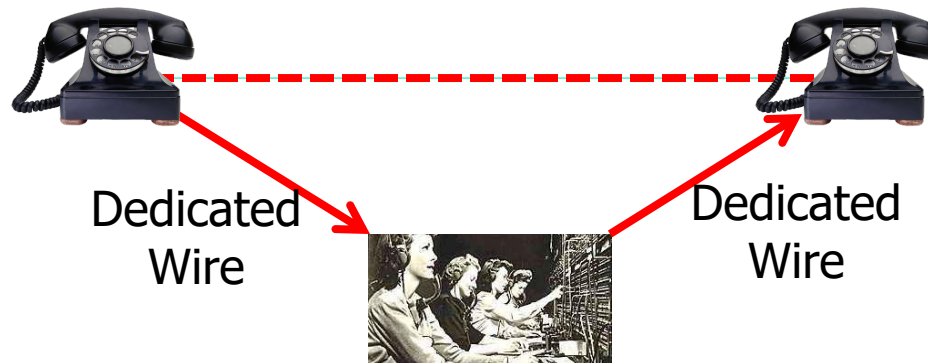
- Mesh of interconnected routers
 - Usually by fiber optics
- Decentralized autonomous systems (ASs)
- Peering between ASs of different ISPs



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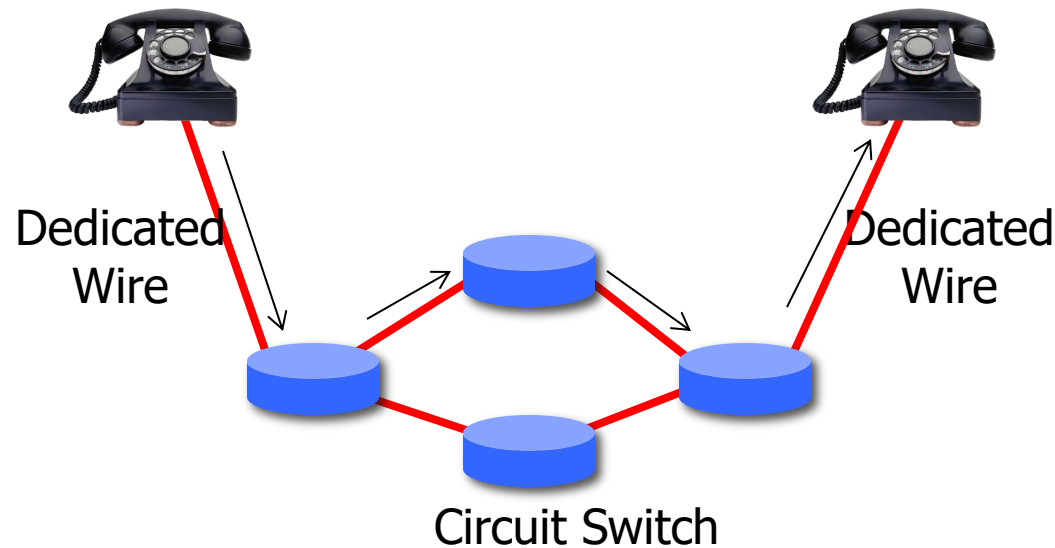
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Circuit Switching: Telephone Network



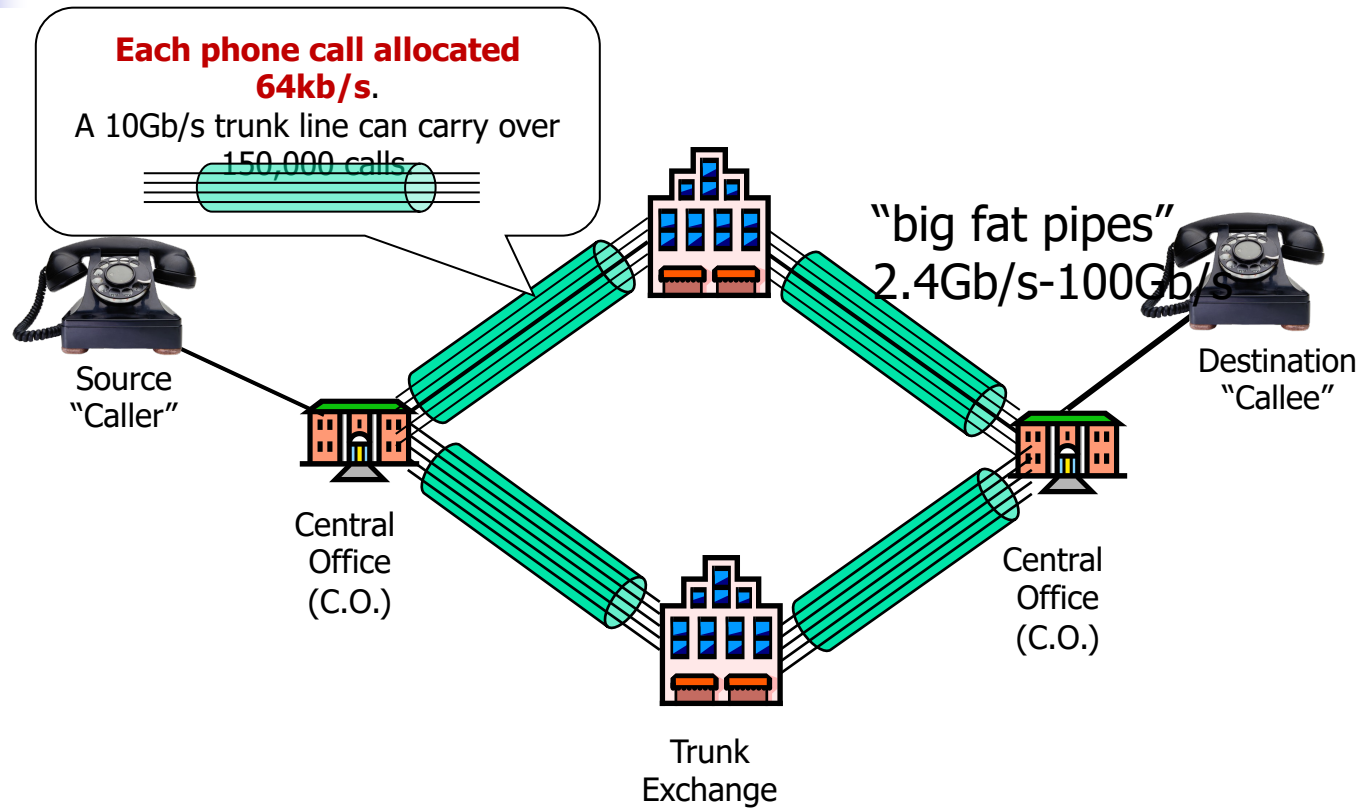
- telephones are connected by a dedicated wire to a local exchange
- Early days, switchboard operators used a big patch-panel to manually connect
- the wire is dedicated to the phone conversation from the start to the end of the phone call.

Circuit Switching: Telephone Network



1. **Dial a number**, which creates a dedicated circuit between the two phones. Each switch maintains state to map the incoming circuit to the correct outgoing circuit.
2. **Talk**: digital phone system, our voice is **sampled and digitized**, and sent over the dedicated circuit, typically **64kb/s for voice**. Our phone conversation has a dedicated circuit
3. **Hang up**: the circuit is removed, and any state is removed at the switches along the path.

Circuit Switching: Telephone Network

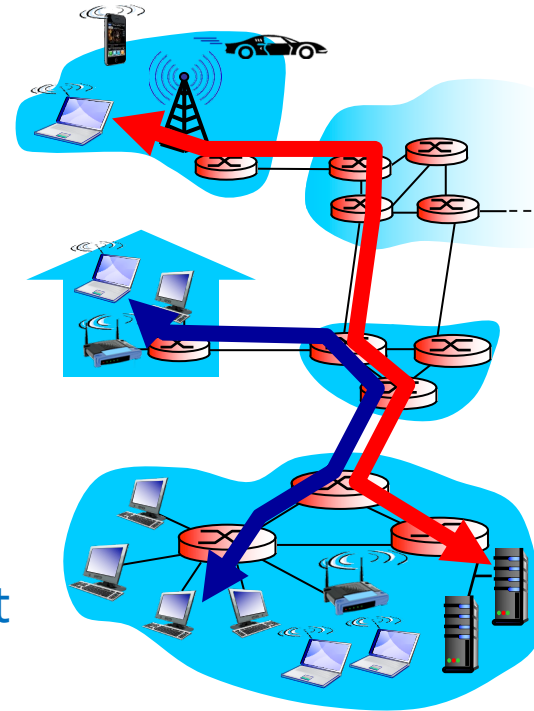


Circuit switching for the Internet's core

end-end resources reserved for, “call” between source, destination:

- dedicated resources: **no** sharing
 - circuit-like (guaranteed) performance
- resource piece idle if not used by owning call (*no sharing*)
- dividing link bandwidth into “pieces”
 - frequency division
 - time division

Restaurant Analogy

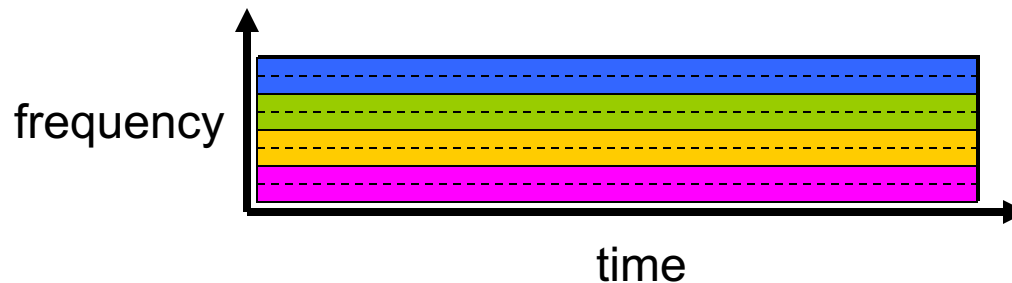


Circuit switching: FDM versus TDM

FDM

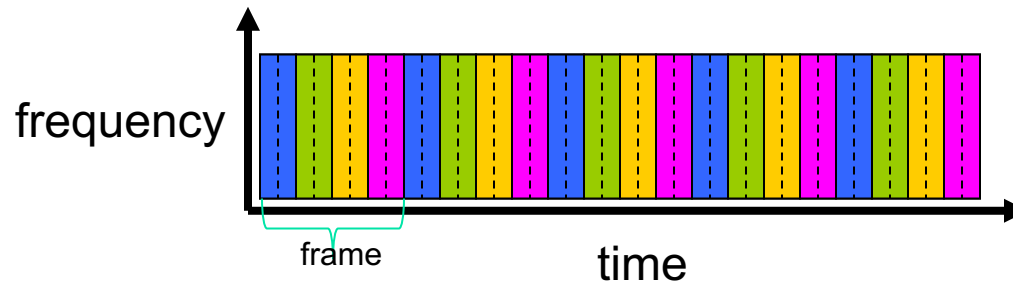
Example:

4 users



FM radio stations use FDM
88MHz to 108 MHz

TDM



time is divided into frames of certain duration

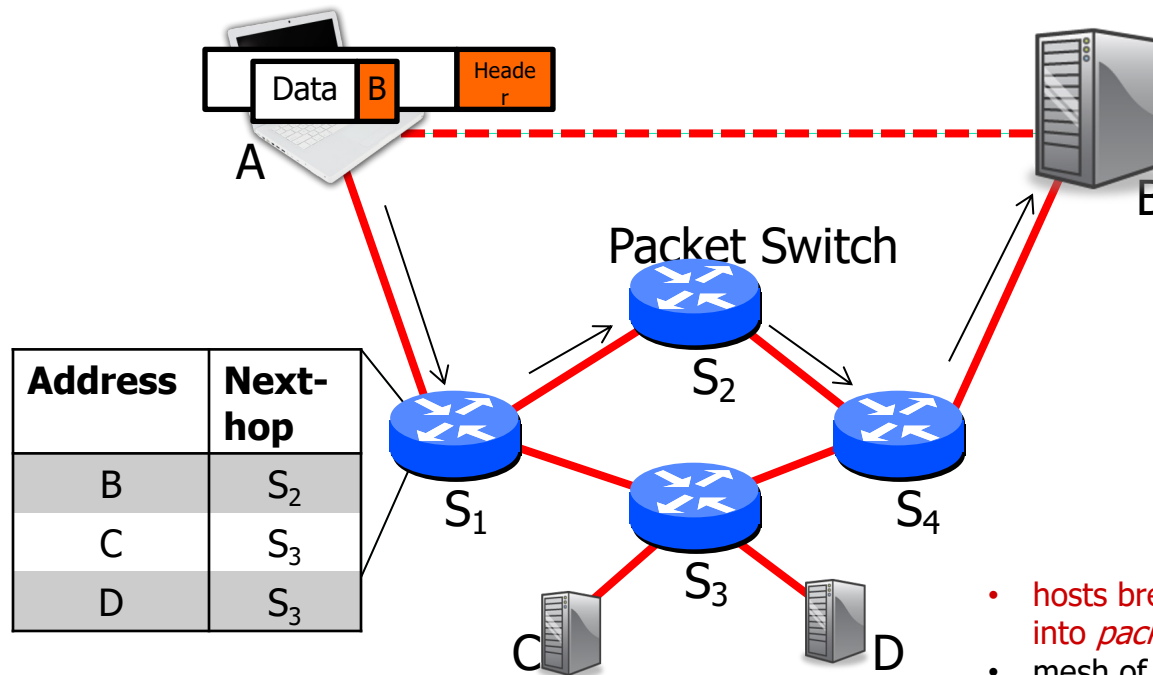
e.g 8000 frames/sec, each slot 8bits -> 64Kbps circuit



Circuit Switching Issues for the Internet

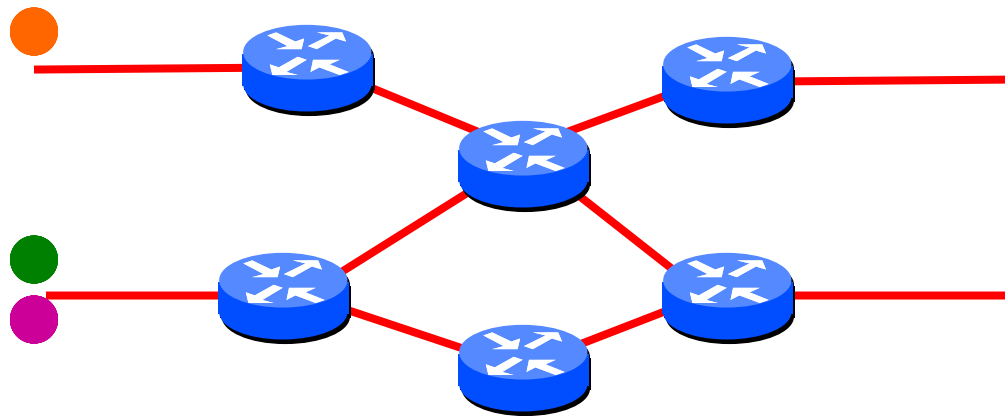
1. **Inefficient.** Computer communication tends to be very **bursty**. *e.g.* typing over an ssh connection, or viewing a sequence of web pages. If each communication has a dedicated circuit, it will be used very inefficiently.
2. **Diverse Rates.** Computers communicate at many different rates. *e.g.* a web server streaming video at 6Mb/s, or typing at 1 character per second over ssh. A fixed rate circuit will not be much use!
3. **State Management.** Circuit switches maintain per-communication state, which must be managed.

Packet Switching



- hosts break application-layer messages into *packets*
- mesh of interconnected routers
- forward packets from one router to the next, across links on path from source to destination
- The routers maintain **no per-communication state**.

Packet Switching

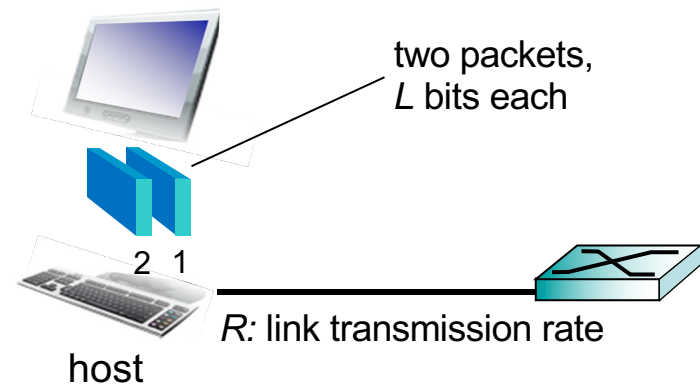


Packet Switch

Host: sends *packets* of data

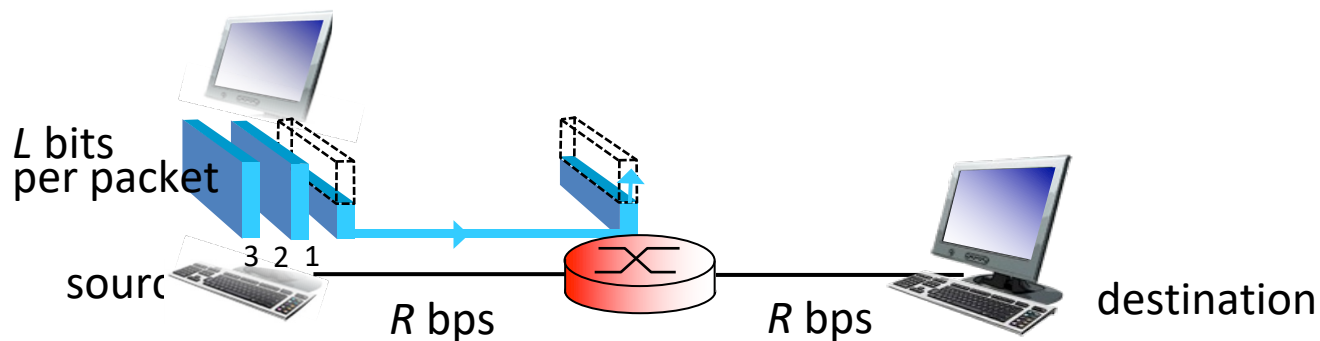
host sending function:

- takes application message
- breaks into smaller chunks, known as *packets*, of length L bits
- transmits packet into access network at *transmission rate* R
 - link transmission rate, aka link *capacity*, aka *link bandwidth*



$$\text{packet transmission delay} = \text{time needed to transmit } L\text{-bit packet into link} = \frac{L \text{ (bits)}}{R \text{ (bits/sec)}}$$

Packet-switching: store-and-forward

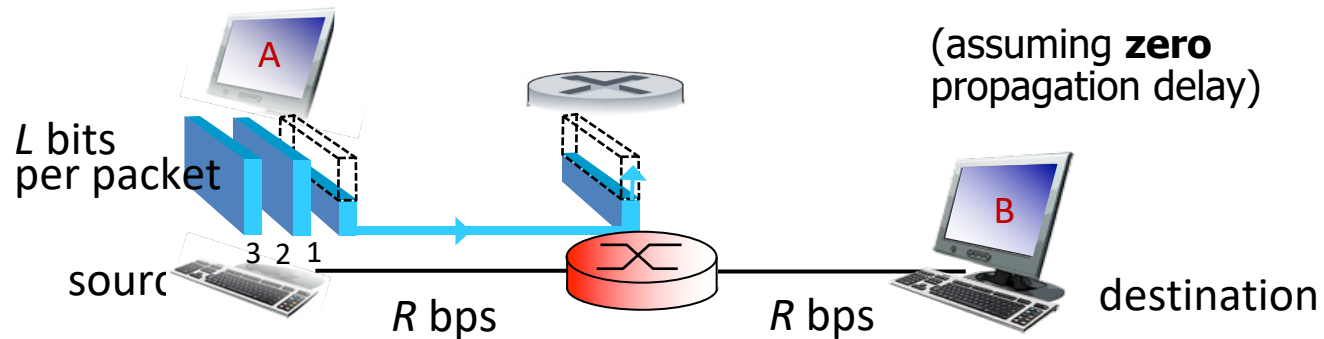


- takes L/R seconds to transmit (push out) L -bit packet into link at R bps
- *store and forward*: entire packet must arrive at router before it can be transmitted on next link

(assuming **zero** propagation, and processing delay)

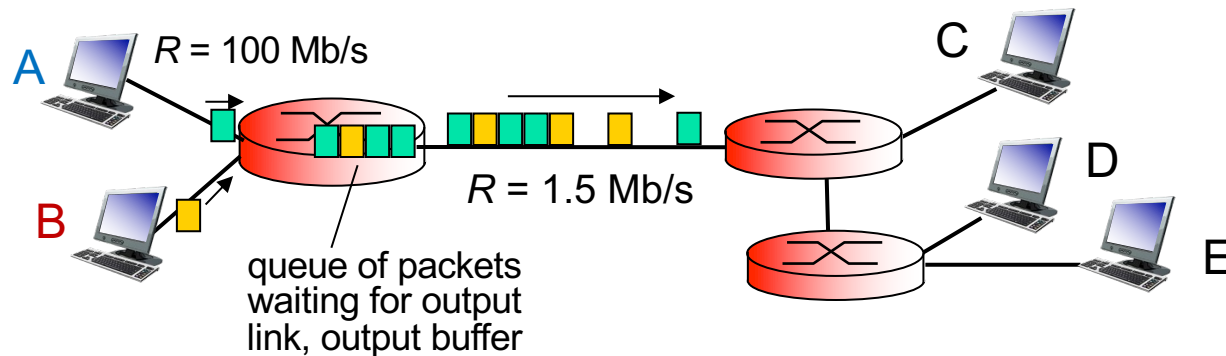
$$\text{delay} = \frac{2L}{R}$$

Packet-switching: store-and-forward



Q: What is the time between when the source begins to send the first packet until the destination has received all three packets?

Packet Switching: queueing delay, loss



resource contention:

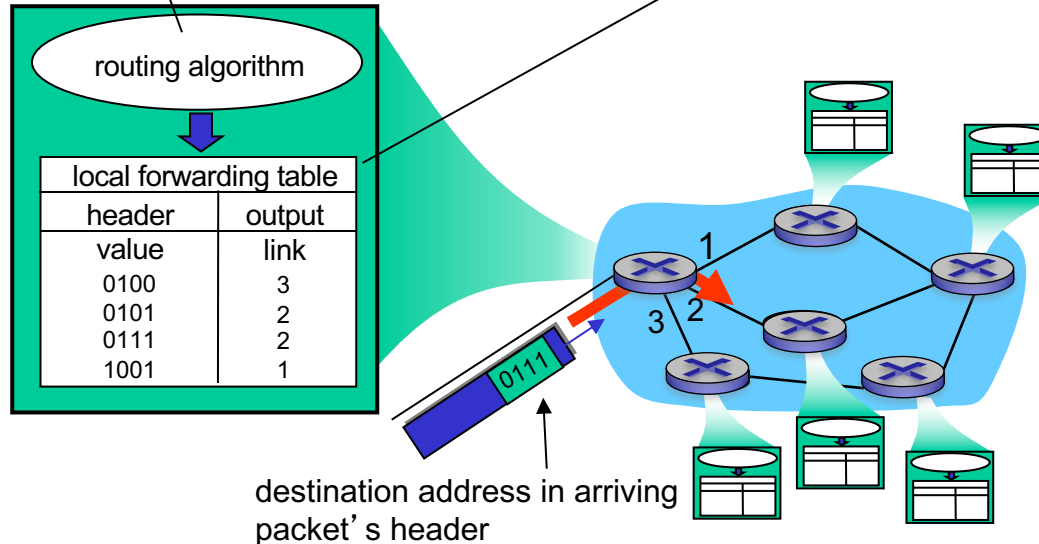
- aggregate resource demand (use of transmission link) can exceed amount available
- *congestion:*
 - packets will **queue**, wait for an available link use
 - packets can be dropped (**lost**) if no memory to store them

Forwarding Tables and Routing Protocols

routing: determines source-destination route taken by packets

- *routing algorithms*

forwarding: move packets from router's input to appropriate router output

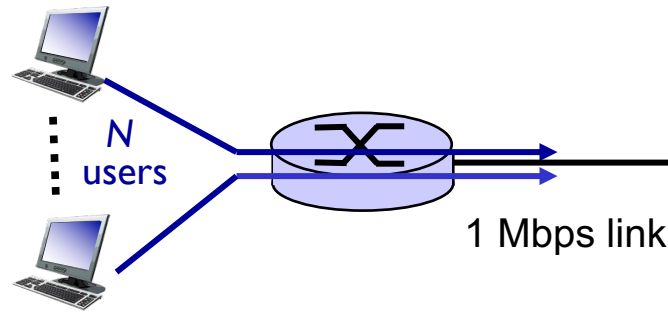


Packet Switching versus Circuit Switching

packet switching allows more users to use network!

Example:

- 1 Mb/s link
- each user:
 - 100 kb/s when “active”
 - active 10% of time



circuit-switching:

10 users

packet switching:

with more than 10 users



Packet switching vs. circuit switching

- great for **bursty** data
 - resource sharing
 - simpler, no call setup
- Resilience to failure of links & routers
- **excessive congestion possible:** packet delay and loss
 - protocols needed for reliable data transfer, congestion control

Q: How to provide circuit-like behavior?

- bandwidth guarantees needed for audio/video apps
- still an unsolved problem (will be revisited later in this course)



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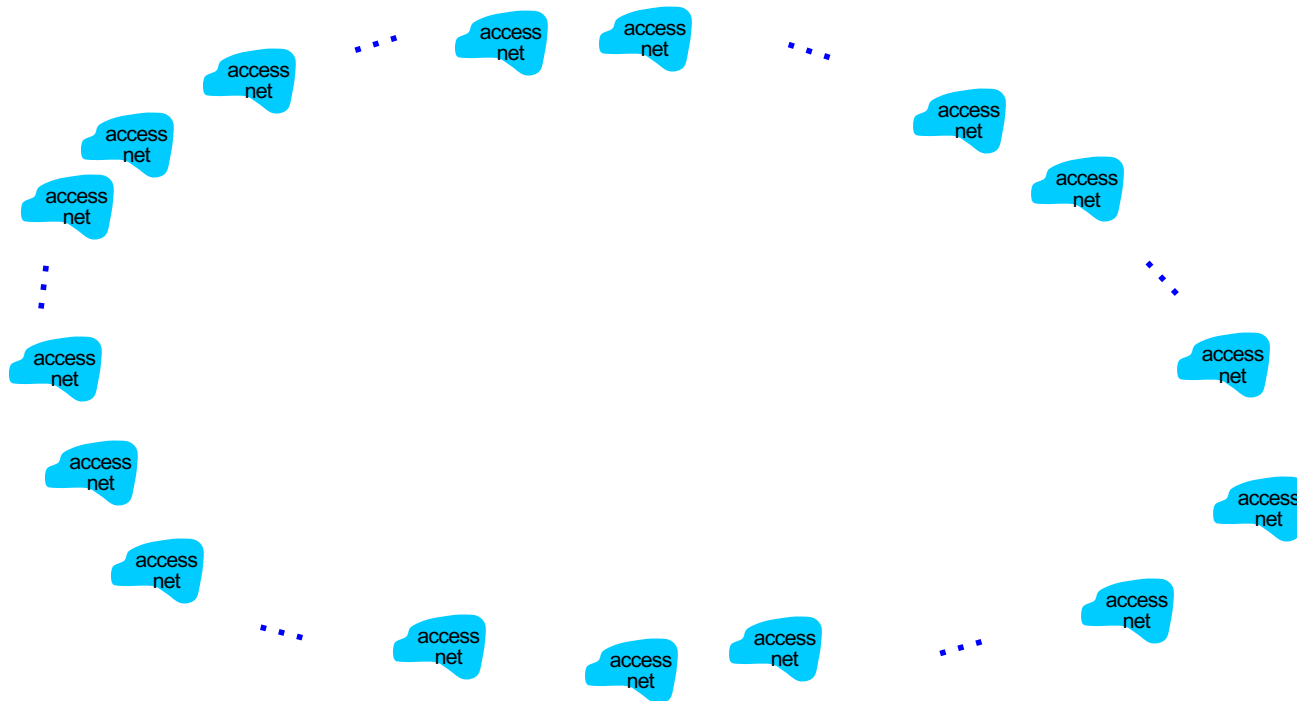


Internet structure: network of networks

- End systems connect to Internet via **access ISPs** (Internet Service Providers)
 - Residential, company and university ISPs
- Access ISPs in turn must be interconnected.
 - So that any two hosts can send packets to each other
- Resulting network of networks is very complex
 - Evolution was driven by **economics** and **national policies**
- Let's take a stepwise approach to describe current Internet structure

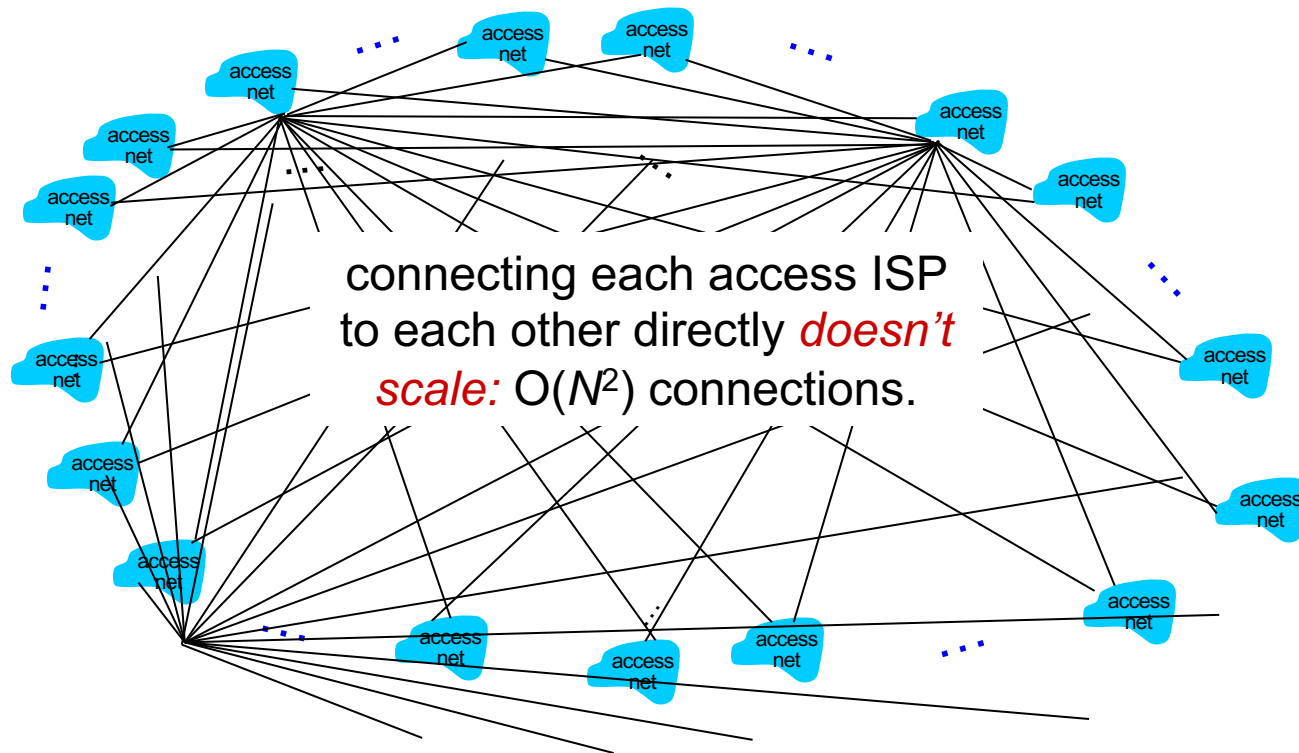
Internet structure: network of networks

Question: given *millions* of access ISPs, how to connect them together?



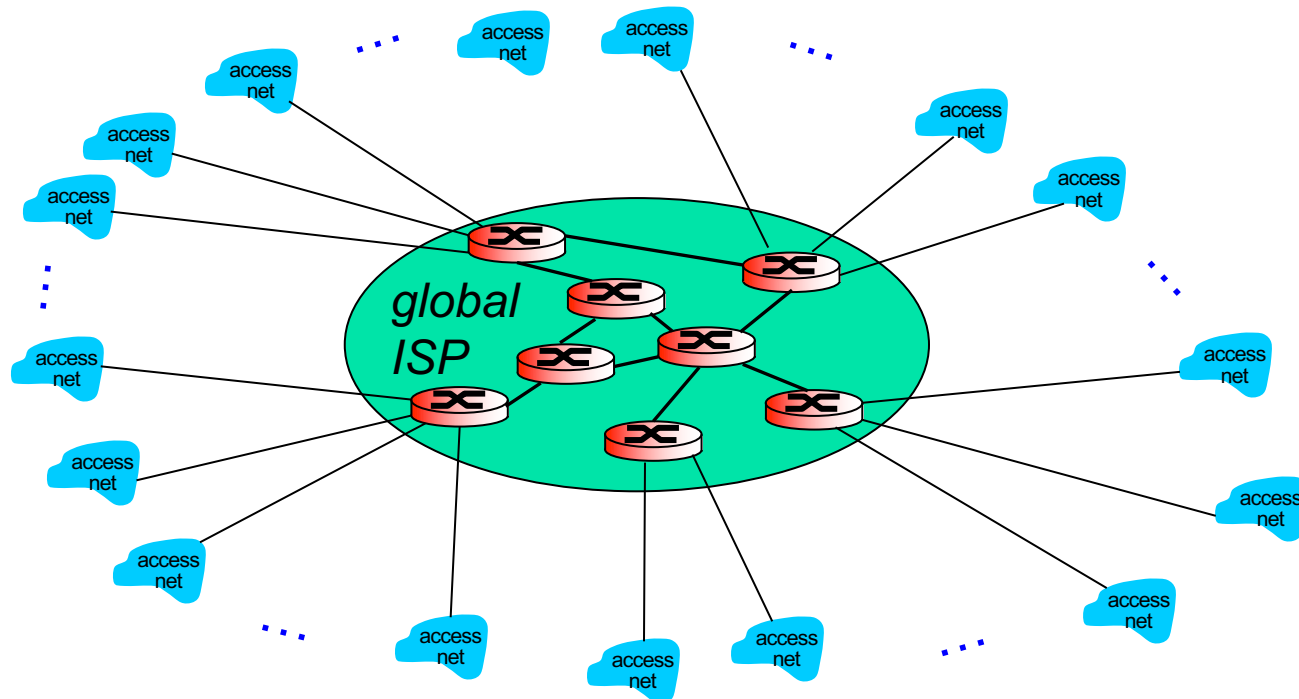
Internet structure: network of networks

Option: connect each access ISP to every other access ISP?



Internet structure: network of networks

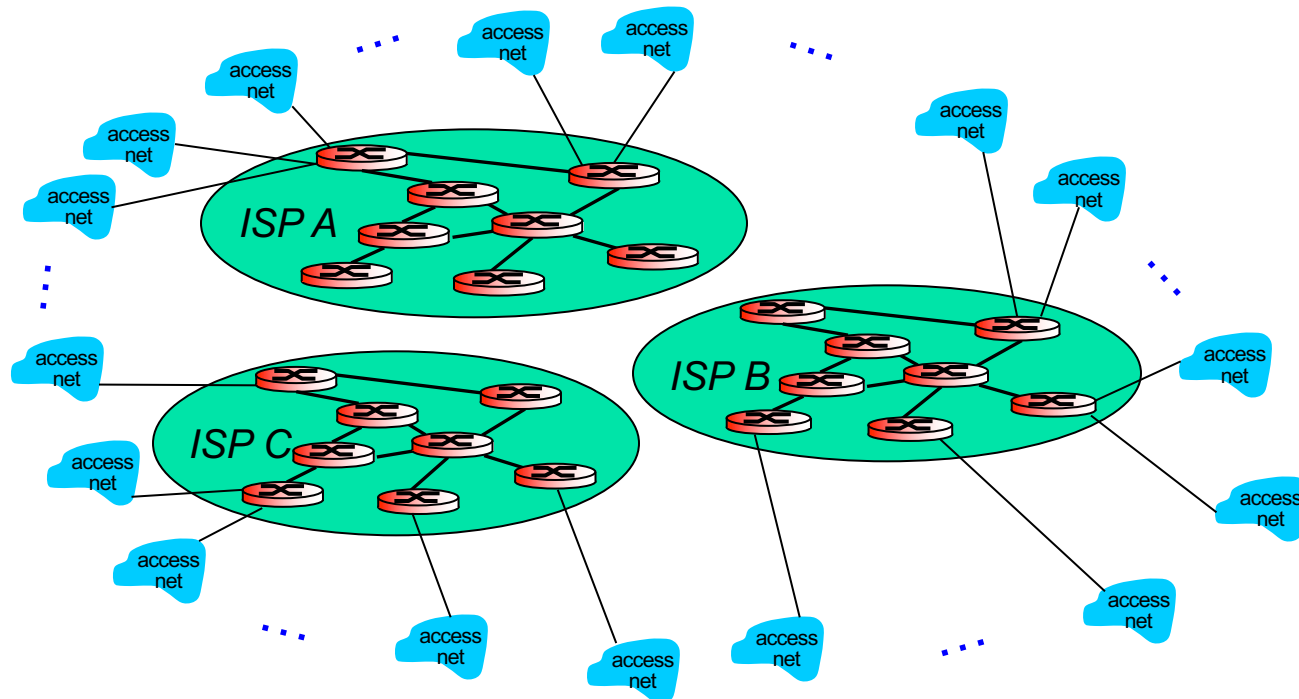
*Option: connect each access ISP to a global transit ISP?
Customer and provider ISPs have economic agreement.*



Internet structure: network of networks

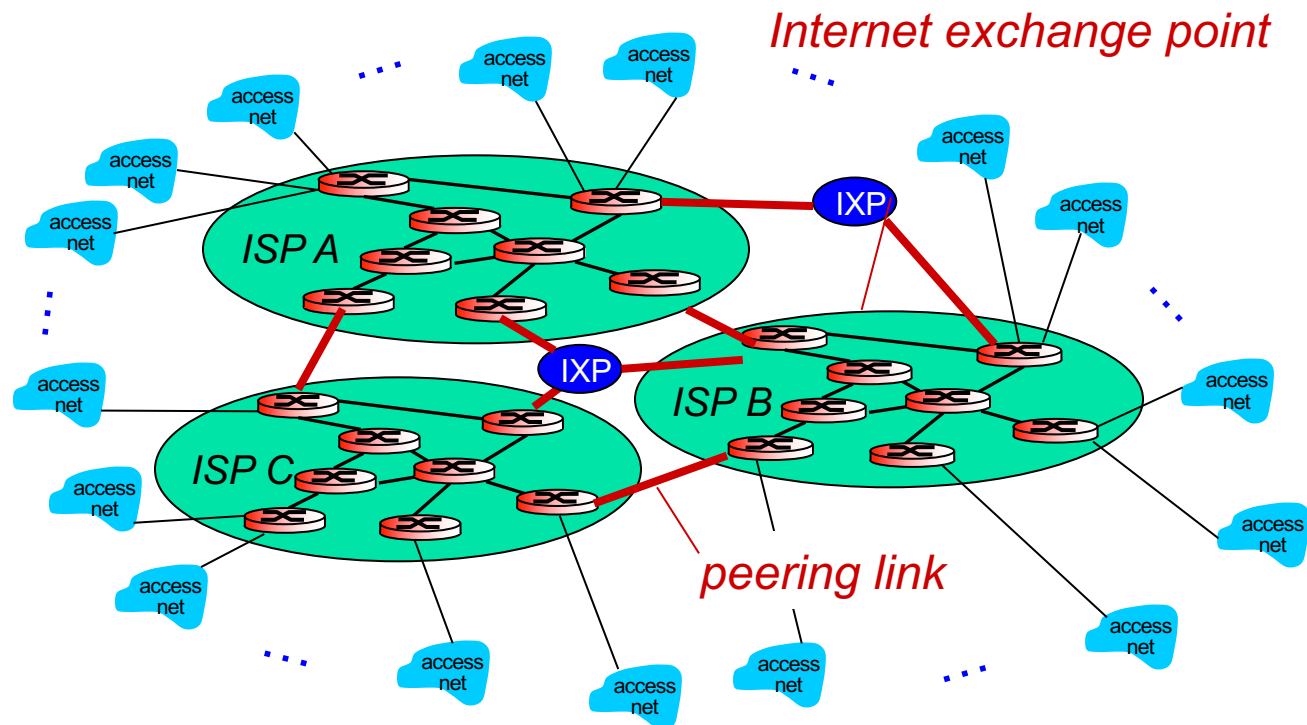
But if one global ISP is viable business, there will be competitors

....



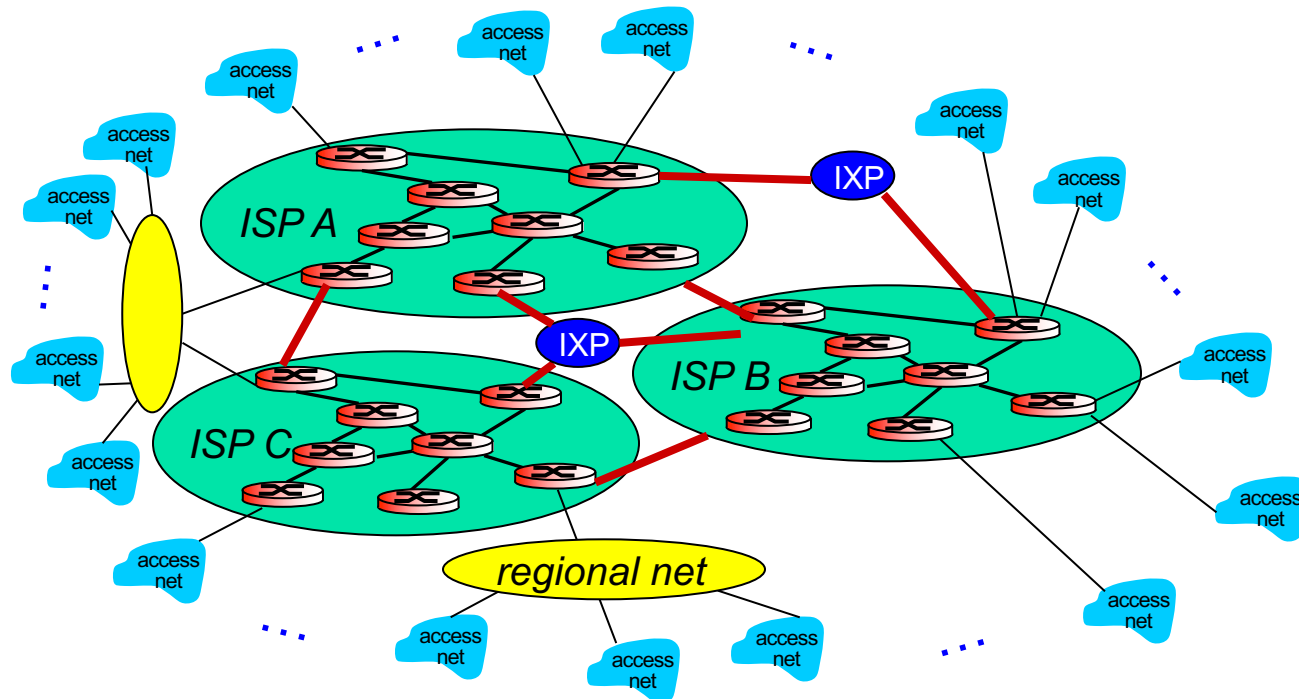
Internet structure: network of networks

But if one global ISP is viable business, there will be competitors
.... which must be interconnected



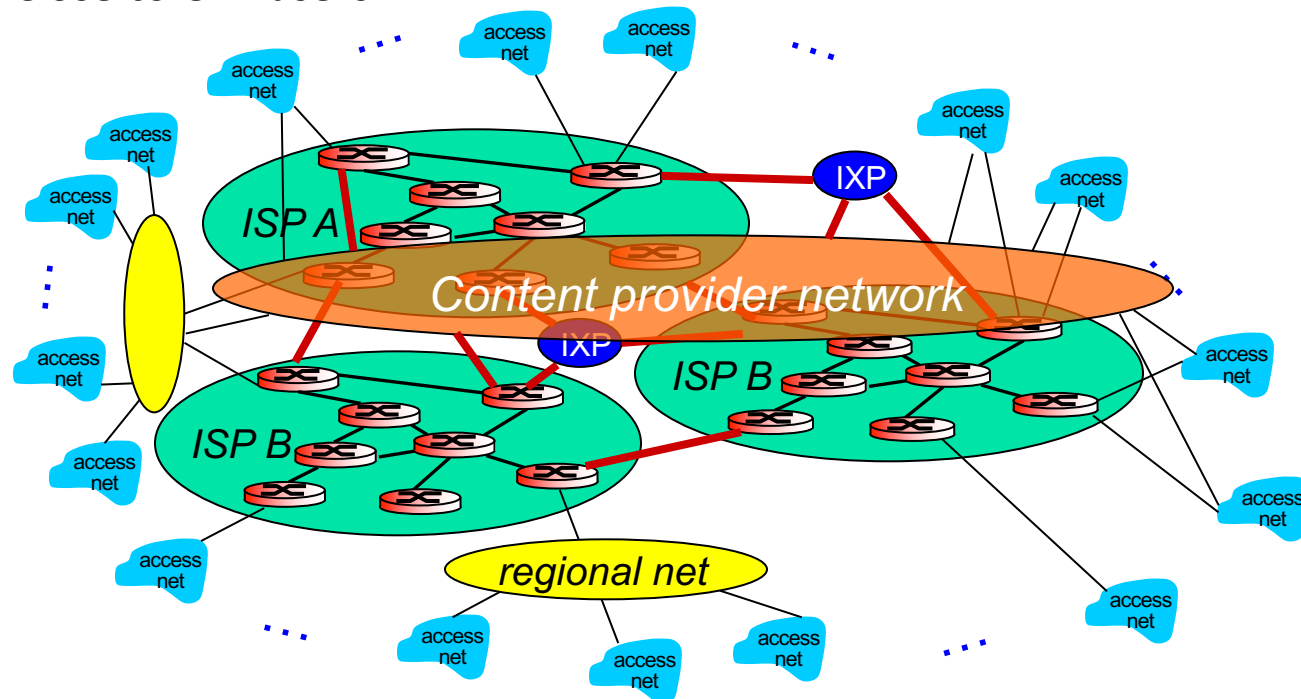
Internet structure: network of networks

... and regional networks may arise to connect access nets to ISPS

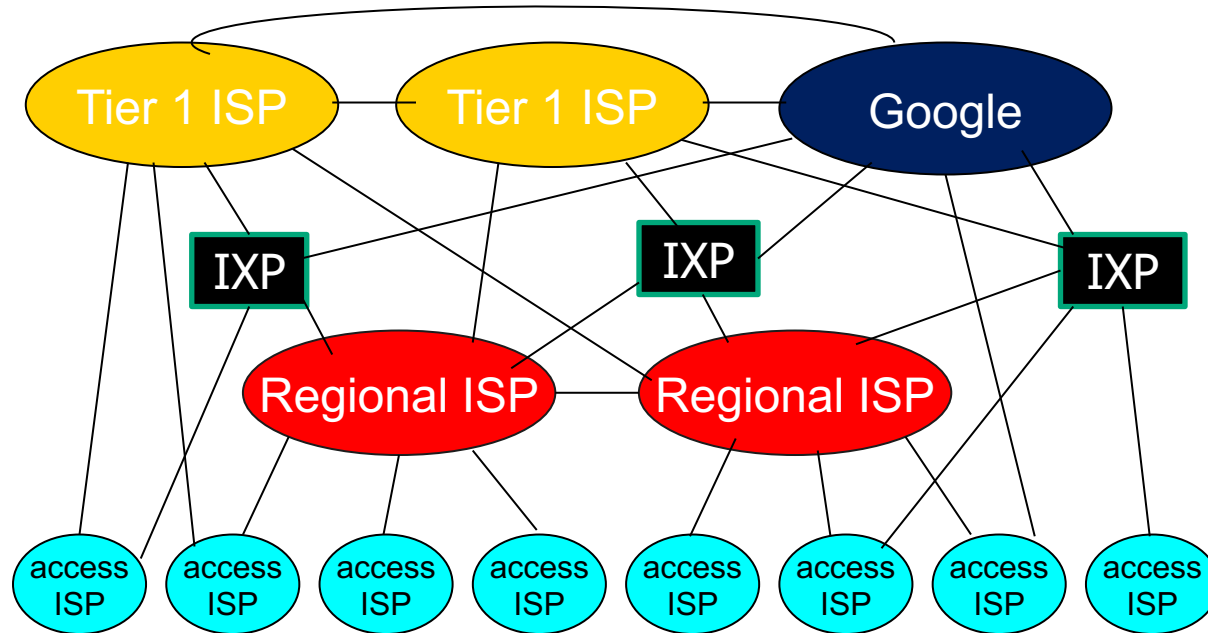


Internet structure: network of networks

... and content provider networks (e.g., Google, Microsoft, Akamai) may run their own network, to bring services, content close to end users



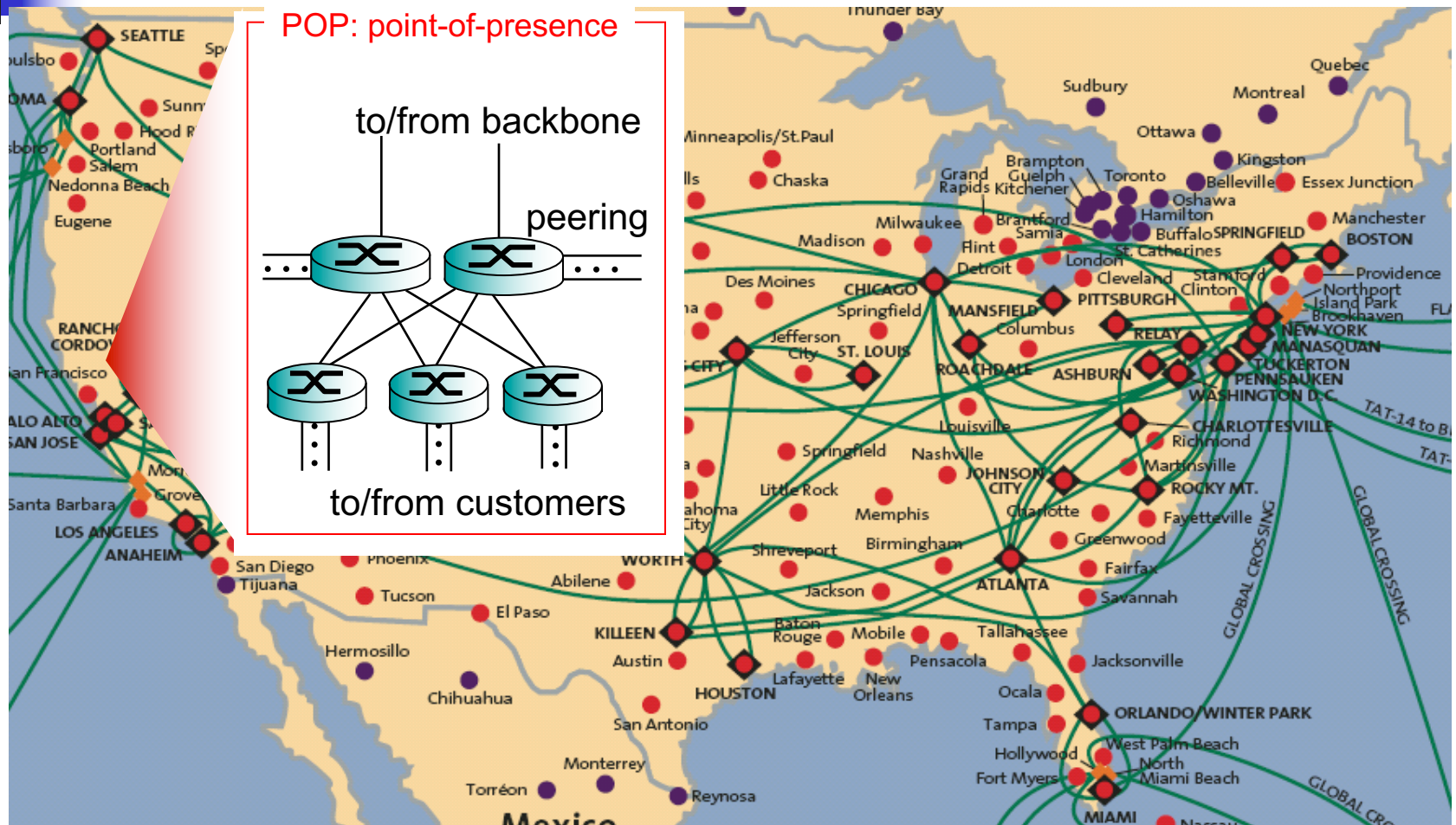
Internet structure: network of networks



at center: small # of well-connected large networks

- “tier-1” commercial ISPs (e.g., Level 3, Sprint, AT&T, NTT), national & international coverage
- content provider network (e.g., Google): private network that connects its data centers to Internet, often bypassing tier-1, regional ISPs

Sprint Network Structure





Next Lecture

1. The Layers!