

# ChIP - Data processing

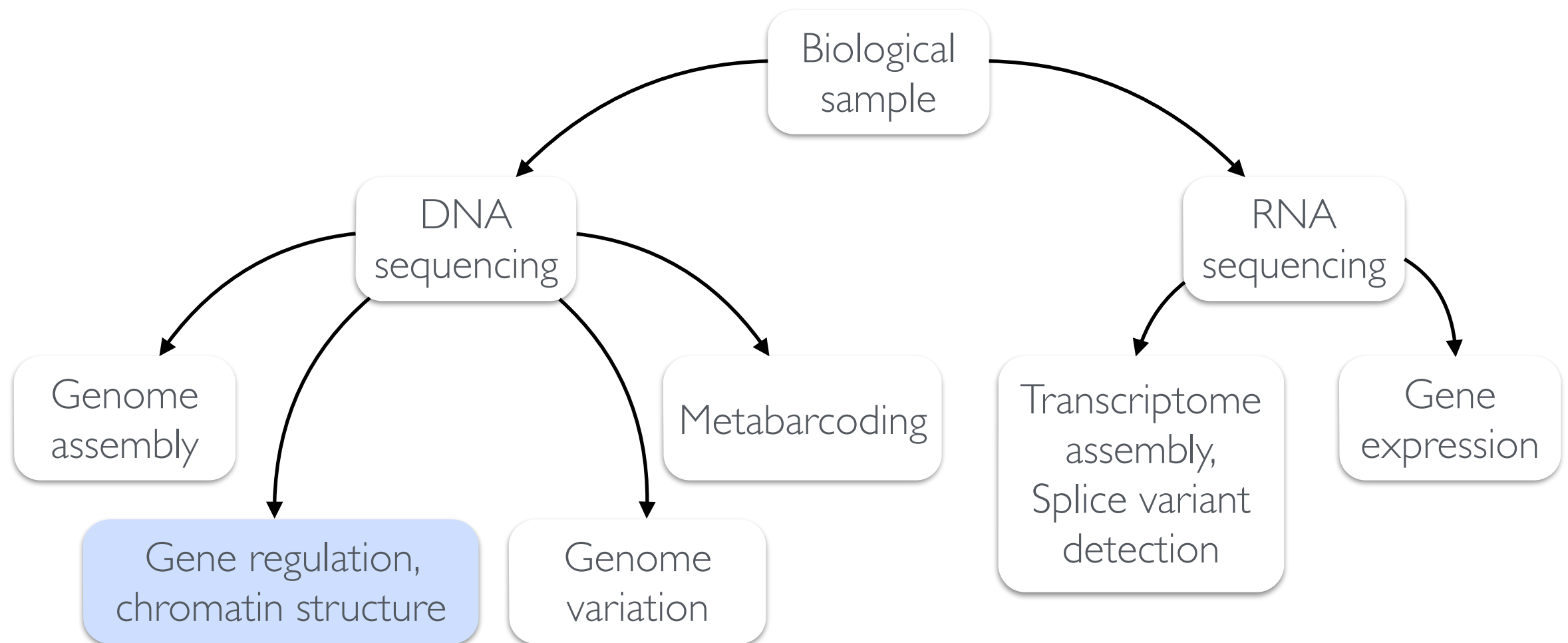
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[http://sschmeier.github.io/bioinf-workshop/  
2015](http://sschmeier.github.io/bioinf-workshop/2015)



# Common analyses overview

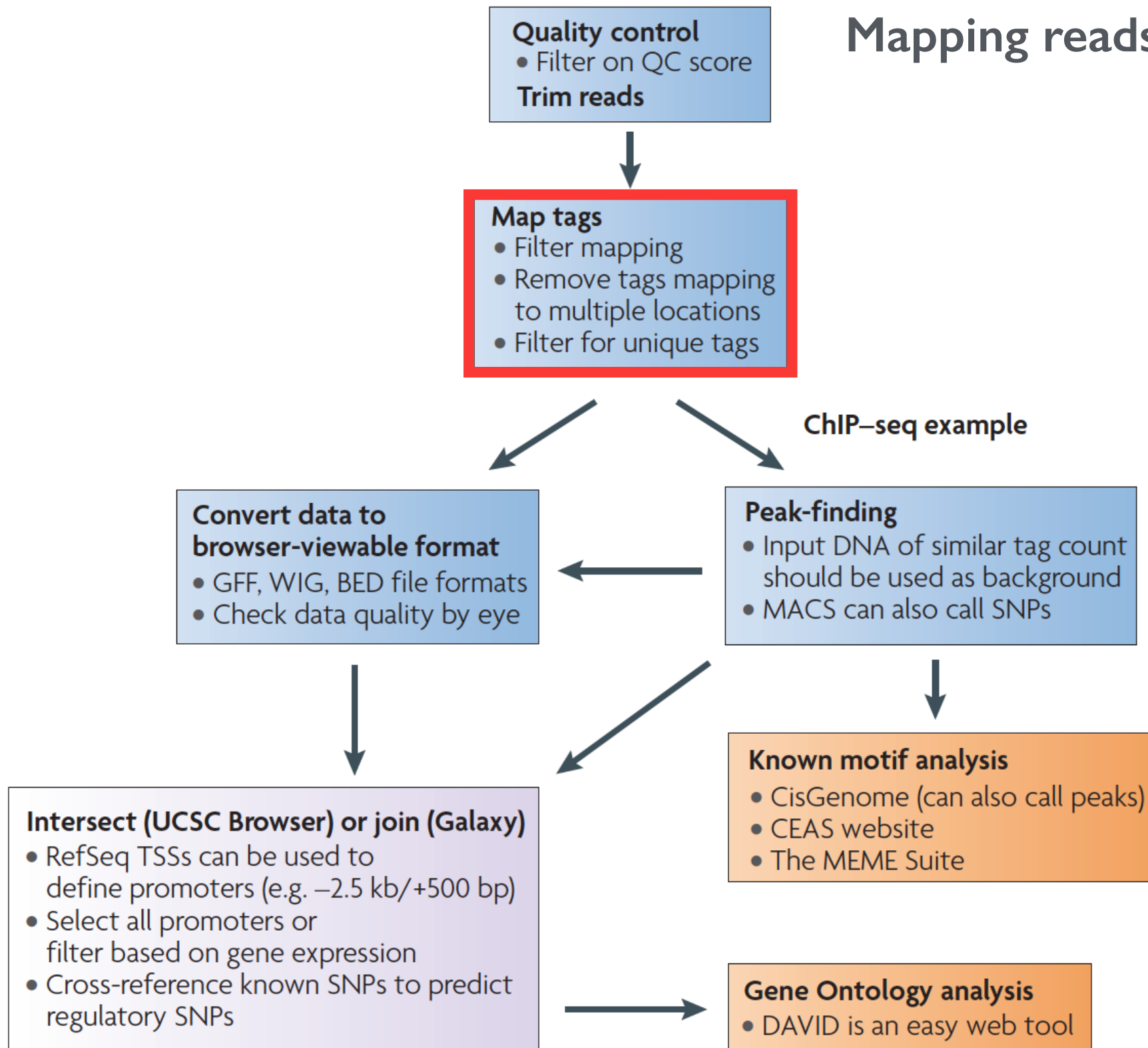


# Gene regulation, chromatin structure

- How do we analyse it?
  - Mapping reads to a reference genome
  - Calling peaks

# Chromatin immunoprecipitation (ChIP)

## Mapping reads



# Mapping reads

- Challenges
- Approximate String Matching Problem
- Burrows-Wheeler transform
- Bowtie

# Challenges of mapping short reads

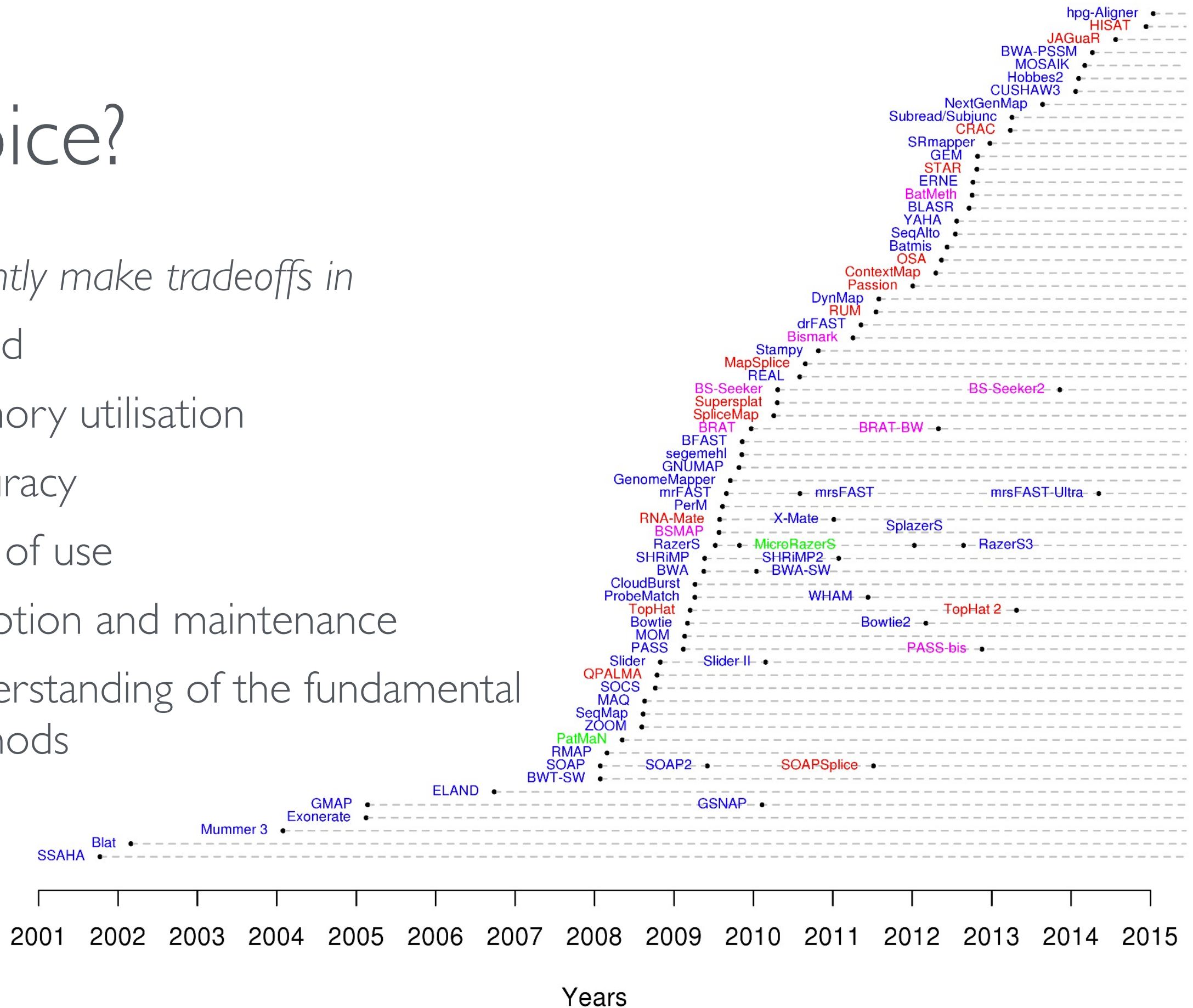
- If the reference genome is very large, and if we have billions of reads, how quickly can we align the reads to the genome?
- The task of mapping billions of sequences to a mammalian-sized genome calls for extraordinarily efficient algorithms, in which every bit of memory is used optimally or near optimally.

# Challenges of mapping short reads

- If a read comes from a repetitive element in the reference, a program must pick which copy of the repeat the read belongs to
  - The program may choose to report multiple possible locations or to pick a location heuristically
  - Sequencing errors or variations between the sequenced chromosomes and the reference genome exacerbate this problem, because the alignment between the read and its true source in the genome may actually have more differences than the alignment between the read and some other copy of the repeat

# Choice?

- *Intelligently make tradeoffs in*
  - Speed
  - Memory utilisation
  - Accuracy
  - Ease of use
  - Adoption and maintenance
  - Understanding of the fundamental methods





# Mapping algorithms

- One could find the true locations using **exact matching**, assuming:
  - a genome had no repeats and a sequencing experiment introduced no errors
  - a sufficient read length relative to the genome size
- Assumption do **NOT** hold

# Mapping algorithms

- **Approximate String Matching Problem**
  - Searching for occurrences of the read sequence within the reference sequence but allowing for *some* mismatches and gaps between the two
- Standard algorithm: *dynamic programming*
  - Too slow
  - Too much memory required

# Mapping algorithms

- **Approximate String Matching Problem**
- Two main ideas for addressing large input sizes (in # of reads and size of the reference):
  - **filtering**
    - quickly exclude large regions of the reference where no approximate match can be found
  - **indexing**
    - *Preprocessing* the reference sequence and/or the set of reads to establish string indices
    - Benefit of preprocessing into string indices is that it typically does not require scanning the whole reference, and it can therefore conduct queries much faster at the expense of larger memory consumption.
    - The string indices that are currently used are:
      - Suffix array
      - Enhanced suffix array
      - **FM-index (Full-text index in Minute space) + Burrows-Wheeler transform**

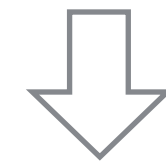
# Burrows-Wheeler transform (BWT)

## Creation

- Write down all rotation of the string
- Sort the matrix lexicographically
- Last column is the BWT(T)
  - The rows in the matrix are essentially the sorted suffixes of the text
- $SA(T)$  is the start offset in the original string

$T = abaaba\$$

|    |    |    |    |    |    |    | $SA(T)$ |
|----|----|----|----|----|----|----|---------|
| \$ | a  | b  | a  | a  | b  | a  | 6       |
| a  | \$ | a  | b  | a  | a  | b  | 5       |
| b  | a  | \$ | a  | b  | a  | a  | 4       |
| a  | b  | a  | \$ | a  | b  | a  | 3       |
| a  | a  | b  | a  | \$ | a  | b  | 2       |
| b  | a  | a  | b  | a  | \$ | a  | 1       |
| a  | b  | a  | a  | b  | a  | \$ | 0       |



|    |    |    |    |    |    |    | $SA(T)$ |
|----|----|----|----|----|----|----|---------|
| \$ | a  | b  | a  | a  | b  | a  | 6       |
| a  | \$ | a  | b  | a  | a  | b  | 5       |
| a  | a  | b  | a  | \$ | a  | b  | 2       |
| a  | b  | a  | \$ | a  | b  | a  | 3       |
| a  | b  | a  | a  | b  | a  | \$ | 0       |
| b  | a  | \$ | a  | b  | a  | a  | 4       |
| b  | a  | a  | b  | a  | \$ | a  | 1       |

**BWT(T) = abba\$aa**

# Burrows-Wheeler transform (BWT)

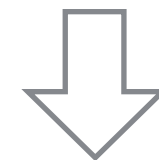
## LF mapping

$T = \text{abaaba}\$$

$\text{BWT}(T) = \text{abba}\$aa$

- We *rank* according to how many times the same character occurred previously in  $\text{BWT}(T)$
- We keep an array of positions in the rotation  $\text{SA}(T)$
- We keep an index of occurrences starting at zero

|    |    |    |    |    |    |    | $\text{SA}(T)$ |
|----|----|----|----|----|----|----|----------------|
| \$ | a  | b  | a  | a  | b  | a  | 6              |
| a  | \$ | a  | b  | a  | a  | b  | 5              |
| a  | a  | b  | a  | \$ | a  | b  | 2              |
| a  | b  | a  | \$ | a  | b  | a  | 3              |
| a  | b  | a  | a  | b  | a  | \$ | 0              |
| b  | a  | \$ | a  | b  | a  | a  | 4              |
| b  | a  | a  | b  | a  | \$ | a  | 1              |



|   | $F$ | $L$ | $rank$ | $\text{SA}(T)$ |
|---|-----|-----|--------|----------------|
|   | \$  | a   | 0      | 6              |
| 0 | a   | b   | 0      | 5              |
| 1 | a   | b   | 1      | 2              |
| 2 | a   | a   | 1      | 3              |
| 3 | a   | \$  | 0      | 0              |
| 0 | b   | a   | 2      | 4              |
| 1 | b   | a   | 3      | 1              |

# Burrows-Wheeler transform (BWT)

*Exact matching*

$T = \text{abaaba\$}$

$\text{BWT}(T) = \text{abba\$aa}$

$P = \text{aba}$

|   | $F$ | $L$ | $rank$ |
|---|-----|-----|--------|
|   | \$  | $a$ | 0      |
| 0 | $a$ | $b$ | 0      |
| 1 | $a$ | $b$ | 1      |
| 2 | $a$ | $a$ | 1      |
| 3 | $a$ | \$  | 0      |
| 0 | $b$ | $a$ | 2      |
| 1 | $b$ | $a$ | 3      |

$P = \text{aba}$

|   | $F$ | $L$ | $rank$ |
|---|-----|-----|--------|
|   | \$  | $a$ | 0      |
| 0 | $a$ | $b$ | 0      |
| 1 | $a$ | $b$ | 1      |
| 2 | $a$ | $a$ | 1      |
| 3 | $a$ | \$  | 0      |
| 0 | $b$ | $a$ | 2      |
| 1 | $b$ | $a$ | 3      |

$P = \text{aba}$

|   | $F$ | $L$ | $rank$ | $\text{SA}(T)$ |
|---|-----|-----|--------|----------------|
|   | \$  | $a$ | 0      | 6              |
| 0 | $a$ | $b$ | 0      | 5              |
| 1 | $a$ | $b$ | 1      | 2              |
| 2 | $a$ | $a$ | 1      | 3              |
| 3 | $a$ | \$  | 0      | 0              |
| 0 | $b$ | $a$ | 2      | 4              |
| 1 | $b$ | $a$ | 3      | 1              |

# Burrows-Wheeler transform (BWT)

*Exact matching*

$T = \text{abaaba\$}$

$\text{BWT}(T) = \text{abba\$aa}$

$P = \text{aba}$

|   | <i>F</i> | <i>L</i> | <i>rank</i> |
|---|----------|----------|-------------|
|   | \$       | <i>a</i> | 0           |
| 0 | <i>a</i> | <i>b</i> | 0           |
| 1 | <i>a</i> | <i>b</i> | 1           |
| 2 | <i>a</i> | <i>a</i> | 1           |
| 3 | <i>a</i> | \$       | 0           |
| 0 | <i>b</i> | <i>a</i> | 2           |
| 1 | <i>b</i> | <i>a</i> | 3           |

$P = \text{aba}$

|   | <i>F</i> | <i>L</i> | <i>rank</i> |
|---|----------|----------|-------------|
|   | \$       | <i>a</i> | 0           |
| 0 | <i>a</i> | <i>b</i> | 0           |
| 1 | <i>a</i> | <i>b</i> | 1           |
| 2 | <i>a</i> | <i>a</i> | 1           |
| 3 | <i>a</i> | \$       | 0           |
| 0 | <i>b</i> | <i>a</i> | 2           |
| 1 | <i>b</i> | <i>a</i> | 3           |

$P = \text{aba}$

|   | <i>F</i> | <i>L</i> | <i>rank</i> | SA( <i>T</i> ) |
|---|----------|----------|-------------|----------------|
|   | \$       | <i>a</i> | 0           | 6              |
| 0 | <i>a</i> | <i>b</i> | 0           | 5              |
| 1 | <i>a</i> | <i>b</i> | 1           | 2              |
| 2 | <i>a</i> | <i>a</i> | 1           | 3              |
| 3 | <i>a</i> | \$       | 0           | 0              |
| 0 | <i>b</i> | <i>a</i> | 2           | 4              |
| 1 | <i>b</i> | <i>a</i> | 3           | 1              |

# Burrows-Wheeler transform (BWT)

*Exact matching*

$T = \text{abaaba}\$$

$\text{BWT}(T) = \text{abba}\$aa$

| P=ab <b>a</b> |          |          |             | P=ab <b>a</b> |          |          |             | P=aba |          |          |                   |
|---------------|----------|----------|-------------|---------------|----------|----------|-------------|-------|----------|----------|-------------------|
|               | <i>F</i> | <i>L</i> | <i>rank</i> |               | <i>F</i> | <i>L</i> | <i>rank</i> |       | <i>F</i> | <i>L</i> | <i>rank</i> SA(T) |
|               | \$       | a        | 0           |               | \$       | a        | 0           |       | \$       | a        | 0 6               |
| 0             | a        | b        | 0           | 0             | a        | b        | 0           | 0     | a        | b        | 0 5               |
| 1             | a        | b        | 1           | 1             | a        | b        | 1           | 1     | a        | b        | 1 2               |
| 2             | a        | a        | 1           | 2             | a        | a        | 1           | 2     | a        | a        | 1 3               |
| 3             | a        | \$       | 0           | 3             | a        | \$       | 0           | 3     | a        | \$       | 0 0               |
| 0             | b        | a        | 2           | 0             | b        | a        | 2           | 0     | b        | a        | 2 4               |
| 1             | b        | a        | 3           | 1             | b        | a        | 3           | 1     | b        | a        | 3 1               |



# Burrows-Wheeler transform (BWT)

*Exact matching*

$T = \text{abaaba}\$$

$\text{BWT}(T) = \text{abba}\$aa$

P=aba

|   | <i>F</i> | <i>L</i> | <i>rank</i> |
|---|----------|----------|-------------|
|   | \$       | a        | 0           |
| 0 | a        | b        | 0           |
| 1 | a        | b        | 1           |
| 2 | a        | a        | 1           |
| 3 | a        | \$       | 0           |
| 0 | b        | a        | 2           |
| 1 | b        | a        | 3           |

P=a**ba**

|   | <i>F</i> | <i>L</i> | <i>rank</i> |
|---|----------|----------|-------------|
|   | \$       | a        | 0           |
| 0 | a        | b        | 0           |
| 1 | a        | b        | 1           |
| 2 | a        | a        | 1           |
| 3 | a        | \$       | 0           |
| 0 | b        | a        | 2           |
| 1 | b        | a        | 3           |

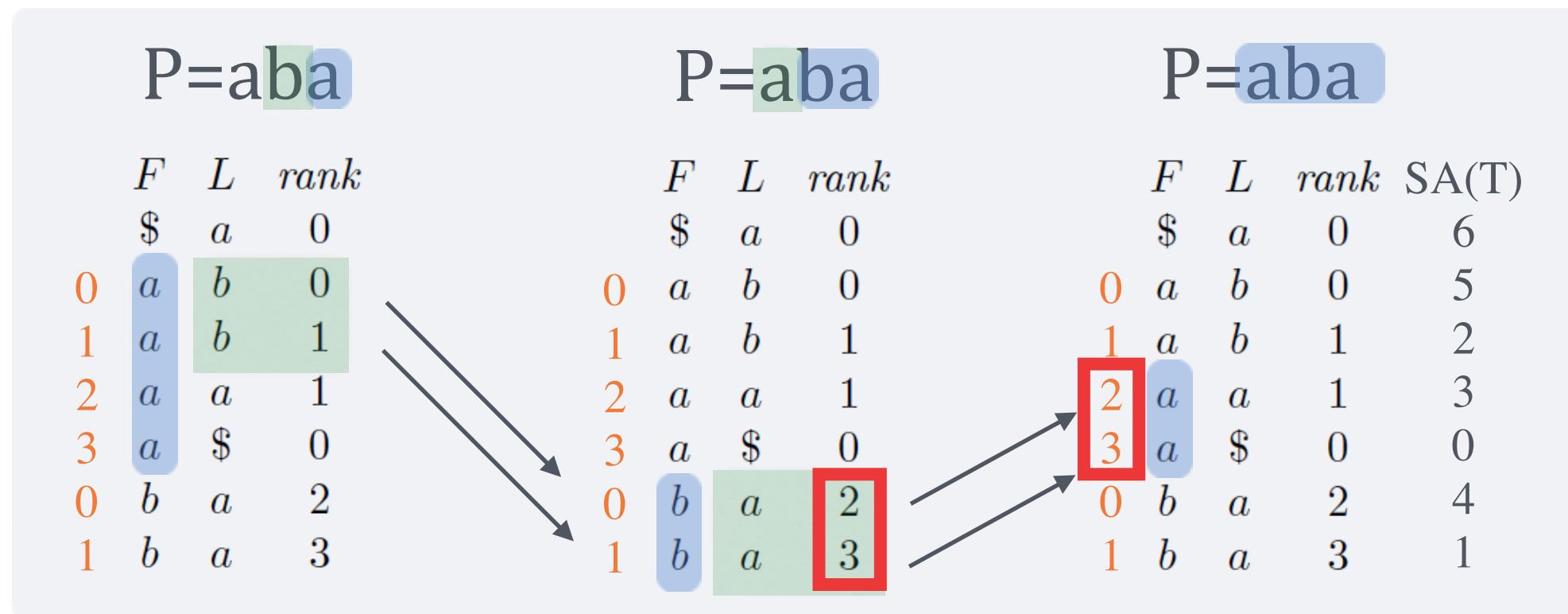
P=aba

|   | <i>F</i> | <i>L</i> | <i>rank</i> | SA(T) |
|---|----------|----------|-------------|-------|
|   | \$       | a        | 0           | 6     |
| 0 | a        | b        | 0           | 5     |
| 1 | a        | b        | 1           | 2     |
| 2 | a        | a        | 1           | 3     |
| 3 | a        | \$       | 0           | 0     |
| 0 | b        | a        | 2           | 4     |
| 1 | b        | a        | 3           | 1     |

# Burrows-Wheeler transform (BWT)

*Exact matching*

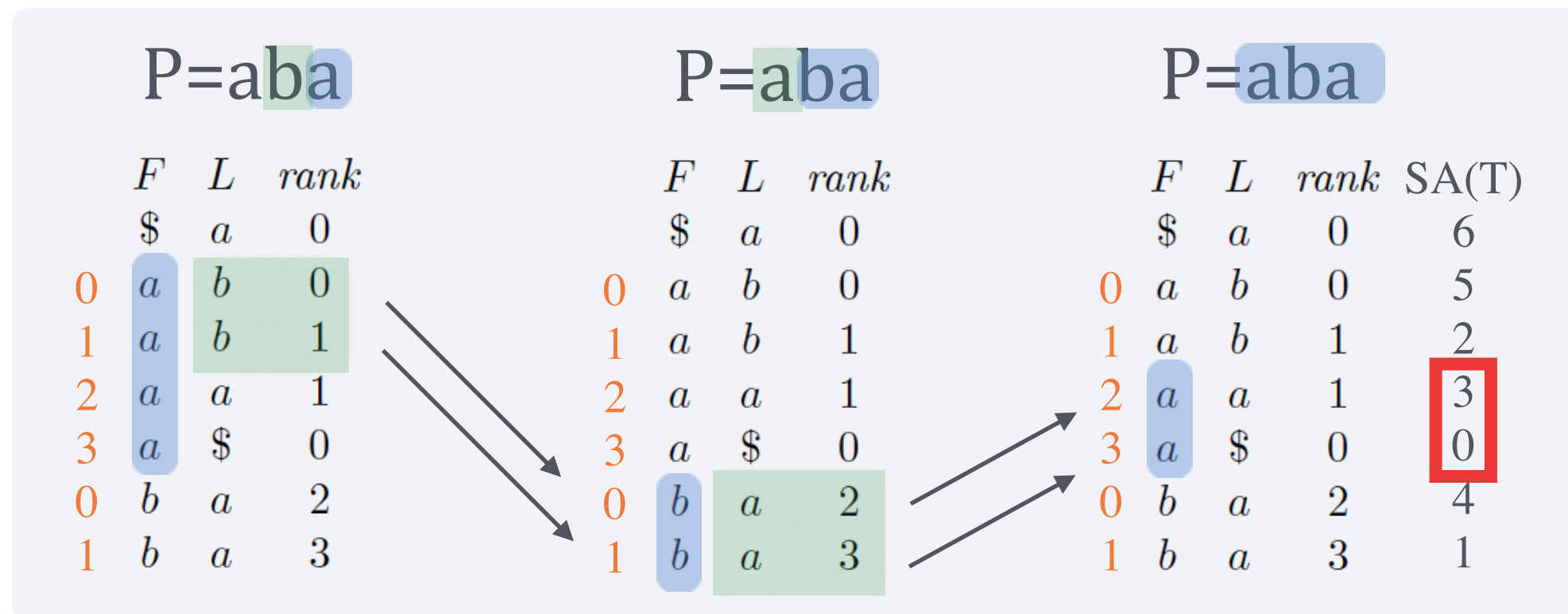
$T = \text{abaaba}\$$   
 $\text{BWT}(T) = \text{abba}\$aa$



# Burrows-Wheeler transform (BWT)

## Exact matching

$T = \text{abaaba}\$$   
 $\text{BWT}(T) = \text{abba}\$aa$



- Allows for matching in constant time

$T = \text{abaaba}\$$   
0 3

## Resource

# Mapping short DNA sequencing reads and calling variants using mapping quality scores

Heng Li,<sup>1</sup> Jue Ruan,<sup>2</sup> and Richard Durbin<sup>1,3</sup>

<sup>1</sup>The Wellcome Trust Sanger Institute, Hinxton CB10 1SA, United Kingdom; <sup>2</sup>Beijing Genomics Institute, Chinese Academy of Science, Beijing 100029, China

## Software

Open Access

# Ultrafast and memory-efficient alignment of short DNA sequences to the human genome

Ben Langmead, Cole Trapnell, Mihai Pop and Steven L Salzberg

Address: Center for Bioinformatics and Computational Biology, Institute for Advanced Computer Studies, University of Maryland, College Park, MD 20742, USA.

Correspondence: Ben Langmead. Email: langmead@cs.umd.edu

Published: 4 March 2009

Genome **Biology** 2009, **10**:R25 (doi:10.1186/gb-2009-10-3-r25)

Received: 21 October 2008

Revised: 19 December 2008

Accepted: 4 March 2009

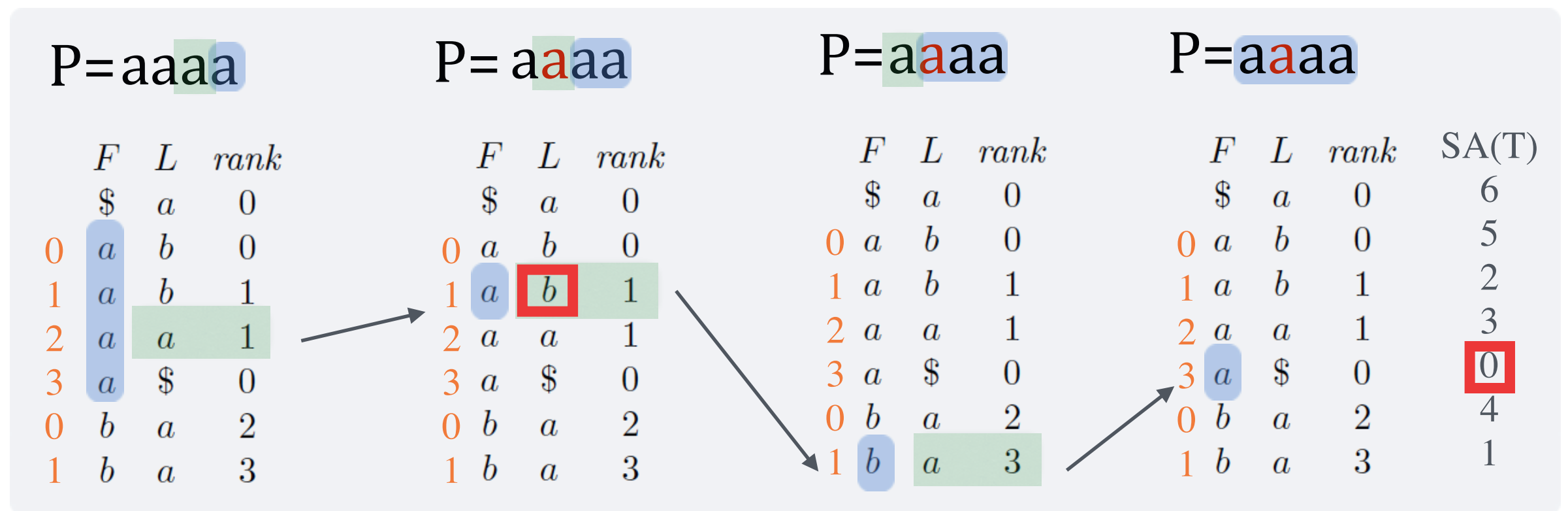
# Bowtie

- FM Index finds exact sequence matches quickly in small memory, but short read alignment demands more:
  - Allowances for mismatches
  - Consideration of quality values

# Bowtie

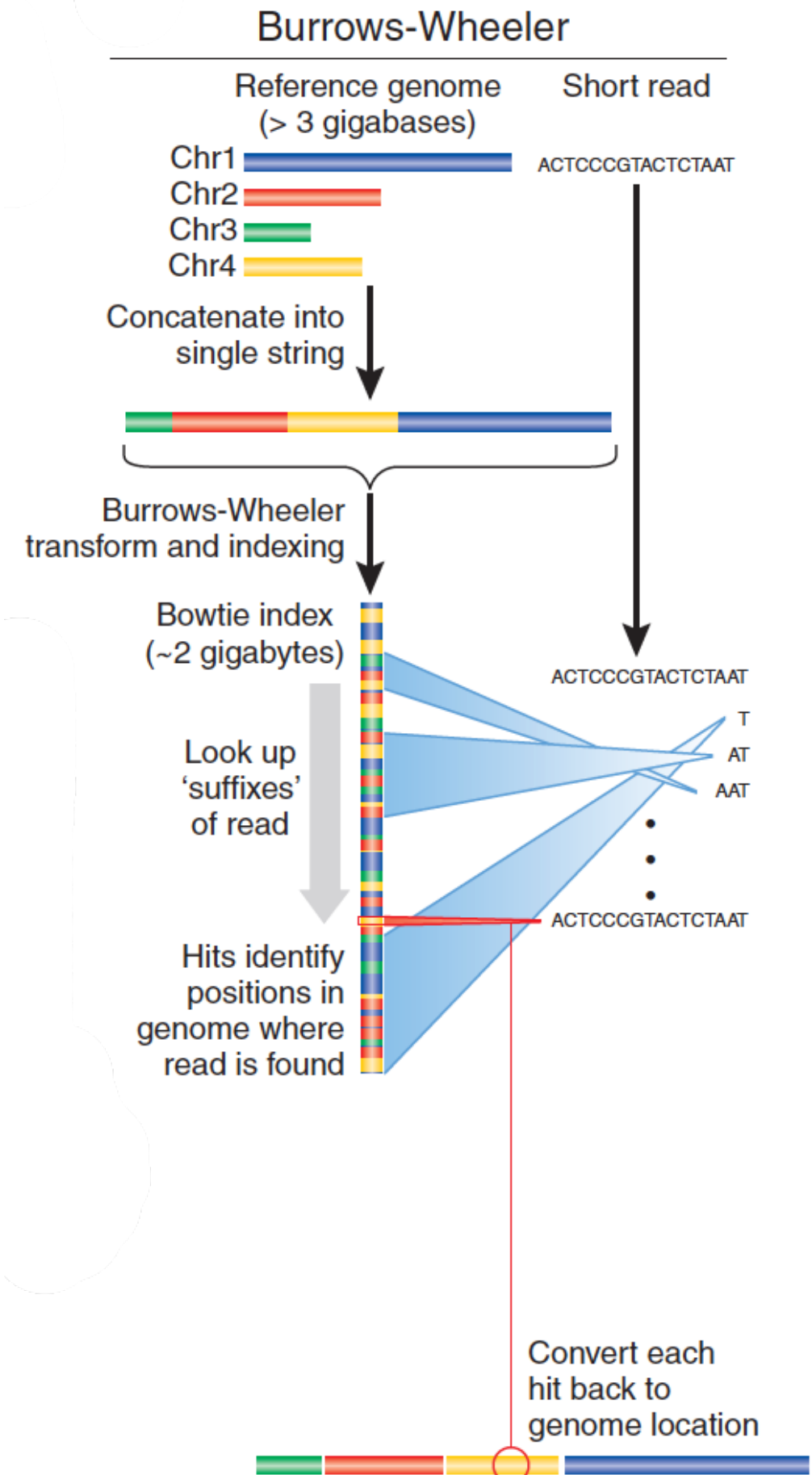
- Bowtie's solution: *backtracking quality-aware* search
  - if a particular base is not found in the index, while traversing the matrix, backtrack and try another “base” based on quality and continue with the search string

$T = \text{abaaba}\$$



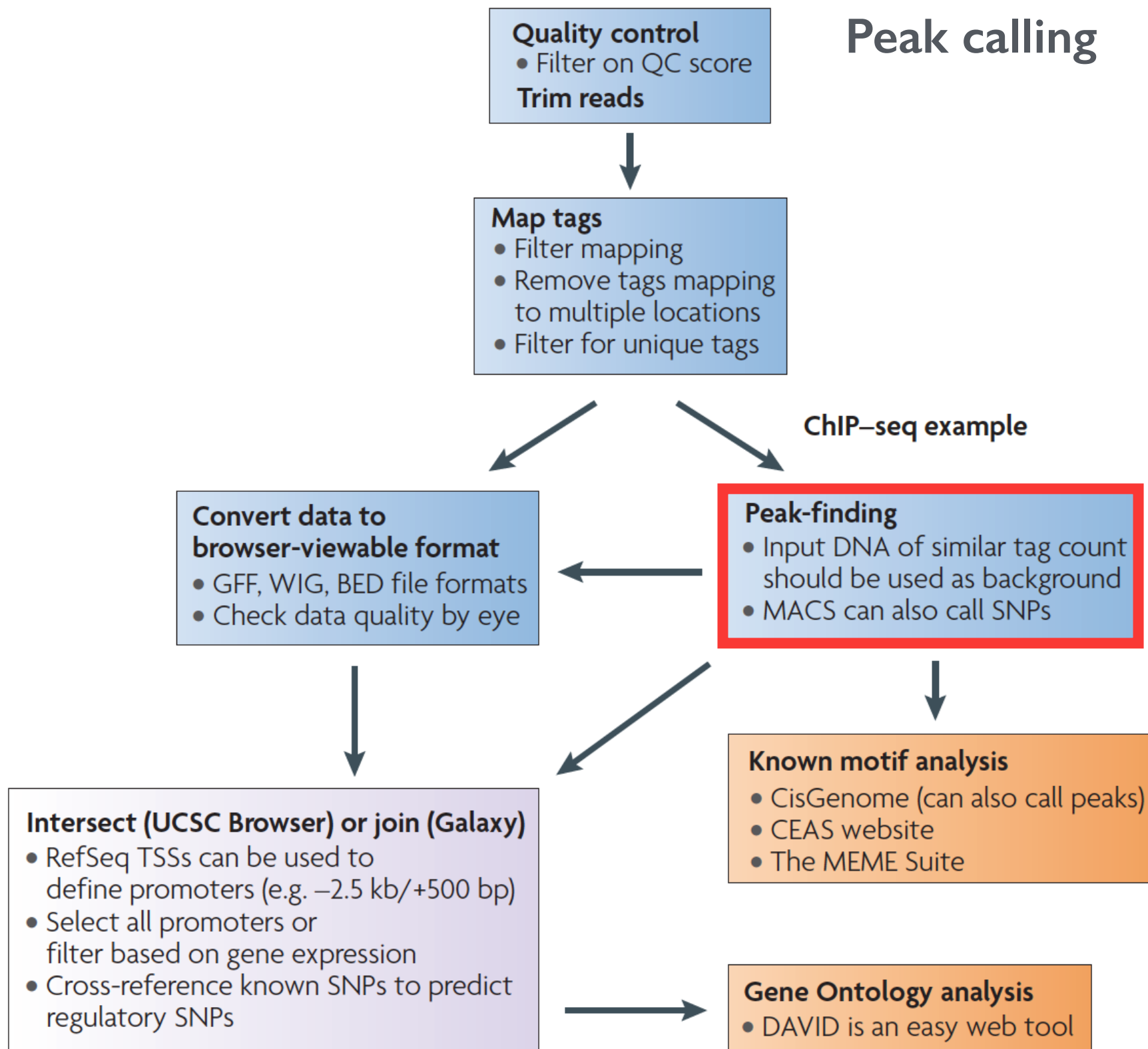
# Burrows-Wheeler transform *genome scale*

- Some clever tricks involved to achieve more compression of the data structures (FM-Index\*)
- Use BWT on the reference genome to build the index
- Look up each read
- Convert to genome locations



# Chromatin immunoprecipitation (ChIP)

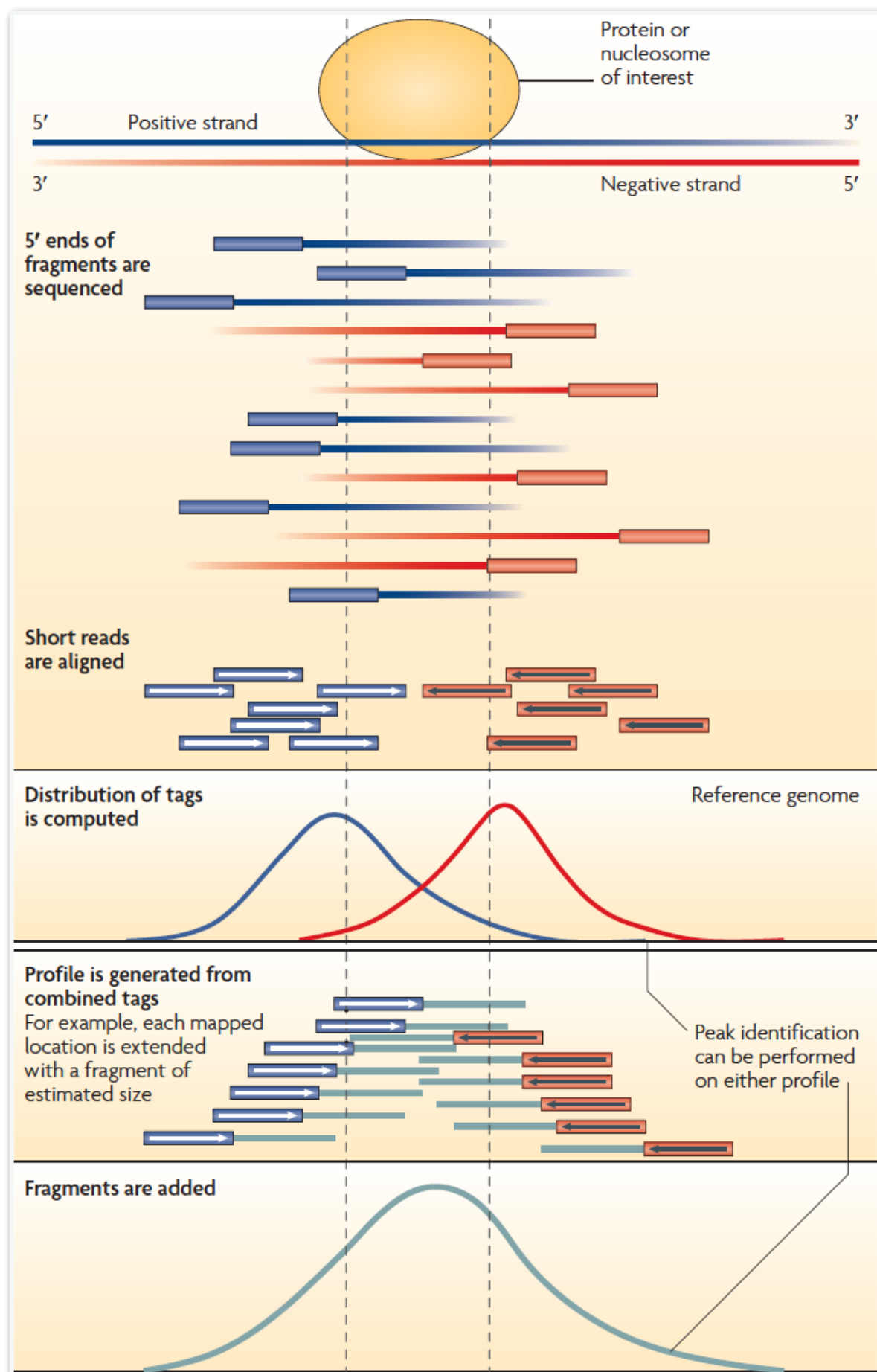
## Peak calling





# Peak calling

- ChIP profile
- Challenges
- MACS



## ChIP profile

- Only 5' ends of ChIPed fragments are sequenced
  - Shifted read distribution
  - Expected symmetry between Watson/Crick read distributions

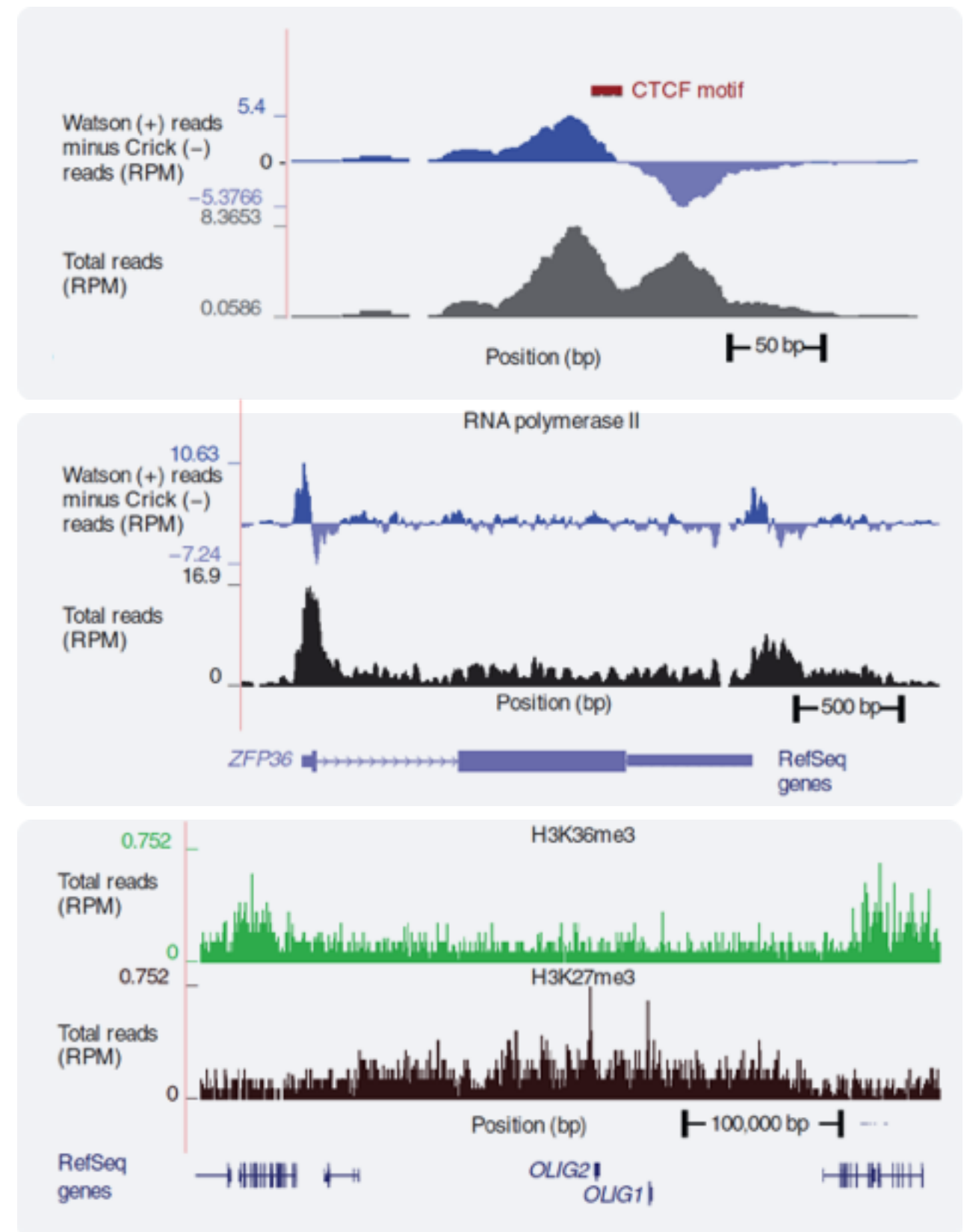
# Peak calling challenges

- Adjust for sequence mappability - regions that contain repetitive elements have different expected tag count

| Organism                       | Genome size (Mb) | Nonrepetitive sequence |            | Mappable sequence |            |
|--------------------------------|------------------|------------------------|------------|-------------------|------------|
|                                |                  | Size (Mb)              | Percentage | Size (Mb)         | Percentage |
| <i>Caenorhabditis elegans</i>  | 100.28           | 87.01                  | 86.8%      | 93.26             | 93.0%      |
| <i>Drosophila melanogaster</i> | 168.74           | 117.45                 | 69.6%      | 121.40            | 71.9%      |
| <i>Mus musculus</i>            | 2,654.91         | 1,438.61               | 54.2%      | 2,150.57          | 81.0%      |
| <i>Homo sapiens</i>            | 3,080.44         | 1,462.69               | 47.5%      | 2,451.96          | 79.6%      |

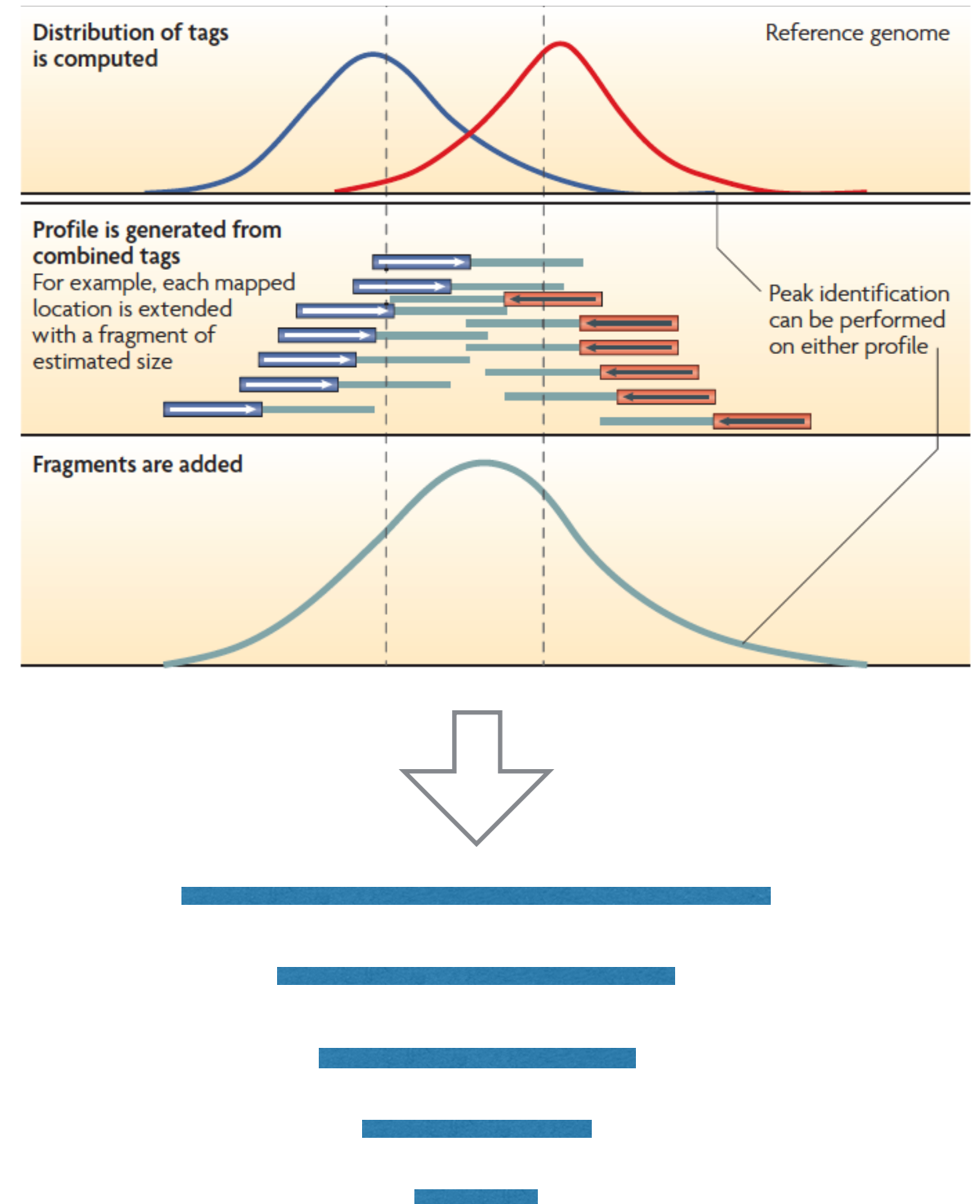
# Peak calling challenges

- Different ChIP-seq applications produce different type of peaks.
- Most current tools have been designed to detect sharp peaks (TF binding, histone modifications at regulatory elements)



# Peak calling challenges

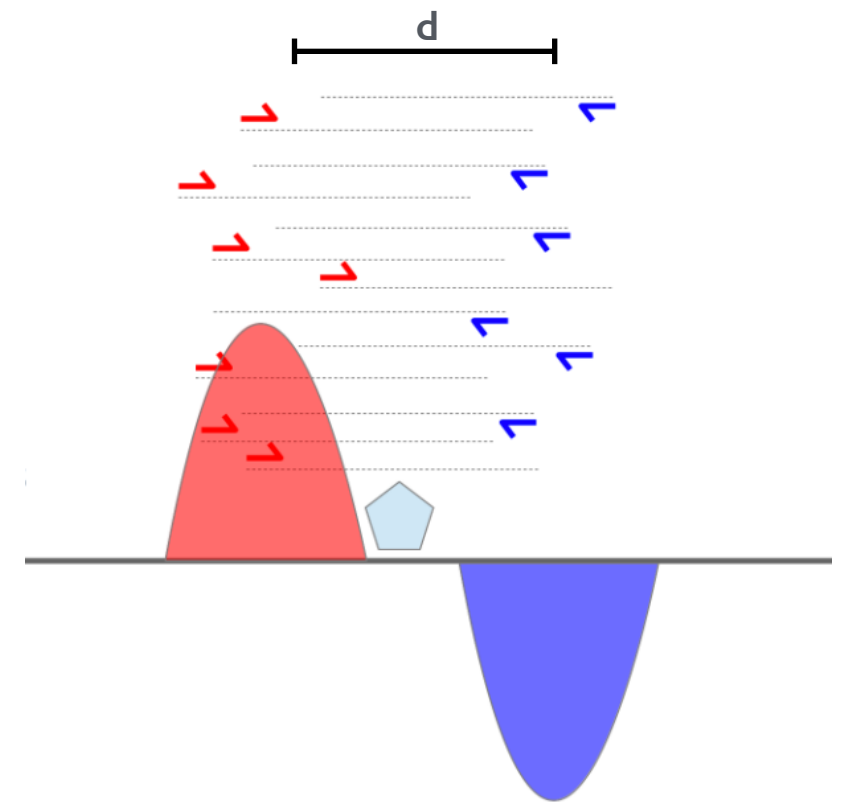
- Definition of enriched regions/peaks:
  - Which statistic to use?
  - What boundaries should be reported?
  - What score to use (ratio, p-val, q-val)?
  - Compute/estimate a FDR?



# MACS

## Step 1: Modelling the tag shift

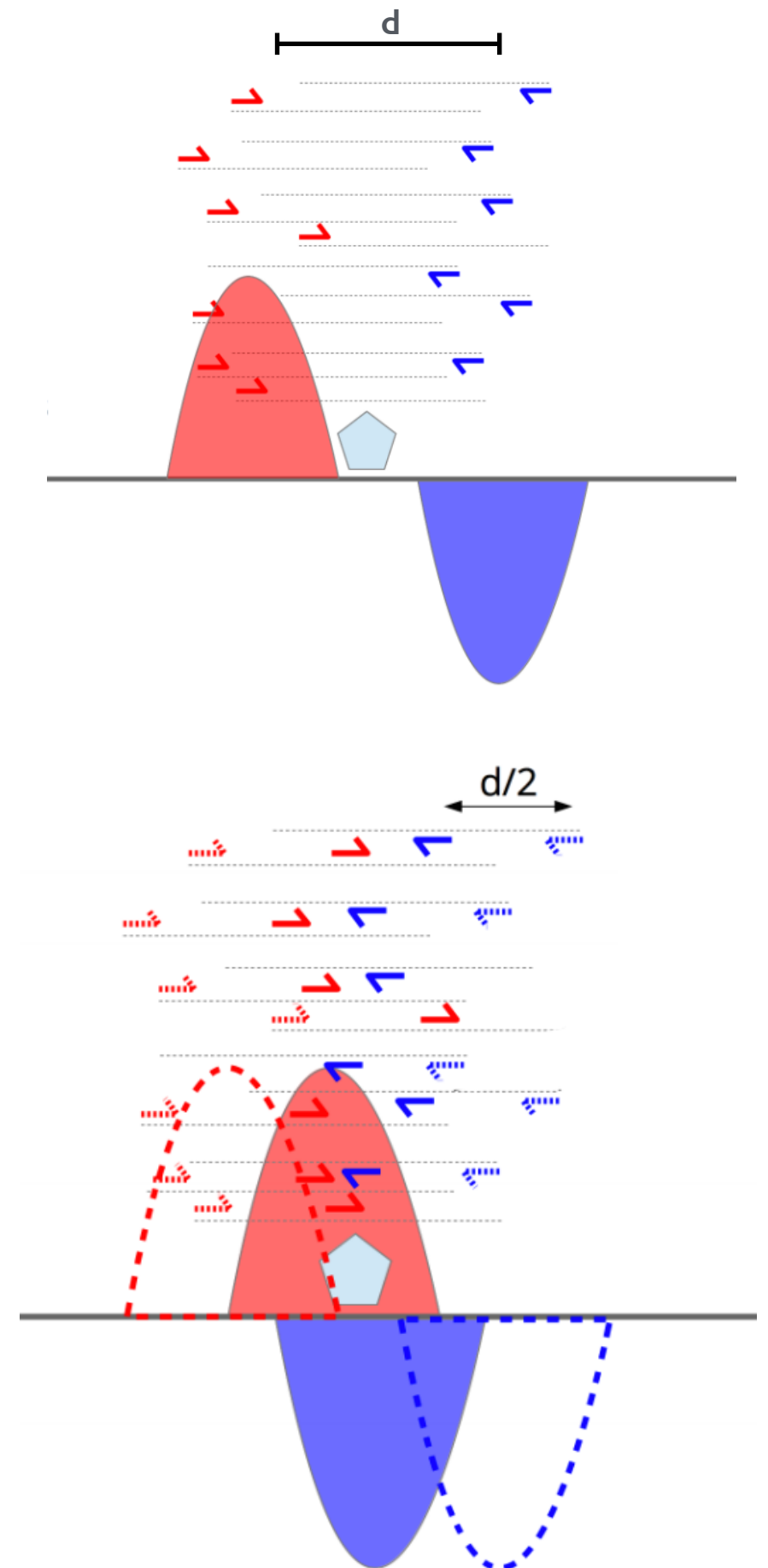
1. Scan genome with a window of user-defined sonication size
2. Keep the best 1000 (or less) peaks having a fold enr.  $>$  mfold (default 32, relative to random model)
3. Separate Watson/Crick tags
4. Shift size is modelled as the distance  $d$  between the modes of the Watson and Crick peaks



# MACS

## Step 2: Peak detection

1. Shift every tag by  $d/2$
2. Slide a  $2d$  window across the genome to find candidate peaks with significant tag enrichment (according to Poisson distribution, default p-value =  $10^{-5}$ )
3. Merge overlapping peaks
4. Report:
  - fold enrichment for called peaks: ratio between tag counts and expected using Poisson distribution (using input data if provided)
  - Position with highest pile-up is defined as the summit of peak
  - Empiric FDR if control sample is provided (sample swap),  
 $\text{FDR} = \# \text{control peaks} / \# \text{ChIP peaks}$





# Again lots of choice

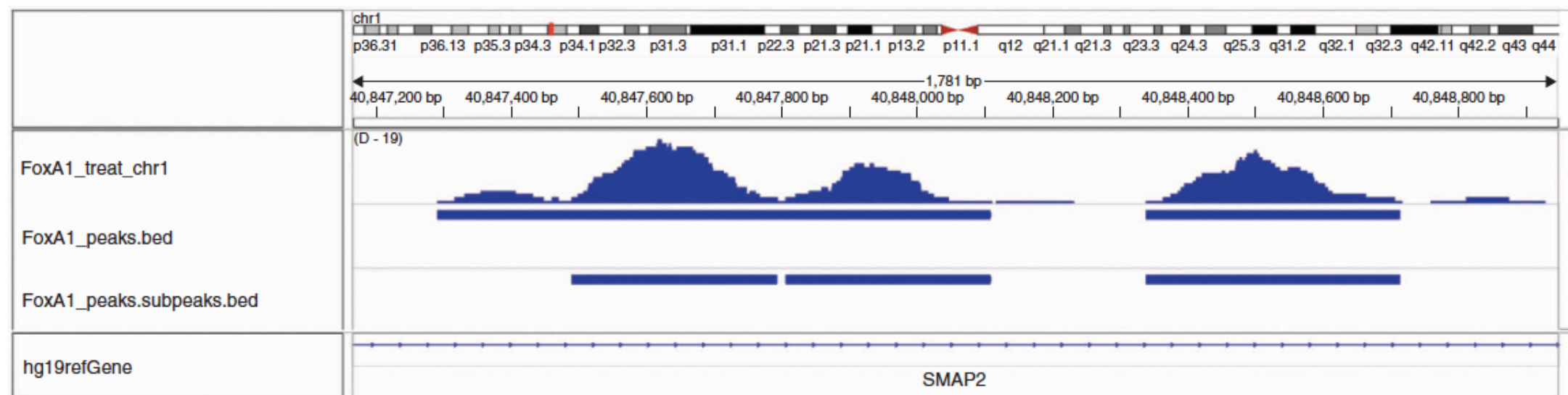
|                    | Profile                         | Peak criteria <sup>a</sup>  | Tag shift   | Control data <sup>b</sup>   | Rank by                         | FDR <sup>c</sup>   | User input parameters <sup>d</sup>                              | Artifact filtering: strand-based/duplicate <sup>e</sup> | Refs. |
|--------------------|---------------------------------|---|---|---|---------------------------------|--|---|---|-------|
| CisGenome v1.1     | Strand-specific window scan     | 1: Number of reads in window<br>2: Number of ChIP reads minus control reads in window | Average for highest ranking peak pairs                    | Conditional binomial used to estimate FDR                             | Number of reads under peak      | 1: Negative binomial<br>2: conditional binomial                              | Target FDR, optional window width, window interval              | Yes / Yes   | 10    |
| ERANGE v3.1        | Tag aggregation                 | 1: Height cutoff<br>High quality peak estimate, per-region estimate, or input         | High quality peak estimate, per-region estimate, or input | Used to calculate fold enrichment and optionally <i>P</i> values      | <i>P</i> value                  | 1: None<br>2: $\frac{\# \text{ control}}{\# \text{ ChIP}}$                   | Optional peak height, ratio to background                       | Yes / No  | 4,18  |
| FindPeaks v3.1.9.2 | Aggregation of overlapped tags  | Height threshold  | Input or estimated  | NA  | Number of reads under peak      | 1: Monte Carlo simulation<br>2: NA   | Minimum peak height, subpeak valley depth                       | Yes / Yes   | 19    |
| F-Seq v1.82        | Kernel density estimation (KDE) | <i>s</i> s.d. above KDE for 1: random background, 2: control                          | Input or estimated  | KDE for local background  | Peak height                     | 1: None<br>2: None   | Threshold s.d. value, KDE bandwidth                             | No / No   | 14    |
| GLITR              | Aggregation of overlapped tags  | Classification by height and relative enrichment                                      | User input tag extension                                  | Multiply sampled to estimate background class values                  | Peak height and fold enrichment | 2: $\frac{\# \text{ control}}{\# \text{ ChIP}}$                              | Target FDR, number nearest neighbors for clustering             | No / No   | 17    |
| MACS v1.3.5        | Tags shifted then window scan   | Local region Poisson <i>P</i> value   | Estimate from high quality peak pairs                     | Used for Poisson fit when available                                   | <i>P</i> value                  | 1: None<br>2: $\frac{\# \text{ control}}{\# \text{ ChIP}}$                   | <i>P</i> -value threshold, tag length, mfold for shift estimate | No / Yes  | 13    |
| PeakSeq            | Extended tag aggregation        | Local region binomial <i>P</i> value  | Input tag extension length                                | Used for significance of sample enrichment with binomial distribution | <i>q</i> value                  | 1: Poisson background assumption<br>2: From binomial for sample plus control | Target FDR  | No / No   | 5     |

and more...



# Visualise to assess quality

- Assess the data quality e.g. positive controls, background
- Determine cutoffs (looking at positive controls)
- Compare different peak finder outputs
- Integration of data / co-visualization



## References

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