

# Contents

<b>1</b>	<b>Chapter 1</b>	<b>1</b>
<b>2</b>	<b>Chapter 9</b>	<b>3</b>

# 1 Chapter 1

## §1.1

finite automaton - finite no. of states, like CPU  
Only current state can be viewed.

Notation

$A$  finite alphabet

Elements of  $A$ : letters

Finite seq. of letters is a word.

Empty sequence/word symbolized by  $\epsilon$ ,  $\Lambda$ , or  $1$

Set of nonempty words symbolized by  $A^+$

Concat non-commutative

$$|uv| = |u| + |v|$$

$$u\epsilon = \epsilon u = u$$

### §1.2.1

Concat product:  $KL = \{uv | u \in K \text{ and } v \in L\}$

Power notation:  $L^n$ , where  $L^0 = \{\epsilon\}$

morphism

For alphabets  $A$  and  $B$ , a morphism from  $A^*$  to  $B^*$  is a mapping.

$\phi: A^* \rightarrow B^*$  s.t. :

$\phi(\epsilon) = \epsilon$

$\forall u, v \in A^*, \phi(u)\phi(v)$

Rational (Regular) Languages:  $\text{Rat}A^*$  is least class of languages over the alphabet,  $A$ , s.t. :

1. the languages  $\emptyset$  and  $\{a\}$  are rational  $\forall a \in A$ ,
2. if  $K$  and  $L$  are rational languages, then  $K \cup L, KL$  and  $L^*$  are also rational.

Extended rational operations: rational ops + intersection, compliment, morphic image (?)

Class of extended rational languages over  $A$  is  $\text{X-Rat}A^*$

§1.2.2

## 2 Chapter 9