

Data Structures

SYBTech(CSE)

Unit – 1

Introduction

PPT-5

Prof. Ms. Manisha A. Bhusa
Dept. Of CSE
COE Ambajogai

Unit 1 [6 hrs]

Introduction:

Data,

Data types,

Data structure,

Abstract Data Type (ADT),

Representation of Information,

Characteristics of algorithm,

Program,

Analyzing programs.

Objective

**Be familiar with analysis of Algorithm
(Time Complexity).**

Time Complexity:

- **$T(P) = \text{Compile Time}(P) + \text{Execution Time}(P)$**

Where P: Program

- **Run Time is denoted by t_p : Instance characteristics**
- **$T_p(n)$ can be obtained experimentally:**
 - 1. Type a program**
 - 2. Compile**
 - 3. Execute**

- **Problem in experimental approach:**

In a multiuser system, the execution time depends on:

- 1. System load, no. of other programs running on computer at the time program P is running.**
 - 2. The characteristics of these other programs.**
- so on.**

Obtain total number of program steps:

Two ways.

1. Introduce a new variable(global), **count**, into the program.

2. Tabular method.

```
Count:=0;  
Algorithm Swap(a, b)  
{  
    a := a+b;  
    count ++;  
    b := a-b;  
    count ++;  
    a := a-b;  
    count ++;  
    Display a, b;  
    count ++;  
}  
Total steps := 4
```

1. Introduce a new variable(global), count:

Algorithm Sum (a, n)

{

s := 0.0;

count:=count + 1;

for i:= 1 to n step 1 do

{

count:=count + 1;

s:= s + a[i];

count:=count + 1;

}

count:=count + 1;

count:=count + 1;

return s;

}

Algorithm Sum (a, n)

{

s := 0.0;

for i:= 1 to n step 1 do

s := s + a[i];

return s

}

```

Algorithm Sum (a, n)
{
    s := 0.0;
    count:=count + 1;
    for i:= 1 to n step 1 do
    {
        count:=count + 1;
        s:= s + a[i];
        count:=count + 1;
    }
    count:=count + 1;
    count:=count + 1;
    return s;
}

```

Simplified Algorithm:

```

Algo Sum(a, n)
{
    for i:= 1 to n step 1 do
        count := count +2;
        count := count +3;
    }
    count= 2n + 3

```

Therefore,

Total steps: $2n + 3$

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If $n=10$

Then Total steps = 23

If 1 step require 1 Nanosecond for execution

then total time required = 23 Nanoseconds

Algorithm RSum (a, n)

```
{  
    count:=count +1  
    if (n<=0) then  
    {  
        count:=count +1  
        return 0.0;  
    }  
    else  
    {  
        count:=count +1  
        return a[n] + RSum(a, n-1);  
    }  
}
```

Recurrence Relation:

$$TRSum(n) = \begin{cases} 2 & ; \text{ if } n=0 \\ 2 + tRSum(n-1) & ; \text{ if } n>0 \end{cases}$$

Recurrence Relation:

$$\text{TRSum}(n) \begin{cases} = 2 & \text{if } n=0 \\ = 2 + \text{tRSum}(n-1) & \text{if } n>0 \end{cases}$$

Solve recurrence relation by repeated substitution:

$$\begin{aligned} \text{tRSum}(n) &= 2 + \text{tRSum}(n-1) \\ &= 2 + 2 + \text{tRSum}(n-2) \\ &= 2(2) + \text{tRSum}(n-2) \\ &= 3(2) + \text{tRSum}(n-3) \\ &\quad \dots \\ &= n(2) + \text{tRSum}(n-n) \\ &= 2n + 2 \end{aligned}$$

$$\text{Step Count} = 2n + 2$$

This is called Linear Time Algorithm

as run time grows linearly in n.