

DIGITAL LOGIC & CIRCUIT DESIGN

Unit: 1

**DIGITAL LOGIC & CIRCUIT
DESIGN**

SUBJECT CODE: ACSE0304

B. Tech: 3rd Sem.

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Associate Professor



NOIDA INSTITUTE OF ENGINEERING AND TECHNOLOGY, GREATER NOIDA

Evaluation Scheme

Sl. No.	Subject Codes	Subject Name	Periods			Evaluation Schemes				End Semester		Total	Credit
			L	T	P	CT	TA	TOTAL	PS	TE	PE		
WEEKS COMPULSORY INDUCTION PROGRAM													
1	AAS0301A	Engineering Mathematics-III	3	1	0	30	20	50		100		150	4
2	ACSE0306	Discrete Structures	3	0	0	30	20	50		100		150	3
3	ACSE0304	Digital Logic & Circuit Design	3	0	0	30	20	50		100		150	3
4	ACSE0301	Data Structures	3	1	0	30	20	50		100		150	4
5	ACS0301	Introduction to Cloud Computing	3	0	0	30	20	50		100		150	3
6	ACSE0305	Computer Organization and Architecture	3	0	0	30	20	50		100		150	3
7	ACSE0354	Digital Logic & Circuit Design Lab	0	0	2				25		25	50	1
8	ACSE0351	Data Structures Lab	0	0	2				25		25	50	1
9	ACS0351	Cloud Computing lab	0	0	2				25		25	50	1
10	ACSE0359	Internship Assessment-I	0	0	2				50			50	1
11	ANC0301/ ANC0302	Cyber Security*/ Environmental Science*(Non Credit)	2	0	0	30	20	50		50		100	0
12		MOOCs (For B.Tech. Hons. Degree)											
		GRAND TOTAL										1100	24

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Course Contents / Syllabus

UNIT-I: Digital System and Binary Numbers: Number System and its arithmetic, Signed binary numbers, Binary codes, Cyclic codes, Hamming Code, Simplification of Boolean Expression: K-map method up to five variable, SOP and POS Simplification Don't Care Conditions, NAND and NOR implementation, Quine Mc-Clusky Method (Tabular Method).

UNIT II : Combinational Logic: Combinational Circuits: Analysis Procedure, Design Procedure, Code Converter, Binary Adder-Subtractor, Decimal Adder, Binary Multiplier, Magnitude Comparator, Decoders, Encoders Multiplexers, Demultiplexers.

UNIT III: Sequential Logic and Its Applications: Storage elements: Latches & Flip Flops, Characteristic Equations of Flip Flops, Excitation Table of Flip Flops, Flip Flop Conversion, Registers, Shift Registers, Ripple Counters, Synchronous Counters, Other Counters: Johnson & Ring Counter.

UNIT IV: Synchronous & Asynchronous Sequential Circuits: Analysis of clocked Sequential Circuits with State Machine Designing, State Reduction and Assignments, Design Procedure. Analysis procedure of Asynchronous Sequential Circuits, Circuit with Latches, Design Procedure, Reduction of State and flow Table, Race-free State Assignment, Hazards.

UNIT-V: Memory & Programmable Logic Devices: Basic concepts and hierarchy of Memory, Memory Decoding, RAM: SRAM, DRAM, ROM: PROM, EPROM, Auxiliary Memories, PLDs: PLA, PAL; Circuit Implementation using ROM, PLA and PAL; CPLD and FPGA..

Branch Wise Application

- The aim of **Digital Logic and Circuit Design** is to provide basic idea about Digital Logic Design or Digital Electronics or Digital. It contains description about Number Systems, Circuit Minimization, Logic Gates, Sequential circuits, Number system conversion, Sum of products and product of sum, binary number operations.
- Digital Logic and Circuit Design used in Mobile Phones, Calculators and Digital Computers, Radios and communication Devices, Signal Generator, Smart Card, Cathode Ray Oscilloscope(CRO), Analog to digital converters (ADC), Digital to analog converters (DAC), etc.

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CONTENT

- Course Objective
- Unit Objective
- Course Outcome
- Co and Po Mapping
- Topic Objective
- Prerequisite
- Introduction to Number Systems
- Code conversion
- Boolean algebra
- SOP & POS form
- K-Map
- QM Algorithm
- Daily quiz & MCQs
- Old Question Papers
- Recap
- Video Links
- Weekly Assignments
- References

COURSE OBJECTIVE

Course Objective: The student will be able to learn about

To Apply concepts of Digital Binary System and implementation of Gates

To Analyze and design of Combinational logic circuits

To Analyze and design of Sequential logic circuits with their applications

To Implement the Design procedure of Synchronous & Asynchronous Sequential Circuits

To Apply the concept of Programmable Logic devices with circuit implementation

COURSE OUTCOME

Course Outcomes: At the end of this course students will able to:

CO1	Apply concepts of Digital Binary System and implementation of Gates
CO2	Design and analyze combinational circuits with MUX / DEMUX, Decoder & Encoder
CO3	Design and analyze of Sequential logic circuits with their applications
CO4	Implement the Design procedure of Synchronous & Asynchronous Sequential Circuits
CO5	Apply the concept of Programmable Logic devices with circuit implementation

Program Outcomes

S. No.	Description
1.	Engineering knowledge: Apply the knowledge of mathematics, science, engineering fundamentals, and an engineering specialization for the solution of complex engineering problems.
2.	Problem analysis: Identify, formulate, research literature, and analyse complex engineering problems reaching substantiated conclusions using first principles of mathematics, natural sciences, and engineering sciences.
3.	Design/development of solutions: Design solutions for complex engineering problems and design system components or processes that meet the specified needs with appropriate consideration for public health and safety, and cultural, societal, and environmental considerations.
4.	Conduct investigations of complex problems: Use research based knowledge and research methods including design of experiments, analysis and interpretation of data, and synthesis of information to provide valid conclusions.
5.	Modern tool usage: Create, select, and apply appropriate techniques, resources, and modern engineering and IT tools, including prediction and modeling to complex engineering activities, with an understanding of the limitations.

Program Outcomes

	Description
	<p>6. The engineer and society: Apply reasoning informed by the contextual knowledge to assess societal, health, safety, legal and cultural issues and the consequent responsibilities relevant to the professional engineering practice.</p> <p>7. Environment and sustainability: Understand the impact of the professional engineering solutions in societal and environmental contexts, and demonstrate the knowledge of, and need for sustainable development.</p> <p>8. Ethics: Apply ethical principles and commit to professional ethics and responsibilities and norms of the engineering practice.</p> <p>9. Individual and team work: Function effectively as an individual, and as a member or leader in diverse teams, and in multidisciplinary settings.</p> <p>10. Communication: Communicate effectively on complex engineering activities with the engineering community and with the society at large, such as, being able to comprehend and write effective reports and design documentation, make effective presentations, and give and receive clear instructions.</p> <p>11. Project management and finance: Demonstrate knowledge and understanding of the engineering and management principles and apply these to one's own work, as a member and leader in a team, to manage projects and in multidisciplinary environments.</p> <p>12. Life-long learning: Recognize the need for, and have the preparation and ability to engage in independent and life-long learning in the broadest context of technological change.</p>

COs-POs Mapping

COs	PO1	PO2	PO3	PO4	PO5	PO6	PO7	PO8	PO9	PO10	PO11	PO12
CO1	3	2	-	-	-	-	-	-	-	-	-	2
CO2	3	3	2	-	-	-	-	-	-	-	-	2
CO3	2	3	2	2	-	-	-	-	-	-	-	2
CO4	3	3	3	2	-	-	-	-	-	-	-	2
CO5	3	2	1	-	-	-	-	-	-	-	-	2
AVERAGE	2.8	2.6	2	2	-	-	-	-	-	-	-	2

On successful completion of graduation degree the Electronics and Communication, graduates will be able to:

1. **Engineering knowledge:** Apply the knowledge of mathematics, science and electronics & communication engineering to work effectively in the industry based on same or related area.
2. **Design/development of solutions:** Use their skills to work in modern electronics & communication engineering tools, software and equipment's to design solutions for complex problems in the related field that meet the specified needs of the society.
3. **Individual and team work:** Function effectively as an individual and as a member or leader of a team by qualifying through examinations like GATE, IES, PSUs, TOEFL, GMAT and GRE etc.

COs-PSOs Mapping

CO	PSO1	PSO2	PSO3
CO1	3	3	3
CO2	3	3	2
CO3	3	3	-
CO4	3	3	-
CO5	3	3	-

Program Education Objectives

The Program Educational Objectives (PEOs) of B.Tech (ECE) program are as follows:

PEO-1 To have excellent scientific and engineering breadth so as to comprehend, analyze, design and solve real- life problems using state-of-the-art technology.

PEO-2 To lead a successful career in industries or to pursue higher studies or to understand entrepreneurial endeavors.

PEO-3 To effectively bridge the gap between industry and academics through effective communication skill, professional attitude and a desire to learn.

Result analysis

Not Applicable

Question Paper Template

Printed page:

Subject Code:

Roll No:

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NOIDA INSTITUTE OF ENGINEERING AND TECHNOLOGY, GREATER NOIDA

(An Autonomous Institute Affiliated to AKTU, Lucknow)

B.Tech B.Voc./MBA/MCA/M.Tech (Integrated)

(SEM: **SESSIONAL EXAMINATION - I**)(2021-2022)

Subject Name:

Time: 1.15Hours

Max. Marks: 30

General Instructions:

- All questions are compulsory. Answers should be brief and to the point.
- This Question paper consists of pages & questions.
- It comprises of three Sections, A, B, and C. You are to attempt all the sections.
- **Section A** Question No. 1 is objective type questions carrying 1 mark each, Question No. 2 is very short answer type carrying 2 mark each. You are expected to answer them as directed.
- **Section B** Question No. 3 is Short answer type questions carrying 5 marks each. You need to attempt any two out of three questions given.
- **Section C** Question No. 4 & 5 are Long answer type (within unit choice) questions carrying 6 marks each. You need to attempt any one part a or b.
- Students are instructed to cross the blank sheets before handing over the answer sheet to the invigilator.
- No sheet should be left blank. Any written material after a blank sheet will not be evaluated/checked.

SECTION - A		[8]	
1.	Attempt all parts	(4×1=4)	CO
a.		(1)	
b.		(1)	
c.		(1)	
d.		(1)	
2.	Attempt all parts	(2×2=4)	CO
a.		(2)	
b.		(2)	
SECTION - B			
3.	Answer any two of the following.	(2×5=10)	CO
a.		(5)	
b.		(5)	
c.		(5)	

SECTION - C			
4.	Answer any one of the following (Any one can be applicable if applicable)	(2×6=12)	CO
a.	Question	(6)	
b.	Question	(6)	
5.	Answer any one of the following.		
a.		(6)	
b.		(6)	

Prerequisite and Recap

- Basic knowledge of Binary Digits.
- Basic knowledge of Number System.
- Basic knowledge of logic gates.
- Basic Knowledge of flip-flops.
- Basic knowledge of combinational circuits.
- Basic knowledge of Sequential circuits.

Brief introduction about the subject with videos

- This course is intended to provide the students with a comprehensive understanding of the fundamental of digital logic circuit. The design of circuits and systems whose input and outputs are represented as discrete variables. These variables are commonly binary i.e., two states in nature. Design at the circuit level is usually done with truth table and state tables. Students will be able to analyze design and implement combinational and sequential circuits.

https://www.youtube.com/watch?v=BoIOLczVulQ&list=PLyqSpQzTE6M_dZdF7Bd-UncI5_L_1VkXF

<https://www.youtube.com/watch?v=oNh6V91zdPY&list=PLbRMhDVUMnge4gDT0vBWjCb3Lz0HnYKkX>

<https://www.youtube.com/watch?v=CeD2L6KbtVM&list=PL803563859BF7ED8C>

CO-PO and PSO Mapping

COs	PO1	PO2	PO3	PO4	PO5	PO6	PO7	PO8	PO9	PO10	PO11	PO12
ACSE0304.1	3	2	-	-	-	-	-	-	-	-	-	2
ACSE0304.2	3	3	2	-	-	-	-	-	-	-	-	2
ACSE0304.3	2	3	2	2	-	-	-	-	-	-	-	2
ACSE0304.4	3	3	3	2	-	-	-	-	-	-	-	2
ACSE0304.5	3	2	1	-	-	-	-	-	-	-	-	2
AVERAGE	2.8	2.6	2	2	-	-	-	-	-	-	-	2

Course Outcome	PSO1	PSO2	PSO3
ACSE0304.1	3	-	3
ACSE0304.2	2	-	3
ACSE0304.3	2	-	3
ACSE0304.4	2	-	3
ACSE0304.5	2	-	3
Average	2.2	-	3

UNIT OBJECTIVE

- To learn number system & its conversions
- To learn basic rules of Boolean algebra.
- To learn the concept of logic gates and SOP& POS.
- To minimize function using K-Map & QM method.

Result analysis

Not Applicable

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ROAD MAP

**PRE
REQUISITE**

**INTRODUCTION
(CONCEPTS
/ THEORY)**

**NUMERICAL
PRACTICE**

**QUIZZES
,
ASSIGNMENTS
&
UNIT
TESTS**

**L
A
B
O
R
A
T
O
R
Y**

**PERFORMANCE
EVALUATION
(Mid Term,
End Sem
Exam)**

Number System

Topic Objective	Mapping with CO
To understand different number systems.	CO1
To understand number system conversion.	CO1

Number System

Prerequisite:

- Basics of decimal number system.

INTRODUCTION TO NUMBER SYSTEMS

In digital electronics, the number system is used for representing the information.

The number system has different bases and the most common of them are the decimal, binary, octal, and hexadecimal.

The **base or radix** of the number system is the total number of the digit used in the number system.

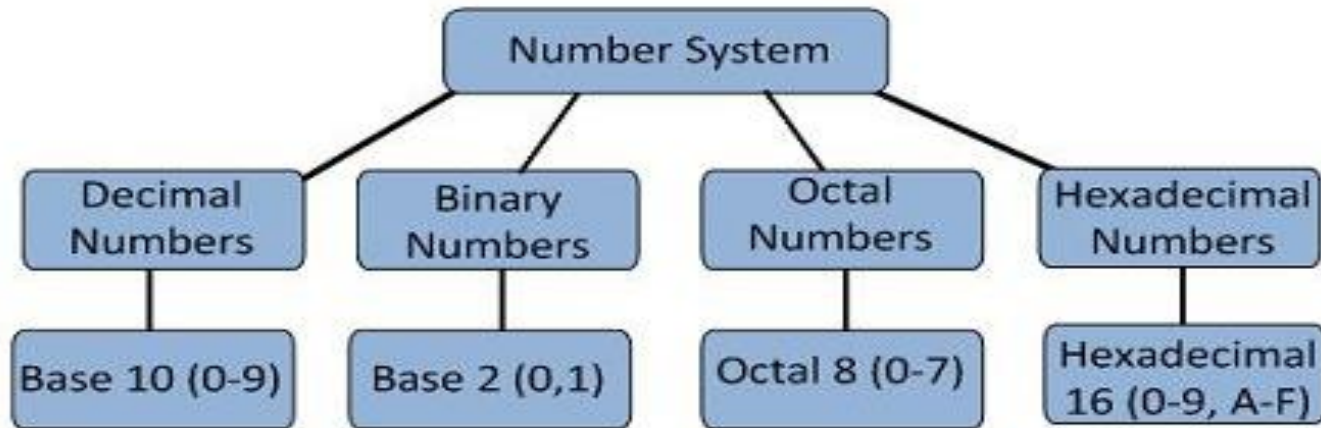
Suppose if the number system representing the digit from 0 – 9 then the base of the system is the 10.

System	Base	Symbols	Used by humans?	Used in computers?
Decimal	10	0, 1, ... 9	Yes	No
Binary	2	0, 1	No	Yes
Octal	8	0, 1, ... 7	No	No
Hexa-decimal	16	0, 1, ... 9, A, B, C, D, E, F	No	No

Types of Number Systems

Some of the important types of number system are:

- Decimal Number System
- Binary Number System
- Octal Number System
- Hexadecimal Number System

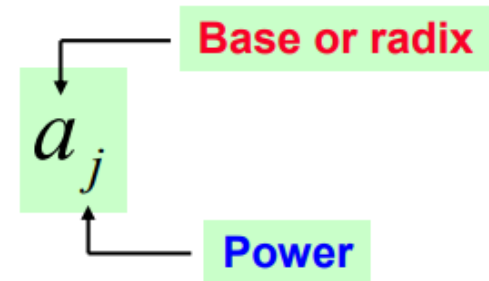


DECIMAL NUMBER SYSTEMS

- The number system is having digit 0, 1, 2, 3, 4, 5, 6, 7, 8, 9.
- This number system is known as a decimal number system because total ten digits are involved.
- The base of the decimal number system is 10.
- Decimal number

$\dots a_5 a_4 a_3 a_2 a_1 \cdot a_{-1} a_{-2} a_{-3} \dots$

↑
Decimal point



⇒ $\dots + 10^5 a_5 + 10^4 a_4 + 10^3 a_3 + 10^2 a_2 + 10^1 a_1 + 10^0 a_0 + 10^{-1} a_{-1} + 10^{-2} a_{-2} + 10^{-3} a_{-3} + \dots$

Example:

$$7,329 = 7 \times 10^3 + 3 \times 10^2 + 2 \times 10^1 + 9 \times 10^0$$

BINARY NUMBER SYSTEMS

- The modern computers do not process decimal number; they work with another number system known as a binary number system which uses only two digits 0 and 1.
- The base of binary number system is 2 because it has only two digit 0 and 1.
- The digital electronic equipment's are works on the binary number system and hence the decimal number system is converted into binary system.
- The table is shown below the decimal, binary, octal, and hexadecimal numbers from 0 to 15 and their equivalent binary number.

BINARY NUMBER SYSTEMS

Decimal	Binary	Octal	Hexa- decimal
0	0	0	0
1	1	1	1
2	10	2	2
3	11	3	3
4	100	4	4
5	101	5	5
6	110	6	6
7	111	7	7
8	1000	10	8
9	1001	11	9
10	1010	12	A
11	1011	13	B
12	1100	14	C
13	1101	15	D
14	1110	16	E
15	1111	17	F

OCTAL NUMBERS

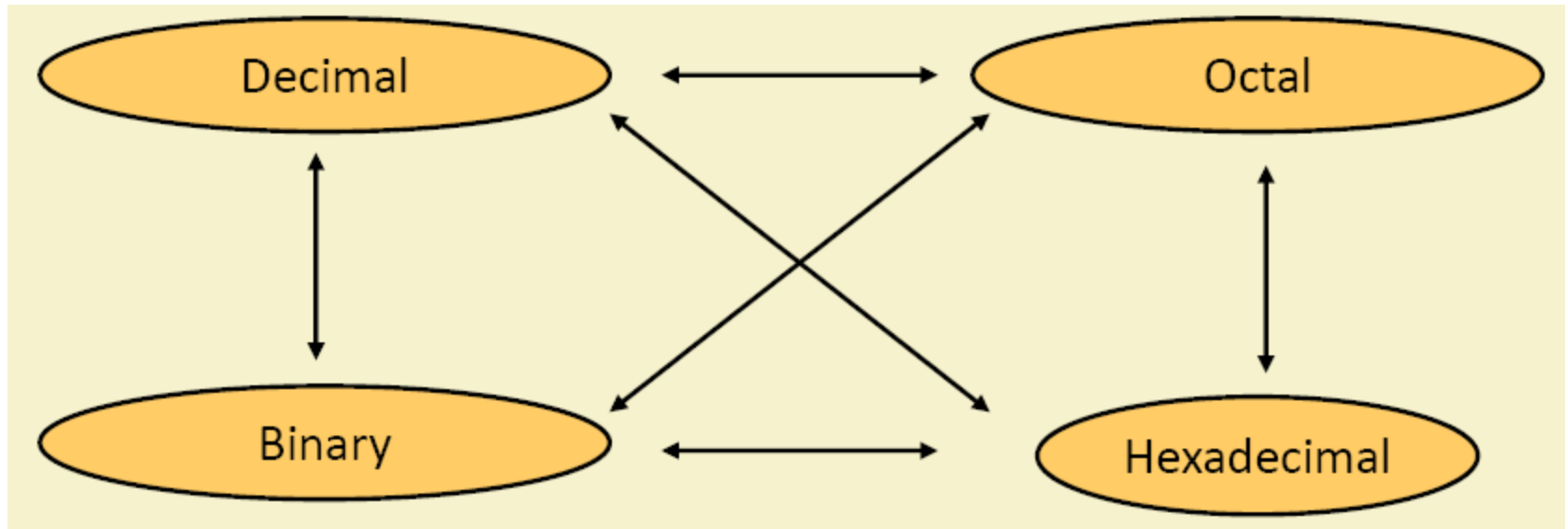
- The base of a number system is equal to the number of digits used, i.e., for decimal number system the base is ten while for the binary system the base is two. The octal system has the base of eight as it uses eight digits 0, 1, 2, 3, 4, 5, 6, 7.
- All these digits from 0 to 7 have the same physical meaning as by decimal symbols, the next digit in the octal number is represented by 10, 11, 12, which are equivalent to decimal digits 8, 9, 10 respectively.
- In this way, the octal number 20 will represent the decimal digit and subsequently, 21, 22, 23.. Octal numbers will represent the decimal number digit 17, 18, 19... etc. and so on.

HEXADECIMAL NUMBERS

- These numbers are used extensively in microprocessor. The hexadecimal number system has a base of 16, and hence it consists of the following sixteen number of digits.
- 0, 1, 2, 3, 4, 5, 6, 7, 8, 9, A, B, C, D, E, F.
- The size of the hexadecimal is much shorter than the binary number which makes them easy to write and remember.
- Let 0000 to 000F representing hexadecimal numbers from zero to fifteen, then 0010, 0011, 0012, ...etc. Will represent sixteen, seventeen, eighteen... etc. till 001F which represent thirty one and so on.

CODE CONVERSION

Conversion among base:



BINARY TO DECIMAL CONVERSION

- Multiply each bit by 2^n , where n is “weight” of the bits.
- The weight is the position of the bit starting from zero from right.
- Add the result
- Example:

Bit “0”

$101011_2 \Rightarrow$

1	x	2^0	=	1
1	x	2^1	=	2
0	x	2^2	=	0
1	x	2^3	=	8
0	x	2^4	=	0
1	x	2^5	=	32
				<hr/>
				43_{10}

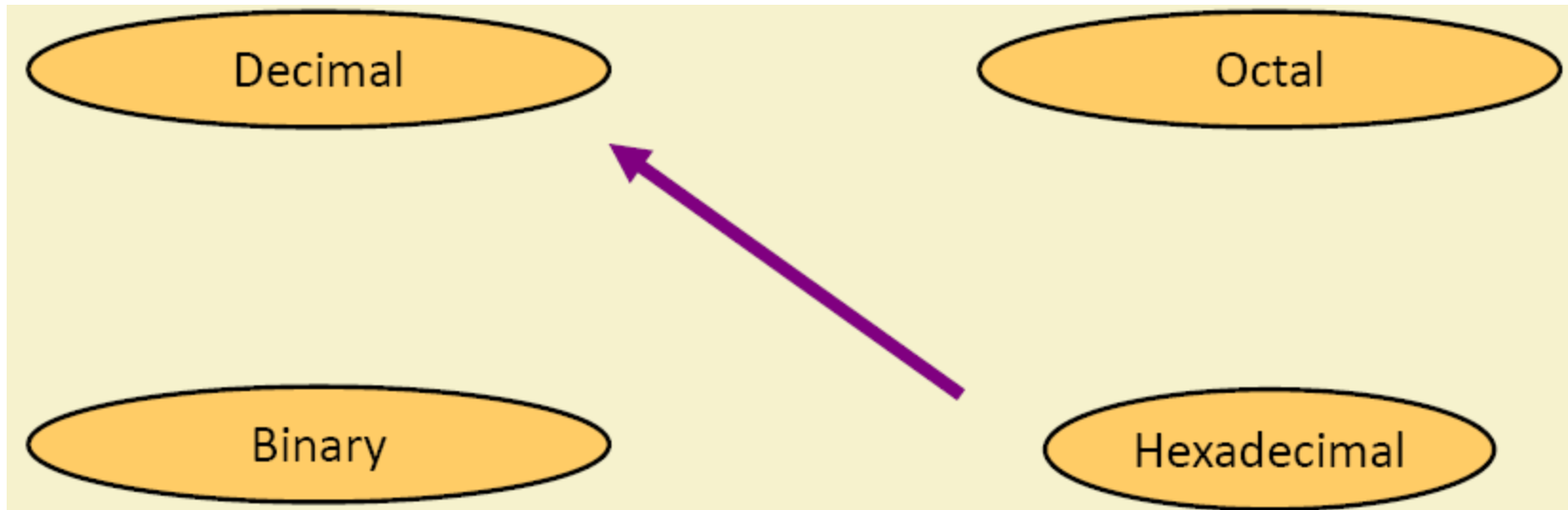
OCTAL TO DECIMAL

- multiplying each digit by 8^n bits where n the weight of the digits.
- the weight is the position of the digit starting from 0 on the right
- add the result
- example:

$724_8 \Rightarrow$

4	$\times 8^0 =$	4
2	$\times 8^1 =$	16
7	$\times 8^2 =$	448
		<hr/>
		468 ₁₀

HEXADECIMAL TO DECIMAL



HEXADECIMAL TO DECIMAL CONVERSION

- Multiplying each digit by 16^n bits where n the weight of the digits.
- The weight is the position of the digit starting from 0 on the right
- Add the result
- Example:

$$\begin{array}{rcll} ABC_{16} \Rightarrow & C \times 16^0 & = 12 \times 1 & = 12 \\ & B \times 16^1 & = 11 \times 16 & = 176 \\ & A \times 16^2 & = 10 \times 256 & = 2560 \\ & & & \hline & & & 2748_{10} \end{array}$$

DECIMAL TO BINARY CONVERSION

Example : convert $(52)_{10}$ to binary.

Sol:

2	52
2	26
2	13
2	6
2	3
2	1
	0

Remainder

0
0
1
0
1
1



$$(52)_{10} = (110100)_2$$

DECIMAL TO OCTAL CONVERSION

Example : Convert $(378.93)_{10}$ to octal Sol:

8	378
8	47
8	5
	0

Remainder

2

7

5



$$0.93 \times 8 = 7.44$$

$$0.44 \times 8 = 3.52$$

$$0.52 \times 8 = 4.16$$

$$0.16 \times 8 = 1.28$$

$$\text{So, } 0.93_{10} = 0.7341_8$$

7

3

4

1



$$\text{Ans: } (378.93)_{10} = (572.7341)_8$$

BINARY TO DECIMAL CONVERSION

Example : Convert $(10101)_2$ to decimal.

$$\begin{aligned} 10101 &= (1 \times 2^4) + (0 \times 2^3) + (1 \times 2^2) + (0 \times 2^1) + (1 \times 2^0) \\ &= 16 + 0 + 4 + 0 + 1 \\ &= 21 \end{aligned}$$

$$(10101)_2 = (21)_{10}$$

Example : Convert $(11011.101)_2$ to decimal.

$$\begin{aligned} 11011.101 &= (1 \times 2^4) + (1 \times 2^3) + (0 \times 2^2) + (1 \times 2^1) + (1 \times 2^0) + (1 \times 2^{-1}) + (0 \times 2^{-2}) + \\ &\quad (1 \times 2^{-3}) \\ &= 16 + 8 + 0 + 2 + 1 + 0.5 + 0 + 0.125 \\ &= 27.625 \end{aligned}$$

$$(11011.101)_2 = (27.625)_{10}$$

Daily Quiz

- What do you mean by radix of any number system?
- Convert $(1101.11)_2$ into decimal.
- Convert $(267.89)_{10}$ into binary.
- The value of $(011010101.110)_2$ in octal and hexadecimal are:
 - a) $(236.6)_8$ and $(D5.B)_{16}$
 - b) $(235.6)_8$ and $(D5.C)_{16}$
 - c) $(325.6)_8$ and $(D5.C)_{16}$
- The value of base x for $(412)_x = (153)_8$ is:
 - a) 9
 - b) 5
 - c) 8
 - d) 4
- Convert $(ACB.8D)_{16}$ into binary and then convert it into octal number system.

FACULTY VIDEO LINKS, YOUTUBE & NPTEL VIDEO LINKS AND ONLINE COURSES DETAILS

Youtube/other Video Links:

- https://www.youtube.com/watch?v=crSGS1uBSNQ&ab_channel=NesoAcademyNesoAcademyVerified
- https://www.youtube.com/watch?v=crSGS1uBSNQ&list=RDCMU_CQYMhOMi_Cdj1CEAU-fv80A&start_radio=1&t=1&ab_channel=NesoAcademyNesoAcademyVerified

Old Questions

- Determine the value of x if $(193)_x = (623)_8$
- Convert $(100000011110)_2$ into hexadecimal and octal number system without converting it to decimal.
- The solution to the quadratic equation $k^2 - 11k + 22 = 0$ are $k = 3$ and $k = 6$. What are the base of number systems?
- Convert $(238.99)_{10}$ into binary, hexadecimal and octal.

Recap

- The number system has different bases and the most common of them are the decimal, binary, octal, and hexadecimal.
- The **base or radix** of the number system is the total number of the digit used in the number system.
- Some of the important types of number system are:

Decimal Number System

Binary Number System

Octal Number System

Hexadecimal Number System

Binary Codes

Topic Objective	Mapping with CO
To understand classification of binary codes.	CO1

Binary Codes

Perquisite:

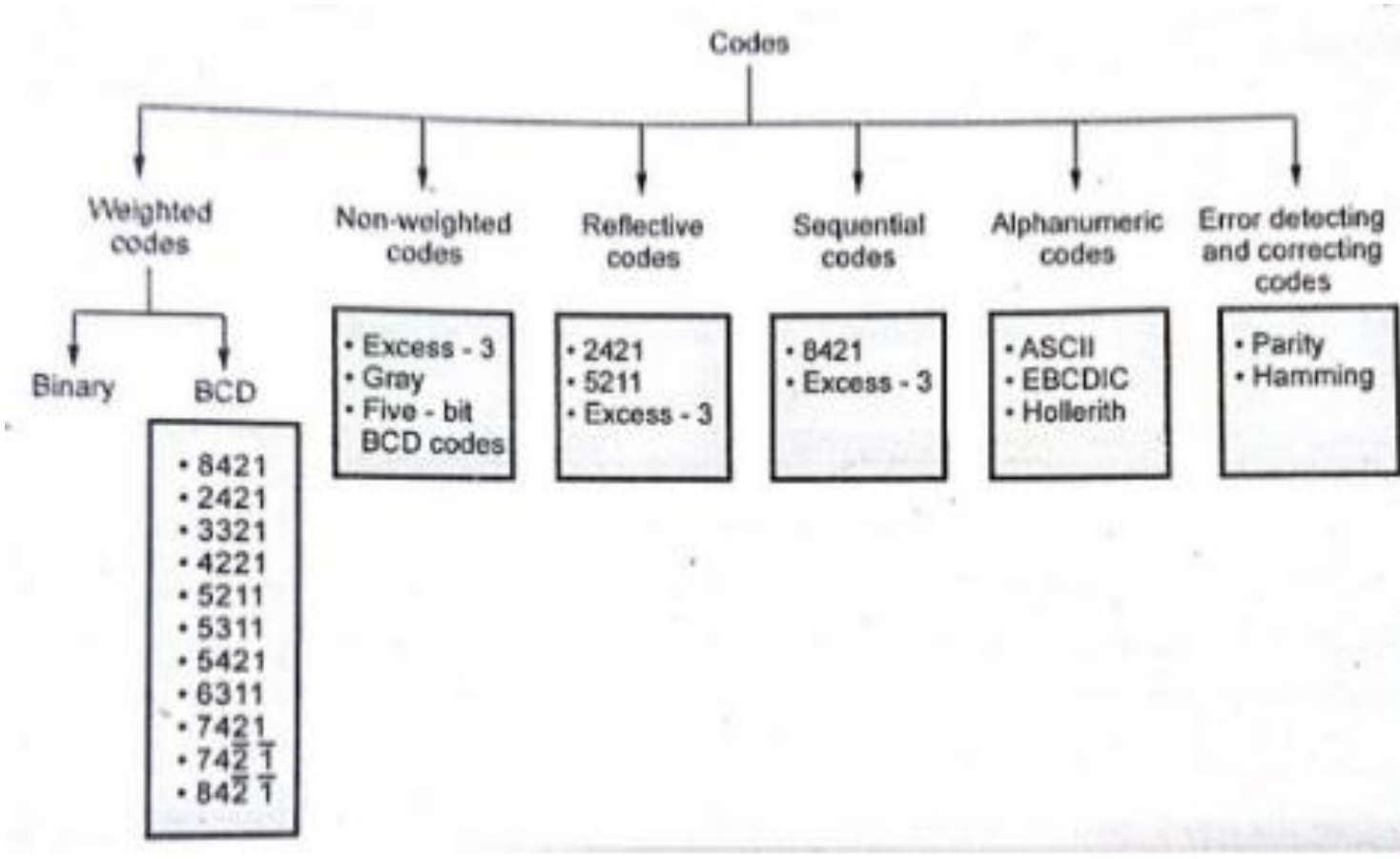
- Knowledge of number system.

Binary Codes

- In the coding, when numbers, letters or words are represented by a specific group of symbols, it is said that the number, letter or word is being encoded. The group of symbols is called as a code.
- The digital data is represented, stored and transmitted as group of binary bits.
- This group is also called as **binary code**. The binary code is represented by the number as well as alphanumeric letter.
- The distinct bit combinations of an n-bit code can be found by counting in binary from 0 to $(2^n - 1)$.

Binary Codes

Classification of Binary Codes



Binary Codes

- **Weighted codes:**
 - In weighted codes, each digit is assigned a specific weight according to its position .
 - Several system of codes are used to express the decimal digits 0 to through 9. These codes have 8421,2421,3321.... All are the weighted codes .
 - In this codes each decimal digit is represented by a group of four bits .
- **Non-weighted codes:**
 - In these codes, positional weights are not assigned.
 - The example of non-weighted codes are excess-3 and gray codes.
- **Reflective codes:**
 - A code is reflective when the code is **self complementing**. In other words, when the code for 9 is the complement the code for 0, 8 for 1, 7 for 2, 6 for 3 and 5 for 4.
 - 2421, 5211 and XS-3 are the examples of reflective codes.

Binary Codes

- **Sequential codes:**
 - In sequential codes, each succeeding 'code is one binary number greater than its preceding code.
 - The 8421 and XS-3 are sequential codes.
- **Alphanumeric codes:**
 - Codes used to represent numbers, alphabetic characters, symbols.
 - Some of these codes are capable to representing some symbols and instructions as well.
 - Example of alphanumeric codes are : ASCII(American Standard Code for Information Interchange), EBCDIC(Extended Binary Coded Decimal Interchange Code).

Binary Codes

- **Error detecting and correcting codes:**
 - When digital data is transmitted from one system to other, unwanted electrical disturbance called “Noise” gets added to it.
 - The noise can force an “error” in digital information. That means a 0 may change to 1 or 1 to 0.
 - To detect and correct such errors we can use some special codes which possess the capacity to detect and correct the error . Such codes are called as error detecting and error correcting codes.

Weighted codes

Binary Coded Decimal (BCD) code:

- BCD codes are weighted codes.
- In this code each decimal digit is represented by a 4-bit binary number. BCD is a way to express each of the decimal digits with a binary code. In the BCD, with four bits we can represent sixteen numbers (0000 to 1111). But in BCD code only first ten of these are used (0000 to 1001). The remaining six code combinations i.e. 1010 to 1111 are invalid in BCD.

Advantages of BCD Codes:

- It is very similar to decimal system.
- We need to remember binary equivalent of decimal numbers 0 to 9 only.

Disadvantages of BCD Codes:

- The addition and subtraction of BCD have different rules.
- The BCD arithmetic is little more complicated.
- BCD needs more number of bits than binary to represent the decimal number. So BCD is less efficient than binary.

BCD Codes

- The **addition of BCD numbers** is slightly different from **binary addition**. Here, the rules of binary addition are partially applicable only to the individual 4-bit groups.
- The **BCD addition**, is thus carried out by individually adding the corresponding 4-bit groups starting from the LSB side.
- If there is a carry to the next group and if the result belongs to any of the 6 illegal states than we add $6_{10}(0110)$ to the sum term of that group and resulting carry is added in the next group.
- **Example: Perform BCD Addition of 5 and 6.**

$$\begin{array}{r}
 0101 \\
 + 0110 \\
 \hline
 1011 \rightarrow \text{Invalid BCD number} \\
 + 0110 \rightarrow \text{Add 6} \\
 \hline
 0001\ 0001 \rightarrow \text{Valid BCD number}
 \end{array}$$

BCD Codes

Example: Perform BCD Addition of 184 and 576

	1	1		
• BCD	0001	1000	0100	184
	+ 0101	0111	0110	+576
Binary sum	<hr/>	<hr/>	<hr/>	
		10000	1010	
Add 6		0110	0110	
BCD sum	<hr/>	<hr/>	<hr/>	<hr/>
	0111	0110	0000	760

Example: Perform BCD Addition of 184 and 976

		1		1	
• BCD	0001	1000	0100	184	
	+ 1001	0111	0110	+976	
Binary sum	<u>1011</u>	<u>10000</u>	<u>1010</u>		
Add 6	<u>0110</u>	<u>0110</u>	<u>0110</u>		
BCD sum	1 0001	0110	0000	1160	

Non-weighted codes

1. Excess-3 code

2. Gray code

Excess-3 Code:

- The excess-3 code is a non-weighted code used to express decimal numbers.
- Excess-3 codes are non-weighted and can be obtained by adding 3 to each decimal digit then it can be represented by using 4 bit binary number for each digit.
- An Excess-3 equivalent of a given binary number is obtained using the following steps:
 1. Find the decimal equivalent of the given binary number.
 2. Add +3 to each digit of decimal number.
 3. Convert the newly obtained decimal number back to binary number to get required excess-3 equivalent.
- Excess-3 code is non-weighted and self complementary code. A self complementary binary codes are always compliment themselves.

Excess-3 Code

- In other words, the 1's complement of an excess-3 code is the excess-3 code for the 9's complement of the corresponding decimal number.
- For example, the excess-3 code for decimal number 5 is 1000 and 1's complement of 1000 is 0111, which is excess-3 code for decimal number 4, and it is 9's complement of number 5.

These are following advantages of Excess-3 codes,

- These are self-complementary codes.
- The codes 0000 and 1111 are not used for any digit which is an advantage for memory organization as these codes can cause fault in transmission line.
- It has no limitation, and it considerably simplifies arithmetic operations.
- It is particularly significant for arithmetic operations as it overcomes shortcoming encountered while using 8421 BCD code to add two decimal digits whose sum exceeds 9.

Excess-3 Code

Example-1 –Convert decimal number 23 to Excess-3 code.

So, according to excess-3 code we need to add 3 to both digit in the decimal number then convert into 4-bit binary number for result of each digit. Therefore,

$= 23+33=56 = 0101\ 0110$ which is required excess-3 code for given decimal number 23.

Example 2– Convert Excess-3 code 1001001 into BCD and decimal number.

So, grouping 4-bit for each group, i.e., 0100 1001 and subtract 0011 0011 from given number. Therefore,

$= 0100\ 1001 - 0011\ 0011 = 0001\ 0110$

So, binary coded decimal number is 0001 0110 and decimal number will be 16.

Binary Codes

Binary codes for the decimal digits

Decimal digit	(BCD) 8421	Excess-3	84-2-1	2421	(Biquinary) 5043210
0	0000	0011	0000	0000	0100001
1	0001	0100	0111	0001	0100010
2	0010	0101	0110	0010	0100100
3	0011	0110	0101	0011	0101000
4	0100	0111	0100	0100	0110000
5	0101	1000	1011	1011	1000001
6	0110	1001	1010	1100	1000010
7	0111	1010	1001	1101	1000100
8	1000	1011	1000	1110	1001000
9	1001	1100	1111	1111	1010000

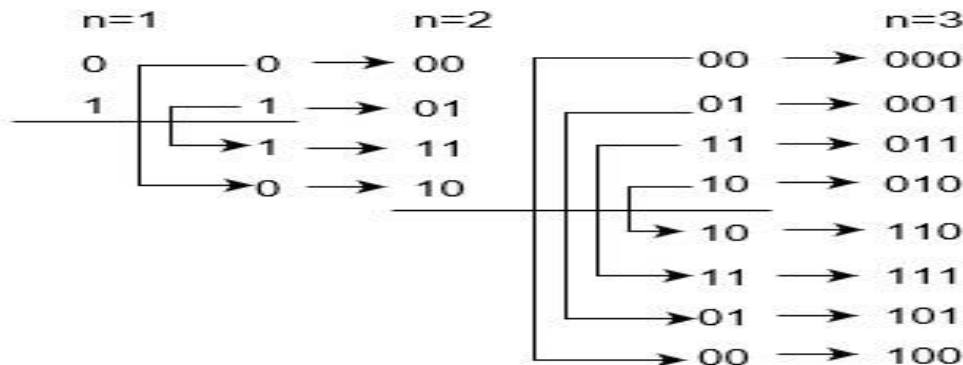
Gray codes:

- The reflected binary code or Gray code is an ordering of the binary numeral system such that two successive values differ in only one bit (binary digit).
- Gray code also known as reflected binary code, because the first $(n/2)$ values compare with those of the last $(n/2)$ values, but in reverse order.
- Gray codes are used in the general sequence of hardware-generated binary numbers.
- The numbers cause ambiguities or errors when the transition from one number to its successive is done. This code simply solves this problem by changing only one bit when the transition is between numbers is done.

Gray codes

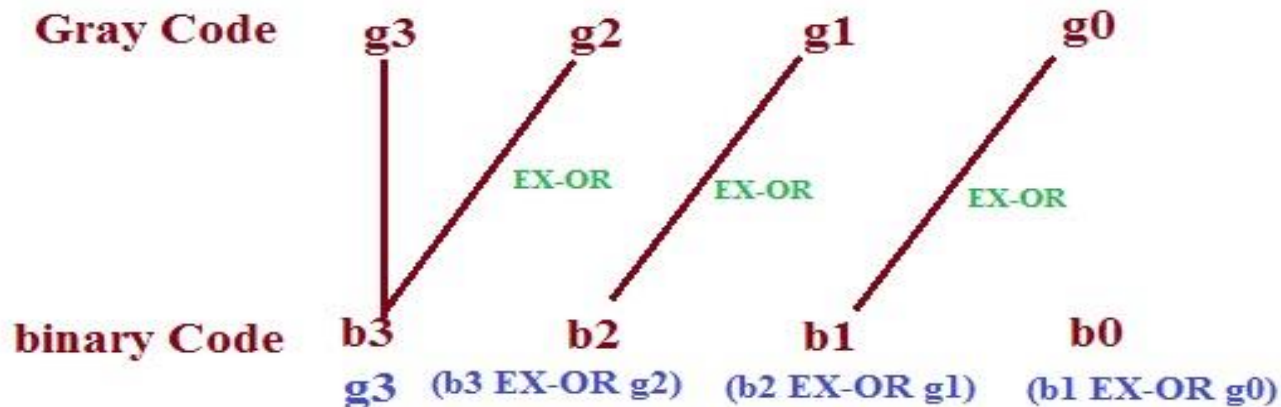
How to generate Gray code?

- The prefix and reflect method are recursively used to generate the Gray code of a number.
- For generating gray code:
 - Generate code for $n=1$: 0 and 1 code.
 - Take previous code in sequence: 0 and 1.
 - Add reversed codes in the following list: 0, 1, 1 and 0.
 - Now add prefix 0 for original previous code and prefix 1 for new generated code: 00, 01, 11, and 10



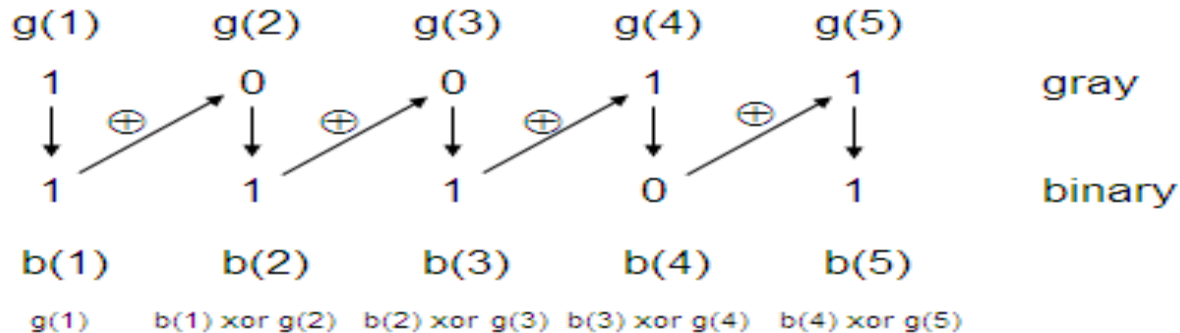
Gray to Binary Code:

- To change gray to binary code, take down the MSB digit of the gray code number, as the primary digit or the MSB of the gray code is similar to the binary digit.
- To get the next straight binary bit, it uses the XOR operation among the primary bit or MSB bit of binary to the next bit of the gray code.

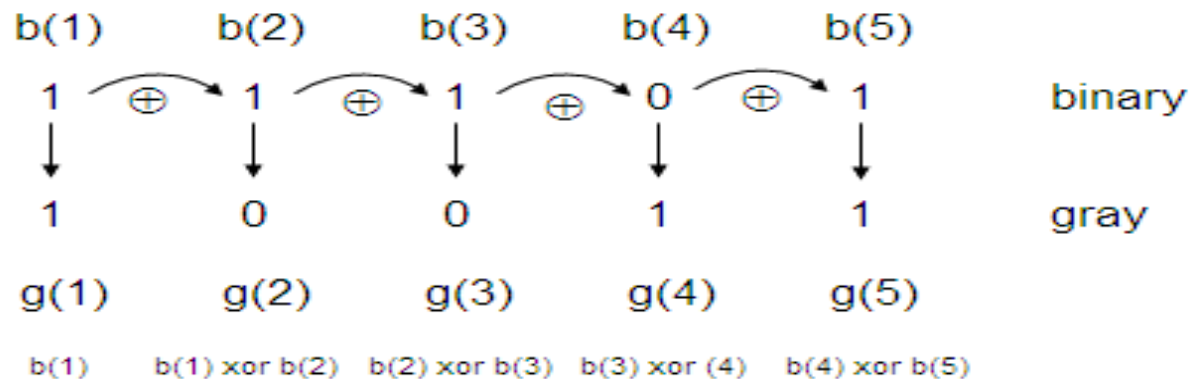


Gray codes

- Convert the 10011 gray code into binary code.



- Convert the 10011 binary code into gray code.



Alphanumeric Codes

ASCII Character Code:

- Many applications of digital computers require the handling of data not only of numbers, but also of letters.
- An alphanumeric character set is a set of elements that includes the 10 decimal digits, the 26 letters of the alphabet, and a number of special characters.
- The standard binary code for the alphanumeric characters is ASCII (American Standard Code for Information Interchange).
- It uses 7 bits to code 128 characters, but it can be considered as an 8-bit code with MSB = 0 always.
- The first 32 ASCII characters are non graphic commands ,these are used only for command purpose.
- The ASCII code set consist of 94 printable characters, SPACE and DELETE characters, and 32 control symbols.

Alphanumeric Codes

American Standard Code for Information Interchange (ASCII)

$b_4b_3b_2b_1$	$b_7b_6b_5$							
	000	001	010	011	100	101	110	111
0000	NUL	DLE	SP	0	@	P	'	p
0001	SOH	DC1	!	1	A	Q	a	q
0010	STX	DC2	"	2	B	R	b	r
0011	ETX	DC3	#	3	C	S	c	s
0100	EOT	DC4	\$	4	D	T	d	t
0101	ENQ	NAK	%	5	E	U	e	u
0110	ACK	SYN	&	6	F	V	f	v
0111	BEL	ETB	'	7	G	W	g	w
1000	BS	CAN	(8	H	X	h	x
1001	HT	EM)	9	I	Y	i	y
1010	LF	SUB	*	:	J	Z	j	z
1011	VT	ESC	+	;	K	[k	{
1100	FF	FS	,	<	L	\	l	
1101	CR	GS	-	=	M]	m	}
1110	SO	RS	.	>	N	^	n	~
1111	SI	US	/	?	O	_	o	DEL

Daily Quiz

- Classify binary codes.
- Differentiate between weighted and non weighted codes.
- Binary equivalent of gray code 101011 is:
 - a) 110010
 - b) 101010
 - c) 110110
 - d) 110011
- Classify binary codes.
- Differentiate between weighted and non weighted codes.
- Binary equivalent of gray code 101011 is:
 - a) 110010
 - b) 101010
 - c) 110110
 - d) 110011

- Excess-3 code of 739 is:
 - a) 1010 0110 1100
 - b) 1101 1001 1000
 - c) 1101 0110 1100
 - d) 1010 1001 1000
- What do you mean by binary codes?
- Classify binary codes.
- Excess-3 codes are also called:
 - a) Non-weighted codes.
 - b) Alphanumeric codes.
 - c) Sequential codes.
 - d) Reflective codes.

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Youtube/other Video Links:

- https://www.youtube.com/watch?v=VBM5XTvJSRQ&ab_channel=JAESCompanyJAESCompany
- https://www.youtube.com/watch?v=1A_NcXxdoCc&ab_channel=NesoAcademyNesoAcademyVerified

Old Questions

- Convert 1110110 binary code into gray code.
- What do you understand by self complementing codes?
- Construct the hamming code, if 4 bit data 1001 is transformed.
- A receiver receives the hamming code 1110101, What is the correct code for even parity is?
- Find binary equivalent of gray code 101011.

Recap

- The digital data is represented, stored and transmitted as group of binary bits, This group is also called as **binary code**.
- In weighted codes, each digit is assigned a specific weight according to its position .
- In weighted codes, positional weights are not assigned.
- A code is reflective when the code is **self complementing**.

Hamming Codes

Topic Objective	Mapping with CO
To understand hamming codes.	CO1

Hamming Codes

Perquisite:

- Knowledge of number system.

Hamming Codes

Hamming Code: When data is transferred from one location to another location, there is always a possibility that an error may occur. Error result in changes to the content of data transferred.

- Hamming code is an error-detection method that can detect some errors, but it is only capable of single error correction.
- Hamming Codes are named after Richard Hamming, the American mathematician.
- It has fix format, there are reserved bit for the parity at the position 2^0 , 2^1 , 2^2 , 2^3 and so on.
- The **minimum value of 'k'** for which the following relation is correct valid is nothing but the required number of parity bits.

$$2^k \geq n+k+1 \quad \text{Where,}$$

‘n’ is the number of bits in the binary code information

‘k’ is the number of parity bits

Therefore, the number of bits in the Hamming code is equal to $n + k$.

Hamming Codes

- Let $n = 4$, bits present is data

$$2^k \geq 4 + k + 1$$

k must be equal to 3.

- Let $n = 8$, bits present is data

$$2^k \geq 8 + k + 1$$

$$k = 4$$

- Table for different data bits and check bits:**

Number of information bits	Number of parity bits
2 to 4	3
5 to 11	4
12 to 26	5
27 to 57	6
58 to 120	7

Parity Bits	Bits to be checked
P1	1,3,5,7,9,11,13,15,.....
P2	2,3,6,7,10,11,14,15,....
P4	4,5,6,7,12,13,14,15,.....
P8	8,9,10,11,12,13,14,15,...

Hamming Codes

Format of Hamming Code for 4 bit:

If D3, D5, D6 and D7 are data bits, (for $n=4$, $k=3$)

7	6	5	4	3	2	1
D7	D6	D5	P4	D3	P2	P1

$P1 = (1, 3, 5, 7)$ should have even number of 1's.

$P2 = (2, 3, 6, 7)$ should have even number of 1's.

$P4 = (4, 5, 6, 7)$ should have even number of 1's.

Format of Hamming Code for 8 bit:

If D3, D5, D6, D7, D9, D10, D11, D12 are data bits, (for $n=8$, $k=4$)

12	11	10	9	8	7	6	5	4	3	2	1
D12	D11	D10	D9	P8	D7	D6	D5	P4	D3	P2	P1

$P1 = (1, 3, 5, 7, 9, 11)$ should have even number of 1's.

$P2 = (2, 3, 6, 7, 10, 11)$ should have even number of 1's.

$P4 = (4, 5, 6, 7, 12)$ should have even number of 1's.

$P8 = (8, 9, 10, 11, 12)$ should have even number of 1's.

Hamming Codes

1. Construct the hamming code, if 4 bit data 1100 is transformed.

Sol: for $n = 4$, $k = 3$

7	6	5	4	3	2	1
1	1	0	P4	0	P2	P1

$P1 = (1, 3, 5, 7)$ ($p1, 0, 0, 1$) should have even number of 1's. $P1 = 1$

$P2 = (2, 3, 6, 7)$ ($P2, 0, 1, 1$) should have even number of 1's. $P2 = 0$

$P4 = (4, 5, 6, 7)$ ($P4, 0, 1, 1$) should have even number of 1's. $P4 = 0$

Hamming Code : **1100001**

2. Construct the hamming code, if 8 bit data 10101011 is transformed.

Sol: for $n = 8$, $k = 4$

12	11	10	9	8	7	6	5	4	3	2	1
1	0	1	0	P8	1	0	1	P4	1	P2	P1

$P1 = (1, 3, 5, 7, 9, 11)$ should have even number of 1's. $P1 = 1$

$P2 = (2, 3, 6, 7, 10, 11)$ should have even number of 1's. $P2 = 1$

$P4 = (4, 5, 6, 7, 12)$ should have even number of 1's. $P4 = 1$

$P8 = (8, 9, 10, 11, 12)$ should have even number of 1's. $P8 = 0$

Hamming Code: **101001011111**

Hamming Codes

- **Correction for the checking of error in 4 bit data with even parity received at the receiver:**

7	6	5	4	3	2	1
D7	D6	D4	P4	D3	P2	P1

- $P1 = (1, 3, 5, 7)$ should have even number of 1's.
If P1 is even means no error so $C1 = 0$
If P1 is not even means error so $C1 = 1$
- $P2 = (2, 3, 6, 7)$ should have even number of 1's.
If P2 is even means no error so $C2 = 0$
If P2 is not even means error so $C2 = 1$
- $P4 = (4, 5, 6, 7)$ should have even number of 1's.
If P4 is even means no error so $C4 = 0$
If P4 is not even means error so $C4 = 1$

So, error bit $C = C_4C_2C_1$

Hamming Codes

- **Correction for the checking of error in 8 bit data with even parity received at the receiver:**

12	11	10	9	8	7	6	5	4	3	2	1
D12	D11	D10	D9	P8	D7	D6	D5	P4	D3	P2	P1

- $P1 = (1, 3, 5, 7, 9, 11)$ should have even number of 1's.
If $P1$ is even means no error so $C1 = 0$
If $P1$ is not even means error so $C1 = 1$
- $P2 = (2, 3, 6, 7, 10, 11)$ should have even number of 1's.
If $P2$ is even means no error so $C2 = 0$
If $P2$ is not even means error so $C2 = 1$
- $P4 = (4, 5, 6, 7, 12)$ should have even number of 1's.
If $P4$ is even means no error so $C4 = 0$
If $P4$ is not even means error so $C4 = 1$
- $P8 = (8, 9, 10, 11, 12)$ should have even number of 1's.
If $P8$ is even means no error so $C8 = 0$
If $P8$ is not even means error so $C8 = 1$
So, error bit $C = C_8C_4C_2C_1$

Hamming Codes

1. Receiver receives the hamming code 1101001, check whether the data received is correct if not, check the error and correct it.

Sol:

7	6	5	4	3	2	1
1	1	0	1	0	0	1

- $P_1 = (1, 3, 5, 7)$ should have even number of 1's.
 P_1 is already even, so no error, $C_1 = 0$
- $P_2 = (2, 3, 6, 7)$ should have even number of 1's.
 P_2 is already even, so no error, $C_2 = 0$
- $P_4 = (4, 5, 6, 7)$ should have even number of 1's.
 P_4 is odd, so there is error, $C_4 = 1$
 So, error bit $C = C_4C_2C_1$, $C = 100 = 4$
 $C = 4$ so the error in 4th bit. The correct code is 1100001

Hamming Codes

2. Receiver receives the hamming code 1111001, check whether the data received is correct if not, check the error and correct it.

Sol:

7	6	5	4	3	2	1
1	1	1	1	0	0	1

- $P_1 = (1, 3, 5, 7)$ should have even number of 1's.
 P_1 is not even, so error, $C_1 = 1$
- $P_2 = (2, 3, 6, 7)$ should have even number of 1's.
 P_2 is even, so no error, $C_2 = 0$
- $P_4 = (4, 5, 6, 7)$ should have even number of 1's.
 P_4 is already even, so no error, $C_4 = 0$
 So, error bit $C = C_4C_2C_1$, $C = 001 = 1$
 $C = 1$ so error is there in the 1st bit. So correct data is 1111000.

Hamming Codes

3. Receiver receives the hamming code 110010001110, check whether the data received is correct if not, check the error and correct it.

Sol:

12	11	10	9	8	7	6	5	4	3	2	1
1	1	0	0	1	0	0	0	1	1	1	0

- $P_1 = (1, 3, 5, 7, 9, 11)$ should have even number of 1's.
 P_1 is even, so no error, $C_1 = 0$
- $P_2 = (2, 3, 6, 7, 10, 11)$ should have even number of 1's.
 P_2 is not even, so error, $C_2 = 1$
- $P_4 = (4, 5, 6, 7, 12)$ should have even number of 1's.
 P_4 is not even, so error, $C_4 = 0$
- $P_8 = (8, 9, 10, 11, 12)$ should have even number of 1's.
 P_8 is not even, so error, $C_8 = 1$

So, error bit $C = C_8 C_4 C_2 C_1$ $C = 1010 = 10$

$C = 10$ so error is there in the 10th bit correct code is: 11**1**010001110

Hamming Codes

4. Receiver receives the hamming code 1101001, check whether the data received is correct if not, check the error and correct it. Use odd parity.

Sol: Odd Parity:

In the case of odd parity, for a given set of bits, the number of 1's are counted. If that count is even, the parity bit value is set to 1, making the total count of occurrences of 1's an odd number. If the total number of 1's in a given set of bits is already odd, the parity bit's value is 0.

7	6	5	4	3	2	1
1	1	0	1	0	0	1

- $P_1 = (1, 3, 5, 7)$ should have odd number of 1's.
 P_1 is even, so error, $C_1 = 1$
- $P_2 = (2, 3, 6, 7)$ should have odd number of 1's.
 P_2 is even, so error, $C_2 = 1$
- $P_4 = (4, 5, 6, 7)$ should have odd number of 1's.
 P_4 is odd, so there is no error, $C_4 = 0$

So, error bit $C = C_4C_2C_1$, $C = 011 = 3$

$C = 3$ so the error in 3rd bit. The correct code is 1101101

Daily Quiz

- Classify binary codes.
- Differentiate between weighted and non weighted codes.
- Binary equivalent of gray code 101011 is:
 - a) 110010
 - b) 101010
 - c) 110110
 - d) 110011
- Construct the hamming code, if 4 bit data 1100 is transformed.(Use odd parity)
 - a) 1101011
 - b) 1101010
 - c) 1101110
 - d) 1011010

- Excess-3 code of 739 is:
 - a) 1010 0110 1100
 - b) 1101 1001 1000
 - c) 1101 0110 1100
 - d) 1010 1001 1000
- What do you mean by binary codes?
- Classify binary codes.
- Excess-3 codes are also called:
 - a) Non-weighted codes.
 - b) Alphanumeric codes.
 - c) Sequential codes.
 - d) Reflective codes.

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Youtube/other Video Links:

- <https://nptel.ac.in/courses/117/106/117106086/>
- <https://www.youtube.com/watch?v=kAezzdEGJ8A>

Old Questions

- Construct the hamming code, if 4 bit data 1001 is transformed.
- A receiver receives the hamming code 1110101, What is the correct code for even parity is?
- Generate the hamming code for the word 11011. Assume that a single error occurs while storing the generated hamming code. Explain how this single error is detected.

Weekly Assignment

- Perform BCD Addition of 999 and 989.
- For 989 and 674, Find BCD Subtraction using 9's complement and 10's complement method.
- Explain hamming code. A receiver receives the hamming code 10101101, check whether the data received is correct if not, check the error and correct it.
- Construct the hamming code, if 4 bit data 1001 is transformed.
- Write a short note on binary codes.

Recap

- Hamming code is an error-detection method that can detect some errors, but it is only capable of single error correction.
- Hamming Codes are named after Richard Hamming, the American mathematician.
- It has fix format, there are reserved bit for the parity at the position 2^0 , 2^1 , 2^2 , 2^3 and so on.
- The **minimum value of 'k'** for which the following relation is correct valid is nothing but the required number of parity bits.

$$2^k \geq n+k+1 \quad \text{Where,}$$

‘n’ is the number of bits in the binary code information

‘k’ is the number of parity bits

Therefore, the number of bits in the Hamming code is equal to $n + k$.

Boolean algebra & Logic Gates

Topic Objective	Mapping with CO
To understand basic rules of Boolean algebra.	CO1
To understand SOP and POS form of Boolean function	CO1

Boolean algebra & Logic Gates

Prerequisite:

- Knowledge of basic logic gates.

Boolean algebra & Logic Gates

- Boolean Algebra is used to analyze and simplify the digital (logic) circuits. It uses only the binary numbers i.e. 0 and 1. It is also called as **Binary Algebra** or **logical Algebra**. Boolean algebra was invented by **George Boole** in 1854.
- Binary logic consists of binary variables and a set of logical operations. The variables are designated by letters of the alphabet, such as A, B, C, x, y, z, etc, with each variable having two and only two distinct possible values: 1 and 0, There are three basic logical operations: AND, OR, and NOT.

Boolean algebra & Logic Gates

Following are the important rules used in Boolean algebra.

- Variable used can have only two values. Binary 1 for HIGH and Binary 0 for LOW.
- Complement of a variable is represented by an overbar (-).
- ORing of the variables is represented by a plus (+) sign between them. For example ORing of A, B, C is represented as $A + B + C$.
- Logical ANDing of the two or more variable is represented by writing a dot between them such as A.B.C. Sometime the dot may be omitted like ABC.

Boolean algebra & Logic Gates

There are six types of Boolean Laws.

1) Commutative law:

- Any binary operation which satisfies the following expression is referred to as commutative operation.

$$(i) A.B = B.A \quad (ii) A + B = B + A$$

Commutative law states that changing the sequence of the variables does not have any effect on the output of a logic circuit.

2) Associative law:

This law states that the order in which the logic operations are performed is irrelevant as their effect is the same.

$$(i) (A.B).C = A.(B.C) \quad (ii) (A + B) + C = A + (B + C)$$

Boolean algebra & Logic Gates

- **Distributive law:**

Distributive law states the following condition.

$$A.(B + C) = A.B + A.C$$

AND law

- These laws use the AND operation. Therefore they are called as **AND** laws.

$$(i) A.0 = 0$$

$$(ii) A.1 = A$$

$$(iii) A.A = A$$

$$(iv) A.\overline{A} = 0$$

OR law

- These laws use the OR operation. Therefore they are called as **OR** laws

$$(i) A + 0 = A$$

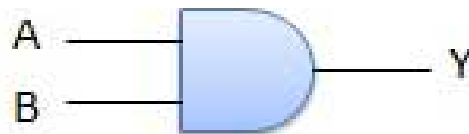
$$(ii) A + 1 = 1$$

$$(iii) A + A = A$$

$$(iv) A + \overline{A} = 1$$

Boolean algebra & Logic Gates

- **AND Gate**
- A circuit which performs an AND operation is shown in figure. It has n input ($n \geq 2$) and one output.
- **Logic diagram**



- **Truth Table**

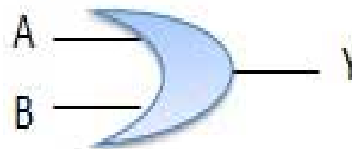
Inputs		Output
A	B	AB
0	0	0
0	1	0
1	0	0
1	1	1

Boolean algebra & Logic Gates

OR Gate

- A circuit which performs an OR operation is shown in figure. It has n input ($n \geq 2$) and one output.

Logic diagram



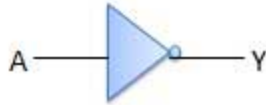
Truth Table

Inputs		Output
A	B	$A + B$
0	0	0
0	1	1
1	0	1
1	1	1

Boolean algebra & Logic Gates

- **NOT Gate**
- NOT gate is also known as **Inverter**. It has one input A and one output Y.

- Symbol



- Truth Table

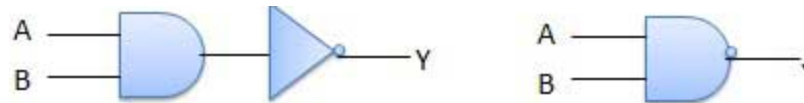
Inputs	Output
A	B
0	1
1	0

Boolean algebra & Logic Gates

NAND Gate

A NOT-AND operation is known as NAND operation. It has n input ($n \geq 2$) and one output.

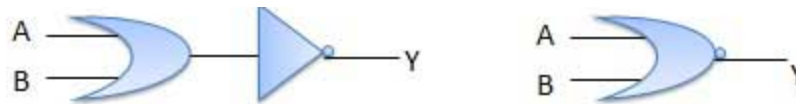
- Logic diagram
- Truth Table



Inputs		Output
A	B	\overline{AB}
0	0	1
0	1	1
1	0	1
1	1	0

Boolean algebra & Logic Gates

- **NOR Gate**
- A NOT-OR operation is known as NOR operation. It has n input ($n \geq 2$) and one output.
- Logic diagram



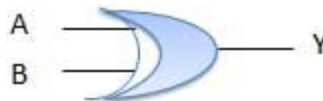
Truth Table

Inputs		Output
A	B	$\overline{A+B}$
0	0	1
0	1	0
1	0	0
1	1	0

Boolean algebra & Logic Gates

- **XOR Gate:**
- XOR or Ex-OR gate is a special type of gate. It can be used in the half adder, full adder and subtractor. The exclusive-OR gate is abbreviated as EX-OR gate or sometime as X-OR gate. It has n input ($n \geq 2$) and one output.

- **Logic diagram**



- **Truth Table**

Inputs		Output
A	B	$A \oplus B$
0	0	0
0	1	1
1	0	1
1	1	0

Boolean algebra & Logic Gates

- **XNOR Gate**
- XNOR gate is a special type of gate. It can be used in the half adder, full adder and subtractor. The exclusive-NOR gate is abbreviated as EX-NOR gate or sometime as X-NOR gate. It has n input ($n \geq 2$) and one output.

- **Logic diagram**

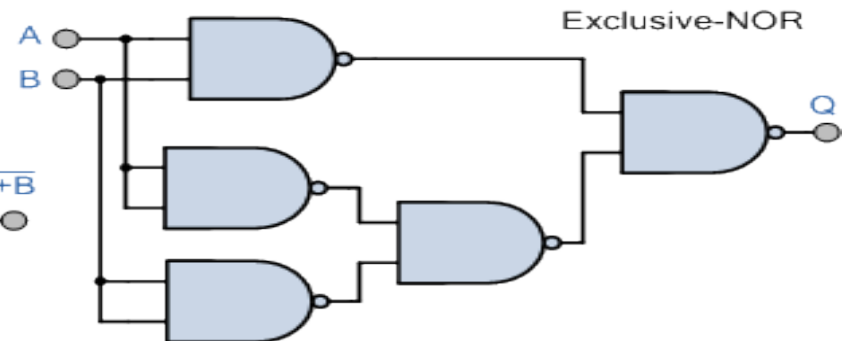
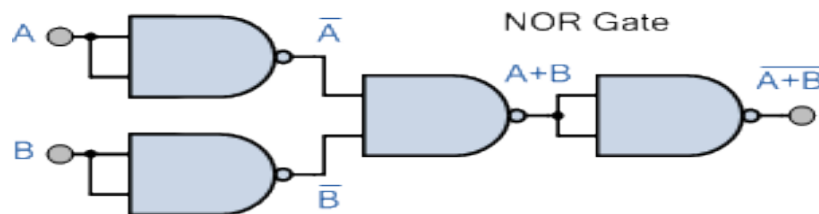
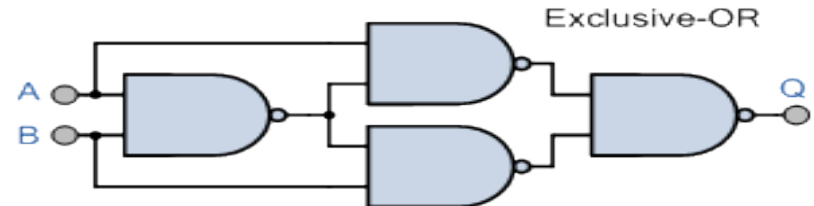
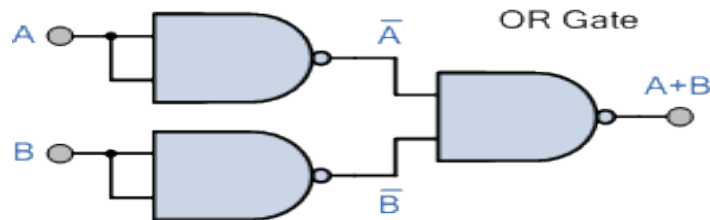
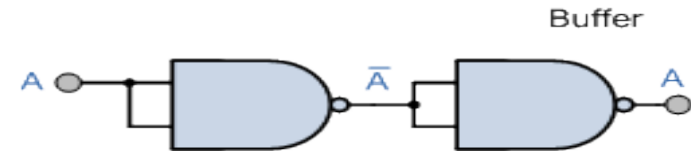
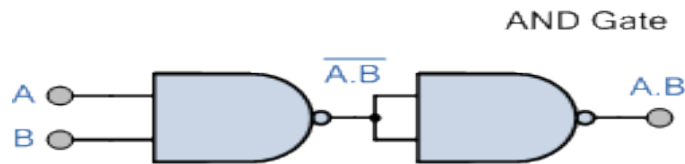
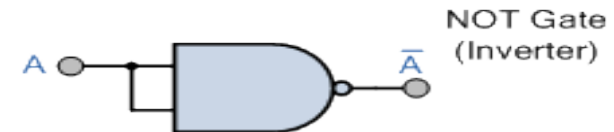
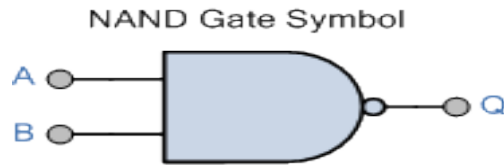


- **Truth Table**

Inputs		Output
A	B	$A \oplus B$
0	0	1
0	1	0
1	0	0
1	1	1

UNIVERSAL GATES

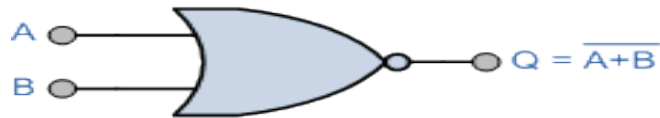
NAND gate as universal gate:



UNIVERSAL GATES

NOR gate as universal gate:

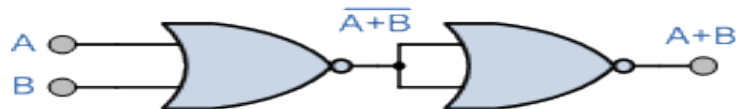
NOR Gate Symbol



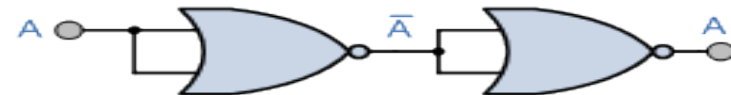
NOT Gate (Inverter)



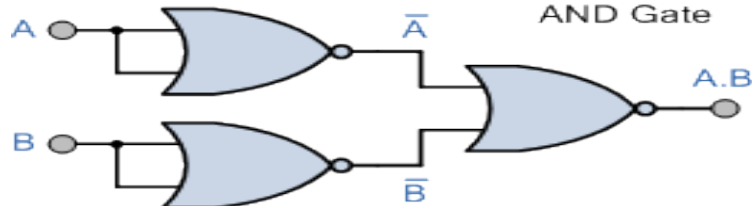
OR Gate



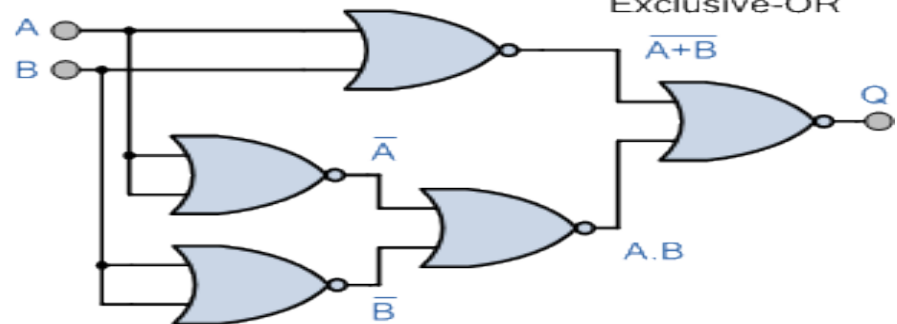
Buffer



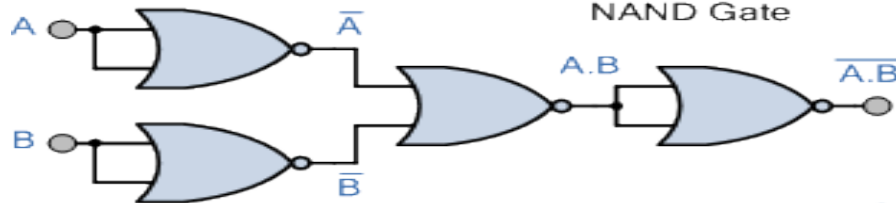
AND Gate



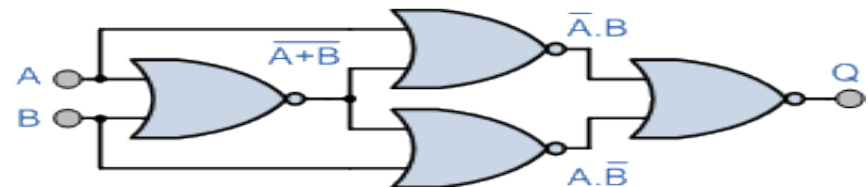
Exclusive-OR



NAND Gate



Exclusive-NOR



Daily Quiz

Q1. The NOR gate output will be high if the two inputs are _____

- a) **00**
- b) 01
- c) 10
- d) 11

Q2. A universal logic gate is one which can be used to generate any logic function. Which of the following is a universal logic gate?

- a) OR
- b) AND
- c) XOR
- d) **NAND**

Q3. Which of following are known as universal gates?

- a) **NAND & NOR**
- b) AND & OR
- c) XOR & OR
- d) EX-NOR & XOR

FACULTY VIDEO LINKS, YOUTUBE & NPTEL VIDEO LINKS AND ONLINE COURSES DETAILS

Youtube/other Video Links:

- https://www.youtube.com/watch?v=VBM5XTvJSRQ&ab_channel=JAESCompanyJAESCompany
- https://www.youtube.com/watch?v=1A_NcXxdoCc&ab_channel=NesoAcademyNesoAcademyVerified

Old Questions

- What do you understand by universal gates?
- Realize a Ex-or gate with the help of NAND gate.
- Realize a OR gate using NAND gates.
- How an Ex-or gate can work as an inverter?
- What do understand by the term SOP and POS?

Recap

- Boolean Algebra is used to analyze and simplify the digital (logic) circuits.
- It uses only the binary numbers i.e. 0 and 1.
- It is also called as **Binary Algebra** or **logical Algebra**.
- Binary logic consists of binary variables and a set of logical operations.
- Here variable having two and only two distinct possible values: 1 and 0.
- There are three basic logical operations: AND, OR, and NOT.

KARNAUGHMAP (K-MAP) REPRESENTATION

Topic Objective	Mapping with CO
To understand rules of K-map simplification	CO1
To understand the logic minimization using K-map.	CO1

Prerequisite:

- Knowledge of Boolean algebra.
- Knowledge of SOP and POS forms.

KARNAUGHMAP (K-MAP) REPRESENTATION

- The complexity of the digital logic gates that implement a Boolean function is directly related to the complexity of the algebraic expression from which the function is implemented.
- The map method provides a simple straightforward procedure for minimizing Boolean functions.
- The map method, first proposed by Veitch and modified by Karnaugh, is also known as the "Veitch diagram" or the "Karnaugh map."
- The map is a diagram made up of squares. Each square represents one minterm.
- Since any Boolean function can be expressed as a sum of minterms, it follows that a Boolean function is recognized graphically in the map from the area enclosed by those squares whose min terms are included in the function.

KARNAUGHMAP (K-MAP) REPRESENTATION

- Two-Variable Karnaugh Map:

		\overline{B} B	
		0	1
A	0	$\overline{A}\overline{B}$ m_0	$\overline{A}B$ m_1
	1	$A\overline{B}$ m_2	AB m_3

$$Y = F(A,B) = \sum m(2,3)$$

A	B	Y
0	0	0
0	1	0
1	0	1
1	1	1

	\overline{B}	B
\overline{A}	0	0
A	1	1

$$Y = A$$

$$Y = F(A,B) = \sum m(1,2,3)$$

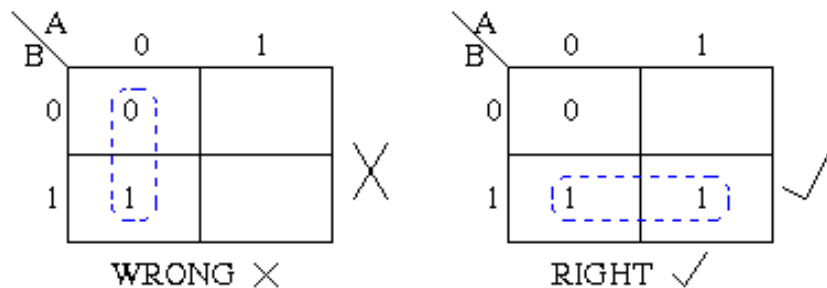
A	B	Y
0	0	0
0	1	1
1	0	1
1	1	1

	\overline{B}	B
\overline{A}	0	1
A	1	1

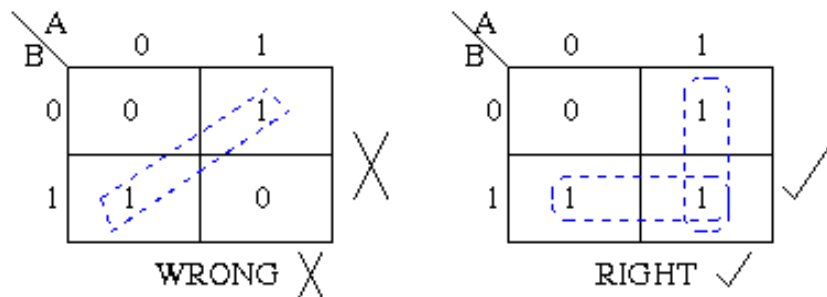
$$Y = A + B$$

KMAP Simplification Rules:

- Groups may not include any cell containing a zero.



- Groups may be horizontal or vertical, but not diagonal.



KMAP

- Groups must contain 1, 2, 4, 8, or in general 2^n cells. That is if $n = 1$, a group will contain two 1's since $2^1 = 2$.
- If $n = 2$, a group will contain four 1's since $2^2 = 4$.

A \ B	0	1
0	1	1
1	0	0

Group of 2

RIGHT ✓

AB \ C	00	01	11	10
0	0	1	1	1
1	0	0	0	0

Group of 3

WRONG ✗

A \ B	0	1
0	1	1
1	1	1

Group of 4

RIGHT ✓

AB \ C	00	01	11	10
0	1	1	1	1
1	0	0	0	1

Group of 5

WRONG ✗

- Each group should be as large as possible.

	AB	00	01	11	10	
C						
0		1	1	1	1	
1		0	0	1	1	✓

RIGHT ✓

	AB	00	01	11	10	
C						
0		1	1	1	1	
1		0	0	1	1	×

WRONG ✗

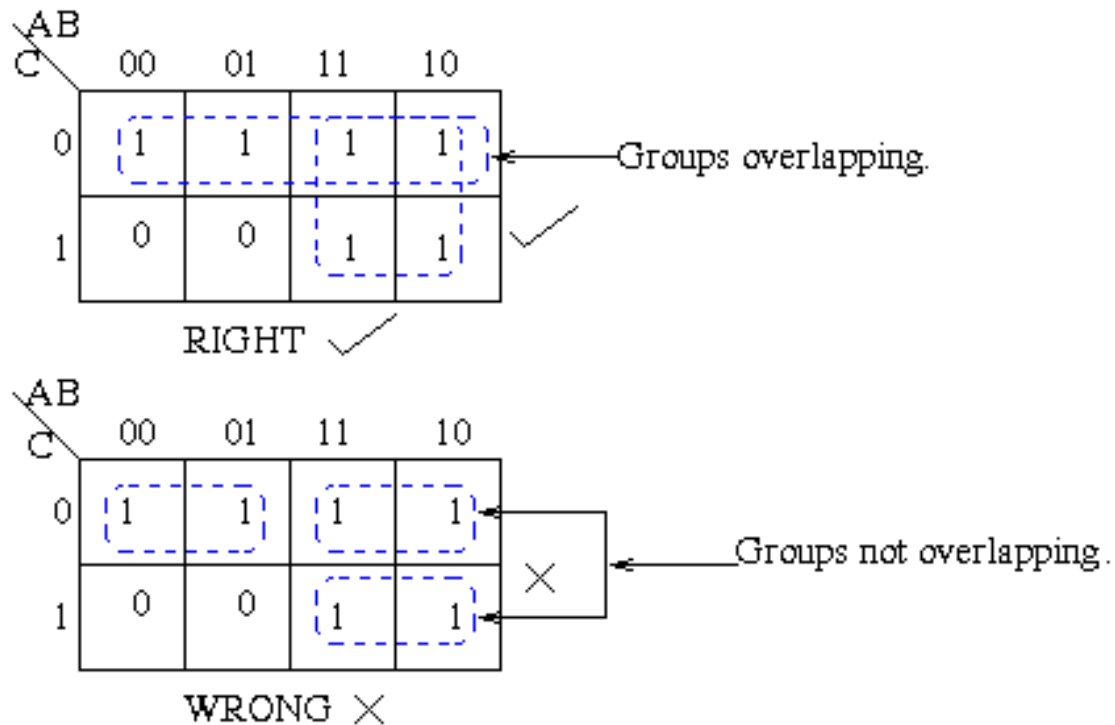
(Note that no Boolean laws broken, but not sufficiently minimal)

- Each cell containing a *one* must be in at least one group.

	AB	00	01	11	10	
C						
0		0	0	1	1	Group I
1		0	0	0	1	Group II

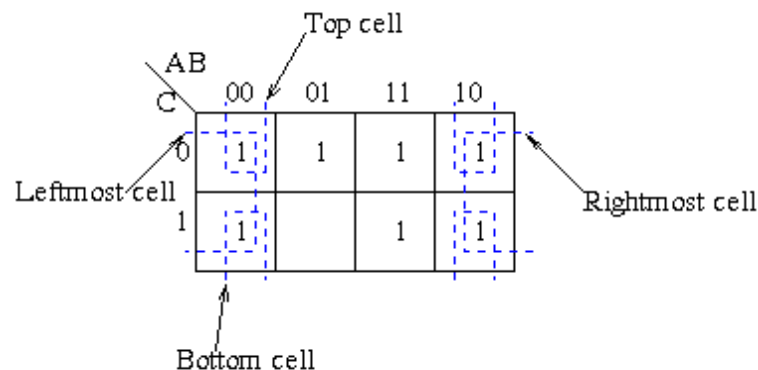
1 present in at least one group.

- Groups may overlap.

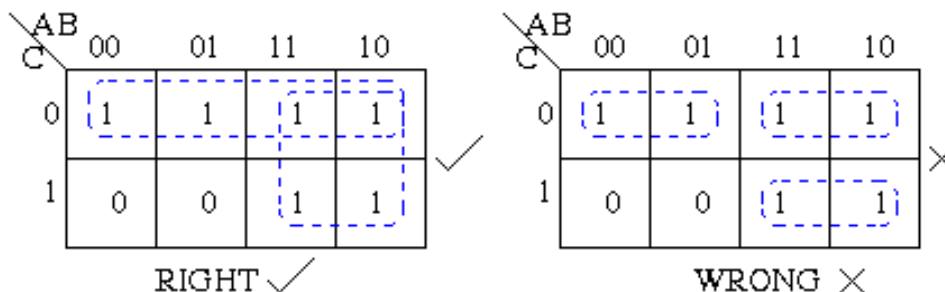


KMAP

- Groups may wrap around the table. The leftmost cell in a row may be grouped with the rightmost cell and the top cell in a column may be grouped with the bottom cell.



- There should be as few groups as possible, as long as this does not contradict any of the previous rules.



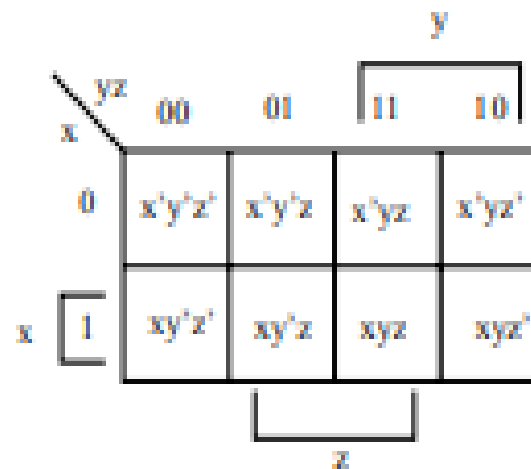
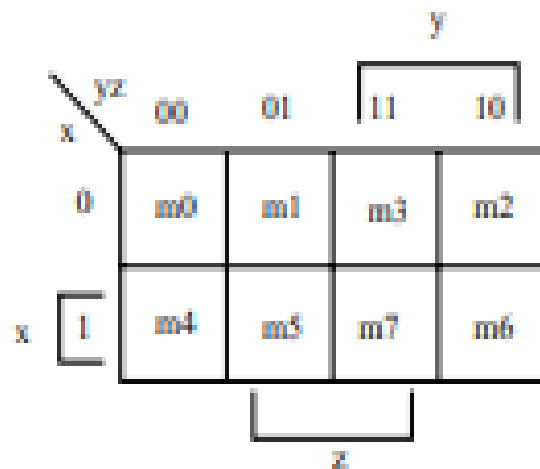
Summary:

1. No zeros allowed.
2. No diagonals.
3. Only power of 2 number of cells in each group.
4. Groups should be as large as possible.
5. Every one must be in at least one group.
6. Overlapping allowed.
7. Wrap around allowed.
8. Fewest number of groups possible.

THREE VARIABLE KARNAUGH MAP

- Three-Variable Karnaugh Map:**

3 variables, $F = f(x, y, z)$



x	y	z	$F(x,y,z)$
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	1
1	0	1	1
1	1	0	1
1	1	1	1

THREE VARIABLE KARNAUGH MAP

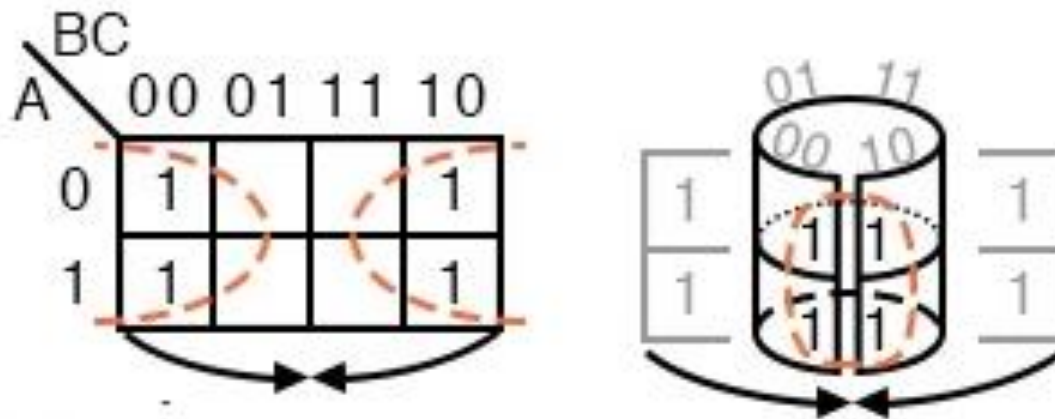
Eg: $F(A,B,C) = \sum(1,2,3,6)$ minimize the given using K-MAP.

A \ BC	BC			
	00	01	11	10
0		1	1	1
1				1

Ans: $F = A'C + BC'$

THREE VARIABLE KARNAUGH MAP

Eg: $F = A'B'C' + AB'C' + A'BC' + ABC'$ minimize the given using K-MAP.



Ans: $F = C'$

THREE VARIABLE KARNAUGH MAP

Eg: $F(x, y, z) = \sum(0, 2, 4, 5, 6)$ minimize the given using K-MAP.

		y			
		z'			
		00	01	11	10
x	0	1			1
x	1	1	1		1
		z			

Ans: $F = z' + x.y'$

THREE VARIABLE KARNAUGH MAP

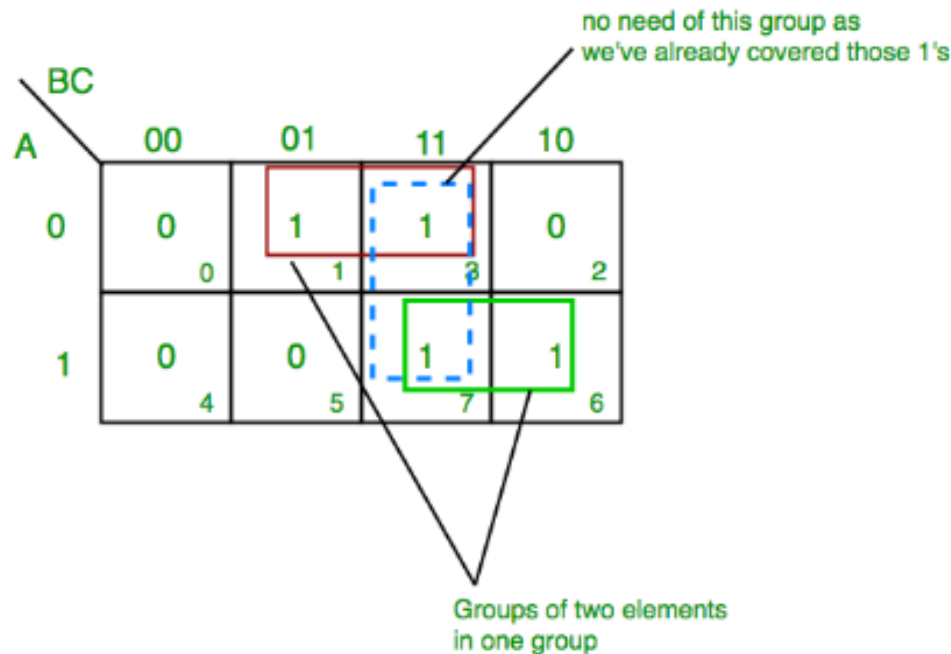
Eg: $F = \sum A, B, C(3, 5, 6, 7)$ minimize the given using K-MAP.

		BC			
		00	01	11	10
A	0			1	
	1		1	1	1

Ans: $F = BC + AC + AB$

THREE VARIABLE KARNAUGH MAP

Eg: $F(A,B,C) = \sum(1,3,6,7)$ minimize the given using K-MAP.



Ans: $F = A'C + AB$

THREE VARIABLE KARNAUGH MAP

Eg: $F(A,B,C) = A'B'C' + A'B + ABC' + AC$ minimize the given using K-MAP.

AB \ C	00	01	11	10
0	1	1	1	
1		1	1	1

Ans: $F = B + AC + A'C'$

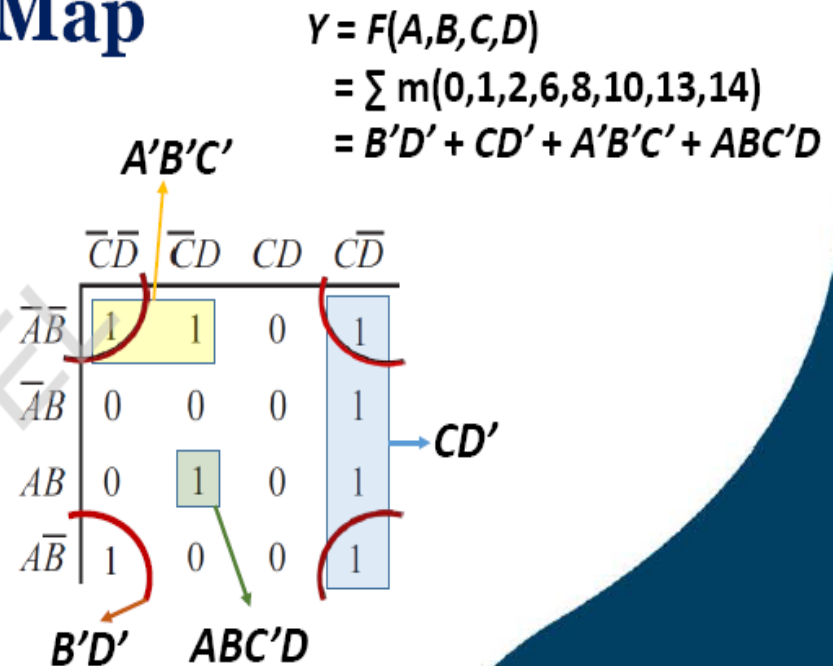
FOUR-VARIABLE KARNAUGH MAP

Four-Variable Karnaugh Map

		$\overline{C}\overline{D}$	$\overline{C}D$	CD	$C\overline{D}$
		00	01	11	10
$\overline{A}\overline{B}$	00	0	1	3	2
$\overline{A}B$	01	4	5	7	6
AB	11	12	13	15	14
$A\overline{B}$	10	8	9	11	10

$F(A,B,C,D)$
minterm numbers

A	B	C	D	Y
0	0	0	0	1
0	0	0	1	1
0	0	1	0	1
0	0	1	1	0
0	1	0	0	0
0	1	0	1	0
0	1	1	1	0
1	0	0	0	1
1	0	0	1	0
1	0	1	0	1
1	0	1	1	0
1	1	0	0	0
1	1	0	1	1
1	1	1	0	1
1	1	1	1	0

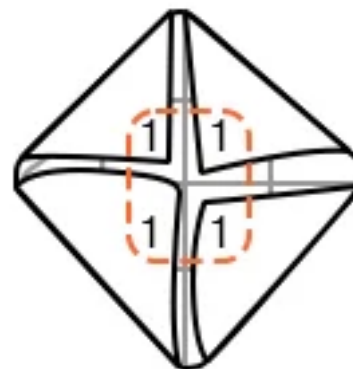


FOUR-VARIABLE KARNAUGH MAP

Eg: $F = \sum m(0, 2, 8, 10)$ minimize the given using K-MAP in SOP form.

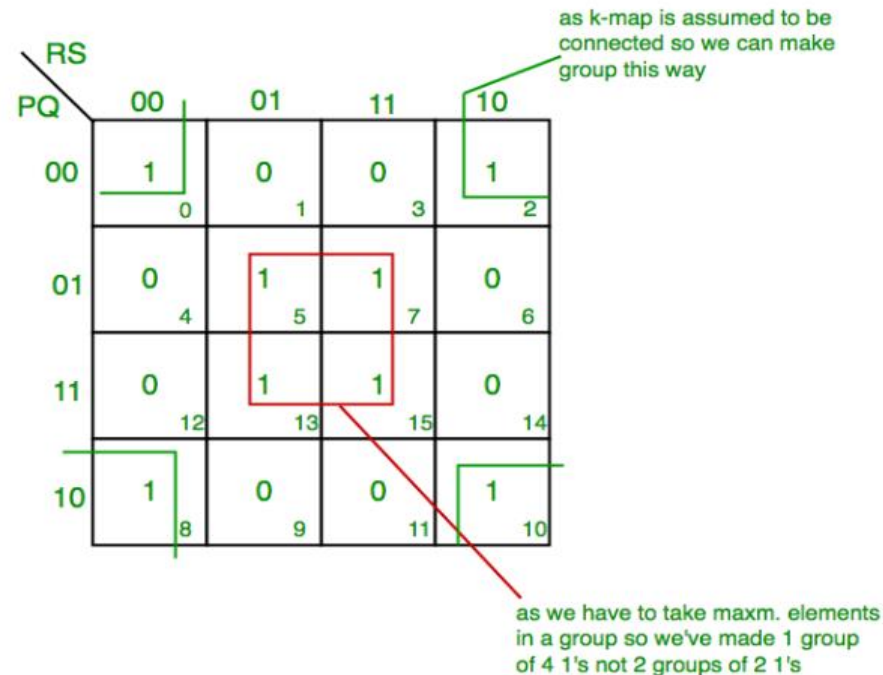
		CD			
		00	01	11	10
A \ B	00	1			1
	01				
	11				
	10	1			1

$$\text{Out} = \overline{B}\overline{D}$$



FOUR-VARIABLE KARNAUGH MAP

- $F(P,Q,R,S) = \sum(0,2,5,7,8,10,13,15)$ minimize the given using K-MAP.

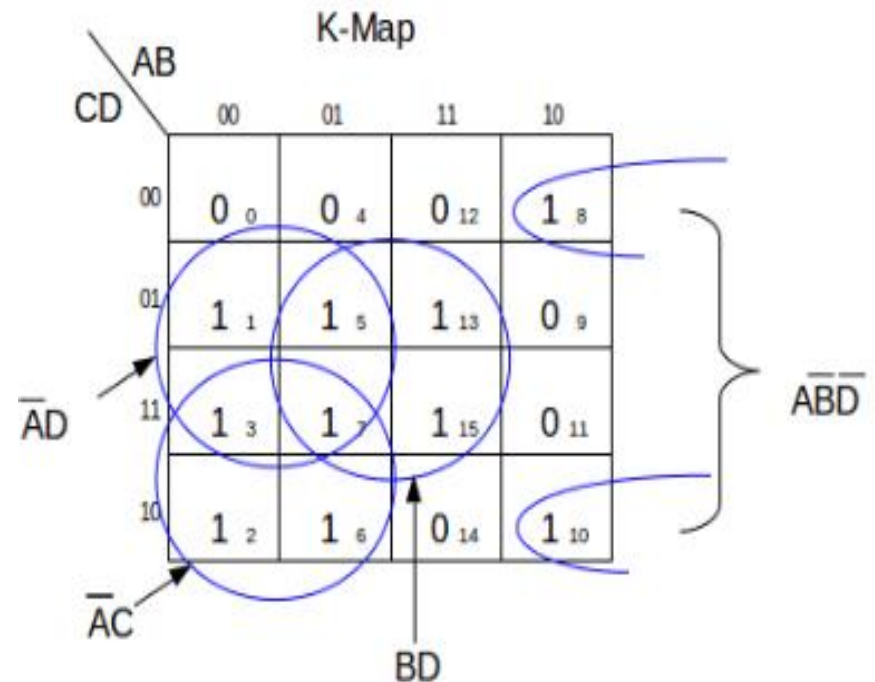


$$F = (QS + Q'S')$$

FOUR-VARIABLE KARNAUGH MAP

- $F(A,B,C,D) = \sum(1,2,3,5,6,7,8,10,13,15)$ minimize the given using K-MAP.

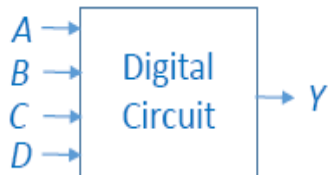
ab \ cd	00	01	11	10
00	0	4	12	8
01	1	4	13	9
11	3	6	15	11
10	2	7	14	10



$$F = \bar{A}\bar{D} + \bar{A}C + BD + \bar{A}\bar{B}\bar{D}$$

DON'T CARE IN KARNAUGH MAP

Design: Y is H when
BCD (Binary Coded
Decimal) input is odd.



$$Y = F(A, B, C, D)$$

$$= \sum m(1, 3, 5, 7, 9) +$$

$$d(10, 11, 12, 13, 14, 15)$$

A	B	C	D	Y
0	0	0	0	0
0	0	0	1	1
0	0	1	0	0
0	0	1	1	1
0	1	0	0	0
0	1	0	1	1
0	1	1	0	0
0	1	1	1	1
1	0	0	0	0
1	0	0	1	1
1	0	1	0	X
1	0	1	1	X
1	1	0	0	X
1	1	0	1	X
1	1	1	0	X
1	1	1	1	X

Not considering X

	$\overline{C}\overline{D}$	$\overline{C}D$	CD	$C\overline{D}$
$\overline{A}\overline{B}$	0	1	1	0
$\overline{A}B$	0	1	1	0
AB	X	X	X	X
$A\overline{B}$	0	1	X	X

$$Y = A'D + B'C'D$$

- If not considered, X = 0.
- If considered, X = 1.
- Consideration wherever helps.

Considering X

	$\overline{C}\overline{D}$	$\overline{C}D$	CD	$C\overline{D}$
$\overline{A}\overline{B}$	0	1	1	0
$\overline{A}B$	0	1	1	0
AB	X	X	X	X
$A\overline{B}$	0	1	X	X

$$Y = D$$

DON'T CARE IN KARNAUGH MAP

EXAMPLES OF DON'T CARE CONDITIONS (1/2)

$$F = \sum m(1, 3, 7) + \sum d(0, 5)$$

Circle the x's that help get bigger groups of 1's (or 0's if POS).

Don't circle the x's that don't help.

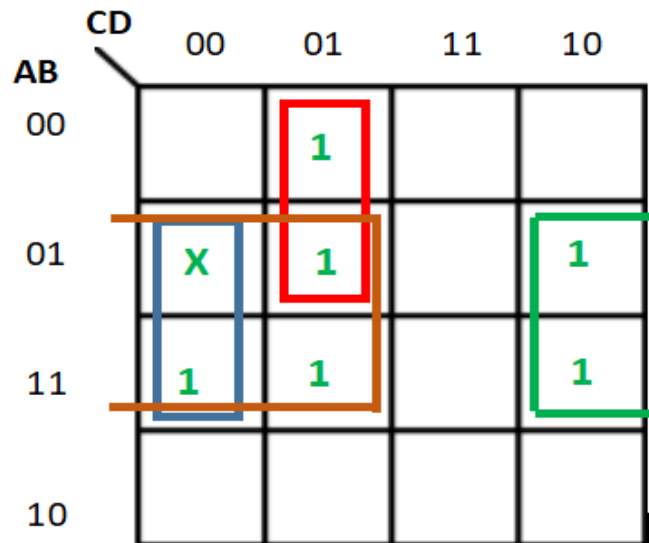
		BC			
		00	01	11	10
A	0	0 X	1 1	3 1	2
	1	4	5 X	7 1	6

Reduced form : $F = C$

DON'T CARE IN KARNAUGH MAP

1. Minimise the following function in SOP minimal form using K-Maps: $f(A,B,C,D) = m(1, 5, 6, 12, 13, 14) + d(4)$

- From **green & blue** group we find terms : BD'
- From **red** group we find terms : $A'C'D$
- From **brown** group we find terms : BC'
- Therefore, SOP minimal is, $f = BC' + BD' + A'C'D$



DON'T CARE IN KARNAUGH MAP

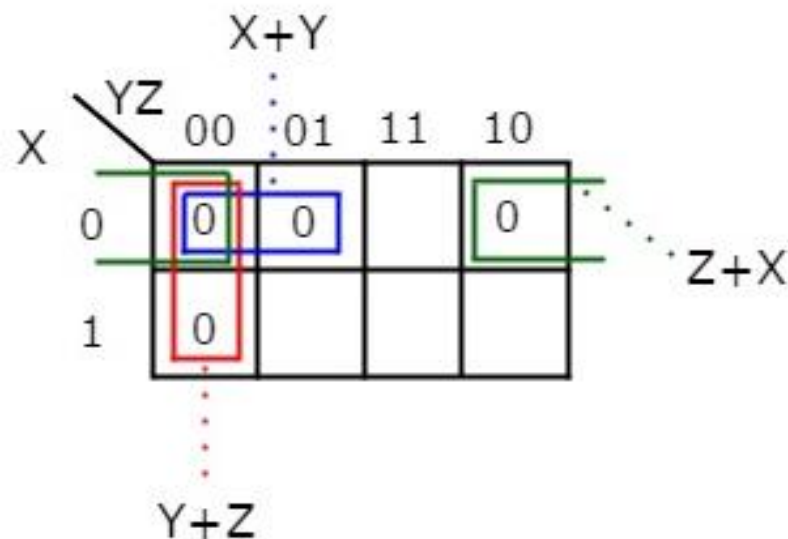
2. Minimise the following function in SOP minimal form using K-Maps:
Maps: $F(A, B, C, D) = m(1, 2, 6, 7, 8, 13, 14, 15) + d(3, 5, 12)$

- From **green** group we find terms : AB
- From **red** group we find terms : $AC'D'$
- From **brown** group we find terms : $A'D$
- From **blue** group we find terms : $A'C$
- $f = AC'D' + A'D + A'C + AB$

		CD			
		00	01	11	10
AB	00		1	X	1
	01		X	1	1
	11	X	1	1	1
	10	1			

KARNAUGH MAP: POS

Eg: $F(X,Y,Z) = \prod M(0,1,2,4)$ minimize the given using K-MAP in POS form.



Ans: $F' = X'Z' + Y'Z' + X'Y'$

$$(F')' = F = (X'Z' + Y'Z' + X'Y')'$$

$$F = (X + Z) \cdot (Y + Z) \cdot (X + Y)$$

KARNAUGH MAP: POS

1. K-map of 3 variables-

$$F(A,B,C)=\pi(0,3,6,7)$$

2 elements in one group

BC A	00	01	11	10
0	0 0	1 1	0 3	1 2
1	1 4	1 5	0 7	0 6

- $F = B.C + A'B'C' + AB$
 $(F')' = F = (B.C + A'B'C' + ABC')' = (B' + C')(A + B + C)(A' + B')$
- From **red** group we find terms : $(B' + C')$
- From **brown** group we find terms : $(A' + B' + C)$ & $(A + B)$
- **Final expression** $(A' + B')(B' + C')(A + B + C)$

KARNAUGH MAP: POS

Eg: Minimise the following function in POS minimal form using K-Maps:
Maps: $F(A, B, C, D) = M(6, 7, 8, 9) + d(10, 11, 12, 13, 14, 15)$

- For **red** group we find terms: $(B' + C')$
- For **green** group we find terms: A'
- POS minimal is, $F' = A + BC$

$$\begin{aligned}(F')' &= (A + BC)' \\ &= A'(B' + C')\end{aligned}$$

		CD			
		00	01	11	10
AB	00				
	01			0	0
	11	X	X	X	X
	10	0	0	X	X

DON'T CARE IN KARNAUGH MAP

Eg: $F(A,B,C,D) = \sum m(4,5,7,8,10,11,13,14) + \sum d(0,1,2)$ minimize the given using K-MAP in POS form.

Karnaugh map for the function $F(A, B, C, D) = A + B + C + D$. The map shows four groups of 1s, each highlighted with a different color (blue, red, green, and orange). The simplified expression is $F = A + B + C + D$.

$$\text{Ans: } F' = ABCD + ABC'D' + A'CD' + B'C'D + A'B'$$

$$F = (A+B) (B+C+\bar{D}) (A+\bar{C}+D) (\bar{A}+\bar{B}+C+D) (\bar{A}+\bar{B}+\bar{C}+\bar{D})$$

FIVE-VARIABLE KARNAUGH MAP

- Maps for more than four variables are not as simple to use. A five-variable map needs 32 squares and a six-variable map needs 64 squares.
- When the number of variables becomes large, the number of squares becomes excessively large and the geometry for combining adjacent squares becomes more involved.
- The five-variable map is shown in Figure. It consists of 2 four-variable maps with variables A, B, C, D, and E.
- Variable A distinguishes between the two maps, as indicated on the top of the diagram.
- The left-hand four-variable map represents the 16 squares where $A = 0$, and the other four-variable map represents the squares where $A = 1$. Minterms 0 through 15 belong with $A = 0$ and minterms 16 through 31 with $A = 1$.

FIVE-VARIABLE KARNAUGH MAP

A	B	C	D	E	minterm	Notation
0	0	0	0	0	$A'B'C'D'E'$	m_0
0	0	0	0	1	$A'B'C'D'E$	m_1
...		
1	1	1	1	0	$ABCDE'$	m_{30}
1	1	1	1	1	$ABCDE$	m_{31}

\bar{A}

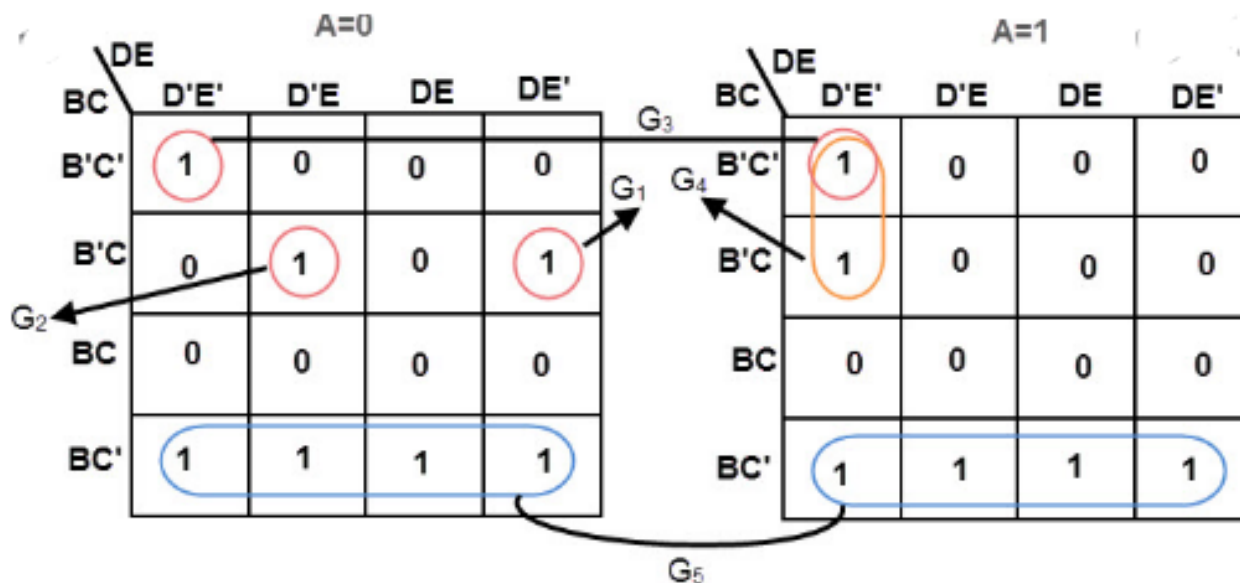
DE \ BC	00	01	11	10
00	0	1	3	2
01	4	5	7	6
11	12	13	15	14
10	8	9	11	10

A

DE \ BC	00	01	11	10
00	16	17	19	18
01	20	21	23	22
11	28	29	31	30
10	24	25	27	26

FIVE-VARIABLE KARNAUGH MAP

Eg: $F(A,B,C,D,E) = \sum m(0, 5, 6, 8, 9, 10, 11, 16, 20, 24, 25, 26, 27)$
minimize the given using K-MAP in SOP form.



Sol: $F = A'B'CDE' + A'B'CD'E + B'C'D'E' + AB'D'E' + BC'$

Daily Quiz

- Complement of the function $F = x'y'z + xy'z'$ is:
 - a) $(x + y + z'). (x' + y + z)$
 - b) $(x' + y' + z). (x + y' + z')$
 - c) $(x + y + z). (x' + y' + z)$
- Minimized function for $F(A, B, C) = \Sigma (2,3,6) + d(0,7)$ using k-map is: (a) $A'C' + BC + BC'$ (b) B (c) AB
- What are the selective prime implicants for $F(w,x,y,z) = \Sigma (1,3,6,7,14)$
 - a) $w'x'z, xyz'$
 - b) $w'x'z, xyz'$
 - c) $w'x'z, w'xy$

FACULTY VIDEO LINKS, YOUTUBE & NPTEL VIDEO LINKS AND ONLINE COURSES DETAILS

Youtube/other Video Links:

- https://www.youtube.com/watch?v=FPrcIhqNPVo&ab_channel=NesoAcademyNesoAcademyVerified
- https://www.youtube.com/watch?v=wjM2RDG5yTI&ab_channel=NesoAcademyNesoAcademyVerified
- https://www.youtube.com/watch?v=BPBiyzc0OBw&ab_channel=IITKharagpurJuly2018IITKharagpurJuly2018

Old Questions

- $F(A,B,C) = \sum m(1,4,5,7)$ minimize the given using K-MAP in SOP form.
- $F(A,B,C,D) = \sum m(0,2,4,6,8,10,12,14)$ minimize the given using K-MAP in SOP form.
- $F(w,x,y,z) = \sum m(4,5,7,8,10,14)$ minimize the given using K-MAP in POS form.
- $F(A,B,C,D,E) = \sum m(4,5,7,8,10,14,27,31) + \sum d(0,1,2,11,19,25)$ minimize the given using K-MAP in SOP form.
- $F(w,x,y,z) = \sum m(0,2,4,5,8,12,14) + \sum d(1,3,5,7,9)$ minimize the given using K-MAP in POS form.

Recap

- The map method provides a simple straightforward procedure for minimizing Boolean functions.
- The map is a diagram made up of squares. Each square represents one minterm.
- K map rule summary:
 1. No zeros allowed.
 2. No diagonals.
 3. Only power of 2 number of cells in each group.
 4. Groups should be as large as possible.
 5. Every one must be in at least one group.
 6. Overlapping allowed.
 7. Wrap around allowed.
 8. Fewest number of groups possible.

QM ALGORITHM

Topic Objective	Mapping with CO
To understand the logic minimization using QM algorithm.	CO1

Prerequisite:

- Knowledge of Boolean algebra.
- Knowledge of SOP and POS forms.

QM ALGORITHM

- The map method of simplification is convenient as long as the number of variables does not exceed five or six. As the number of variables increases, the excessive number of squares prevents a reasonable selection of adjacent squares.
- For functions of six or more variables, it is difficult to be sure that the best selection has been made
- The tabulation method overcomes this difficulty. It is a specific step-by-step procedure that is guaranteed to produce a simplified standard-form expression for a function.
- The tabulation method was first formulated by Quine and later improved by McCluskey. It is also known as the Quine-McCluskey method.

QM ALGORITHM

Procedure of Quine-McCluskey Tabular Method

- Follow these steps for simplifying Boolean functions using Quine-McCluskey tabular method.
- **Step 1** – Arrange the given min terms in an **ascending order** and make the groups based on the number of ones present in their binary representations. So, there will be **at most ‘n+1’ groups** if there are ‘n’ Boolean variables in a Boolean function or ‘n’ bits in the binary equivalent of min terms.
- **Step 2** – Compare the min terms present in **successive groups**. If there is a change in only one-bit position, then take the pair of those two min terms. Place this symbol ‘_’ in the differed bit position and keep the remaining bits as it is.
- **Step 3** – Repeat step2 with newly formed terms till we get all **prime implicants**.

QM ALGORITHM

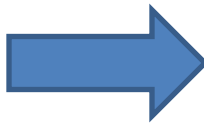
- **Step 4** – Formulate the **prime implicant table**. It consists of set of rows and columns. Prime implicants can be placed in row wise and min terms can be placed in column wise. Place ‘1’ in the cells corresponding to the min terms that are covered in each prime implicant.
- **Step 5** – Find the essential prime implicants by observing each column. If the min term is covered only by one prime implicant, then it is **essential prime implicant**. Those essential prime implicants will be part of the simplified Boolean function.
- **Step 6** – Reduce the prime implicant table by removing the row of each essential prime implicant and the columns corresponding to the min terms that are covered in that essential prime implicant. Repeat step 5 for Reduced prime implicant table. Stop this process when all min terms of given Boolean function are over.

QM ALGORITHM

E.g Minimize the $F(W,X,Y,Z)=\sum m(2,6,8,9,10,11,14,15)$ using Quine-McCluskey method

- Step1 : The given min terms are 2, 6, 8, 9, 10, 11, 14 and 15. The ascending order of these min terms based on the number of ones present in their binary equivalent is 2, 8, 6, 9, 10, 11, 14 and 15.

Min terms	W	X	Y	Z
2	0	0	1	0
8	1	0	0	0
6	0	1	1	0
9	1	0	0	1
10	1	0	1	0
11	1	0	1	1
14	1	1	1	0
15	1	1	1	1



Min terms	W	X	Y	Z
2,6	0	-	1	0
2,10	-	0	1	0
8,9	1	0	0	-
8,10	1	0	-	0
6,14	-	1	1	0
9,11	1	0	-	1
10,11	1	0	1	-
10,14	1	-	1	0
11,15	1	-	1	1
14,15	1	1	1	-

QM ALGORITHM

Step 2 : The given min terms are arranged into 4 groups based on the number of ones present in their binary equivalents. The following table shows the possible **merging of min terms** from adjacent groups.

Min terms	W	X	Y	Z
2,6	0	-	1	0
2,10	-	0	1	0
8,9	1	0	0	-
8,10	1	0	-	0
6,14	-	1	1	0
9,11	1	0	-	1
10,11	1	0	1	-
10,14	1	-	1	0
11,15	1	-	1	1
14,15	1	1	1	-



Min terms	W	X	Y	Z
2,6,10,14	-	-	1	0
2,10,6,14	-	-	1	0
8,9,10,11	1	0	-	-
8,10,9,11	1	0	-	-
10,11,14,15	1	-	1	-
10,14,11,15	1	-	1	-

QM ALGORITHM

Step 3 : The min terms, which are differed in only one-bit position from adjacent groups are merged. That differed bit is represented with this symbol, ‘-’. In this case, there are three groups and each group contains combinations of two min terms. The following table shows the possible **merging of min term pairs** from adjacent groups.

Min terms	W	X	Y	Z
2,6,10,14	-	-	1	0
2,10,6,14	-	-	1	0
8,9,10,11	1	0	-	-
8,10,9,11	1	0	-	-
10,11,14,15	1	-	1	-
10,14,11,15	1	-	1	-



Min terms	W	X	Y	Z
2,6,10,14	-	-	1	0
8,9,10,11	1	0	-	-
10,11,14,15	1	-	1	-

QM ALGORITHM

- These combinations of 4 min terms are available in two rows. So, we can remove the repeated rows. The reduced table after removing the redundant rows is shown below.

Min terms	W	X	Y	Z
2,6,10,14	-	-	1	0
8,9,10,11	1	0	-	-
10,11,14,15	1	-	1	-

- There are three rows in the above table. So, each row will give one prime implicant. Therefore, the **prime implicants** are YZ' , WX' & WY .

Min terms / Prime Implicants	2	6	8	9	10	11	14	15
YZ'	1	1			1		1	
WX'			1	1	1	1		
WY					1	1	1	1

QM ALGORITHM

Min terms / Prime Implicants	2	6	8	9	10	11	14	15
YZ'	1	1			1		1	
WX'			1	1	1	1		
WY					1	1	1	1

- In this example problem, we got three prime implicants and all the three are essential. Therefore, the **simplified Boolean function** is

$$f(W,X,Y,Z) = YZ' + WX' + WY.$$

QM ALGORITHM

Eg: $F(A,B,C,D) = \sum m(5,7,11,12,27,29) + \sum d(14,20,21,22,23)$
minimize the given using QM method.

Sol: At step 1 we have to search for prime Implicants.

2-1s	5 ✓	5,7(2) ✓	<div>20,21,22,23(1,2) PI</div> <div>5,7,21,23(2,16) PI</div> <div>20,22,21,23(2,1)</div> <div>5,21,7,23(16,2)</div>
	12 ✓	5,21(16) ✓	
	20 ✓	12,14(2) PI	
		20,21(1) ✓	
3-1s		20,22(2) ✓	<p>Don't forget that if it did not get checked off then it has to go into the PI table (unless it turns out to be completely made up of Don't Cares)</p>
	7 ✓	7,23(16) ✓	
	11 ✓	11,27(16) PI	
	14 ✓	21,23(2) ✓	
	21 ✓	21,29(8) PI	
4-1s	22 ✓	22,23(1) ✓	
	23 ✓	X	
	27 ✓		
	29 ✓		

QM ALGORITHM

PI Table

			✓ 5	✓ 7	✓ 11	✓ 12	✓ 27	✓ 29
Don't Cares Only	20,21,22,23	(1,2)						
EPI	5,7,21,23	(2,16)	⊗	⊗				
EPI	12,14	(2)				⊗		
EPI	11,27	(16)			⊗		⊗	
EPI	21,29	(8)						⊗

			16 A	8 B	4 C	2 D	1 E	Boolean
EPI	5,7,21,23	(2,16)	—	0	1	—	1	$\overline{B}CE$
EPI	12,14	(2)	0	1	1	—	0	$\overline{A}BC\overline{E}$
EPI	11,27	(16)	—	1	0	1	1	$B\overline{C}DE$
EPI	21,29	(8)	1	—	1	0	1	$AC\overline{D}E$

$$f(A,B,C,D,E) = \overline{B}CE + \overline{A}BC\overline{E} + B\overline{C}DE + AC\overline{D}E$$

Recap

- The tabulation method is a specific step-by-step procedure that is guaranteed to produce a simplified standard-form expression for a function.
- **Step 1** – Arrange the given min terms in an **ascending order** and make the groups based on the number of ones present in their binary representations.
- **Step 2** – Compare the min terms present in **successive groups**.
- **Step 3** – Repeat step2 with newly formed terms till we get all **prime implicants**.
- **Step 4** – Formulate the **prime implicant table**.
- **Step 5** – Find the essential prime implicants by observing each column.

Daily Quiz

- What are the advantages of tabulation methods over K-Map method?
- Write the steps to minimize function using QM method.
- $F(A,B,C,D) = \sum m(1,3,7,11,15) + \sum d(0,2,5)$ minimize the given using QM method.
- What minimization methods are used when number of variables are greater than 6:
 - a) K-Map method
 - b) QM method
 - c) (a) or (b) both

Weekly Assignment

- Convert the following numbers
 - i) $(163.789)_{10}$ to Octal number
 - ii) $(11001101.0101)_2$ to base-8 and base-4
 - iii) $(4567)_{10}$ to base 2
 - iv) $(4D.56)_{16}$ to Binary
- Subtract $(111001)_2$ from $(101011)_2$ using 1's complement?
- Represent the decimal number 3452 in
 - i) BCD
 - ii) Excess-3
- perform $(-50) - (-10)$ in binary using the signed-2's complement
- Determine the value of base x if $(211)_x = (152)_8$ (L4) (2M)
- Convert the following numbers
 - i) $(250.5)_{10} = ()_2$
 - ii) $(673.23)_{10} = ()_8$
 - iii) $(101110.01)_2 = ()_8$
- Convert the following to binary and then to gray code $(AB33)_{16}$
- Perform the following Using BCD arithmetic $(7129)_{10} + (7711)_{10}$
- Explain the Binary codes with examples?
- What is Digital System? Characteristics of digital systems. Explain the difference between analog and digital systems.

Weekly Assignment

- Design the circuit by Using NAND gates $F = ABC' + DE + AB'D'$
- Simplify and implementation the following SOP function using NOR gates $F(A,B,C,D) = \sum m(0,1,4,5,10,11,14,15)$
- Convert the following
 - a) $(1AD)_{16} = ()_{10}$
 - b) $(453)_8 = ()_{10}$
 - c) $(10110011)_2 = ()_{10}$
 - d) $(5436)_{10} = ()_3$
- Explain binary to Gray & Gray to binary conversion with example?
- State and Explain the DeMorgan's Theorem and Consensus Theorem
- Convert the following numbers
 - a) $(615)_{10} = ()_{16}$
 - b) $(214)_{10} = ()_8$
 - c) $(0.8125)_{10} = ()_2$
 - d) $(658.825)_{10} = ()_8$ v) $(54)_{10} = ()_2$ 10)
- Explain the Excess-3 code? Write about Error correction & Detection?

FACULTY VIDEO LINKS, YOUTUBE & NPTEL VIDEO LINKS AND ONLINE COURSES DETAILS

Youtube/other Video Links:

- https://www.youtube.com/watch?v=11jgq0R5EwQ&t=2s&ab_channel=NesoAcademy
- https://www.youtube.com/watch?v=Shj-u66gdE8&ab_channel=EkeedaEkeedaVerified

1. The number system which uses alphabets as well as numerals is Binary number system ?
 - Octal number system
 - Decimal number system
 - Hexadecimal number system

2. The number that come immediately after and before hex number (FFEF)₁₆ are respectively
 - (FFF0)₁₆ and (FFE0)₁₆
 - (FFFE)₁₆ and (FFEE)₁₆
 - (FFF0)₁₆ and (FFFE)₁₆
 - (FFF0)₁₆ and (FFEE)₁₆

- 4 Two input exclusive NOR gate gives high output
- When one input is high and the output is low
 - Only when both the inputs are low
 - When both the inputs are same
 - Only when both the inputs are high
5. What is 2's complement or 0011 0101 1001 1100 number?
- 1100 1010 1100 1011
 - 1100 1010 0110 0011
 - 1100 1010 0110 0100
 - None of these

5. The maxterm designator of the term $A' + B' + C + D'$ is

- 2
- 13
- 10
- None of these

6. The minterm, designator of the term $AB'CD$ is

- 4
- 15
- 11
- None of these

Glossary questions

- The minterm corresponding to decimal number 15 is -----
- Decimal digit 5 is represented by -----using 7312 weighted code.
- Represent the decimal number $(0.875)_{10}$ into binary form.
- A decimal number 6 in excess-3 code is-----
- $(1111)_2 - (1111)_2$ is equal to-----
- 9's complement of 5436 is-----
- $(28CE5)_H + (AB2C3)_H = (-----)_H$
- Characteristic of Gray code is -----,-----and it is unit -----code.
- For the identity $AB + A'C + BC = AB + A'C$, the dual form is-----
- The minimum number of NAND gates required to implement $A + AB' + AB'C$ is equal to-----

Cont...

- The code used for labelling the cells of the K-map is-----
- Switches connected in parallel behave as-----
- Output of gate -----is $(A'B)'$
- Positive logic in a logic circuit is one in which logic 0 voltage level is----- than logic 1 voltage level.
- AND-OR realization is equivalent to -----realization.
- 2-input EX-OR gate can also be performed as inverter if its one of the input is fixed as logic-----
- The output of a logic gate is '1' when all its input are at logic '0' the gate is either----- or an -----
- Given the logic function of Four variables $f(A,B,C,D) = (A' + BC)(B + CD)$. The function as a sum of product will be-----

Old Questions

- Minimize the following function by Quine McClusky method and also perform the NAND implementation of the simplified function.

$$F(w,x,y,z) = \sum m(1,4,8,9,13,14,15) + d(2,3,11,12)$$

- $F(A,B,C,D) = \sum m(0,1,2,3,10,11,12,13,14,15)$ minimize the given using QM method.
- $F(A,B,C,D) = \sum m(1,3,7,11,15) + \sum d(0,2,5)$ minimize the given using QM method.

OLD QUESTION PAPERS

B. TECH.
(SEM-III) THEORY EXAMINATION 2019-20
DIGITAL SYSTEM DESIGN

Time: 3 Hours

Total Marks: 100

Note: Attempt all Sections. If require any missing data; then choose suitably.

SECTION A

1. Attempt all questions in brief.

2 x 10 = 20

Qno.	Question	Marks	CO
a.	The solution to the quadratic equation $k^2 - 11k + 22 = 0$ are $x = 3$ and $x = 6$. What is the base of the number system?	2	1
b.	Simplify the expression $F(A, B, C, D) = ACD + \bar{A}B + \bar{D}$ by K- Map.	2	1
c.	Construct half subtractor using logic gates.	2	2
d.	Implement a 4:1 multiplexer using 2:1 multiplexer.	2	2
e.	What do you mean by race around condition in JK Flip Flop?	2	3
f.	Distinguish between Leach and Flip Flop.	2	3
g.	What is logic family? Give the classification of logic families in brief.	2	4
h.	Describe figure of merit & noise immunity of TTL & CMOS ICs.	2	4
i.	What are the advantages and disadvantages of flash type ADC?	2	5
j.	The basic step of a 9-bit DAC is 10.3 mV. If 000000000 represents 0Volts, what is the output for an input of 101101111?	2	5

OLD QUESTION PAPERS

SECTION B

2. Attempt any *three* of the following:

3 x 10 = 30

Qno.	Question	Marks	CO
a.	Design an excess-3 to BCD code converter.	10	1
b.	Implement a full adder by using 8:1 multiplexer.	10	2
c.	Design a sequential circuit with two Flip Flops, A & B and one input x. When x=0, the State of the circuit remains the same when x=1 the circuit passes through the state transitions from 00 to 01 to 11 to 10 back to 00 & repeat.	10	3
d.	Compare TTL and CMOS logic families and also draw CMOS NOR gate.	10	4
e.	Explain the operation of successive approximation ADC. Discuss its merits and demerits.	10	5

SECTION C

3. Attempt any *one* part of the following:

1 x 10 = 10

Qno.	Question	Marks	CO
a.	Minimize the logic function using Quine-McCluskey Method $F(A, B, C, D, E) = \sum m(8, 9, 10, 11, 13, 15, 16, 18, 21, 24, 25, 26, 27, 30, 31)$	10	1
b.	Simplify the logic expression using K-Map $F(A, B, C, D, E, F) = \sum m(0, 5, 7, 8, 9, 12, 13, 23, 24, 25, 28, 29, 37, 40, 42, 44, 46, 55, 56, 57, 60, 61)$	10	1

OLD QUESTION PAPERS

4. Attempt any *one* part of the following:

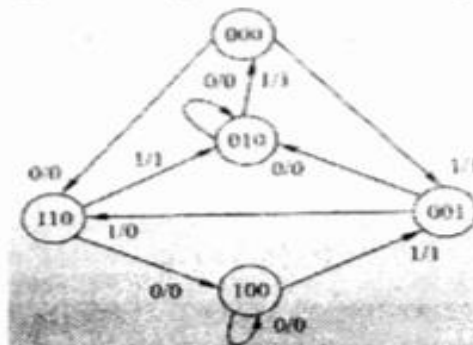
1 x 10 = 10

Qno.	Question	Marks	CO
a.	Design a 4-bit parallel binary Adder/Subtractor circuit.	10	2
b.	Design a 4-bit comparator circuit using logic gates.	10	2

5. Attempt any *one* part of the following:

1 x 10 = 10

Qno.	Question	Marks	CO
a.	Discuss Mealy and Moore FSM. What do you mean by excitation table?	10	3
b.	For the given state diagram design the circuit using T flip flop	10	3



6. Attempt any *one* part of the following:

1 x 10 = 10

Qno.	Question	Marks	CO
a.	Draw three input standard TTL NAND gate circuit and explain its operation.	10	4
b.	Implement the following function using PLA $F_1 = \sum m(0,3,4,7)$ $F_2 = \sum m(1,2,5,7)$	10	4

Expected Questions for University Exam

- Simplify the Boolean expression using K-MAP:
 $F(A,B,C,D,E) = \sum m(0,1,4,5,16,17,21,25,29)$
- Simplify the Boolean expression using K-MAP:
 $F(A,B,C,D) = \sum m(1,2,3,8,9,10,11,14) + d(7,15)$
- Simplify the Boolean expression using K-map and implement using NAND gates
 $F(A,B,C,D) = \sum m(0,2,3,8,10,11,12,14)$
- Simplify the Boolean expressions to minimum number of literals
 - $(A + B)(A + C')(B' + C') (L3) (3M)$
 - $AB + (AC)' + AB'C (AB + C) (L3)(4M)$
 - $(A+B)' (A'+B')' (L5) (3M)$
- Reduce the expression $f(x,y,z,w) = \pi M(0,2,7,8,9,10,11,15) .d (3,4)$ using K-Map?
- Simplify the Boolean expression using K-map?:
 $F(A,B,C,D,E) = \sum m(0,2,4,6,9,11,13,15,17,21,25,27,29,31)$
- Obtain the a) SOP b) POS expression for the function given below:
 $F(A,B,C,D) = \sum m(0,1,2,5,8,9,10)$
- Simplify the Boolean expressions to minimum number of literals:
 - $X' + XY + XZ' + XYZ'$
 - $(X+Y) (X+Y')$
- Obtain the Complement of Boolean Expression (L4) (5M)
 - $A+B+A'B'C$
 - $AB + A(B+C) + B'(B+D)$
- Determine the minimal sum of product form of (L4) (5M,5M)
 - $f(w,x,y,z) = \sum m(4,5,7,12,14,15) + d(3,8,10)$
 - $F(A,B,C,D) = \pi M(0,3,5,6,8,12,15)$

REFERENCE

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Thank You