## I/O-efficient Manipulation of Binary Decision Diagrams

Steffan Christ Sølvsten

S. C. Sølvsten, J. van de Pol, A. B. Jakobsen, and M. W. B. Thomasen. *Adiar: Binary Decision Diagrams in External Memory.* 2022



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What are Binary Decision Diagrams?

Why do they break?

How can we fix it?

CountPaths

Apply

**Equality Checking** 

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What are Binary Decision Diagrams?

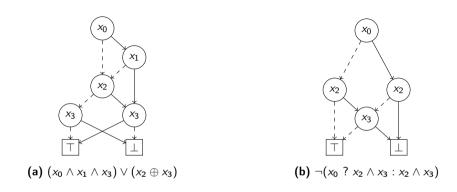
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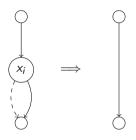
Apply

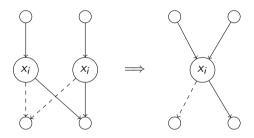
**Equality Checking** 



Examples of (Reduced Ordered) Binary Decision Diagrams.

**Theorem (Bryant '86)**For a fixed variable order, if one exhaustively applies the two rules below, then one obtains the Reduced OBDD, which is a unique canonical form of the function.





(1) Remove redundant nodes

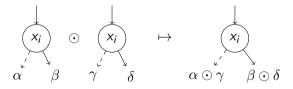
(2) Merge duplicate nodes

 $bdd_apply(f,g,\odot)$ 

### Base Case $(f, g \in \mathbb{B})$ :



#### Inductive Case:

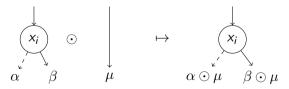


 $bdd_apply(f,g,\odot)$ 

#### Base Case $(f, g \in \mathbb{B})$ :



#### Inductive Case:



$$bdd_apply(f,g,\odot)$$

Let  $N_f$ ,  $N_g$  be the size of the BDDs for f and g.

Let T be the  $O(N_f \cdot N_g)$  size of the BDD for  $f \odot g$ .

#### Theorem

 $bdd_apply(f,g,\odot)$  runs in  $O(N_f + N_g + T)$  time

- Memoisation (*Computation Cache*) ensures each  $(t_f, t_g)$  is only computed once.
- Reduction Rules can be maintained with a make\_node(i, t, e) in O(1) time.
  - 1 Redundancy is resolved with an if-statement.
  - 2 Duplication is avoided with a hash table (Unique Node Table).

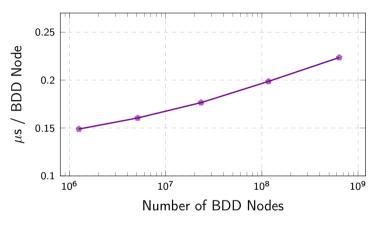
#### Corollary

 $bdd_apply(f,g,\odot)$  runs in O(1) time per BDD node.

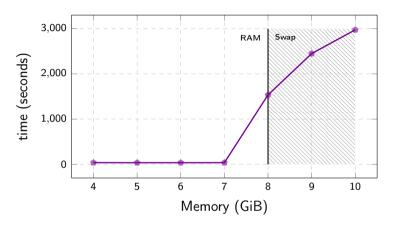
# Adiar

I/O-efficient Decision Diagrams

github.com/ssoelvsten/adiar



Running time of BuDDy for the N-Queens problem.



Running time of BuDDy for 3D Tic-Tac-Toe with N=21.

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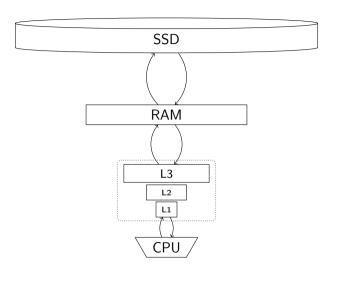
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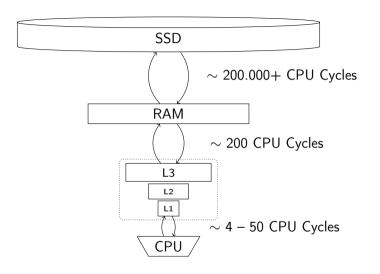
How can we fix it

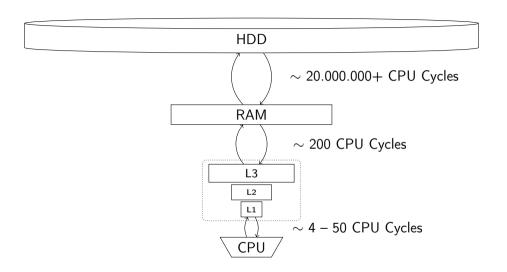
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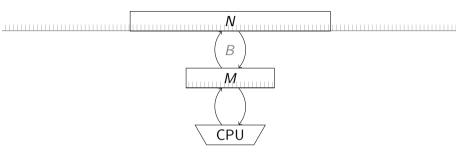
Apply

**Equality Checking** 









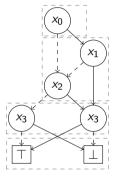
The I/O model by Aggarwal and Vitter '87

For any realistic values of N, M, and B we have that

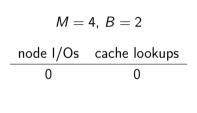
$$N/B < \operatorname{sort}(N) \triangleq N/B \cdot \log_{M/B} N/B \ll N$$
,

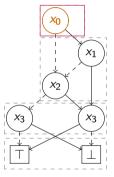
Theorem (Aggarwal and Vitter '87) N elements can be sorted in  $\Theta(sort(N))$  I/Os.

**Theorem (Arge '95)** A Priority Queue can do N insertions and extractions in  $\Theta(sort(N))$  I/Os.

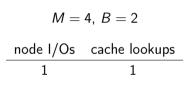


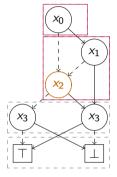
(a) 
$$(x_0 \land x_1 \land x_3) \lor (x_2 \oplus x_3)$$





(a) 
$$(x_0 \wedge x_1 \wedge x_3) \vee (x_2 \oplus x_3)$$



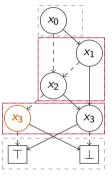


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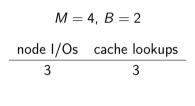
$$M = 4$$
,  $B = 2$ 

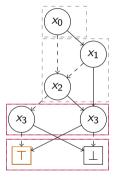
node I/Os cache lookups

2 2

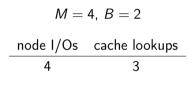


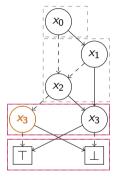
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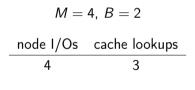


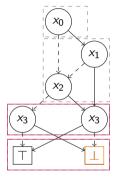
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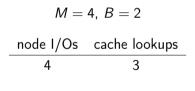


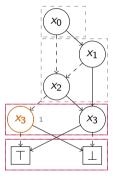
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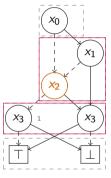
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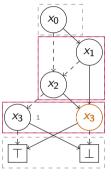
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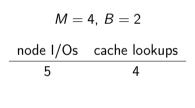


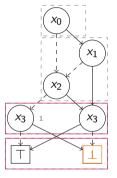
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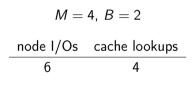


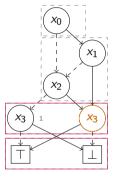
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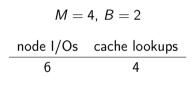


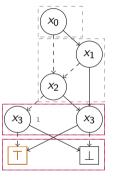
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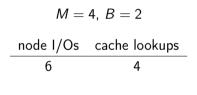


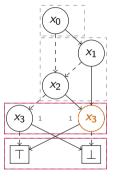
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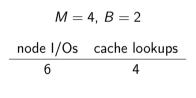


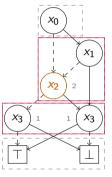
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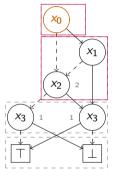
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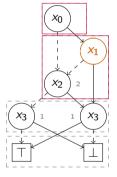
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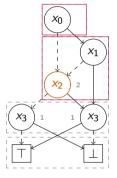
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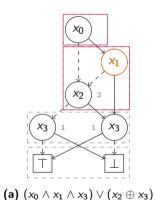


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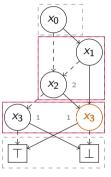
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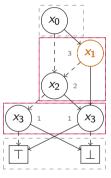
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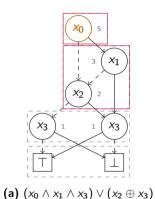
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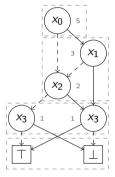
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10 7



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node I/Os cache lookups

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Algorithm	Time Complexity
bdd_pathcount	$O(N_f)$
bdd_not	$O(N_f)$
bdd_restrict	$O(N_f)$
bdd_apply	$O(N_f \cdot N_g)$
bdd_equal	O(1)

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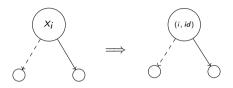
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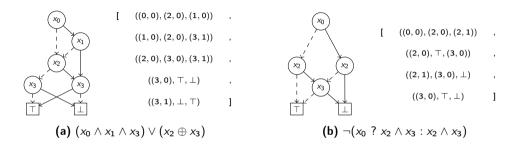
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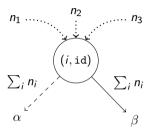


$$(i_1, id_1) < (i_2, id_2) \equiv i_1 < i_2 \lor (i_1 = i_2 \land id_i < id_j)$$



Node-based representation of prior shown BDDs

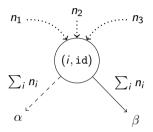
#### **CountPaths**



# Idea

Count the number of in-going paths to each node.

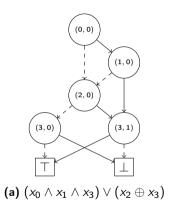
#### **CountPaths**

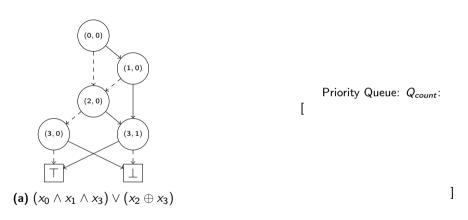


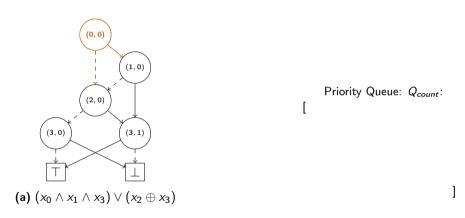
## **Time-Forward Processing**

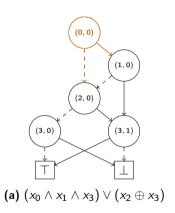
Defer work with  $Q_{count}$ : PriorityQueue $\langle (s \to t, \mathbb{N}) \rangle$  sorted on t in ascending order.

$$((i, \mathtt{id}) \xrightarrow{\perp} \alpha, \quad \sum_{i} n_{i}), \qquad ((i, \mathtt{id}) \xrightarrow{\top} \beta, \quad \sum_{i} n_{i})$$

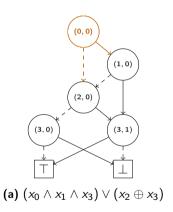


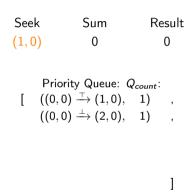


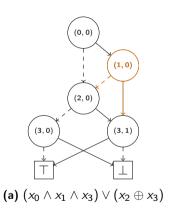


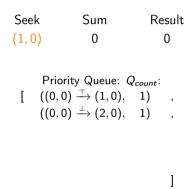


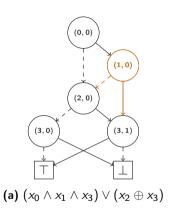
Priority Queue: 
$$Q_{count}$$
: [  $((0,0) \xrightarrow{\top} (1,0), 1)$  ,  $((0,0) \xrightarrow{\bot} (2,0), 1)$  ,

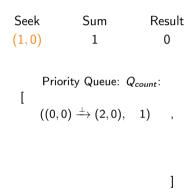


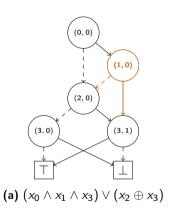


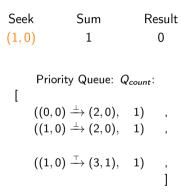


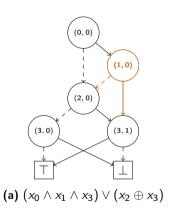


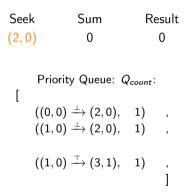


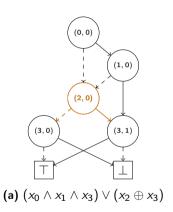


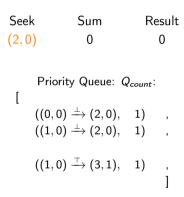


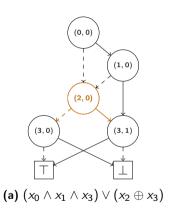


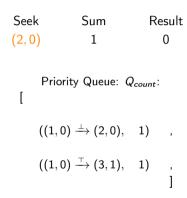


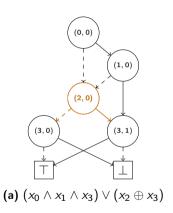


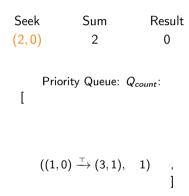


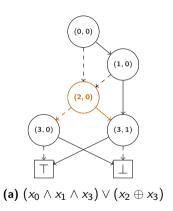




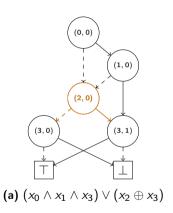




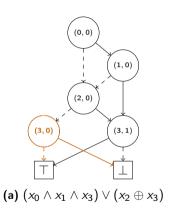




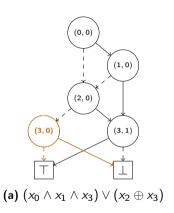
```
Seek
        Sum
                Result
(2,0)
         2
                 0
   Priority Queue: Qcount:
```

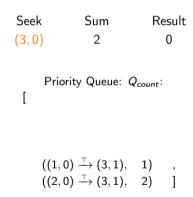


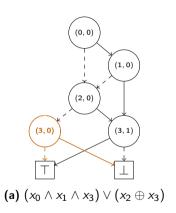
```
Seek
        Sum
                Result
(3,0)
         0
                 0
   Priority Queue: Qcount:
```

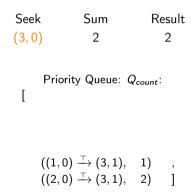


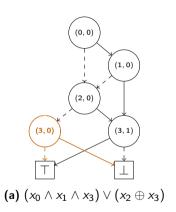
```
Seek
        Sum
                Result
(3,0)
         0
                 0
   Priority Queue: Qcount:
```

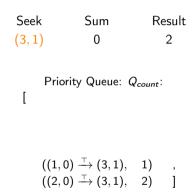


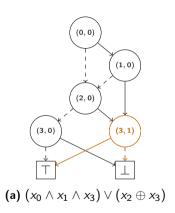


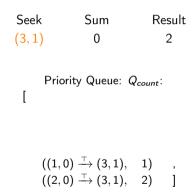


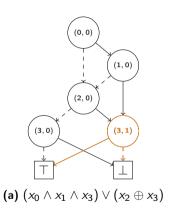


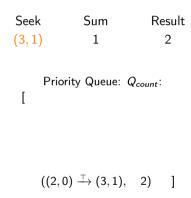


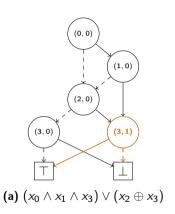


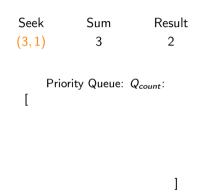


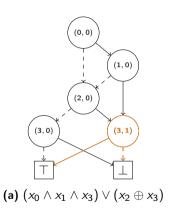


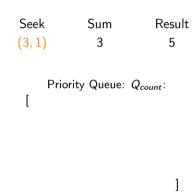




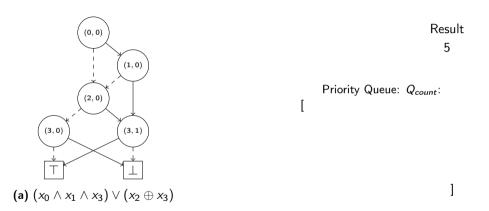








# CountPaths : Example



#### **Contents**

What are Binary Decision Diagrams?

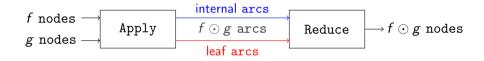
Why do they break?

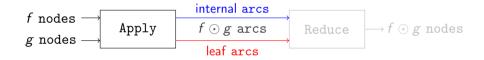
How can we fix it?

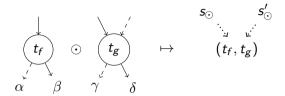
CountPaths

Apply

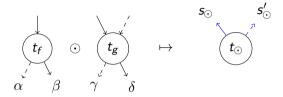
**Equality Checking** 







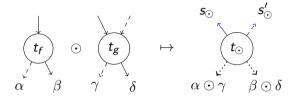
# **Time-Forward Processing**



#### Observation (semi-tranposition)

 $\leftarrow$  :  $s \rightarrow t$  (Internal Arcs) are output at time t and hence sorted by t.

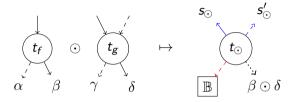
# **Time-Forward Processing**



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# **Time-Forward Processing**



#### Observation (semi-tranposition)

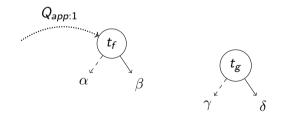
- $\leftarrow$  :  $s \rightarrow t$  (Internal Arcs) are output at time t and hence sorted by t.
- $\rightarrow$  :  $s \rightarrow \mathbb{B}$  (Terminal Arcs) are output at time s.

# **Time-Forward Processing**

 $Q_{app:1}$ : PriorityQueue $\langle (s o (t_f, t_g)) 
angle$  sorted on  $\min(t_f, t_g)$  in ascending order.

#### Case 1

 $t_f.var() \neq t_g.var()$ 



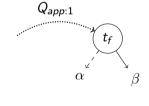
 $Q_{app:1}$ : PriorityQueue $\langle (s \rightarrow (t_f, t_g)) \rangle$  sorted on  $\min(t_f, t_g)$  in ascending order.

#### Case 1

 $t_f.var() 
eq t_g.var()$ 

# Case 2(a):

 $t_f.var() = t_g.var() \land t_f.id() = t_g.id()$ 

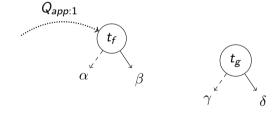




 $Q_{app:1}$ : PriorityQueue $\langle (s \rightarrow (t_f, t_g)) \rangle$  sorted on min $(t_f, t_g)$  in ascending order.

#### Case 1

$$t_f.var() \neq t_g.var()$$



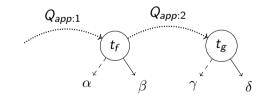
 $Q_{app:2}$ : PriorityQueue $\langle (s \to (t_f, t_g), (\alpha, \beta)) \rangle$  sorted on max $(t_f, t_g)$  in ascending order.

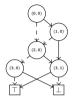
#### Case 2(a):

$$t_f.var() = t_g.var() \land t_f.id() = t_g.id()$$

#### Case 2(b):

$$t_f.var() = t_g.var() \land t_f.id() \neq t_g.id()$$



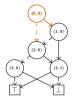


(a)  $(x_0 \wedge x_1 \wedge x_3) \vee (x_2 \oplus x_3)$ 



(b) 
$$\neg (x_0 ? x_2 \lor x_3 : x_2 \land x_3)$$

(c)  $(a) \wedge (b)$ 

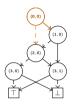


(a) 
$$(x_0 \wedge x_1 \wedge x_3) \vee (x_2 \oplus x_3)$$



(b) 
$$\neg (x_0 ? x_2 \lor x_3 : x_2 \land x_3)$$

**(c)** (a) ∧ (b)



Priority Queue: Qapp:1:

 $[ (0,0) \xrightarrow{\top} ((1,0),(2,1)) ,$  $(0,0) \xrightarrow{\bot} ((2,0),(2,0)) ,$  (0,0)

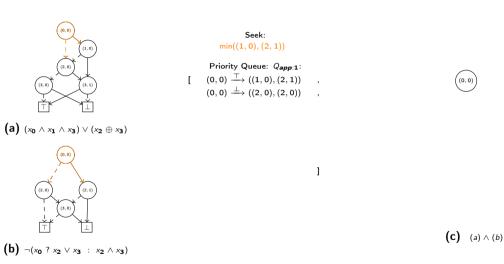
(a) 
$$(x_0 \wedge x_1 \wedge x_3) \vee (x_2 \oplus x_3)$$

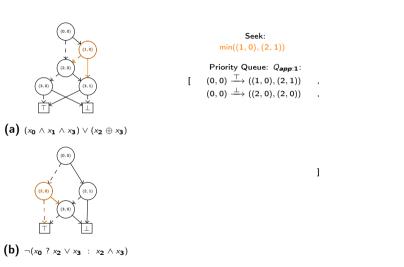


(b) 
$$\neg (x_0 ? x_2 \lor x_3 : x_2 \land x_3)$$

]

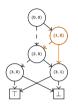
(c) 
$$(a) \wedge (b)$$





(0,0)

**(c)** (a) ∧ (b)



(a) 
$$(x_0 \wedge x_1 \wedge x_3) \vee (x_2 \oplus x_3)$$



**(b)**  $\neg (x_0 ? x_2 \lor x_3 : x_2 \land x_3)$ 

# $\begin{array}{c} \text{Seek:} \\ \min((1,0),(2,1)) \end{array}$

Priority Queue: Qapp:1:

 $(0,0) \xrightarrow{\top} ((1,0),(2,1))$ 

 $(0,0) \xrightarrow{\perp} ((2,0),(2,0)) ,$ 

 $(1,0) \xrightarrow{\perp} ((2,0),(2,1))$ 

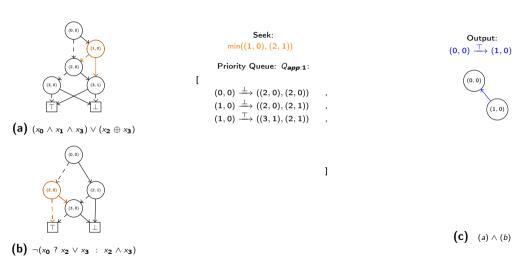
 $(1,0) \xrightarrow{\top} ((3,1),(2,1))$ 

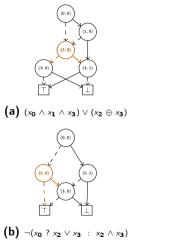
(0,0)

(1,0)

J

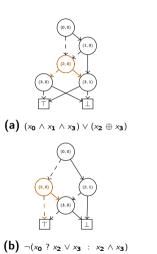
(c)  $(a) \wedge (b)$ 





Seek: min((2,0),(2,0))Priority Queue: Qapp:1:  $(0,0) \xrightarrow{\perp} ((2,0),(2,0))$  $(1,0) \xrightarrow{\perp} ((2,0),(2,1))$  $(1,0) \xrightarrow{\top} ((3,1),(2,1))$ 

Output: (c)  $(a) \wedge (b)$ 

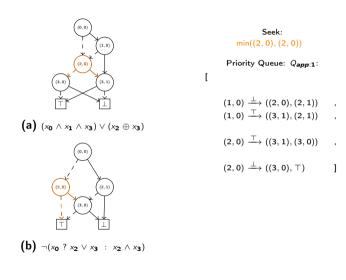


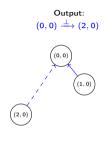
Seek: min((2,0),(2,0))Priority Queue: Qapp:1:  $(0,0) \xrightarrow{\perp} ((2,0),(2,0))$  $(1,0) \xrightarrow{\perp} ((2,0),(2,1))$  $(1,0) \xrightarrow{\top} ((3,1),(2,1))$  $(2,0) \xrightarrow{\top} ((3,1),(3,0))$  $(2,0) \xrightarrow{\perp} ((3,0),\top)$  ]

Output:

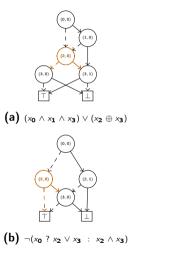
(c)  $(a) \wedge (b)$ 

(2,0)





**(c)** (a) ∧ (b)

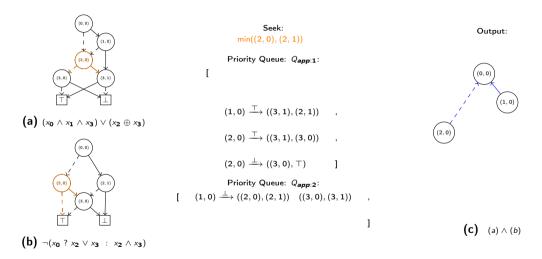


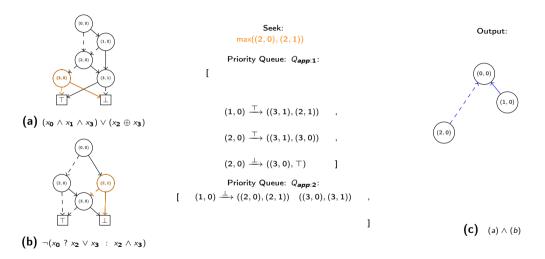
Seek: min((2,0),(2,1))Priority Queue: Qapp:1:  $(1,0) \xrightarrow{\perp} ((2,0),(2,1))$  $(1,0) \xrightarrow{\top} ((3,1),(2,1))$  $(2,0) \xrightarrow{\top} ((3,1),(3,0))$  $(2,0) \xrightarrow{\perp} ((3,0),\top)$  ]

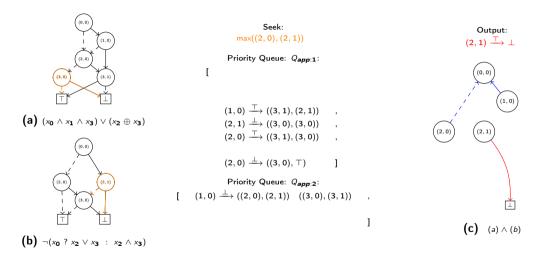
(2,0)

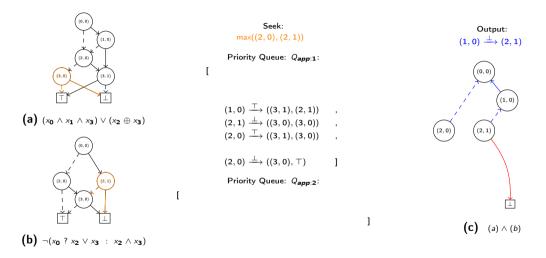
Output:

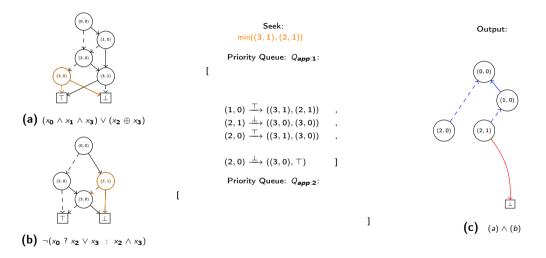
(c)  $(a) \wedge (b)$ 

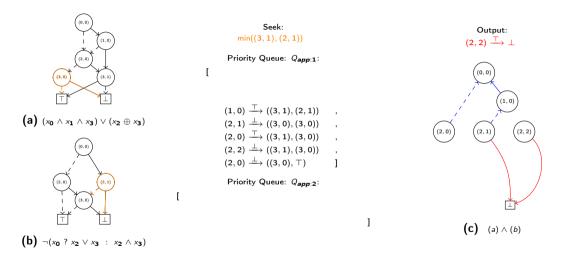


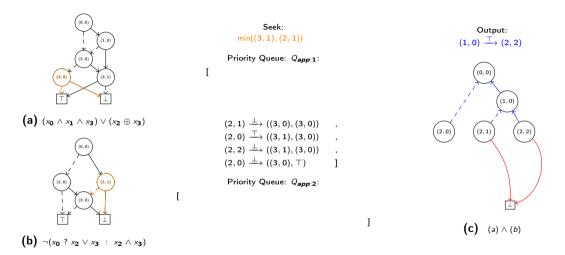


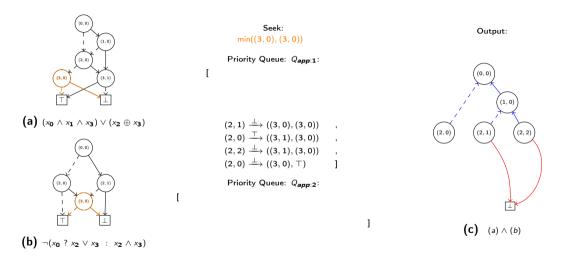


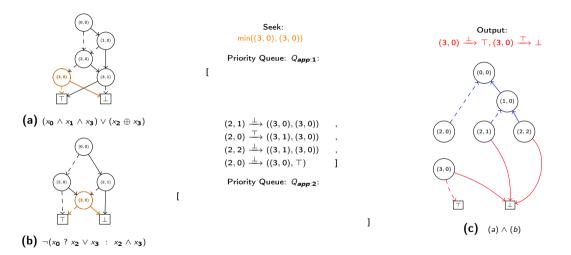


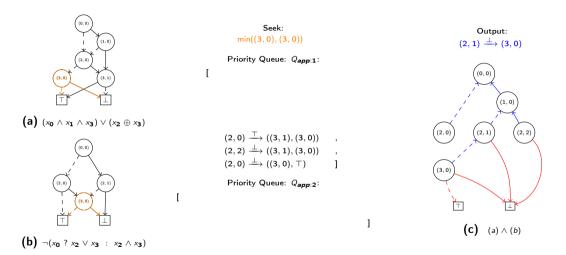


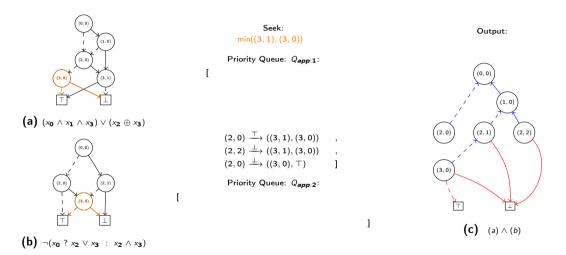


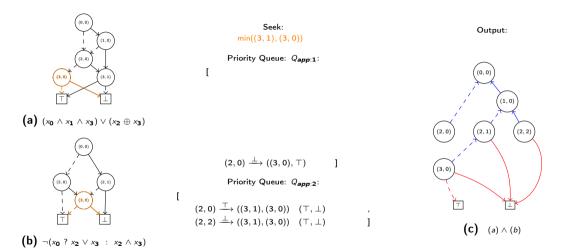


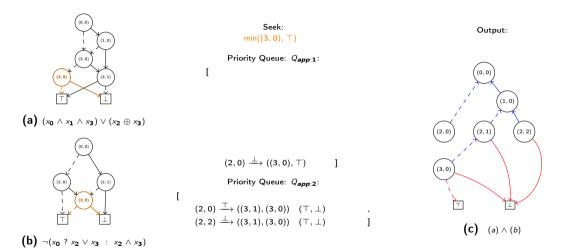


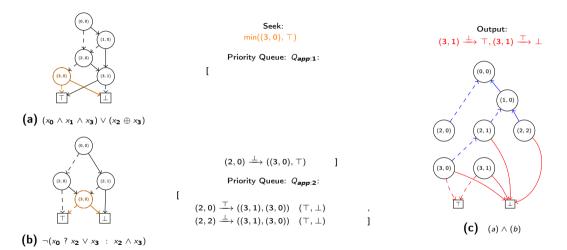


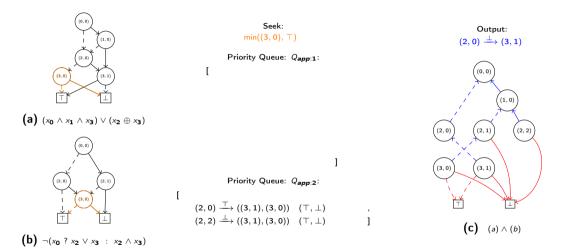


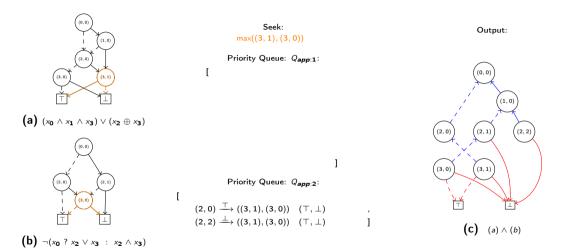


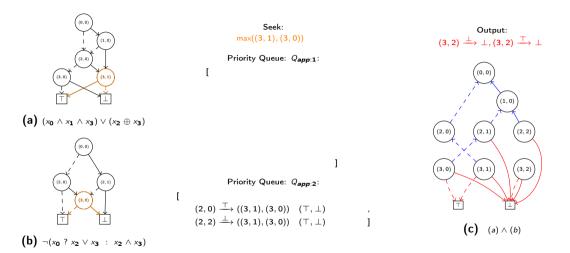


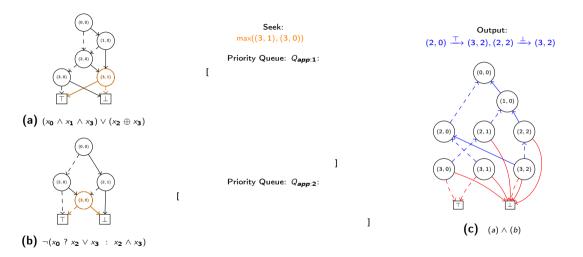


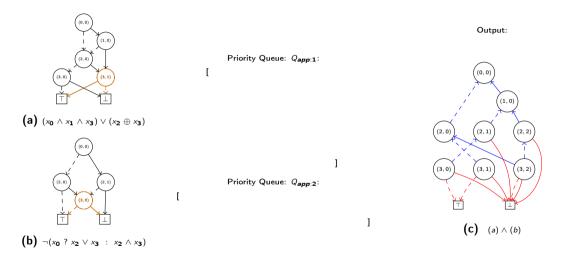




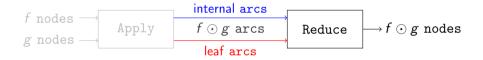


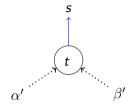






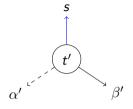
#### **Apply**





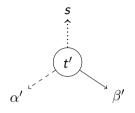
## **Time-Forward Processing**

Send reduction t' with  $Q_{red}$  : PriorityQueue $\langle (s o t') 
angle$  descending on parent s.



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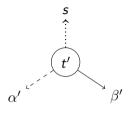


### **Time-Forward Processing**

Send reduction t' with  $Q_{red}$  : PriorityQueue $\langle (s \to t') \rangle$  descending on parent s.

#### Observation (semi-tranposition)

 $\leftarrow$  :  $s \rightarrow t$  (Internal Arcs) provide parents of unreduced node t.

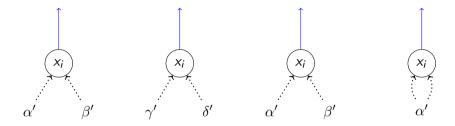


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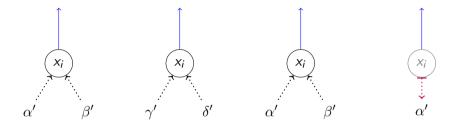
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- $\leftarrow$  :  $s \rightarrow t$  (Internal Arcs) provide parents of unreduced node t.
- ightarrow:  $s 
  ightarrow \mathbb{B}$  (Terminal Arcs) are reduced and already sorted as per  $Q_{red}$ .



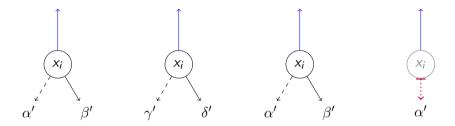
#### Reduce Level i:

1 Obtain nodes from  $Q_{red}$  and terminal arcs. Filter and remember redundant nodes.



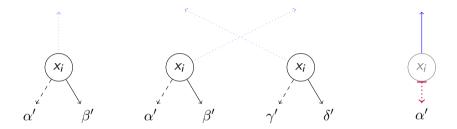
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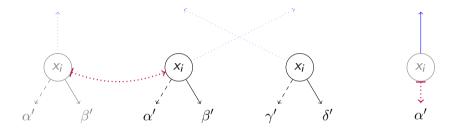


#### Reduce Level i:

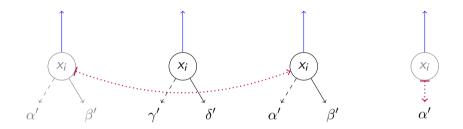
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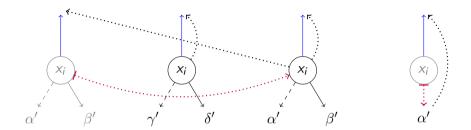
- 1 Obtain nodes from  $Q_{red}$  and terminal arcs. Filter and remember redundant nodes.
- Sort remaining nodes by children, output unique nodes, and remember duplications.



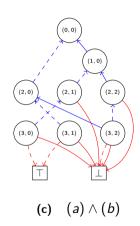
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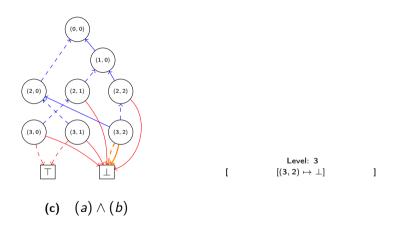


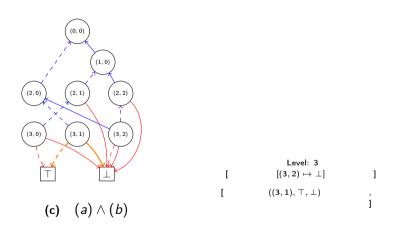
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- Sort remaining nodes by children, output unique nodes, and remember duplications.
- 3 Sort back to match internal arcs and forward to parents with  $Q_{red}$ .

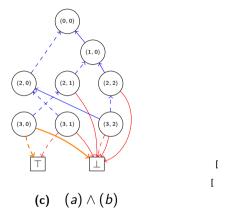


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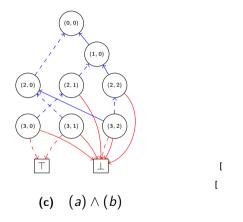


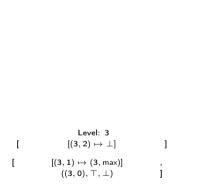




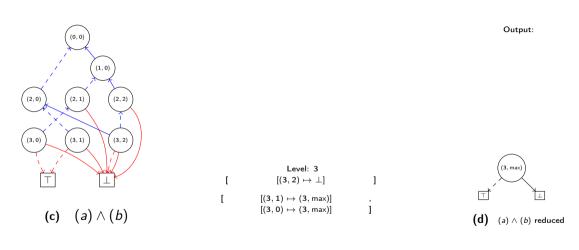


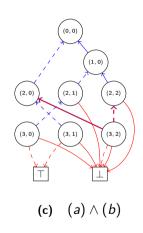
Level: 3 
$$[(3,2)\mapsto \bot]$$
 
$$[ \qquad \qquad ((3,1),\top,\bot) \qquad \qquad ((3,0),\top,\bot) \qquad \qquad ]$$

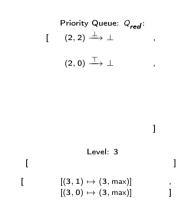


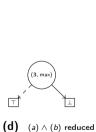


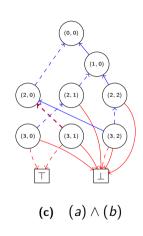


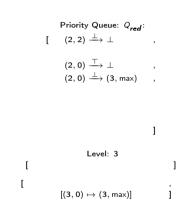


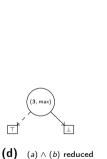


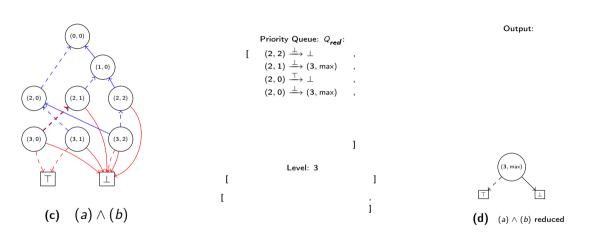


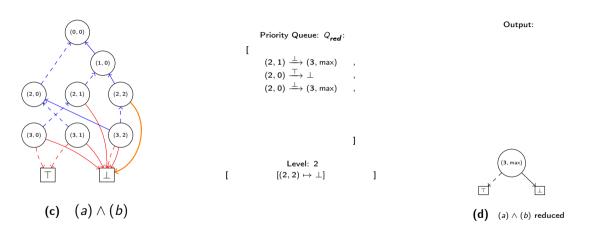


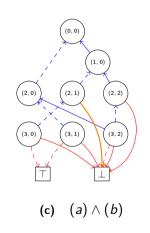


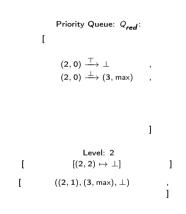




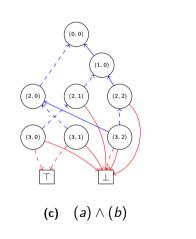


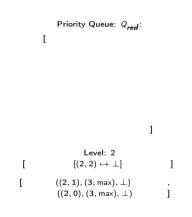




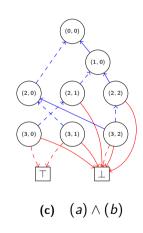


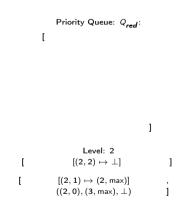


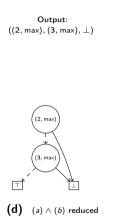


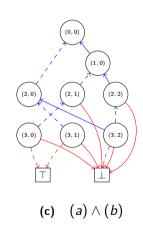


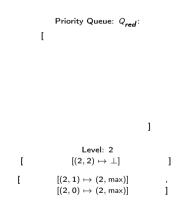


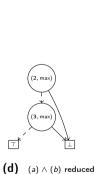


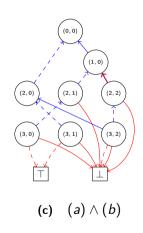


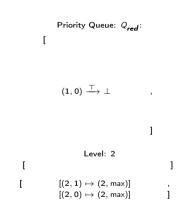


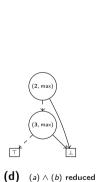


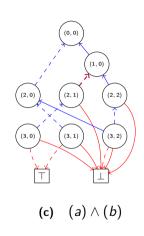


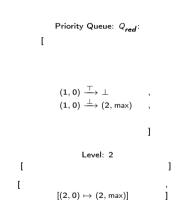


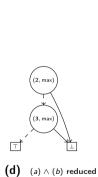


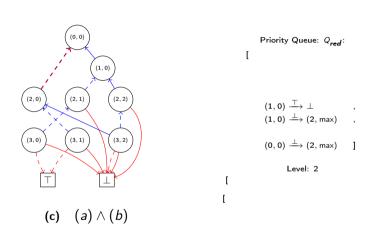


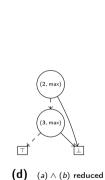


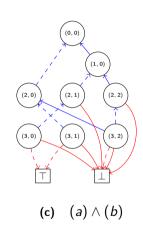


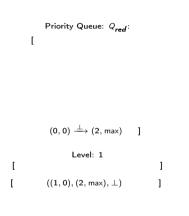


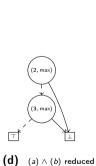


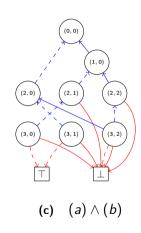


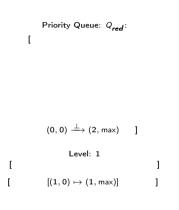


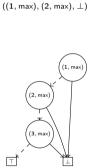




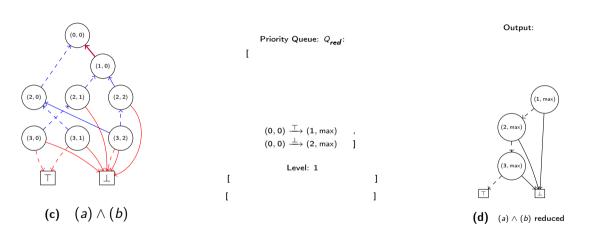


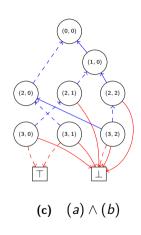


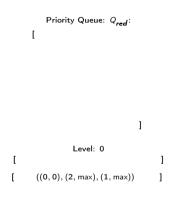


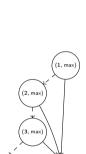


Output:

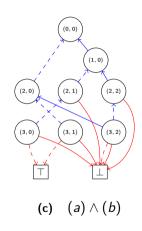


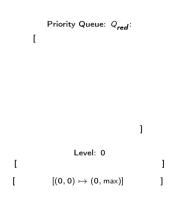


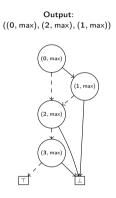


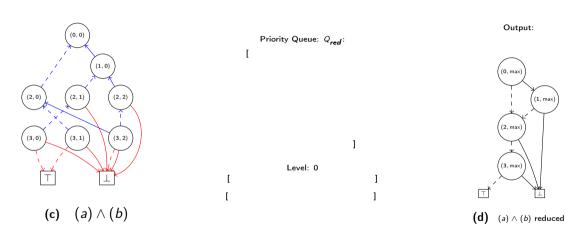


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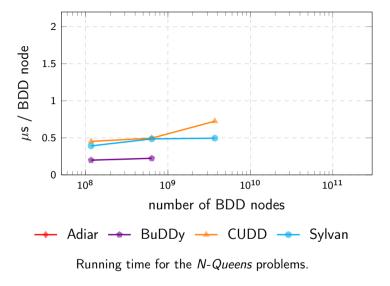


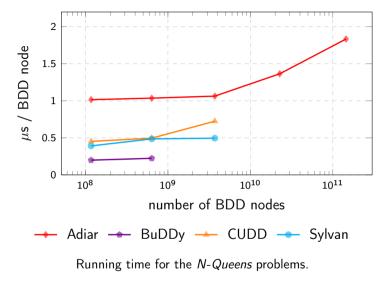






Algorithm	I/O-Complexity
bdd_pathcount	$O(\operatorname{sort}(N_f))$
bdd_not	$2N_f/B$
bdd_restrict	$O(\operatorname{sort}(N_f))$
bdd_apply	$O(\operatorname{sort}(N_f \cdot N_g))$





#### **Contents**

What are Binary Decision Diagrams?

Why do they break?

How can we fix it?

CountPaths

Apply

**Equality Checking** 

Algorithm	I/O-Complexity
bdd_pathcount	$O(\operatorname{sort}(N_f))$
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bdd_equal	?

$$f\leftrightarrow g\equiv \top$$

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$$\underbrace{\textit{O}(\mathsf{sort}(\textit{N}^2))}_{\texttt{Apply}} + \underbrace{\textit{O}(\mathsf{sort}(\textit{N}^2))}_{\texttt{Reduce}} + \underbrace{\textit{O}(1))}_{\texttt{check is }\top} = \textit{O}(\mathsf{sort}(\textit{N}^2))$$

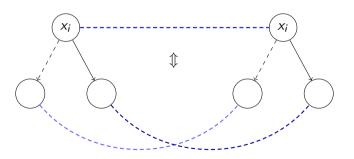
#### Theorem (Bryant '86)

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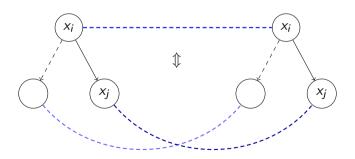
Let  $\pi$  be a variable order and  $f: \mathbb{B}^n \to \mathbb{B}$  then there exists a unique (up to isomorphism) Reduced Ordered Binary Decision Diagram representing f with ordering  $\pi$ .

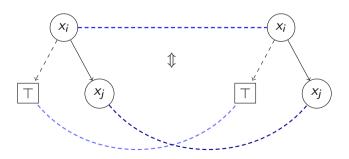
## Trivial cases: $f \not\equiv g$ if there is a mismatch in

	$N_f  eq N_g$	Number of nodes	O(1) I/Os
•	$L_f  eq L_g$	Number of levels	<i>O</i> (1) I/Os
•	$N_{f,i}  eq N_{g,i}$	Number of nodes on a level	O(L/B) I/Os
•	$L_{f,i} \neq L_{g,i}$	Label of an <i>i</i> th level	O(L/B) I/Os

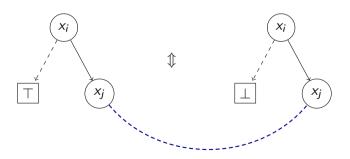


## Theorem (Bryant '86)

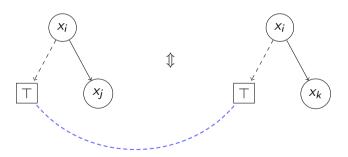




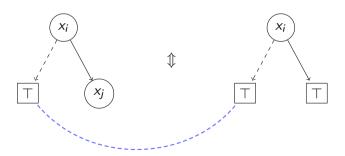
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IsIsomorphic(f, g)

- Check whether root  $v_f$  of f and root  $v_g$  of g have a local violation.
- Check  $low(v_f) \sim low(v_g)$  and  $high(v_f) \sim high(v_g)$  "recursively".

Return false on first violation. If there are no violations then return true.

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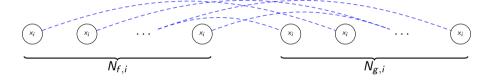
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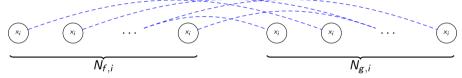
$$\underbrace{O(\mathsf{sort}(\mathit{N}^2))}_{\mathtt{Apply'}} + \underbrace{O(\mathsf{sort}(\mathit{N}^2))}_{\mathtt{Reduce}} + \underbrace{O(1))}_{\mathtt{check is} \ \top} = O(\mathsf{sort}(\mathit{N}^2))$$

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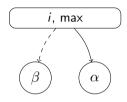
Let  $\pi$  be a variable order and  $f: \mathbb{B}^n \to \mathbb{B}$  then there exists a unique (up to isomorphism) Reduced Ordered Binary Decision Diagram representing f with ordering  $\pi$ .



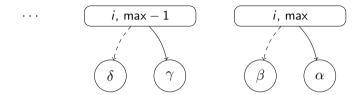
Return false if more than  $N_{f,i} = N_{g,i}$  pairs of nodes are checked on level i.

$$\underbrace{O(\mathsf{sort}(\Sigma_i \ N_{f,i}))}_{\mathsf{Apply''}} = O(\mathsf{sort}(N))$$

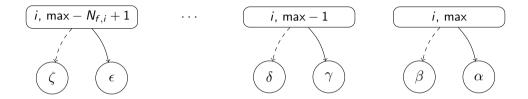
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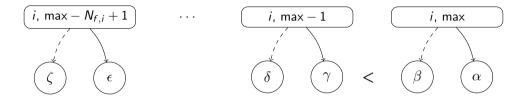
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Each level output by the Reduce algorithm has the following properties:

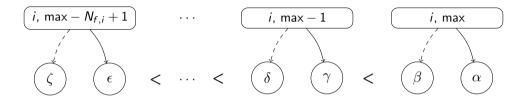
■ Nodes on level *i* have their identifiers *consecutively* numbered.



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#### Observation

- Nodes on level *i* have their identifiers *consecutively* numbered.
- Nodes on level *i* are output sorted by their children.

## Theorem

If  $G_f$  and  $G_g$  are outputs of Reduce.

 $G_f \sim G_g \iff For \ all \ i \in [0; N_f) \ the \ node \ G_f[i] \ matches \ G_g[i] \ numerically.$ 

#### Proof.

← : Must describe the exact same graph.

 $\Rightarrow$ : Strong induction on BDD levels bottom-up.

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## Corollary

If  $G_f$  and  $G_g$  are outputs of Reduce then  $f \equiv g$  is computable using  $2 \cdot N/B$  I/Os.

$$\begin{array}{c|c} & \text{Algorithm} & \text{Time (s)} \\ \hline f \leftrightarrow g \equiv \top & 0.38 \end{array}$$

Checking the (EPFL Benchmark) voter circuit's single output gate ( $|N_f| = |N_g| = 5.76$  MiB).

Algorithm Time (s)
$$f \leftrightarrow g \equiv \top \quad 0.38$$

$$O(\operatorname{sort}(N)) \quad 0.058$$

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Algorithm	Time (s)
$f\leftrightarrow g\equiv \top$	0.38
O(sort(N))	0.058
2N/B	0.006

Checking the (EPFL Benchmark) voter circuit's single output gate ( $|N_f| = |N_g| = 5.76$  MiB).

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✓ soelvsten@cs.au.dk

ssoelvsten.github.io

# Adiar

github.com/ssoelvsten/adiar

ssoelvsten.github.io/adiar



Algorithm
bdd_pathcount

Depth-First









Time-Forwarding







bdd not

bdd\_apply

bdd\_restrict





 $O(N_f)$ 

 $O(N_f)$ 













 $O(\operatorname{sort}(N_f))$