

Rel alg - procedural

SQL - non-procedural
- declarative



Rel alg: elem operations

$$(\sigma, \pi, \rho, \delta, \Delta, \dots)$$

- many of them
output relations

= inputs:

(a) core rel alg:
sets of tuples

(b) extended: bags of
tuples



SQL: UNION ALL
 $R \cup S$ - "bag union"
 (bag version of
 set union)

| $R:$ | <u>A</u> | <u>B</u> |
|------|----------|----------|
| | 1 | 2 |
| | 3 | 4 |
| | 1 | 2 |
| | 5 | 6 |
| | 3 | 4 |

| $S:$ | <u>A</u> | <u>B</u> |
|------|----------|----------|
| | 3 | 4 |
| | 3 | 4 |
| | 5 | 6 |
| | 5 | 6 |



Bag union: output the unchanged inputs

| $R \oplus S:$ | <u>A</u> | <u>B</u> |
|---------------|----------|---------------|
| 1 | 2 | - twice |
| 3 | 4 | - four times |
| 5 | 6 | - three times |



$$R \Delta S$$

$$\equiv$$

bag version of set intersection
 [min on # of copies in the inputs]

$$R \Delta S:$$

| | A | B |
|--|---|---|
| | 3 | 4 |
| | 3 | 4 |
| | 5 | 6 |



$R \pm \setminus$ - bag version of
set difference

$$R \pm \setminus : \begin{array}{c} A \quad B \\ \hline / \quad \backslash \\ / \quad \backslash \end{array}$$



δ - eliminating duplicates

τ - sorting - potentially on several attributes
 \Rightarrow the output is a list

π generalized - arithmetic, duplicate columns, etc

ρ - grouping and aggregation.
 $=$ standard fns: COUNT, SUM, MIN, MAX, AVG,
~~to~~ [CNT'D]



$$R : \begin{array}{c} A \quad B \\ \hline \end{array}$$

$$\begin{array}{cc} \rightarrow & 0 \quad 1 \\ \rightarrow & 2 \quad 3 \\ \rightarrow & 0 \quad 4 \\ \rightarrow & 2 \quad 3 \\ \rightarrow & 3 \quad 4 \end{array}$$

$$\therefore \begin{array}{c} B \quad C \\ \hline \end{array}$$

$$\begin{array}{cc} \rightarrow & 0 \quad 1 \\ \rightarrow & 2 \quad 5 \\ \rightarrow & 2 \quad 4 \\ \rightarrow & 3 \quad 4 \\ \rightarrow & 2 \\ \rightarrow & 4 \end{array}$$



$R \bowtie S$: natural join

$R \bowtie^o S$: natural full outerjoin

$\underline{R} \bowtie_L S$: natural left outerjoin

$R \bowtie_R \underline{S}$: natural right outerjoin

$R \dot{\times} S = R \times S$

$R \bowtie_{R.B \leq S.B} S$ full theta
outerjoin



$R \bowtie S$, $R \bowtie_L S$, $R \bowtie_R S$

| | A | B | C |
|-----------------|------|---|------|
| | 2 | 3 | 4 |
| | 2 | 3 | 4 |
| $R \bowtie_L S$ | 0 | 1 | NULL |
| | 0 | 4 | NULL |
| | 3 | 4 | NULL |
| | NULL | 0 | 1 |
| | NULL | 2 | 5 |
| | NULL | 2 | 4 |
| | NULL | 0 | 2 |

$R \bowtie S$

$R \bowtie_L S$

$R \bowtie_R S$



T: A B

| | |
|---|---|
| 1 | 2 |
| 1 | 2 |
| 3 | 4 |
| 5 | 6 |
| 5 | 6 |

U: A B

| | |
|---|---|
| 1 | 2 |
| 3 | 4 |
| 3 | 4 |
| 3 | 4 |
| 7 | 8 |

$\delta(u) : \frac{A}{1} \frac{B}{2}$

| | |
|---|---|
| 3 | 4 |
| 7 | 8 |



$\tau(\text{Sorting})$: ASC [default]
DESC

$\tilde{\tau}(W)$: list:

| | A | B |
|---|---|---|
| 4 | 1 | . |
| 1 | 2 | . |
| 1 | 3 | . |
| 2 | 3 | . |
| 2 | 3 | . |
| 2 | 5 | . |

W: A B

| | | |
|---|---|---|
| 1 | 3 | |
| 2 | 3 | ← |
| 4 | 1 | . |
| 2 | 3 | ← |
| 2 | 5 | . |
| 1 | 2 | . |



$\tilde{\pi}_A \text{ DESC } (u)$

$\pi_{A, A-B, A} (\tau) :$

| A | $A-B$ | A |
|-----|-------|-----|
| 1 | -1 | 1 |
| 1 | -1 | 1 |
| 3 | -1 | 3 |
| 5 | -1 | 5 |
| 5 | -1 | 5 |



$\mathcal{N}(\rho)$

individual attr names
and/or aggregation fn's
applied to attribute
names



P: A B C

| | | | |
|---|---|---|---|
| | 1 | 2 | 3 |
| . | 4 | 5 | 1 |
| . | 4 | 6 | 2 |
| | 1 | 2 | 7 |
| . | 4 | 5 | 3 |
| | 1 | 2 | 7 |



$$\mathcal{N}_{\min(B)}(P) = \mathbb{Z}$$

$$\underline{\mathcal{N}_{A,B}(P)}$$

$$\underline{\mathcal{N}_{A,B,\text{sum}(C)}(P)}$$



Evaluating a $\prod_L R$

E] Partition R by values of
the unaggregated attributes
in the subscript L of \prod

$$(1, 2) \rightarrow \begin{array}{c|ccc} & A & B & C \\ \hline 1 & 1 & 2 & 3 \\ 2 & & 2 & 7 \\ \hline & 1 & 2 & 7 \end{array}$$

$$(4, 5) \rightarrow \begin{array}{c|ccc} & A & B & C \\ \hline 4 & 4 & 5 & 1 \\ 5 & & 5 & 3 \\ \hline & 4 & 5 & 3 \end{array}$$

$$(4, 6) \rightarrow \begin{array}{c|ccc} & A & B & C \\ \hline 4 & 4 & 6 & 2 \\ 6 & & & \\ \hline & 4 & 6 & 2 \end{array}$$



[2] Within each individual partition, apply all the aggregation functions in the subscript L of π^L

[3] For each individual partition, form and output a single tuple, with its values of all the unaggregated attributes in L^U , and with its value for each aggregated output.



| $\nabla_{A, B, \text{sum}(C)}$ | (P) | $=$ | $\begin{array}{c ccc} & A & B & \cancel{\text{sum}(C)} \\ \hline 1 & 2 & & 17 \\ 4 & 5 & & 4 \\ 4 & 6 & & 2 \end{array}$ |
|--------------------------------|-------|-----|--|
|--------------------------------|-------|-----|--|

$\nabla_{A, \text{MIN}(B), \text{COUNT}(C)}$ $\underbrace{(\sum_{C>3} (P))}_{(P)}$

(1) only non-NULL values
are considered

(2) If all the values are NULL
 \Rightarrow the aggregation will

(E)

θ, A, \pm - different
atom set
ops

$$R \cup R = R$$

$$R \cup R \neq R$$



SQL is declarative

- you are describing, in your query, a relation* that you want to see in the answer



* All student info:

SELECT [list of all the
attr.]
FROM Student;

SELECT *
FROM Student;



All students with GPA > 3.7:

(3) SELECT *
(1) FROM Student
(2) WHERE GPA > 3.7;

π_* ($\sigma_{GPA > 3.7}$ Student))

