

Data Structures

SYBTech(CSE)

Unit – 1

Introduction

PPT-4

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Unit 1 [6 hrs]

Introduction:

Data,

Data types,

Data structure,

Abstract Data Type (ADT),

Representation of Information,

Characteristics of algorithm,

Program,

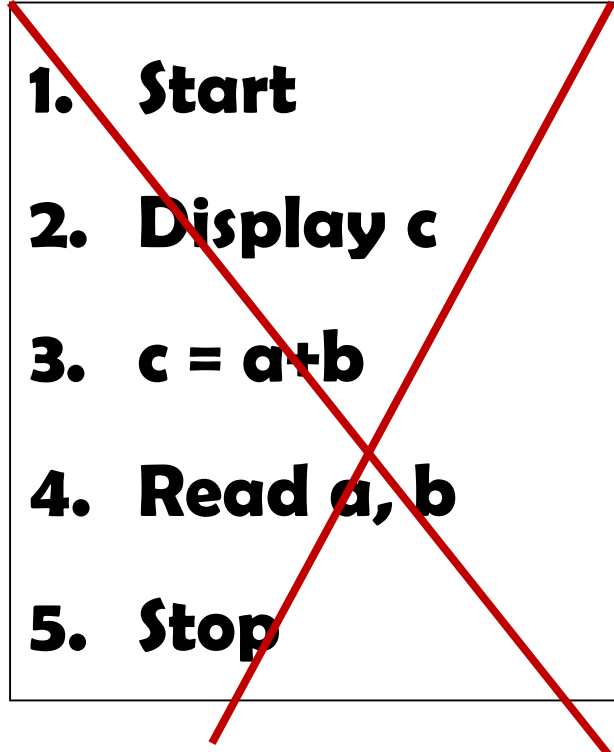
Analyzing programs.

Objective

Be familiar with basics of Algorithm and its criteria.

Algorithm:

This word comes from name of the Persian author, who wrote a book on Mathematics.

- 
- 1. Start**
 - 2. Display c**
 - 3. $c = a + b$**
 - 4. Read a, b**
 - 5. Stop**

- 1. Start**
- 2. Read a, b**
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Definition:

An Algorithm is a set of instructions if accomplishes, perform a specified task.

All algorithms should satisfy following **criteria:**

1. Input:

Zero or more

2. Output:

At least one

3. Definiteness:

Instruction should be clear & unambiguous.

Ex: Add 6 or 7 to x

Compute 5/0

4. Finiteness:

Terminate after a finite number of steps.

Ex: Chess

5. Effectiveness:

Instruction must be very basic enough so that it can be solved by a person using only pencil and paper.

Data: integer

Program Analysis:

1. Space Complexity

2. Time Complexity

1. Space complexity: Amount of memory algorithm needs to run for completion.

2. Time complexity: Amount of computer time algorithm needs to run for completion.

Space Complexity:

- **Space = Fixed part + Variable part**
of an algorithm.
- 1. **Fixed part:** Independent of characteristics of the I/Ps and O/Ps.
It includes space for
 - a. **Code.**
 - b. **Simple variables and**
fixed size component variables.
 - c. **Constants.**

2. Variable part consists of space needed by component variables whose size is dependent on:

- a. Particular problem instance** being solved.
- b. Space needed by referenced variables.**
- c. Recursion stack space.**

Space requirement for algorithm P is:

$$S(P) = C + S_p$$

where, C: Constant

S_p : Instance characteristics

Algorithm abc (a, b, c)

```
{  
    return a + b + b * c + (a + b - c) / ( a + b) + 4.0;  
}
```

- Characterized by 1 constant, **4.0** & values of **a, b & c**
- Assume each value require **1 word** space.
- No **instance characteristics**.

$$S(abc) = C + S_{abc}$$

Therefore $S_{abc} = 0$

$$C = 4$$

$$S(abc) \geq 4 + 0$$

Sum of n numbers. (Iterative Method)

Algorithm Sum (a, n)

{

 s := 0.0;

 for i:= 1 to n step 1 do

 s:= s + a[i];

 return s

}

- **Characterized by n and a.**
 n: number of elements
 a: array of floating point numbers.
- **Variables:** n, i, s (3 words)
- **n words** for a[]
 S(Sum) >= n + 3

Sum of n numbers. (Recursive)

Algorithm RSum (a, n)

{

if (n<=0) then

return 0.0;

else

return a[n] + RSum(a, n-1);

}

n = 3 a = {3, 5, 4}

Rsum(a,3)

```
{  
  if(3<=0)  
    False  
  else  
    return a[3]+Rsum(a,2)  
}
```

Diagram illustrating the recursive call **Rsum(a,3)**. The condition **if(3<=0)** is **False**. The **else** branch is executed, returning **a[3] + Rsum(a,2)**. The values **4** and **8** are shown in red, with a blue bracket indicating their sum is **12**.

Rsum(a,1)

```
{  
  if(1<=0)  
    False  
  else  
    return a[1]+Rsum(a,0)  
}
```

Diagram illustrating the recursive call **Rsum(a,1)**. The condition **if(1<=0)** is **False**. The **else** branch is executed, returning **a[1] + Rsum(a,0)**. The values **3** and **0** are shown in red, with a blue bracket indicating their sum is **3**.

Rsum(a,2)

```
{  
  if(2<=0)  
    False  
  else  
    return a[2]+Rsum(a,1)  
}
```

Diagram illustrating the recursive call **Rsum(a,2)**. The condition **if(2<=0)** is **False**. The **else** branch is executed, returning **a[2] + Rsum(a,1)**. The values **5** and **3** are shown in red, with a blue bracket indicating their sum is **8**.

Rsum(a,0)

```
{  
  if(0<=0)  
    return 0  
}
```

- The recursion stack space includes:
 Space for
 1. formal parameters.
 2. local variables.
 3. return address.
- Assume **return addr.** requires **1 word** of **memory**.
- RSum requires 3 words:
 1. Space for the value of n
 2. Return address
 3. A pointer to a[]
- The **depth** of the **recursion** is **n+1**.

Therefore, $S(\text{RSum}) \geq 3(n+1)$