

Networking

3-point

* Idea \rightarrow Reality

- Drive
- tenacity
- persistence

Random Access protocol

Controlled Access prot.

Channelization protocol

Multiple Access protocol

* The Data link layer is responsible for transmission of data between two nodes.

Funcⁿ:-

- Data link control
- Multiple Access control

* Data link - The data link control is responsible for reliable transmission of message over transmission channel by using techniques like framing, error control and flow control. For data link control 'refer to Stop and Wait ARQ'. (Back)

* Multiple Access control - If there is a dedicated link b/w the sender and the receiver then data link control layer is sufficient, however if there is not dedicated link present then multiple stations can access the channel simultaneously. Hence Multiple Access protocol required to decrease collision and avoid cross talk.

Ex:- If Teacher asked a question & all student started answering simultaneously so then data/answer is overlapped and collision occurred. So, Teacher control the students to 'Answer one at a time'.

Computer Network

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Random Access protocol :-

↳ ALOHA ↳ CSMA ↳ CSMA/CD ↳ CSMA/CA

⇒ In Random Access:- all stations have same superiority that is no station has more priority than another station. Any station can send data depending on medium's state (idle or busy).

- Features:-
- ① There is no fixed time for sending data.
 - ② There is no fixed sequence of stations sending data.

> ALOHA:- Designed for Wireless LAN, But also applicable for shared medium. In this, Multiple stations can transmit data at the same time and can hence lead to collision and data being garbled.

> Pure ALOHA:- When a station sends data it waits for an acknowledgement. If Acknowledgement doesn't come within the allotted time then the station waits for a random amount of time called back-off time (T_b) and re-sends the data. Since, different stations wait for different amount of time, the probability of further collision decreases.

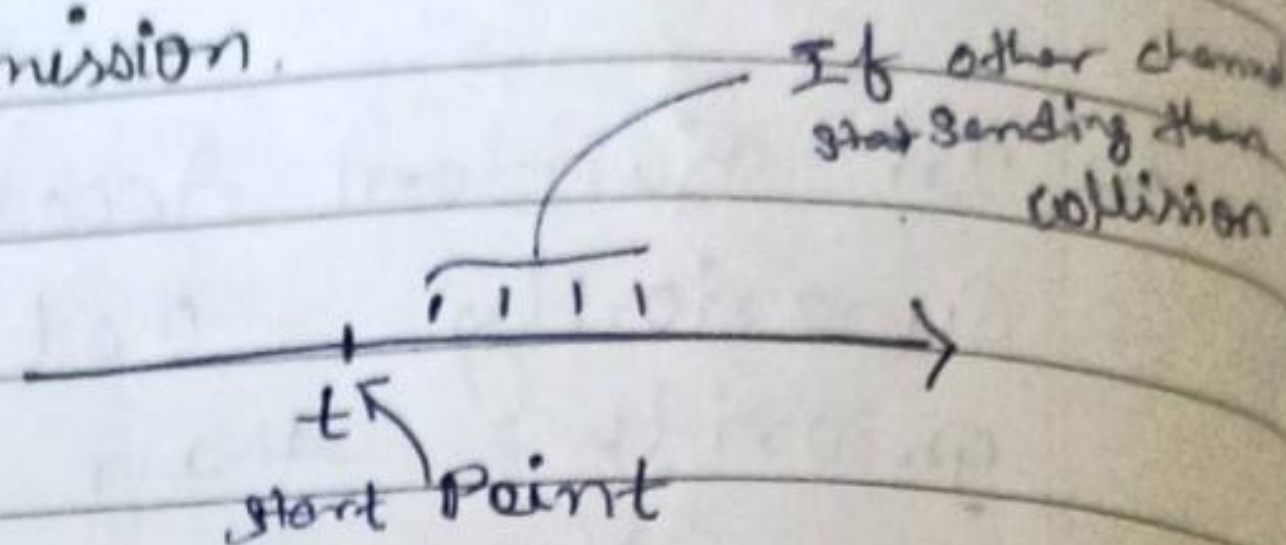
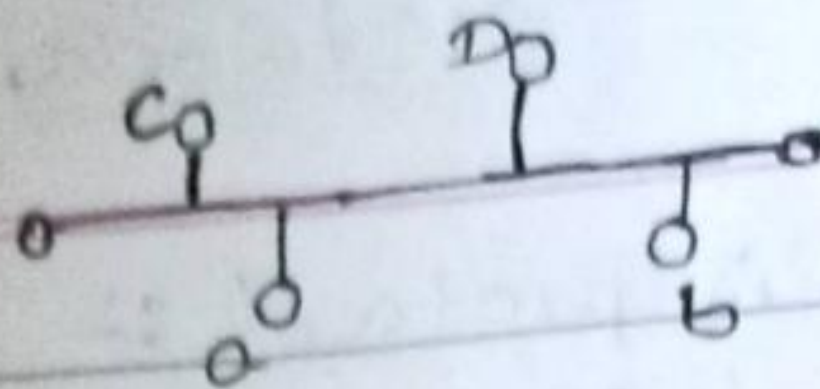
Pure, ALOHA

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Features:-

- Random Access prot.
- Ack. is there
- LAN Based / Any time transmission.
- only transmission time,
No propagation time



→ Vulnerable time? (V_t)

→ Efficiency $\eta = ?$

$$[e = 2.71]$$

$$[\text{Efficiency } \eta = G \times e^{-2G}]$$

Example :- $TT: \text{Assumed some tol All}$

$$[TT = \frac{M}{B.W} = \frac{100 \text{ bits} \times 10^{-3}}{100 \text{ Kbps}} =$$

$$[= 10 \text{ msec}]$$

$$\frac{d\eta}{dG} = G \times e^{-2G} (-2) + e^{-2G} (1) = 0$$

$$= [e^{-2G} (-2G + 1) = 0]$$

$$= [-2G + 1 = 0]$$

Max left $G = \frac{1}{2}$

G :- No. of station looking for particular T_t

Note

If 2 station available & 1 (one) is only transmitting message. So efficiency 'High'

$$[\text{Vulnerable time} = 2 \times T_t]$$

TT/T_t :- Transmission Time

$$[M = \text{message size}]$$

$$[B.W = \text{Bandwidth}]$$

$$V_t = 2 \times 10 \text{ ms} = 20 \text{ ms}$$

Note: If V_t time is given and in that

time no other channel should send the message in past and also after transmission of station A. Then collision will not occur.

$$[\text{max efficiency } \eta = \frac{1}{2} \times e^{-2 \times \frac{1}{2}} = \frac{1}{2} \times e^{-1}]$$

$$= 0.184 = [18.4\%] \text{ for } [G \approx 0.5]$$

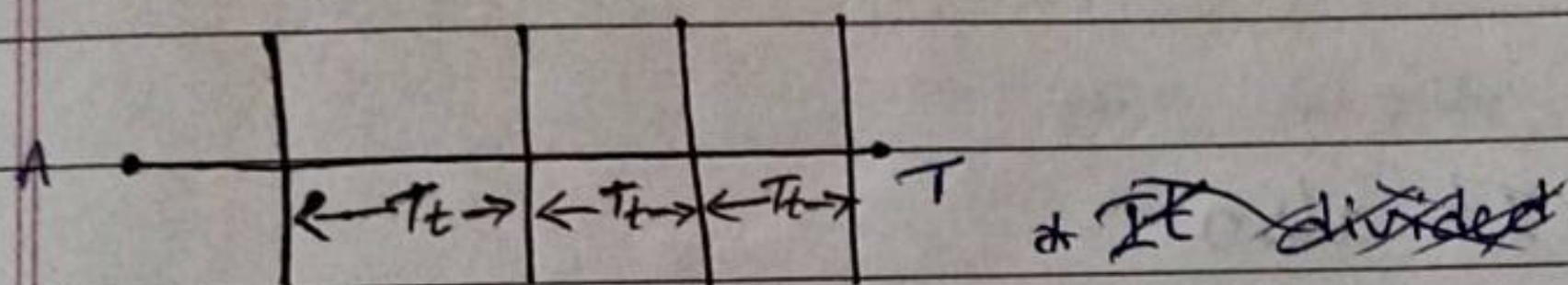
Slotted ALOHA

⇒ It is similar to pure ALOHA, except that we divide time into 'slots' and sending of data is allowed only at the beginning of these slots.

If a station misses out the allowed time, it must wait for the next 'slot'. This reduces the probability of collision.

Note:- Max throughput/efficiency $\boxed{= 0.368 = 36.8\%}$ for $G=1$

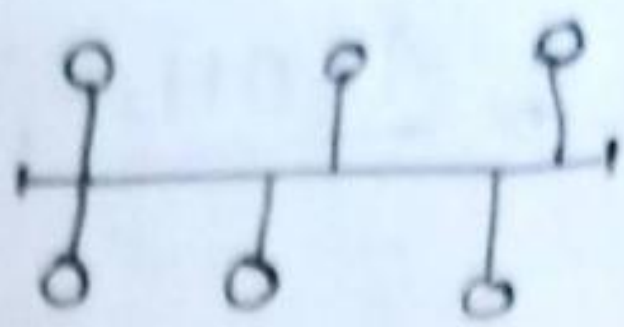
$$(ii) \boxed{VT = T_t}$$



It is Fixed for All

Rule:- In this topic Any station can start sending message 'at beginning', Not after begin. or (Inmid). But If No. of stations started at beginning then collision occurred.

$$\begin{aligned} \frac{d\eta}{dG} &= G \times e^{-G}(-1) + e^{-G}(1) = 0 \\ &= e^{-G}(-G+1) = 0 \\ &= -G+1 = 0 \\ &= \boxed{G=1} \end{aligned}$$

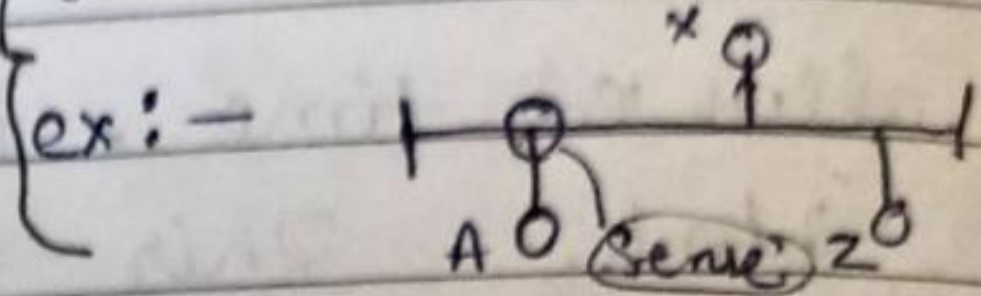


Carrier - Sense Multiple Access (CSMA)

[1-persistent 0-persistent P-persistent.]

→ It sense only own path and point.

* 1-persistent:- It sense continuously that path & point is idle or Not. And if idle then immediately send start sending message.



* 0-persistent:- It also sense the channel if it get busy than it fix a timer, to wait till that. Some time it can be in wait state for too long period.

(Hybrid)
* P-persistent:- Some working as a 1-persistent, But if channel is free then it 'P' value & follow them to start transmitting message.

* Wifi = follow the P-persistent concept.

- * Key Points *
1. persistent:- दृढ़/निश्चि
 2. Carrier:- वाहक
 3. vulnerable:- उपेक्षित/गैर
 4. Propagation:- प्रसारण / प्रचार
 5. efficiency:- क्षमता
 6. Superiority:- श्रेष्ठता/प्रधानता
 7. garbled:- विकृत
 8. Acknowled:- स्वीकृति / प्राप्ति सूचना

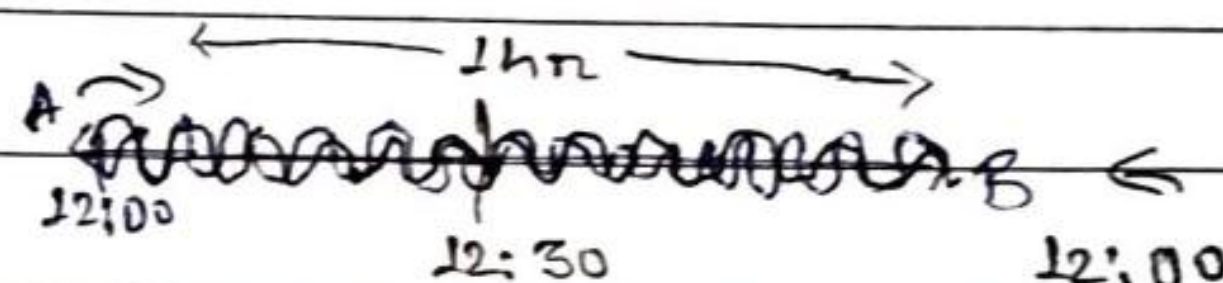
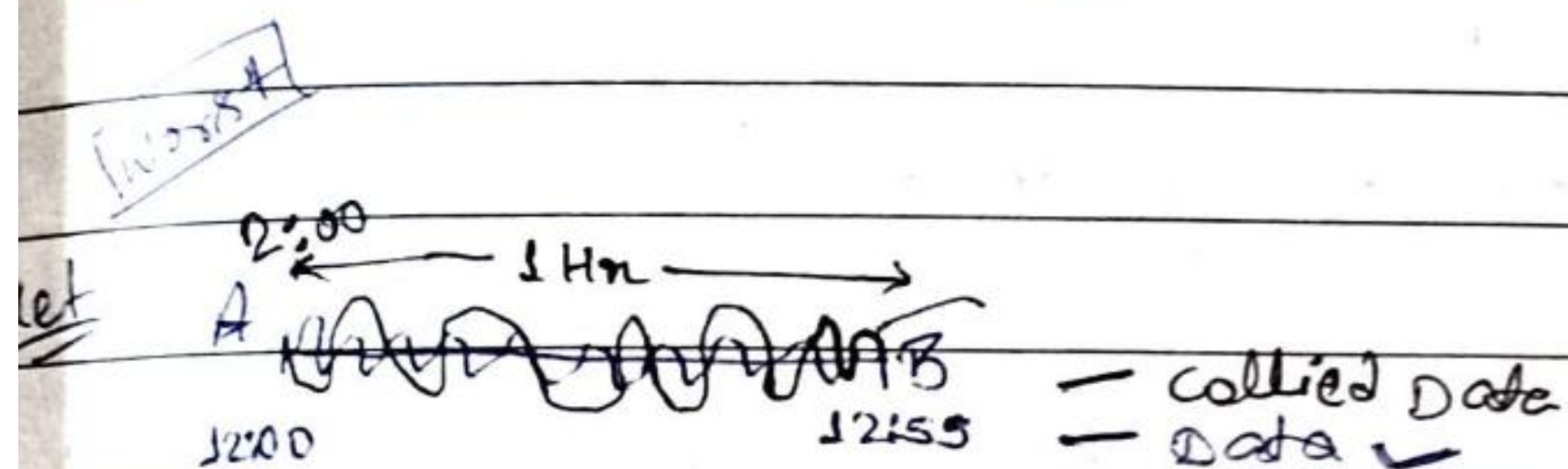
CSMA/CD (wire medium).

* Carrier Sense Multiple Access / Collision Detection

↳ Three Energy Level :-

① E

* No Ack. is There ⊗



* In above fig. A is started transmission at 12:00. but 'B' started at 12:59 before getting A's Data. so, this is "worst case" for A because A got collided data at 2:00 and if A is transmitting data at 2:00 then get that their data collided but 'A' is idle then they will not understand.

$$\begin{aligned} & \text{IMP.} \rightarrow \begin{cases} TT \geq 2 * PD \\ L \geq 2 * PD * BW \end{cases} \end{aligned}$$

TT: Transmission Time
PD: Propagation delay
L: Length of the message.
BW: Bandwidth

* efficiency $\eta = \frac{1}{1 + 6.44a}$

$$a = \frac{PD}{TT}$$

$$TT = \frac{L}{B}$$

CSMA/CD

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* → It used in Ethernet/LAN Technology.

* Bus Topology.

* Not used in Star Topo. or point to point.

Efficiency conclusion:-

1. If distance ↑ increases, the efficiency of CSMA Decreases ↓
2. ^{CSMA} It is not suitable for long-distance N/w like WAN.
3. If Length of packet is bigger, the efficiency of CSMA also 'increase'
But Max limit for L. is 1500 Bytes.

CSMA/CA

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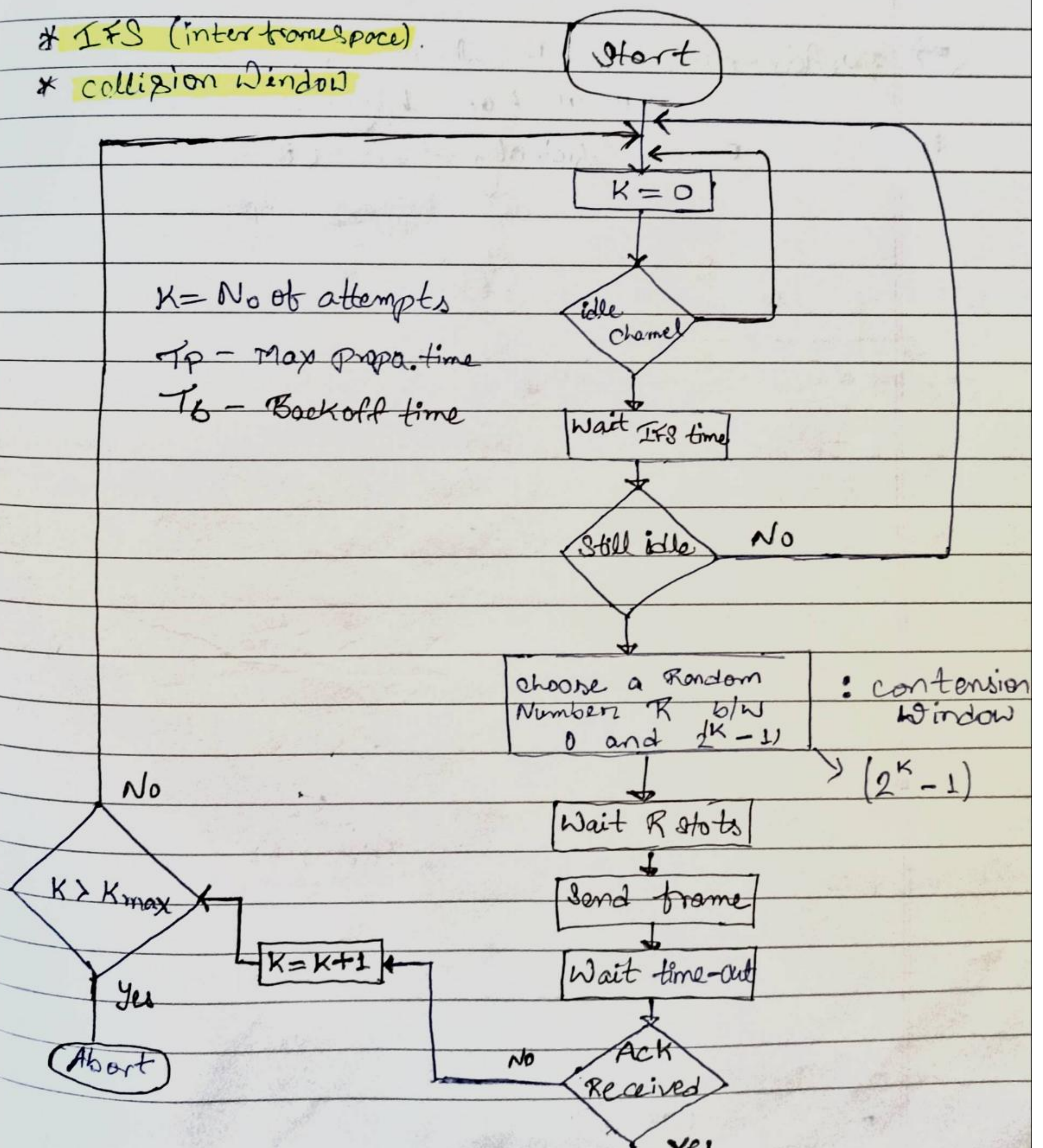
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⇒ It used in 'Wireless'.

- In wireless N/w, much of the sent energy is lost in transmission. The received signal has very little energy. Therefore, A collision may add up 5 to 10% additional energy. This is not effective collision detection.

* IFS (inter frame space).

* collision Window



⇒ Advantages of CDMA/CA :-

- CDMA/CA prevents collision.
- Due to acknowledgements, data is not lost unnecessarily.
- It avoids wasteful transmission.
- It is very much suited for wireless transmissions

- ## ⇒ Disadvantage
- The Algo. calls for long wait.
 - It has high consumption.
 - Less efficiently than CDMA/CD.