

Task Scheduler

You are given an array of CPU tasks, each represented by letters A to Z, and a cooling time, n. Each cycle or interval allows the completion of one task. Tasks can be completed in any order, but there's a constraint: **identical** tasks must be separated by at least n intervals due to cooling time.

Return the *minimum number of intervals* required to complete all tasks.

Example 1:

Input: tasks = ["A","A","A","B","B","B"], n = 2

Output: 8

Explanation: A possible sequence is: A -> B -> idle -> A -> B -> idle -> A -> B.

After completing task A, you must wait two cycles before doing A again. The same applies to task B. In the 3rd interval, neither A nor B can be done, so you idle. By the 4th cycle, you can do A again as 2 intervals have passed.

Example 2:

Input: tasks = ["A","C","A","B","D","B"], n = 1

Output: 6

Explanation: A possible sequence is: A -> B -> C -> D -> A -> B.

With a cooling interval of 1, you can repeat a task after just one other task.

Example 3:

Input: tasks = ["A","A","A","B","B","B"], n = 3

Output: 10

Explanation: A possible sequence is: A -> B -> idle -> idle -> A -> B -> idle -> idle -> A -> B.

There are only two types of tasks, A and B, which need to be separated by 3 intervals. This leads to idling twice between repetitions of these tasks.

Constraints:

- $1 \leq \text{tasks.length} \leq 10^4$
- `tasks[i]` is an uppercase English letter.
- $0 \leq n \leq 100$

```
class Solution { public: int leastInterval(vector& tasks, int n) { } };
```

Minimum Number of Arrows to Burst Balloons

There are some spherical balloons taped onto a flat wall that represents the XY-plane. The balloons are represented as a 2D integer array `points` where `points[i] = [xstart, xend]` denotes a balloon whose **horizontal diameter** stretches between `xstart` and `xend`. You do not know the exact y-coordinates of the balloons.

Arrows can be shot up **directly vertically** (in the positive y-direction) from different points along the x-axis. A balloon with `xstart` and `xend` is **burst** by an arrow shot at `x` if `xstart <= x <= xend`. There is **no limit** to the number of arrows that can be shot. A shot arrow keeps traveling up infinitely, bursting any balloons in its path.

Given the array `points`, return *the **minimum** number of arrows that must be shot to burst all balloons.*

Example 1:

Input: `points = [[10,16],[2,8],[1,6],[7,12]]`

Output: 2

Explanation: The balloons can be burst by 2 arrows:

- Shoot an arrow at `x = 6`, bursting the balloons `[2,8]` and `[1,6]`.
- Shoot an arrow at `x = 11`, bursting the balloons `[10,16]` and `[7,12]`.

Example 2:

Input: `points = [[1,2],[3,4],[5,6],[7,8]]`

Output: 4

Explanation: One arrow needs to be shot for each balloon for a total of 4 arrows.

Example 3:

Input: `points = [[1,2],[2,3],[3,4],[4,5]]`

Output: 2

Explanation: The balloons can be burst by 2 arrows:

- Shoot an arrow at `x = 2`, bursting the balloons `[1,2]` and `[2,3]`.
- Shoot an arrow at `x = 4`, bursting the balloons `[3,4]` and `[4,5]`.

Constraints:

- `1 <= points.length <= 105`
- `points[i].length == 2`
- `-231 <= xstart < xend <= 231 - 1`

```
class Solution { public: int findMinArrowShots(vector<vector<int>>& points) { } };
```

Insert Interval

You are given an array of non-overlapping intervals `intervals` where `intervals[i] = [starti, endi]` represent the start and the end of the i^{th} interval and `intervals` is sorted in ascending order by `starti`. You are also given an interval `newInterval = [start, end]` that represents the start and end of another interval.

Insert `newInterval` into `intervals` such that `intervals` is still sorted in ascending order by `starti` and `intervals` still does not have any overlapping intervals (merge overlapping intervals if necessary).

Return `intervals` *after the insertion*.

Note that you don't need to modify `intervals` in-place. You can make a new array and return it.

Example 1:

Input: `intervals = [[1,3],[6,9]]`, `newInterval = [2,5]`

Output: `[[1,5],[6,9]]`

Example 2:

Input: `intervals = [[1,2],[3,5],[6,7],[8,10],[12,16]]`, `newInterval = [4,8]`

Output: `[[1,2],[3,10],[12,16]]`

Explanation: Because the new interval `[4,8]` overlaps with `[3,5]`, `[6,7]`, `[8,10]`.

Constraints:

- $0 \leq \text{intervals.length} \leq 10^4$
- `intervals[i].length == 2`
- $0 \leq \text{start}_i \leq \text{end}_i \leq 10^5$
- `intervals` is sorted by `starti` in **ascending** order.
- `newInterval.length == 2`
- $0 \leq \text{start} \leq \text{end} \leq 10^5$

```
class Solution { public: vector<> insert(vector<>& intervals, vector& newInterval) { } };
```

Peak Index in a Mountain Array

You are given an integer **mountain** array `arr` of length `n` where the values increase to a **peak element** and then decrease.

Return the index of the peak element.

Your task is to solve it in $O(\log(n))$ time complexity.

Example 1:

Input: `arr = [0,1,0]`

Output: 1

Example 2:

Input: `arr = [0,2,1,0]`

Output: 1

Example 3:

Input: `arr = [0,10,5,2]`

Output: 1

Constraints:

- $3 \leq \text{arr.length} \leq 10^5$
- $0 \leq \text{arr}[i] \leq 10^6$
- `arr` is **guaranteed** to be a mountain array.

```
class Solution { public: int peakIndexInMountainArray(vector& arr) { } };
```

Find Minimum in Rotated Sorted Array

Suppose an array of length n sorted in ascending order is **rotated** between 1 and n times. For example, the array `nums = [0,1,2,4,5,6,7]` might become:

- `[4,5,6,7,0,1,2]` if it was rotated 4 times.
- `[0,1,2,4,5,6,7]` if it was rotated 7 times.

Notice that **rotating** an array `[a[0], a[1], a[2], ..., a[n-1]]` 1 time results in the array `[a[n-1], a[0], a[1], a[2], ..., a[n-2]]`.

Given the sorted rotated array `nums` of **unique** elements, return *the minimum element of this array*.

You must write an algorithm that runs in $O(\log n)$ time.

Example 1:

Input: `nums = [3,4,5,1,2]`

Output: 1

Explanation: The original array was `[1,2,3,4,5]` rotated 3 times.

Example 2:

Input: `nums = [4,5,6,7,0,1,2]`

Output: 0

Explanation: The original array was `[0,1,2,4,5,6,7]` and it was rotated 4 times.

Example 3:

Input: `nums = [11,13,15,17]`

Output: 11

Explanation: The original array was `[11,13,15,17]` and it was rotated 4 times.

Constraints:

- $n == \text{nums.length}$
- $1 \leq n \leq 5000$
- $-5000 \leq \text{nums}[i] \leq 5000$
- All the integers of `nums` are **unique**.
- `nums` is sorted and rotated between 1 and n times.

```
class Solution { public: int findMin(vector& nums) { } };
```

Find Peak Element

A peak element is an element that is strictly greater than its neighbors.

Given a **0-indexed** integer array `nums`, find a peak element, and return its index. If the array contains multiple peaks, return the index to **any of the peaks**.

You may imagine that `nums[-1] = nums[n] = -∞`. In other words, an element is always considered to be strictly greater than a neighbor that is outside the array.

You must write an algorithm that runs in $O(\log n)$ time.

Example 1:

Input: `nums = [1,2,3,1]`

Output: `2`

Explanation: 3 is a peak element and your function should return the index number 2.

Example 2:

Input: `nums = [1,2,1,3,5,6,4]`

Output: `5`

Explanation: Your function can return either index number 1 where the peak element is 2, or index number 5 where the peak element is 6.

Constraints:

- $1 \leq \text{nums.length} \leq 1000$
- $-2^{31} \leq \text{nums}[i] \leq 2^{31} - 1$
- $\text{nums}[i] \neq \text{nums}[i + 1]$ for all valid i .

```
class Solution { public: int findPeakElement(vector& nums) { } };
```

Search in Rotated Sorted Array

There is an integer array `nums` sorted in ascending order (with **distinct** values).

Prior to being passed to your function, `nums` is **possibly rotated** at an unknown pivot index k ($1 \leq k < \text{nums.length}$) such that the resulting array is `[nums[k], nums[k+1], ..., nums[n-1], nums[0], nums[1], ..., nums[k-1]]` (**0-indexed**). For example, `[0,1,2,4,5,6,7]` might be rotated at pivot index 3 and become `[4,5,6,7,0,1,2]`.

Given the array `nums` **after** the possible rotation and an integer `target`, return *the index of target if it is in nums, or -1 if it is not in nums*.

You must write an algorithm with $O(\log n)$ runtime complexity.

Example 1:

Input: `nums = [4,5,6,7,0,1,2]`, `target = 0`

Output: 4

Example 2:

Input: `nums = [4,5,6,7,0,1,2]`, `target = 3`

Output: -1

Example 3:

Input: `nums = [1]`, `target = 0`

Output: -1

Constraints:

- $1 \leq \text{nums.length} \leq 5000$
- $-10^4 \leq \text{nums}[i] \leq 10^4$
- All values of `nums` are **unique**.
- `nums` is an ascending array that is possibly rotated.
- $-10^4 \leq \text{target} \leq 10^4$

```
class Solution { public: int search(vector& nums, int target) { } };
```

Search in Rotated Sorted Array II

There is an integer array `nums` sorted in non-decreasing order (not necessarily with **distinct** values).

Before being passed to your function, `nums` is **rotated** at an unknown pivot index `k` ($0 \leq k < \text{nums.length}$) such that the resulting array is `[nums[k], nums[k+1], ..., nums[n-1], nums[0], nums[1], ..., nums[k-1]]` (**0-indexed**). For example, `[0,1,2,4,4,4,5,6,6,7]` might be rotated at pivot index 5 and become `[4,5,6,6,7,0,1,2,4,4]`.

Given the array `nums` **after** the rotation and an integer `target`, return `true` *if target is in nums*, or `false` *if it is not in nums*.

You must decrease the overall operation steps as much as possible.

Example 1:

Input: `nums = [2,5,6,0,0,1,2]`, `target = 0`

Output: `true`

Example 2:

Input: `nums = [2,5,6,0,0,1,2]`, `target = 3`

Output: `false`

Constraints:

- $1 \leq \text{nums.length} \leq 5000$
- $-10^4 \leq \text{nums}[i] \leq 10^4$
- `nums` is guaranteed to be rotated at some pivot.
- $-10^4 \leq \text{target} \leq 10^4$

Follow up: This problem is similar to [Search in Rotated Sorted Array](#), but `nums` may contain **duplicates**. Would this affect the runtime complexity? How and why?

```
class Solution { public: bool search(vector& nums, int target) { } };
```


Search a 2D Matrix

You are given an $m \times n$ integer matrix `matrix` with the following two properties:

- Each row is sorted in non-decreasing order.
- The first integer of each row is greater than the last integer of the previous row.

Given an integer `target`, return `true` *if target is in matrix* or `false` *otherwise*.

You must write a solution in $O(\log(m * n))$ time complexity.

Example 1:

1	3	5	7
10	11	16	20
23	30	34	60

Input: `matrix = [[1,3,5,7],[10,11,16,20],[23,30,34,60]]`, `target = 3`

Output: `true`

Example 2:

1	3	5	7
10	11	16	20
23	30	34	60

Input: `matrix = [[1,3,5,7],[10,11,16,20],[23,30,34,60]]`, `target = 13`

Output: `false`

Constraints:

- `m == matrix.length`
- `n == matrix[i].length`
- `1 <= m, n <= 100`

- $-10^4 \leq \text{matrix}[i][j], \text{target} \leq 10^4$

```
class Solution { public: bool searchMatrix(vector<vector<int>>& matrix, int target) { } };
```

Search a 2D Matrix II

Write an efficient algorithm that searches for a value `target` in an $m \times n$ integer matrix `matrix`. This matrix has the following properties:

- Integers in each row are sorted in ascending from left to right.
- Integers in each column are sorted in ascending from top to bottom.

Example 1:

1	4	7	11	15
2	5	8	12	19
3	6	9	16	22
10	13	14	17	24
18	21	23	26	30

Input: `matrix = [[1,4,7,11,15],[2,5,8,12,19],[3,6,9,16,22],[10,13,14,17,24],[18,21,23,26,30]]`,
`target = 5`
Output: `true`

Example 2:

1	4	7	11	15
2	5	8	12	19
3	6	9	16	22
10	13	14	17	24
18	21	23	26	30

Input: `matrix = [[1,4,7,11,15],[2,5,8,12,19],[3,6,9,16,22],[10,13,14,17,24],[18,21,23,26,30]]`,
`target = 20`
Output: `false`

Constraints:

- `m == matrix.length`
- `n == matrix[i].length`
- `1 <= n, m <= 300`
- `-109 <= matrix[i][j] <= 109`
- All the integers in each row are **sorted** in ascending order.
- All the integers in each column are **sorted** in ascending order.
- `-109 <= target <= 109`

```
class Solution { public: bool searchMatrix(vector<vector<int>>& matrix, int target) { } };
```

Find K Closest Elements

Given a **sorted** integer array `arr`, two integers `k` and `x`, return the `k` closest integers to `x` in the array. The result should also be sorted in ascending order.

An integer `a` is closer to `x` than an integer `b` if:

- $|a - x| < |b - x|$, or
- $|a - x| == |b - x|$ and $a < b$

Example 1:

Input: `arr = [1,2,3,4,5]`, `k = 4`, `x = 3`

Output: `[1,2,3,4]`

Example 2:

Input: `arr = [1,2,3,4,5]`, `k = 4`, `x = -1`

Output: `[1,2,3,4]`

Constraints:

- $1 \leq k \leq \text{arr.length}$
- $1 \leq \text{arr.length} \leq 10^4$
- `arr` is sorted in **ascending** order.
- $-10^4 \leq \text{arr}[i], x \leq 10^4$

```
class Solution { public: vector findClosestElements(vector& arr, int k, int x) { } };
```

Minimum Size Subarray Sum

Given an array of positive integers `nums` and a positive integer `target`, return *the **minimal length** of a subarray whose sum is greater than or equal to target*. If there is no such subarray, return 0 instead.

Example 1:

Input: `target = 7, nums = [2,3,1,2,4,3]`

Output: 2

Explanation: The subarray `[4,3]` has the minimal length under the problem constraint.

Example 2:

Input: `target = 4, nums = [1,4,4]`

Output: 1

Example 3:

Input: `target = 11, nums = [1,1,1,1,1,1,1,1]`

Output: 0

Constraints:

- $1 \leq \text{target} \leq 10^9$
- $1 \leq \text{nums.length} \leq 10^5$
- $1 \leq \text{nums}[i] \leq 10^4$

Follow up: If you have figured out the $O(n)$ solution, try coding another solution of which the time complexity is $O(n \log(n))$.

```
class Solution { public: int minSubArrayLen(int target, vector& nums) { } };
```

Fruit Into Baskets

You are visiting a farm that has a single row of fruit trees arranged from left to right. The trees are represented by an integer array `fruits` where `fruits[i]` is the **type** of fruit the i^{th} tree produces.

You want to collect as much fruit as possible. However, the owner has some strict rules that you must follow:

- You only have **two** baskets, and each basket can only hold a **single type** of fruit. There is no limit on the amount of fruit each basket can hold.
- Starting from any tree of your choice, you must pick **exactly one fruit** from **every** tree (including the start tree) while moving to the right. The picked fruits must fit in one of your baskets.
- Once you reach a tree with fruit that cannot fit in your baskets, you must stop.

Given the integer array `fruits`, return *the **maximum** number of fruits you can pick*.

Example 1:

Input: `fruits = [1,2,1]`

Output: 3

Explanation: We can pick from all 3 trees.

Example 2:

Input: `fruits = [0,1,2,2]`

Output: 3

Explanation: We can pick from trees [1,2,2].

If we had started at the first tree, we would only pick from trees [0,1].

Example 3:

Input: `fruits = [1,2,3,2,2]`

Output: 4

Explanation: We can pick from trees [2,3,2,2].

If we had started at the first tree, we would only pick from trees [1,2].

Constraints:

- $1 \leq \text{fruits.length} \leq 10^5$
- $0 \leq \text{fruits}[i] < \text{fruits.length}$

```
class Solution { public: int totalFruit(vector& fruits) { } };
```

Permutation in String

Given two strings `s1` and `s2`, return `true` if `s2` contains a permutation of `s1`, or `false` otherwise.

In other words, return `true` if one of `s1`'s permutations is the substring of `s2`.

Example 1:

Input: `s1 = "ab", s2 = "eidbaooo"`

Output: `true`

Explanation: `s2` contains one permutation of `s1` ("ba").

Example 2:

Input: `s1 = "ab", s2 = "eidboao"`

Output: `false`

Constraints:

- $1 \leq s1.length, s2.length \leq 10^4$
- `s1` and `s2` consist of lowercase English letters.

```
class Solution { public: bool checkInclusion(string s1, string s2) { } };
```


