



# OpenMP Loop Level Parallelism









# **Outline**

- Work-sharing Constructs
  - loop
- Combined Constructs
- Reduction Clause
- Schedule Clause
- Synchronisation Constructs
  - master, critical, barrier, atomic, flush, ordered
- Monte Carlo Method
- Timing with OpenMP





# **Loop Construct**

- splits up loop iterations among the threads in a team.
- don't create a team of threads;
- there is an implied barrier at the end (unless nowait is used).

```
C/C++:
```

```
#pragma omp parallel [clauses]
   #pragma omp for [clauses] [nowait]
```

### Fortran:

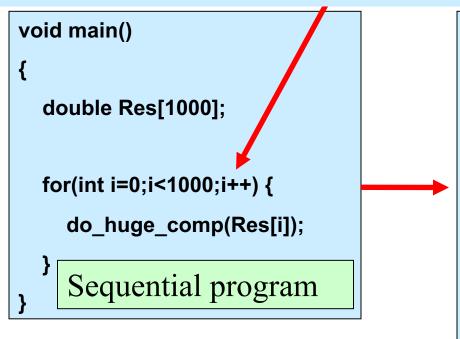
```
!$omp parallel [clauses]
   !$omp do [clauses]
   !$omp end do [nowait]
!$omp end
parallel
```



# **How is OpenMP Typically Used?**

- Find your most time consuming loops.
- Split them up between threads.

### Split-up this loop between multiple threads



```
THREADS
void main()
 double Res[1000];
 #pragma omp parallel
 #pragma omp for
 for(int i=0;i<1000;i++) {
    do_huge_comp(Res[i]);
              Parallel progra
```

**THREADS** 

DO



### **Performance Considerations:**

- Loop dependency
- Cache miss

```
for(i=1; i<n; i++)

for(j=1; j<n; j++)

a[i][j]=2*a[i-1][j];

#pragma omp parallel

#pragma omp for private(i)

for(int j=0; j<n; j++)

for(i=1; i<n; i++)

a[i][j]=2*a[i-1][j];
```





# **Example:** (Parallel matrix vector product y=Ax)

```
C:
#pragma omp parallel
default(none), shared(n, A, x, y),
private(i, j, sum, tid)
tid=omp get thread num();
#pragma omp for
for (i = 0; i < n; i++) {
   sum=0;
   for (j=0; j< n; j++)
      sum+=A[i*n+j]*x[j];
   y[i]=sum;
   printf("%d: y[%d]=%f\n", tid,
i, y[i]);
```

### Fortran: !\$omp parallel default(none), shared(A, x, y, n), private(i, j, sum, tid) tid=omp get thread num() !\$omp do do i=1,nsum=0do j=1,nsum = sum + A((i-1)\*n+j)\*x(j)enddo y(i)=sumprint \*, tid, y(i) enddo !\$omp end do

!\$omp end parallel



```
for(i=0; i<n*n; i++){
    A[i]=i;
}
for(i=0; i<n; i++){
    x[i]=1;
}</pre>
```

### **Compile:** (Intel)

>icc -openmp hello.c -o a.out

>ifort -openmp hello.f90 -o a.out

### **Execute:**

>export OMP\_NUM\_THREADS=4

>./a.out

Thread 0: y[0]=120.000000

Thread 0: y[1]=376.000000

Thread 0: y[2]=632.000000

Thread 0: y[3]=888.000000

Thread 1: y[4]=1144.000000

Thread 1: y[5]=1400.000000

Thread 1: y[6]=1656.000000

Thread 1: y[7]=1912.000000

Thread 2: y[8]=2168.000000

Thread 2: y[9]=2424.000000

Thread 2: y[10]=2680.000000

Thread 2: y[11]=2936.000000

Thread 3: y[12]=3192.000000

Thread 3: y[13]=3448.000000

Thread 3: y[14]=3704.000000

Thread 3: y[15]=3960.000000







# **Combined Constructs**

- The following shortcuts are supported:
  - parallel for / parallel do
  - parallel sections
  - parallel workshare
- Equivalent to a parallel construct followed by a work-sharing construct.

**#pragma omp parallel for** 

Same as

**#pragma omp parallel** #pragma omp for

!\$omp parallel do

Same as

!\$omp parallel !\$omp do





# **Parallel Loop Directives:**

```
C/C++:
#pragma omp parallel for [clause[clause]...]
for ( index = first; index <= last ; index++ ){</pre>
        body of the loop
Fortran:
!$omp parallel do [clause[clause]...]
do index = first, last [, stride]
  body of the loop
enddo
!$omp end parallel do
```





Performs a reduction operation on the variables in the list

C: reduction(operator: list)

<u>Fortran:</u> reduction(*operator*/*intrinsic: list*)

- A local copy of each list variable is created
- Compiler finds the reduction operation and updates local copy
- Combine local copies into a single variable
- for, parallel, sections
- Reduction operations: +, -, \*, ... (C/C++/Fortran)

max, min, ... (Fortran)





# **Example:** (Parallel matrix vector product y=Ax)

```
C:
for (i = 0; i < n; i++) {
  sum=0;
#pragma omp parallel for
default(shared) private(j, tid)
reduction(+:sum)
  for (j=0; j< n; j++){
     sum=sum+(A[i*n+j]*x[j]);
     if(i==0){
        tid=omp get thread num();
        printf("Tid=%d, A[%d]=%f,
sum=%f\n'', tid, i*n+j, A[i*n+j], sum);
  y[i]=sum;
  printf("y[%d]=%f\n", i, y[i]);
```

```
>./a.out
Tid=0, A[0]=0.000000, sum=0.000000
Tid=0, A[1]=1.000000, sum=1.000000
Tid=0, A[2]=2.000000, sum=3.000000
Tid=0, A[3]=3.000000, sum=6.000000
Tid=1, A[4]=4.000000, sum=4.000000
Tid=1, A[5]=5.000000, sum=9.000000
Tid=1, A[6]=6.000000, sum=15.000000
Tid=1, A[7]=7.000000, sum=22.000000
Tid=2, A[8]=8.000000, sum=8.000000
Tid=2, A[9]=9.000000, sum=17.000000
Tid=2, A[10]=10.000000, sum=27.000000
Tid=2, A[11]=11.000000, sum=38.000000
Tid=3, A[12]=12.000000, sum=12.000000
Tid=3, A[13]=13.000000, sum=25
Tid=3, A[14]=14.000000, sum=2
```





# **Schedule Clause**

 Describes how iterations of the loop are divided among the team threads

C: schedule (type [,chunk])

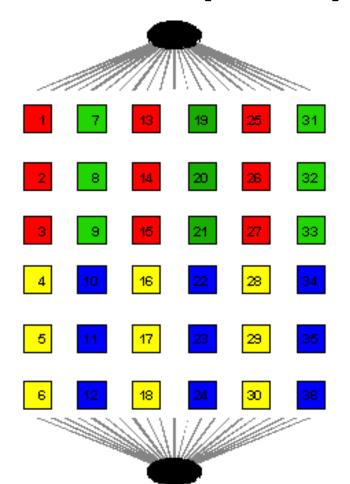
<u>Fortran:</u> schedule (*type [,chunk]*)

- Types:
  - schedule(static [,chunk])
  - schedule(dynamic [,chunk])
  - schedule(guided [,chunk])
  - schedule(runtime)
- More: ordered clause









12/03/2013

- Iterations are divided evenly among threads
- Divides the work into chunk sized parcels
- If there are N threads, each does every N<sup>th</sup> chunk of work

### Fortran:

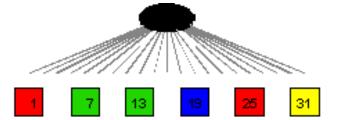
```
!$omp parallel do schedule(static,3)
```

!\$omp end do
OpenMP Loop Level Parallelism





# **SCHEDULE(DYNAMIC):**













12/03/2013

- Divides the workload into chunk sized parcels
- As a thread finishes one chunk, it grabs the next available chunk
- More overhead, but better load balancing

### Fortran:

```
!$omp parallel do schedule(dynamic,1)
```

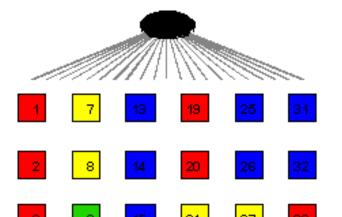
```
do i = 1, 36
     Work (i)
end do
```

!\$omp end do
OpenMP Loop Level Parallelism





# **SCHEDULE(GUIDED):**



- 4 10 16 22 28 34
- 5 11 17 23 29 35
- 6 12 18 24 30 36

- Iterations are divided into chunks such that the size of each successive chunk decreases.
- chunk: the size of the smallest chunk size

### Fortran:

```
!$omp parallel do schedule(guided,1)
```

!\$omp end do
OpenMP Loop Level Parallelism







# **Synchronisation Constructs**

- To control how the execution of each thread proceeds relative to other team threads
- To protect against data races
- Synchronisation Constructs:
  - master
  - critical
  - barrier
  - atomic
  - flush
  - ordered







master

### **Master directive:**

Only master thread can enter the structured block

# Critical directive:

Only one thread at a time can enter a critical section





### **Barrier directive:**

Threads wait each other until all have reached that barrier

Fortran: C/C++:

#pragma omp barrier

!\$omp barrier

### **Atomic directive:**

- Access to a specific memory location atomically
- x binop=expr, x++, --x, ... (C/C++)
- x=x op expr, x=expr op x, ... (Fortran)

C/C++: Fortran:

#pragma omp atomic statement

!\$omp atomic statement







### Flush directive:

A synchronization point at which thread visible variables are flushed

```
Fortran:
C/C++:
                               !$omp flush
#pragma omp flush (list)
                               (list)
```

### **Ordered directive:**

Iterations of the enclosed loop will be executed in the same order as if they were executed sequentially

```
Fortran:
C/C++:
                                   !$omp ordered
#pragma omp ordered
                                   !$omp end ordered
```





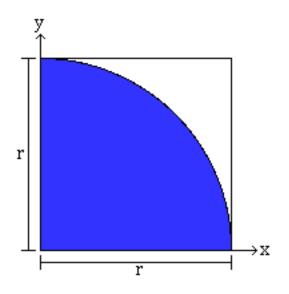
# **Monte Carlo Method**

- Based on the use of random numbers and probability statistics
- Applied in many research areas
  - Financial Analysis
  - Nuclear Physics
  - Population Dynamics
  - Weather Prediction
- Random Number Generation
- Very easy to parallelize





# **Calculation of pi:**



- Area of the shaded circle: πr<sup>2</sup>/4
- Area of the square: r<sup>2</sup>
- Probability=π/4
- r=1 at (0,0)
- distance from the origin: sqrt(x²+y²)
- If the distance  $\leq 1$ , then it is a hit.

•  $\Pi = 4 \times \text{total number of hits/total number of random numbers}$ 





### C:

```
#pragma omp parallel for default(shared), private(i),
reduction(+:hit), schedule(dynamic, 1)
for (i = 0; i < rnums; i++) {
    double x = (double) rand() / RAND_MAX;
    double y = (double) rand() / RAND_MAX;
    if (x * x + y * y <= 1)
        hit++;
}
printf("The number of hits=%d, Estimate of pi = %f\n", hit, 4.0
* hit / rnums);</pre>
```

### True value of pi = 3.141593

rnums=10: The number of hits=8, Estimate of pi = 3.200000 rnums=100: The number of hits=78, Estimate of pi = 3.120000 rnums=1000: The number of hits=783, Estimate of pi = 3.132000 rnums=10000: The number of hits=7909, Estimate of pi = 3.163600 rnums=100000: The number of hits=78524, Estimate of pi = 3.140960



```
C:
#pragma omp parallel default(shared), private(thit, i, tid)
thit=0;
#pragma omp for schedule(dynamic, 1)
for (i = 0; i < rnums; i++) {
   double x = (double) rand() / RAND_MAX;
   double y = (double) rand() / RAND MAX;
   if (x * x + y * y \le 1)
      thit++;
tid=omp get thread num();
printf("Tid=%d thit=%d\n", tid, thit);
#pragma omp critical
   hit+=thit;
```

Tid=0 thit=22484, Tid=3 thit=20419, Tid=2 thit=18034, Tid=1 thit=17728. The number of hits=78665, Estimate of pi = 3.146600







# **Timing with OpenMP**

omp\_get\_wtime

C: double omp\_get\_wtime(void)

<u>Fortran:</u> double precision function omp\_get\_wtime()

 Returns a double-precision floating point value equal to the number of elapsed seconds since some point in the past.

```
double start, end, work;
start = omp_get_wtime();
... work to be timed ...
end = omp_get_wtime();
work=end-start;
```

omp\_get\_wtick

