I/O-efficient Manipulation of Binary Decision Diagrams

Steffan Christ Sølvsten

S. C. Sølvsten, J. van de Pol, A. B. Jakobsen, and M. W. B. Thomasen. *Adiar: Binary Decision Diagrams in External Memory.* 2022



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What are Binary Decision Diagrams?

Why do they break?

How can we fix it?

CountPaths

Apply

Equality Checking

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What are Binary Decision Diagrams?

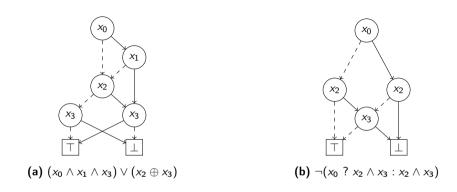
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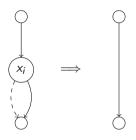
Apply

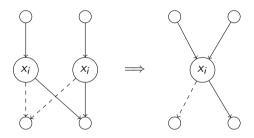
Equality Checking



Examples of (Reduced Ordered) Binary Decision Diagrams.

Theorem (Bryant '86)For a fixed variable order, if one exhaustively applies the two rules below, then one obtains the Reduced OBDD, which is a unique canonical form of the function.





(1) Remove redundant nodes

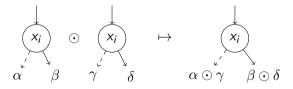
(2) Merge duplicate nodes

 $bdd_apply(f,g,\odot)$

Base Case $(f, g \in \mathbb{B})$:



Inductive Case:

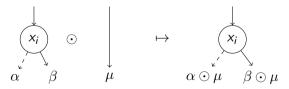


 $bdd_apply(f,g,\odot)$

Base Case $(f, g \in \mathbb{B})$:



Inductive Case:



$$bdd_apply(f,g,\odot)$$

Let N_f , N_g be the size of the BDDs for f and g.

Let T be the $O(N_f \cdot N_g)$ size of the BDD for $f \odot g$.

Theorem

 $bdd_apply(f,g,\odot)$ runs in $O(N_f + N_g + T)$ time

- Memoisation (*Computation Cache*) ensures each (t_f, t_g) is only computed once.
- Reduction Rules can be maintained with a make_node(i, t, e) in O(1) time.
 - 1 Redundancy is resolved with an if-statement.
 - 2 Duplication is avoided with a hash table (Unique Node Table).

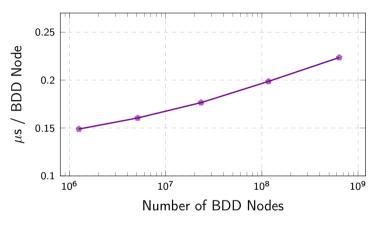
Corollary

 $bdd_apply(f,g,\odot)$ runs in O(1) time per BDD node.

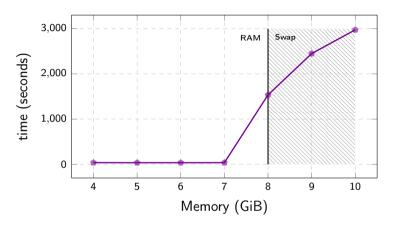
Adiar

I/O-efficient Decision Diagrams

github.com/ssoelvsten/adiar



Running time of BuDDy for the N-Queens problem.



Running time of BuDDy for 3D Tic-Tac-Toe with N=21.

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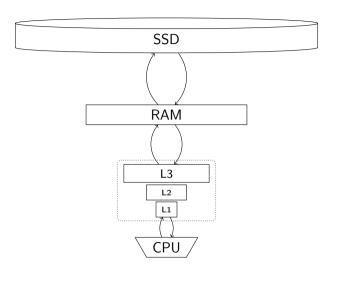
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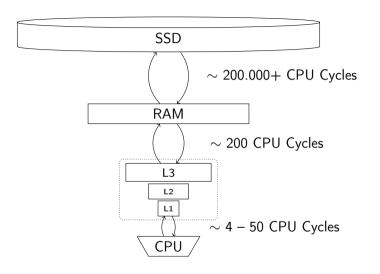
How can we fix it

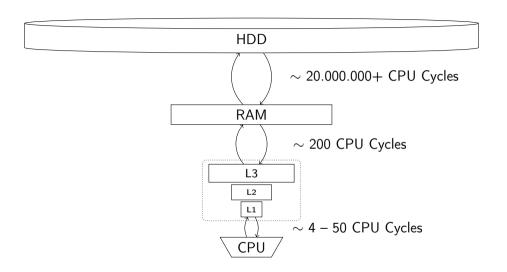
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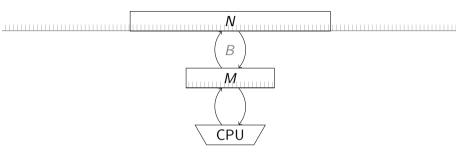
Apply

Equality Checking









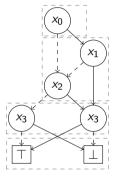
The I/O model by Aggarwal and Vitter '87

For any realistic values of N, M, and B we have that

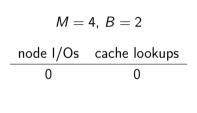
$$N/B < \operatorname{sort}(N) \triangleq N/B \cdot \log_{M/B} N/B \ll N$$
,

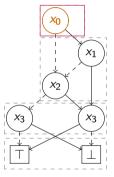
Theorem (Aggarwal and Vitter '87) N elements can be sorted in $\Theta(sort(N))$ I/Os.

Theorem (Arge '95) A Priority Queue can do N insertions and extractions in $\Theta(sort(N))$ I/Os.

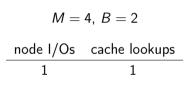


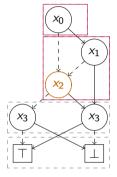
(a)
$$(x_0 \land x_1 \land x_3) \lor (x_2 \oplus x_3)$$





(a)
$$(x_0 \wedge x_1 \wedge x_3) \vee (x_2 \oplus x_3)$$



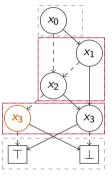


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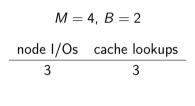
$$M = 4$$
, $B = 2$

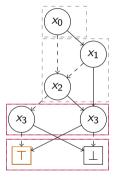
node I/Os cache lookups

2 2

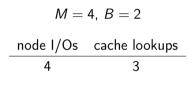


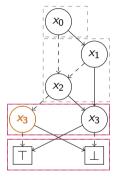
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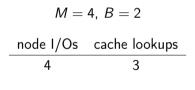


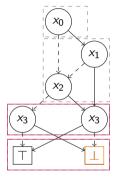
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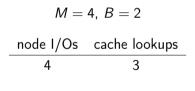


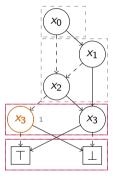
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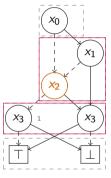
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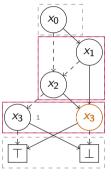
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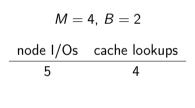


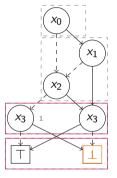
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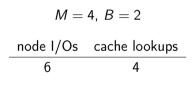


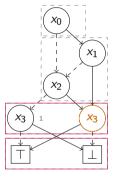
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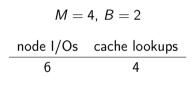


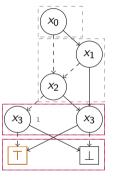
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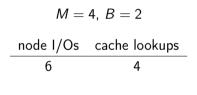


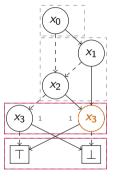
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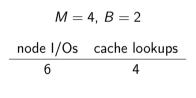


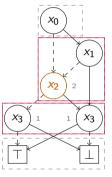
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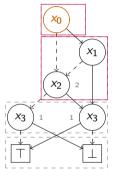
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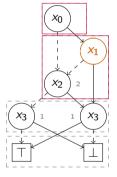
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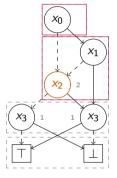
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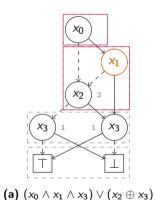


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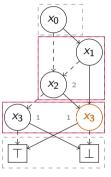
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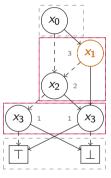
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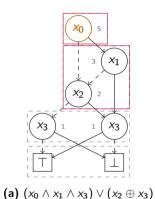
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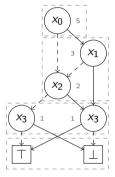
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Algorithm	Time Complexity
bdd_pathcount	$O(N_f)$
bdd_not	$O(N_f)$
bdd_restrict	$O(N_f)$
bdd_apply	$O(N_f \cdot N_g)$
bdd_equal	O(1)

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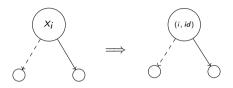
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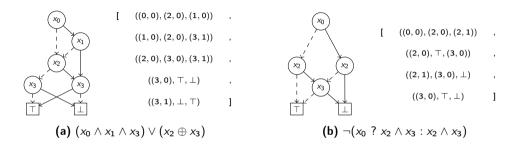
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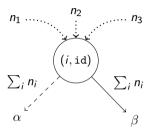


$$(i_1, id_1) < (i_2, id_2) \equiv i_1 < i_2 \lor (i_1 = i_2 \land id_i < id_j)$$



Node-based representation of prior shown BDDs

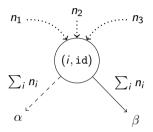
CountPaths



Idea

Count the number of in-going paths to each node.

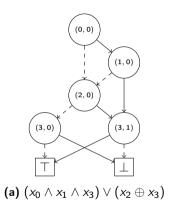
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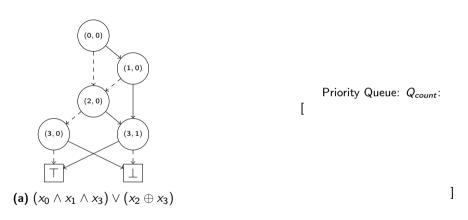


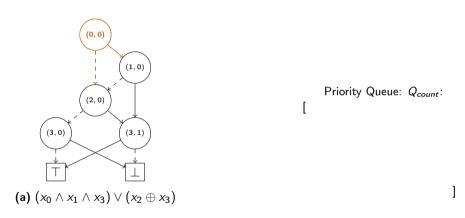
Time-Forward Processing

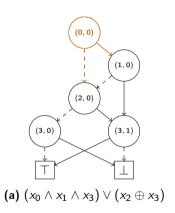
Defer work with Q_{count} : PriorityQueue $\langle (s \to t, \mathbb{N}) \rangle$ sorted on t in ascending order.

$$((i, \mathtt{id}) \xrightarrow{\perp} \alpha, \quad \sum_{i} n_{i}), \qquad ((i, \mathtt{id}) \xrightarrow{\top} \beta, \quad \sum_{i} n_{i})$$

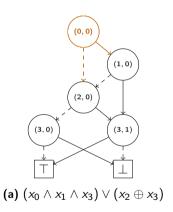


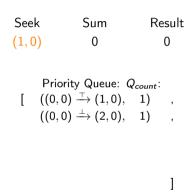


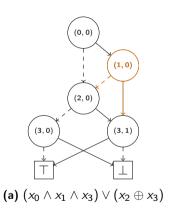


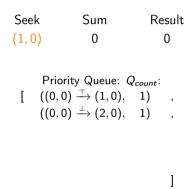


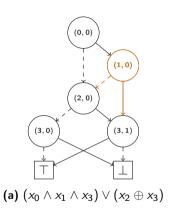
Priority Queue:
$$Q_{count}$$
: [$((0,0) \xrightarrow{\top} (1,0), 1)$, $((0,0) \xrightarrow{\bot} (2,0), 1)$,

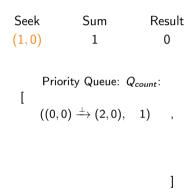


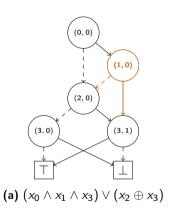


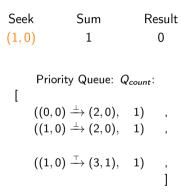


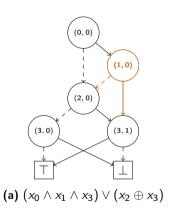


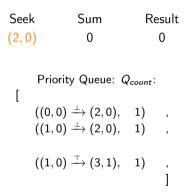


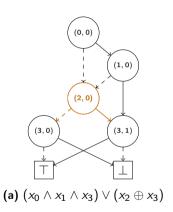


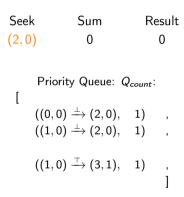


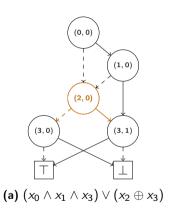


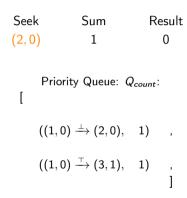


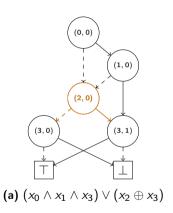


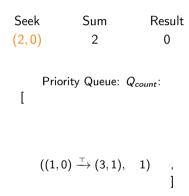


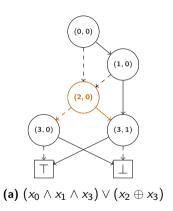




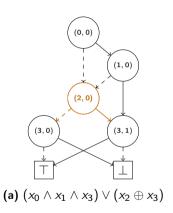




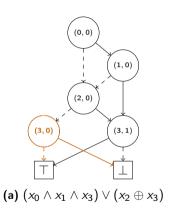




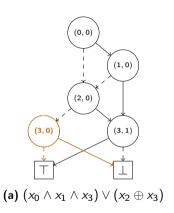
```
Seek
        Sum
                Result
(2,0)
         2
                 0
   Priority Queue: Qcount:
```

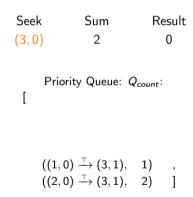


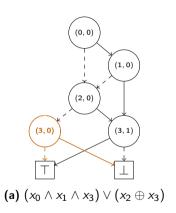
```
Seek
        Sum
                Result
(3,0)
         0
                 0
   Priority Queue: Qcount:
```

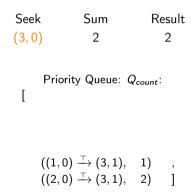


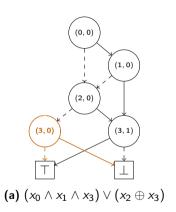
```
Seek
        Sum
                Result
(3,0)
         0
                 0
   Priority Queue: Qcount:
```

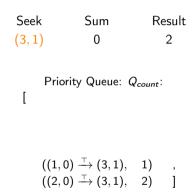


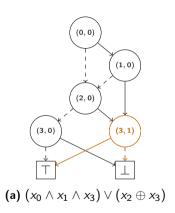


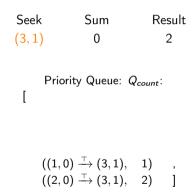


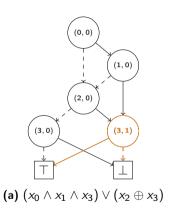


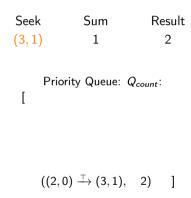


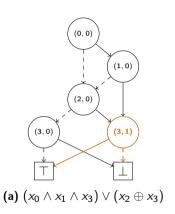


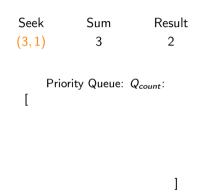


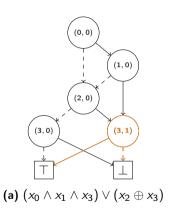


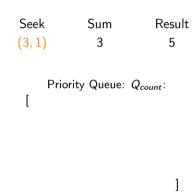




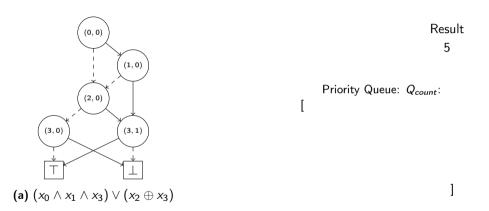








CountPaths : Example



Contents

What are Binary Decision Diagrams?

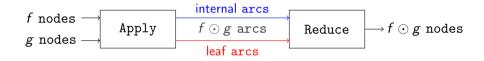
Why do they break?

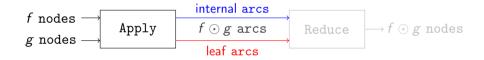
How can we fix it?

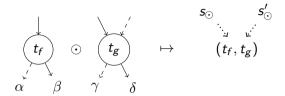
CountPaths

Apply

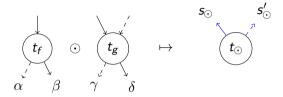
Equality Checking







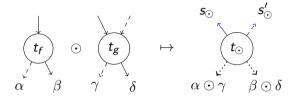
Time-Forward Processing



Observation (semi-tranposition)

 \leftarrow : $s \rightarrow t$ (Internal Arcs) are output at time t and hence sorted by t.

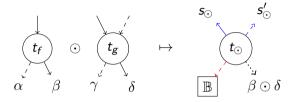
Time-Forward Processing



Observation (semi-tranposition)

 \leftarrow : $s \rightarrow t$ (Internal Arcs) are output at time t and hence sorted by t.

Time-Forward Processing



Observation (semi-tranposition)

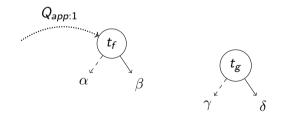
- \leftarrow : $s \rightarrow t$ (Internal Arcs) are output at time t and hence sorted by t.
- \rightarrow : $s \rightarrow \mathbb{B}$ (Terminal Arcs) are output at time s.

Time-Forward Processing

 $Q_{app:1}$: PriorityQueue $\langle (s o (t_f, t_g))
angle$ sorted on $\min(t_f, t_g)$ in ascending order.

Case 1

 $t_f.var() \neq t_g.var()$



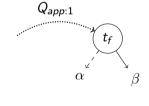
 $Q_{app:1}$: PriorityQueue $\langle (s \rightarrow (t_f, t_g)) \rangle$ sorted on $\min(t_f, t_g)$ in ascending order.

Case 1

 $t_f.var()
eq t_g.var()$

Case 2(a):

 $t_f.var() = t_g.var() \land t_f.id() = t_g.id()$

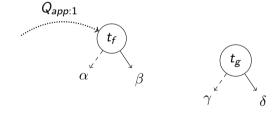




 $Q_{app:1}$: PriorityQueue $\langle (s \rightarrow (t_f, t_g)) \rangle$ sorted on min (t_f, t_g) in ascending order.

Case 1

$$t_f.var() \neq t_g.var()$$



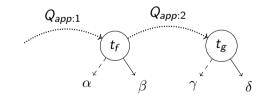
 $Q_{app:2}$: PriorityQueue $\langle (s \to (t_f, t_g), (\alpha, \beta)) \rangle$ sorted on max (t_f, t_g) in ascending order.

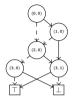
Case 2(a):

$$t_f.var() = t_g.var() \land t_f.id() = t_g.id()$$

Case 2(b):

$$t_f.var() = t_g.var() \land t_f.id() \neq t_g.id()$$



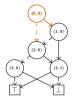


(a) $(x_0 \wedge x_1 \wedge x_3) \vee (x_2 \oplus x_3)$



(b)
$$\neg (x_0 ? x_2 \lor x_3 : x_2 \land x_3)$$

(c) $(a) \wedge (b)$

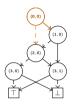


(a)
$$(x_0 \wedge x_1 \wedge x_3) \vee (x_2 \oplus x_3)$$



(b)
$$\neg (x_0 ? x_2 \lor x_3 : x_2 \land x_3)$$

(c) (a) ∧ (b)



Priority Queue: Qapp:1:

 $[(0,0) \xrightarrow{\top} ((1,0),(2,1)) ,$ $(0,0) \xrightarrow{\bot} ((2,0),(2,0)) ,$ (0,0)

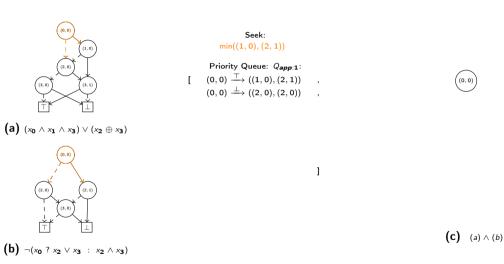
(a)
$$(x_0 \wedge x_1 \wedge x_3) \vee (x_2 \oplus x_3)$$

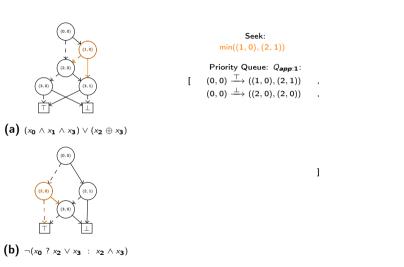


(b)
$$\neg (x_0 ? x_2 \lor x_3 : x_2 \land x_3)$$

]

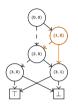
(c)
$$(a) \wedge (b)$$





(0,0)

(c) (a) ∧ (b)



(a)
$$(x_0 \wedge x_1 \wedge x_3) \vee (x_2 \oplus x_3)$$



(b) $\neg (x_0 ? x_2 \lor x_3 : x_2 \land x_3)$

$\begin{array}{c} \text{Seek:} \\ \min((1,0),(2,1)) \end{array}$

Priority Queue: Qapp:1:

 $(0,0) \xrightarrow{\top} ((1,0),(2,1))$

 $(0,0) \xrightarrow{\perp} ((2,0),(2,0)) ,$

 $(1,0) \xrightarrow{\perp} ((2,0),(2,1))$

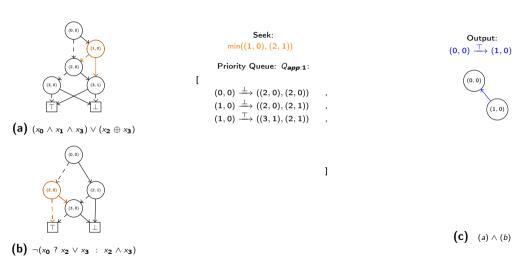
 $(1,0) \xrightarrow{\top} ((3,1),(2,1))$

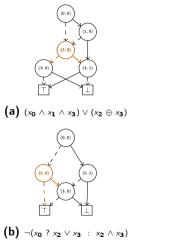
(0,0)

(1,0)

J

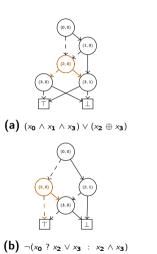
(c) $(a) \wedge (b)$





Seek: min((2,0),(2,0))Priority Queue: Qapp:1: $(0,0) \xrightarrow{\perp} ((2,0),(2,0))$ $(1,0) \xrightarrow{\perp} ((2,0),(2,1))$ $(1,0) \xrightarrow{\top} ((3,1),(2,1))$

Output: (c) $(a) \wedge (b)$

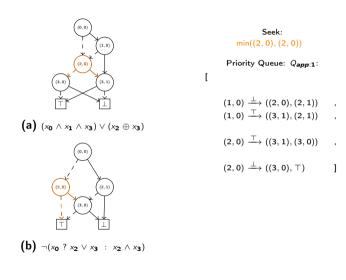


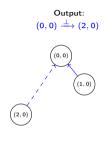
Seek: min((2,0),(2,0))Priority Queue: Qapp:1: $(0,0) \xrightarrow{\perp} ((2,0),(2,0))$ $(1,0) \xrightarrow{\perp} ((2,0),(2,1))$ $(1,0) \xrightarrow{\top} ((3,1),(2,1))$ $(2,0) \xrightarrow{\top} ((3,1),(3,0))$ $(2,0) \xrightarrow{\perp} ((3,0),\top)$]

Output:

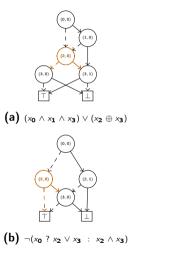
(c) $(a) \wedge (b)$

(2,0)





(c) (a) ∧ (b)

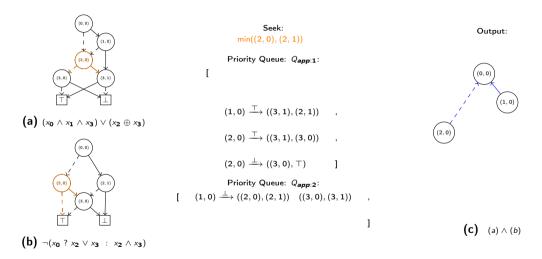


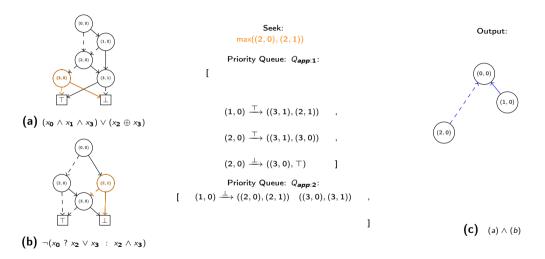
Seek: min((2,0),(2,1))Priority Queue: Qapp:1: $(1,0) \xrightarrow{\perp} ((2,0),(2,1))$ $(1,0) \xrightarrow{\top} ((3,1),(2,1))$ $(2,0) \xrightarrow{\top} ((3,1),(3,0))$ $(2,0) \xrightarrow{\perp} ((3,0),\top)$]

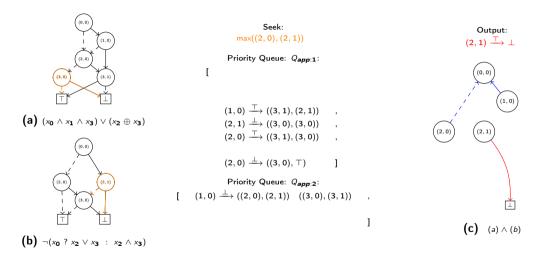
(2,0)

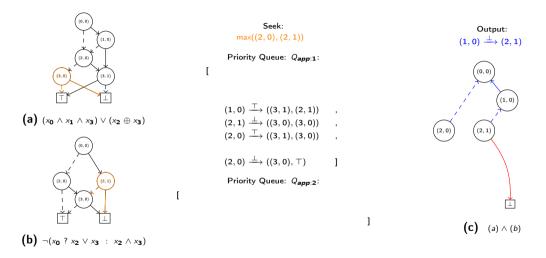
Output:

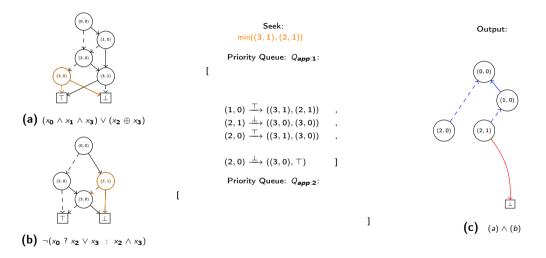
(c) $(a) \wedge (b)$

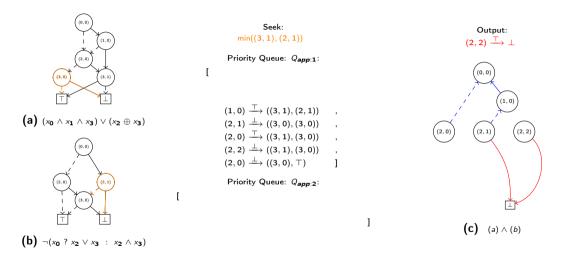


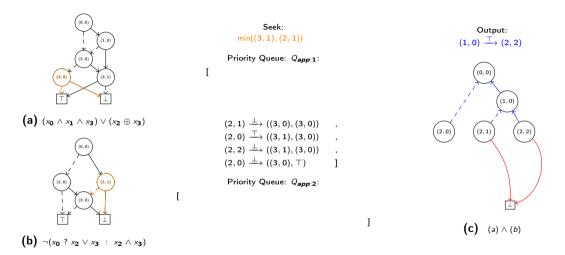


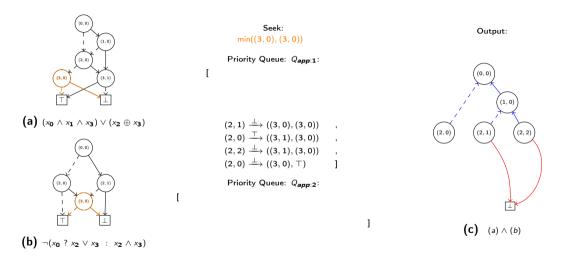


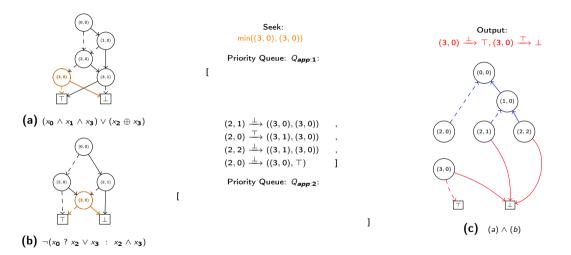


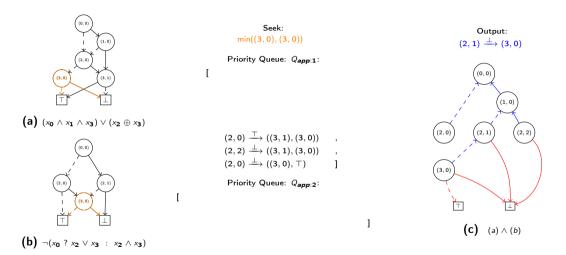


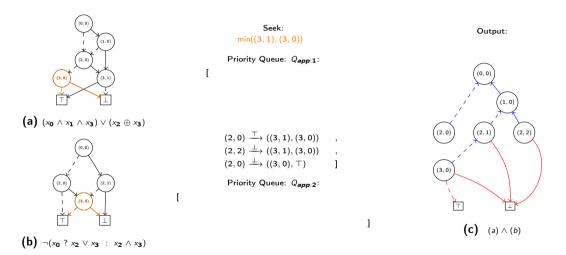


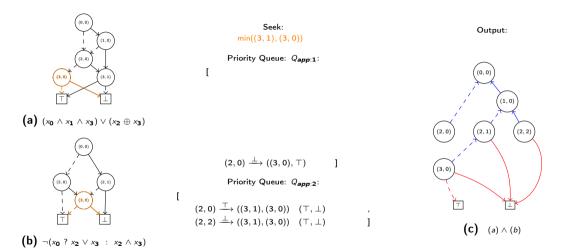


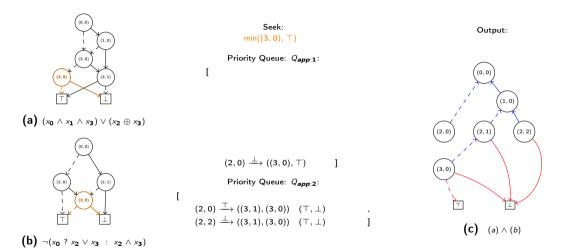


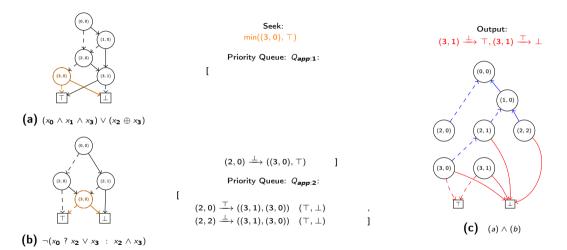


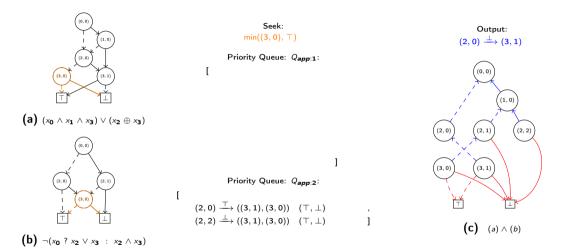


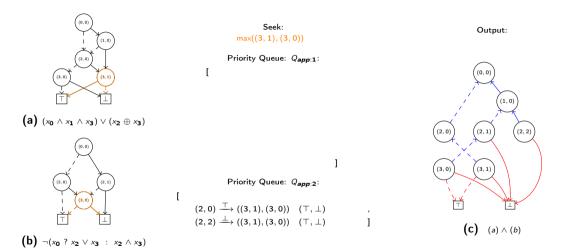


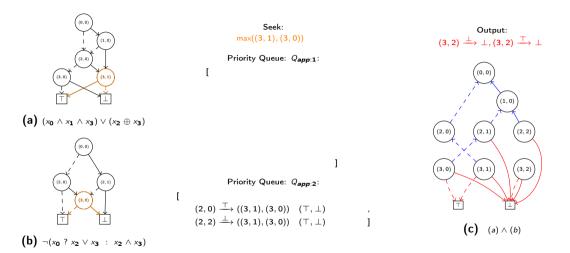


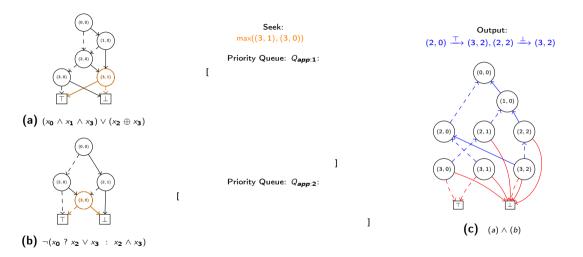


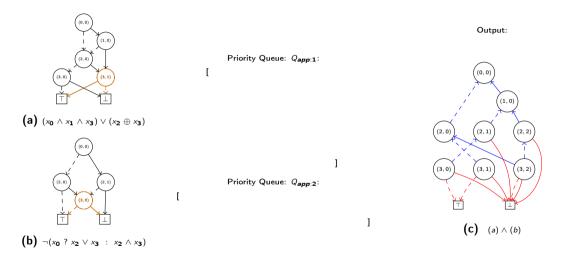




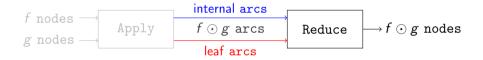


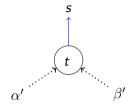






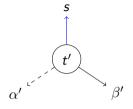
Apply





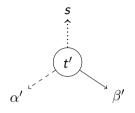
Time-Forward Processing

Send reduction t' with Q_{red} : PriorityQueue $\langle (s o t')
angle$ descending on parent s.



Time-Forward Processing

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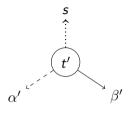


Time-Forward Processing

Send reduction t' with Q_{red} : PriorityQueue $\langle (s \to t') \rangle$ descending on parent s.

Observation (semi-tranposition)

 \leftarrow : $s \rightarrow t$ (Internal Arcs) provide parents of unreduced node t.

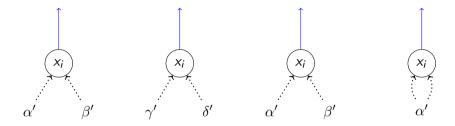


Time-Forward Processing

Send reduction t' with Q_{red} : PriorityQueue $\langle (s o t')
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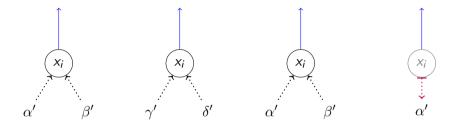
Observation (semi-tranposition)

- \leftarrow : $s \rightarrow t$ (Internal Arcs) provide parents of unreduced node t.
- ightarrow: $s
 ightarrow \mathbb{B}$ (Terminal Arcs) are reduced and already sorted as per Q_{red} .



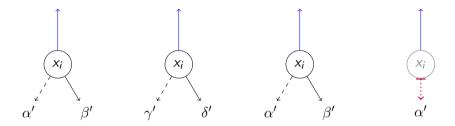
Reduce Level i:

1 Obtain nodes from Q_{red} and terminal arcs. Filter and remember redundant nodes.



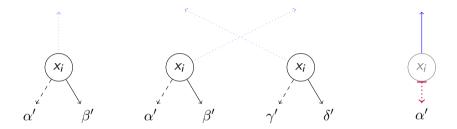
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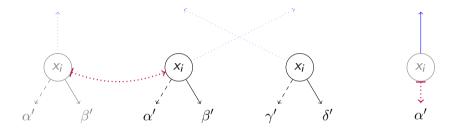


Reduce Level i:

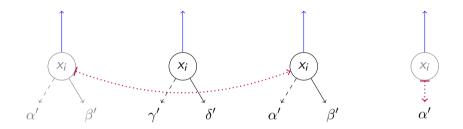
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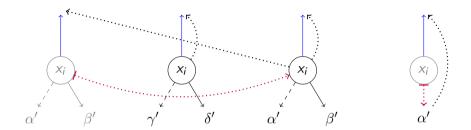
- 1 Obtain nodes from Q_{red} and terminal arcs. Filter and remember redundant nodes.
- Sort remaining nodes by children, output unique nodes, and remember duplications.



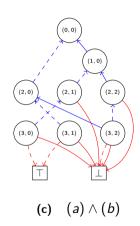
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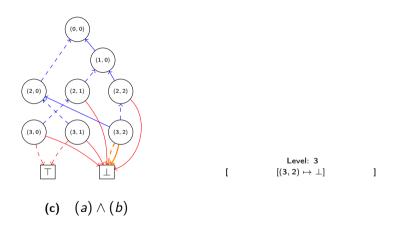


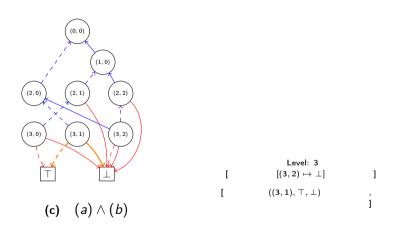
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- 3 Sort back to match internal arcs and forward to parents with Q_{red} .

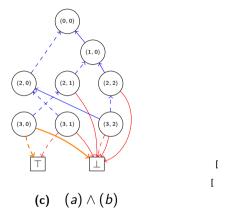


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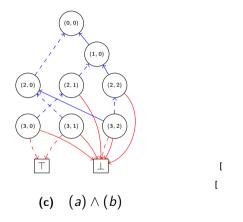


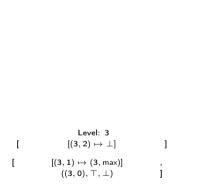




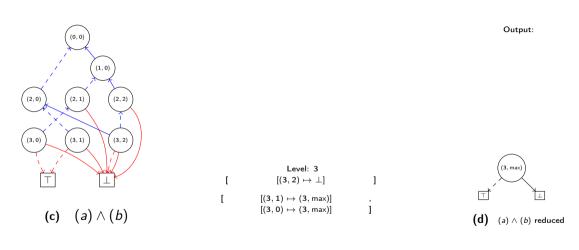
Level: 3
$$[(3,2)\mapsto \bot]$$

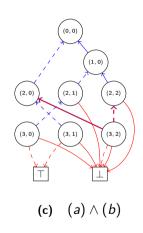
$$[\qquad \qquad ((3,1),\top,\bot) \qquad \qquad ((3,0),\top,\bot) \qquad \qquad]$$

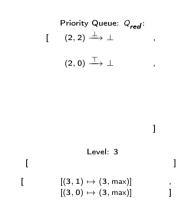


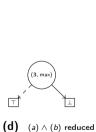


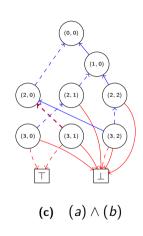


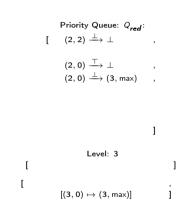


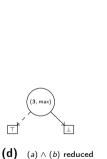


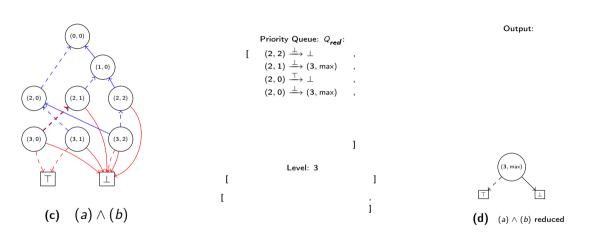


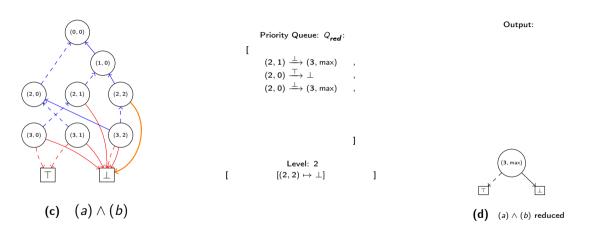


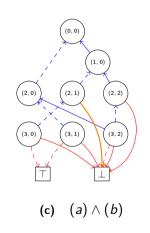


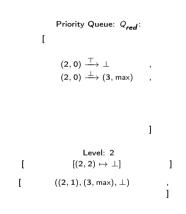




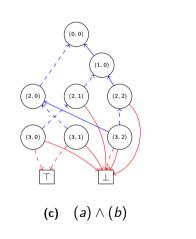


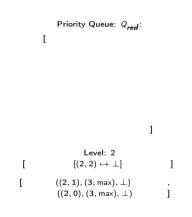




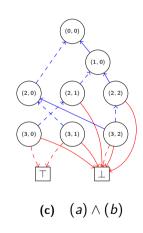


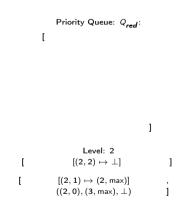


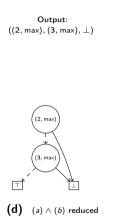


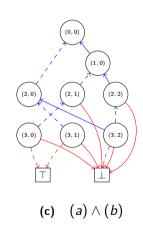


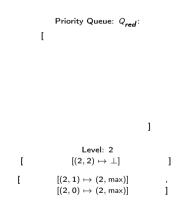


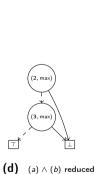


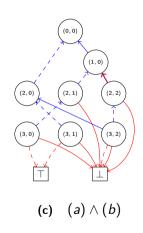


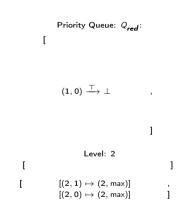


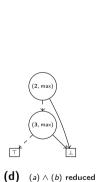


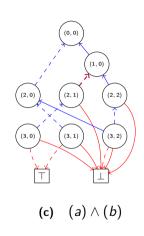


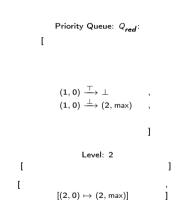


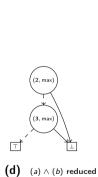


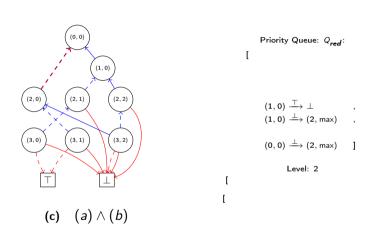


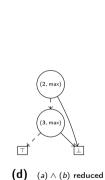


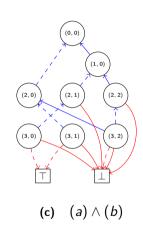


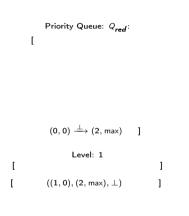


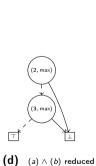


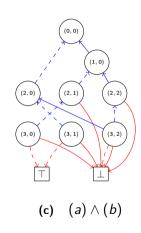


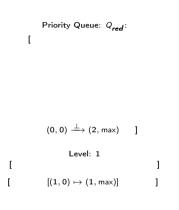


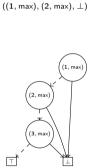




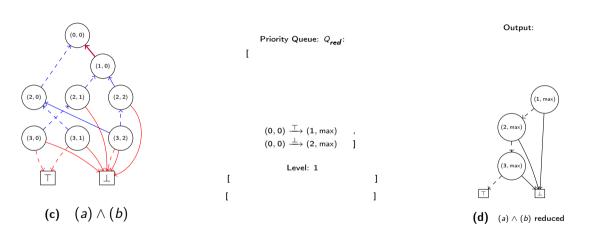


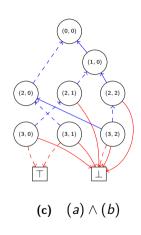


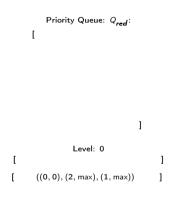


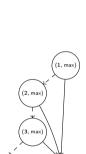


Output:

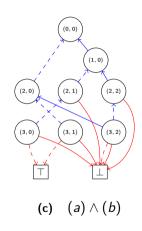


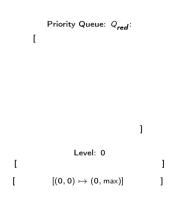


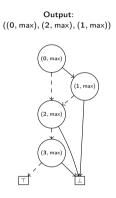


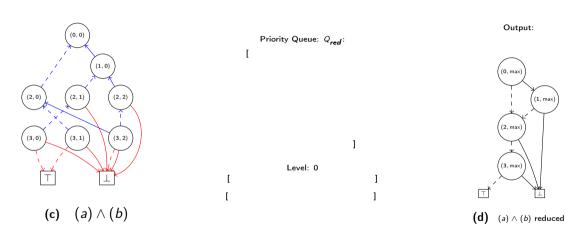


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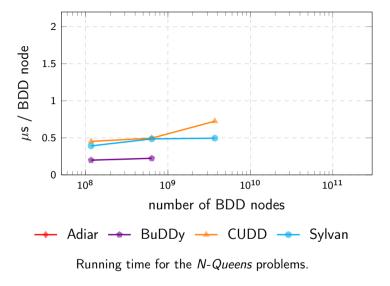


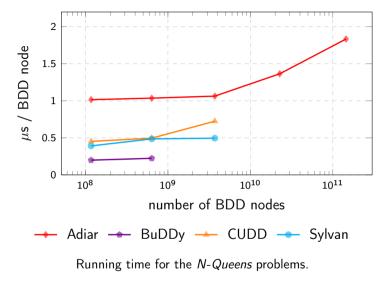






Algorithm	I/O-Complexity	
bdd_pathcount	$O(\operatorname{sort}(N_f))$	
bdd_not	$2N_f/B$	
bdd_restrict	$O(\operatorname{sort}(N_f))$	
bdd_apply	$O(\operatorname{sort}(N_f \cdot N_g))$	





Contents

What are Binary Decision Diagrams?

Why do they break?

How can we fix it?

CountPaths

Apply

Equality Checking

Algorithm	I/O-Complexity
bdd_pathcount	$O(\operatorname{sort}(N_f))$
bdd_not	$2N_f/B$
bdd_restrict	$O(\operatorname{sort}(N_f))$
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bdd_equal	?

$$f\leftrightarrow g\equiv \top$$

$$f \leftrightarrow g \equiv \top$$

$$\underbrace{\textit{O}(\mathsf{sort}(\textit{N}^2))}_{\texttt{Apply}} + \underbrace{\textit{O}(\mathsf{sort}(\textit{N}^2))}_{\texttt{Reduce}} + \underbrace{\textit{O}(1))}_{\texttt{check is }\top} = \textit{O}(\mathsf{sort}(\textit{N}^2))$$

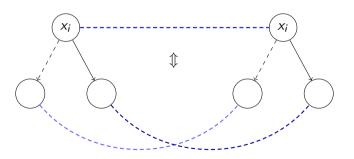
Theorem (Bryant '86)

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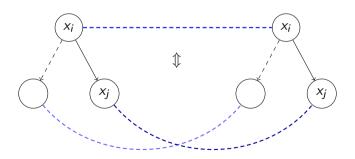
Let π be a variable order and $f: \mathbb{B}^n \to \mathbb{B}$ then there exists a unique (up to isomorphism) Reduced Ordered Binary Decision Diagram representing f with ordering π .

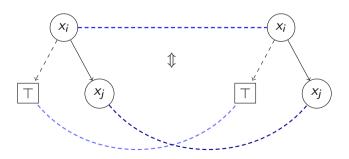
Trivial cases: $f \not\equiv g$ if there is a mismatch in

	$N_f eq N_g$	Number of nodes	O(1) I/Os
•	$L_f eq L_g$	Number of levels	<i>O</i> (1) I/Os
•	$N_{f,i} eq N_{g,i}$	Number of nodes on a level	O(L/B) I/Os
•	$L_{f,i} \neq L_{g,i}$	Label of an <i>i</i> th level	O(L/B) I/Os

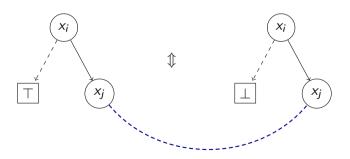


Theorem (Bryant '86)

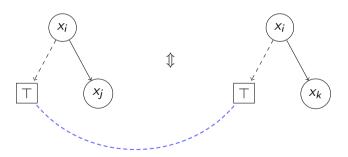




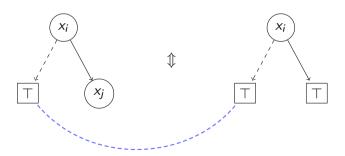
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Let π be a variable order and $f: \mathbb{B}^n \to \mathbb{B}$ then there exists a unique (up to isomorphism) Reduced Ordered Binary Decision Diagram representing f with ordering π .

IsIsomorphic(f, g)

- Check whether root v_f of f and root v_g of g have a local violation.
- Check $low(v_f) \sim low(v_g)$ and $high(v_f) \sim high(v_g)$ "recursively".

Return false on first violation. If there are no violations then return true.

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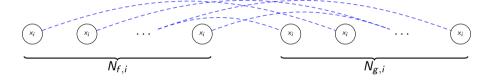
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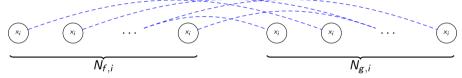
$$\underbrace{O(\mathsf{sort}(\mathit{N}^2))}_{\mathtt{Apply'}} + \underbrace{O(\mathsf{sort}(\mathit{N}^2))}_{\mathtt{Reduce}} + \underbrace{O(1))}_{\mathtt{check is} \ \top} = O(\mathsf{sort}(\mathit{N}^2))$$

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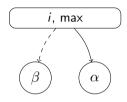
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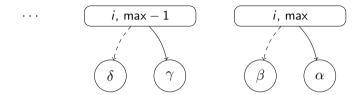
Return false if more than $N_{f,i} = N_{g,i}$ pairs of nodes are checked on level i.

$$\underbrace{O(\mathsf{sort}(\Sigma_i \ N_{f,i}))}_{\mathsf{Apply''}} = O(\mathsf{sort}(N))$$

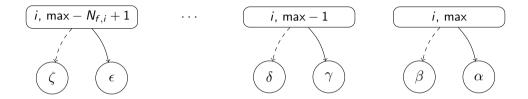
Observation



Observation



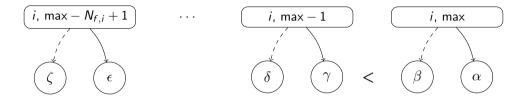
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Each level output by the Reduce algorithm has the following properties:

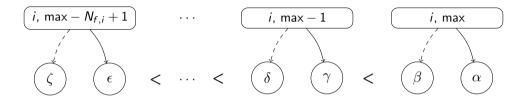
■ Nodes on level *i* have their identifiers *consecutively* numbered.



Observation

Each level output by the Reduce algorithm has the following properties:

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Observation

- Nodes on level *i* have their identifiers *consecutively* numbered.
- Nodes on level *i* are output sorted by their children.

Theorem

If G_f and G_g are outputs of Reduce.

 $G_f \sim G_g \iff For \ all \ i \in [0; N_f) \ the \ node \ G_f[i] \ matches \ G_g[i] \ numerically.$

Proof.

← : Must describe the exact same graph.

 \Rightarrow : Strong induction on BDD levels bottom-up.

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 \Rightarrow : Strong induction on BDD levels bottom-up.

Corollary

If G_f and G_g are outputs of Reduce then $f \equiv g$ is computable using $2 \cdot N/B$ I/Os.

$$\begin{array}{c|c} & \text{Algorithm} & \text{Time (s)} \\ \hline f \leftrightarrow g \equiv \top & 0.38 \end{array}$$

Checking the (EPFL Benchmark) voter circuit's single output gate ($|N_f| = |N_g| = 5.76$ MiB).

Algorithm Time (s)
$$f \leftrightarrow g \equiv \top \quad 0.38$$

$$O(\operatorname{sort}(N)) \quad 0.058$$

Checking the (EPFL Benchmark) voter circuit's single output gate ($|N_f| = |N_g| = 5.76$ MiB).

Algorithm	Time (s)
$f\leftrightarrow g\equiv \top$	0.38
O(sort(N))	0.058
2N/B	0.006

Checking the (EPFL Benchmark) voter circuit's single output gate ($|N_f| = |N_g| = 5.76$ MiB).

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ssoelvsten.github.io

Adiar

github.com/ssoelvsten/adiar

ssoelvsten.github.io/adiar



Algorithm	Depth-First	Time-Forwared
bdd_pathcount	$O(N_f)$	$O(\operatorname{sort}(N_f))$
bdd_not	$O(N_f)$	$2N_f/B$
bdd_restrict	$O(N_f)$	$O(\operatorname{sort}(N_f))$
bdd_apply	$O(N_f N_g)$	$O(\operatorname{sort}(N_f N_g))$
bdd_equal	O(1)	2 N /B