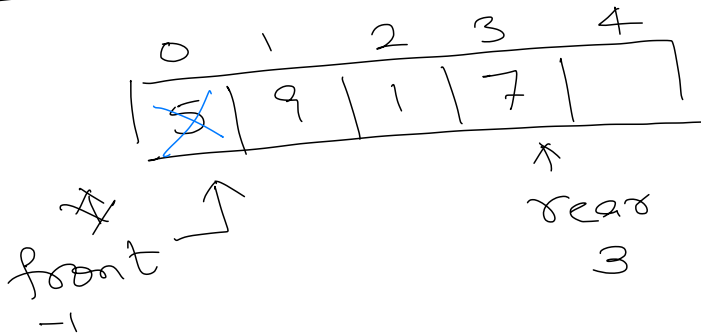


Linear Queue



Delete Q()

- Move front to next element.
- Remove element from front.
- Shift all elements of queue to left by 1 place.
↓
This is done to avoid condition that queue is empty as well as full. In efficient way.

OR

- if queue is empty AND queue is full then
 - Reset front and rear
 $\text{front} = -2, \text{rear} = -1.$

5	7	3
0	1	2

Add Q(5)
Add Q(7)
Add Q(3)

front \rightarrow ~~1~~ ~~0~~ ~~2~~ 2
rear \rightarrow ~~1~~ ~~0~~ ~~2~~ 2 \Leftarrow Queue is full.

\Rightarrow Queue is empty.

Delete Q() \Rightarrow 5
Delete Q() \Rightarrow 7
Delete Q() \Rightarrow 3

```

class C1 {
    data member
    member functions    F1();
}

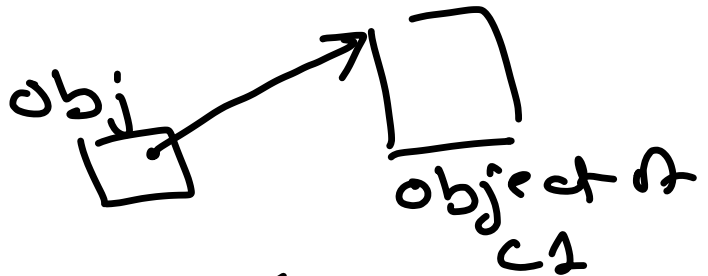
```

C1 obj;

\rightarrow variable that stores reference to object of class.

obj.F1(); ← Null Pointer Exception.

obj = new C1();



obj.F1(); ✓

```
class Stack {  
    public void Push(int)  
    public int Pop() { ... }  
};
```

3

Defining Stack
as ADT



```
interface StackIntf {  
    public void Push(int);  
    public int Pop();  
};
```

7

class Stack Using Array
implements StackIntf

{
:
}

3
Reverse (int [] elements,
StackIntf stack)
{
stack.Push(..);
:
stack.Pop();
}

}

stack
↗

obj = new Stack Using Array();
Reverse (arr, obj);

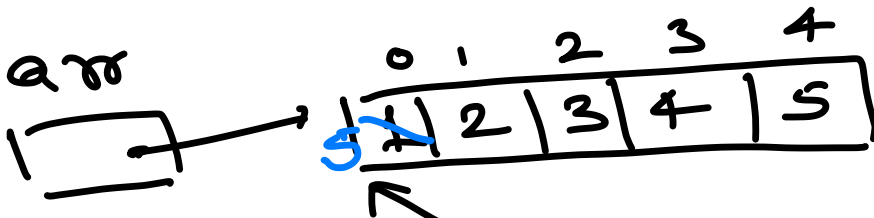
main()

int[] arr = {1, 2, ..., 5};

Reverse(arr, ...);

⋮

}

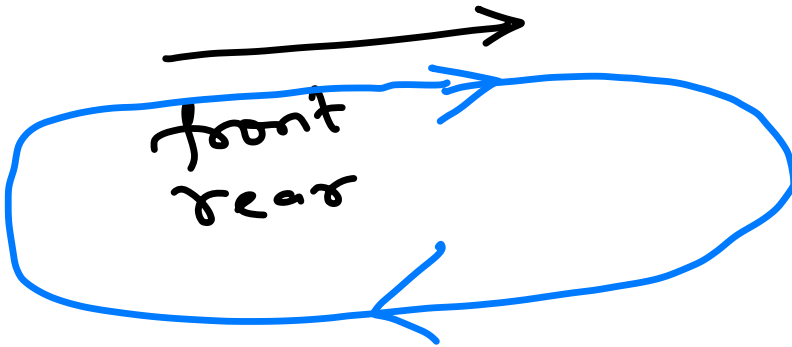
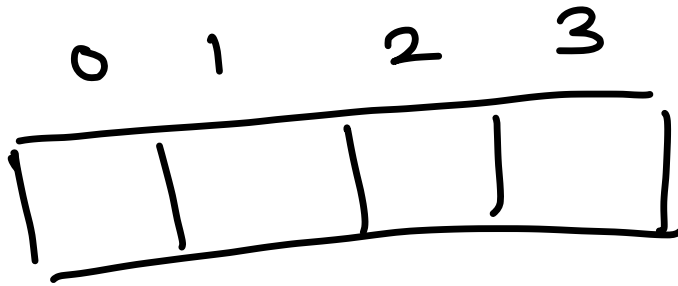


OR

Reverse(^{elements}int[] elements, ...) {
elements[0] = 5;

}

int[] arr;
arr = new int[5];
arr[0] = 1; arr[2] = 3;
arr[1] = 2; arr[3] = 4;
arr[4] = 5;



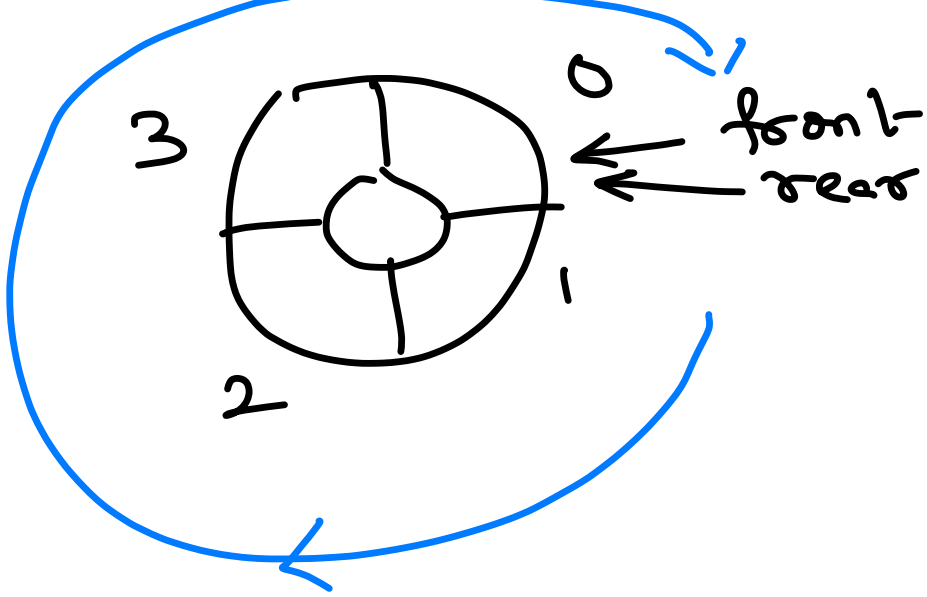
```

++ rear;
if (rear == n)
    rear = 0;

```

OR

$$rear = \frac{(rear + 1) \% n}{0 \dots (n-1)}$$



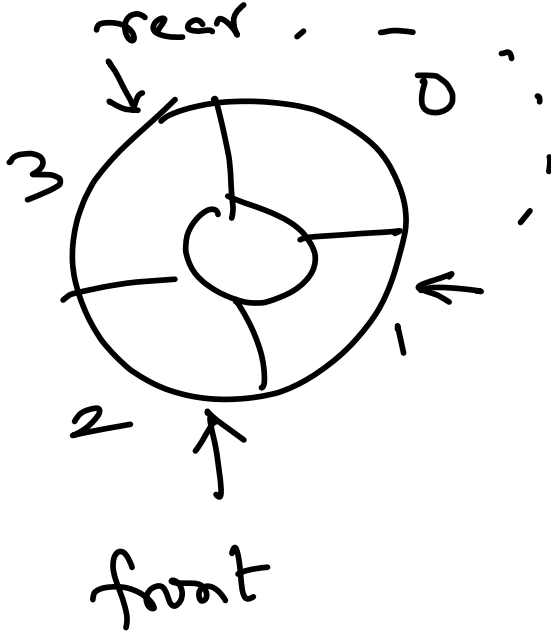
front equals rear \leftarrow Empty

rear just before front \leftarrow Full

\Downarrow
 $(\text{rear} + 1) \% n$ equals front

new with [4].

0	1	2	3	...
1	1	1	1	...



① Check for balanced parenthesis.

$$\begin{array}{c} () () \\ \hline ([]] \end{array} \checkmark$$

$$\begin{array}{c}) (\\ \hline ([]) \end{array} \checkmark$$

Hint: use stack.

boolean IsBalanced (String str)

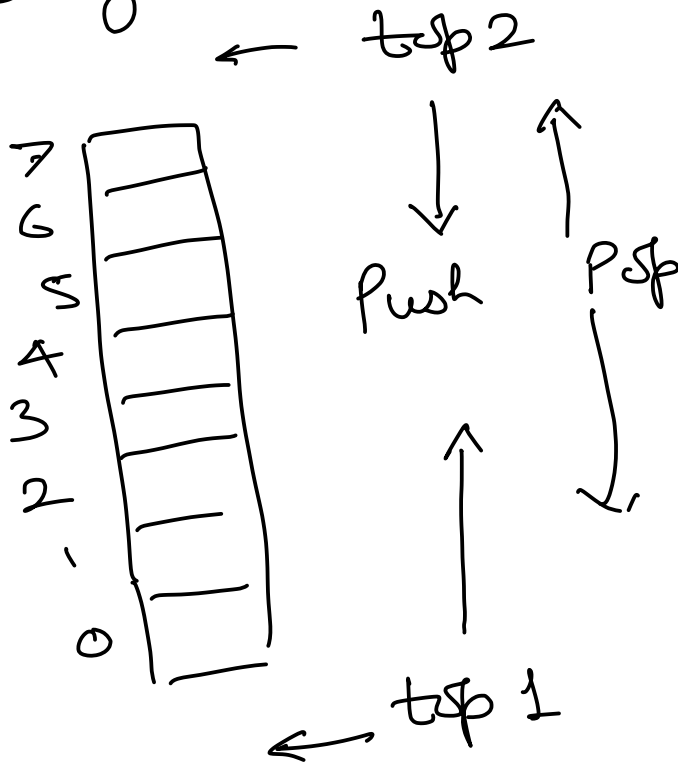
* ② Implement Stack using Queue.

* ③ Implement Queue using Stack.

Hint: ② will need two queues.

③ will need two stacks.

④ Implement 2 Stacks in a single array.



Linked list

Array : Need?

When we need to store multiple elements.
And do same processing on those elements.

Properties of array:

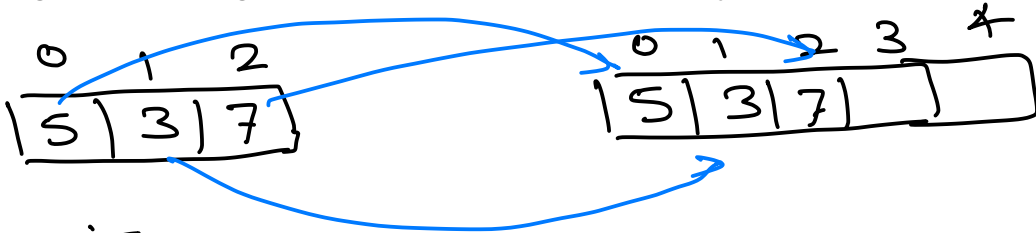
- Data structure that stores multiple elements, all of the **same type**.
- All elements of an array are **stored sequentially** in memory, one after another.

Advantages of array:

- Efficient lookup OR **Random access**.
- Efficient in adding or removing elements **at the end of array**.

Disadvantages of array

- **Fixed size**. Resizing of array is inefficient.
- **Inserting and deleting** of elements, in middle of array is inefficient.



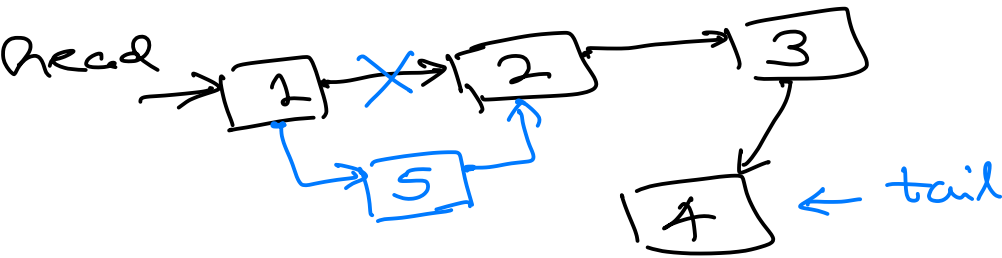
Resize

- Create a larger array
- Copy values from existing array to new one.
- Release old memory.

Initial array Size = 1000 copy 1000 elem.
Resized array Size = 2000
Resized array Size = 4000 copy 2000 elements

Linked list → Do not store elements next to each other.

0	1	2	3	4
1	2	3	4	



Properties of linked list

- Stores data as a chain of **nodes**.
- Each node contains **data** and a **pointer** to next node in chain.
- We need to know where first node is of list - **head**.

→ where data is stored.

Advantages of linked list

- Can easily grow / shrink in size.
- Efficient in insertion and deletion of elements.

Disadvantages of linked list

- Random access is inefficient.

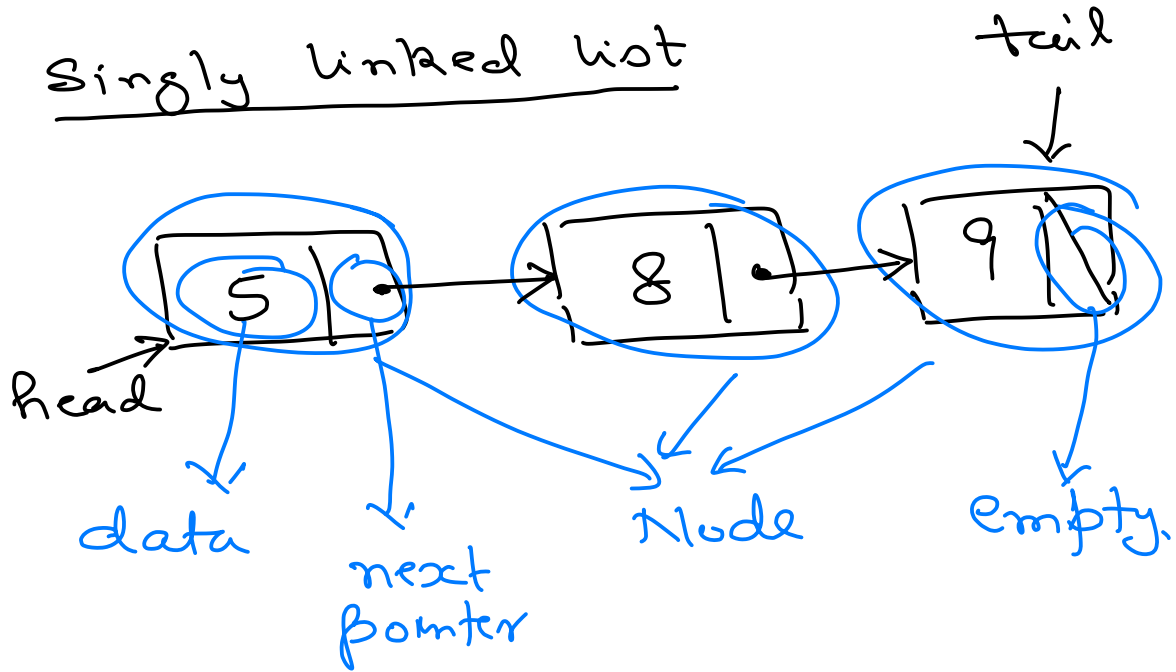
Types of linked list

- Singly linked list (Uni-directional)
- Doubly linked list (Bi-directional)
- Circular list.

One node keeps track of one neighbour only.

Each node keeps track of both of its neighbours

Singly linked list



Traversal

Starting from first element, access each element one at a time, till the last element.

Traversal.

```
for ( i = 0; i < elements.length; ++i)
    stack.Push (elements[i]);
```

- ① what if list is empty?
- ② what if list is not empty?

head \rightarrow empty \Leftarrow list is empty.

head \rightarrow  \Leftarrow list is not empty



\swarrow * current \searrow *
 \downarrow empty

o/p: 1 2 3

Traversal of singly list

- if list is empty then Stop.
- Set current to first node of list.
- while (current is not empty) do
 - Process current node.
 - Move current to current node's next.
- Stop.

Traversal of singly list (Optimised)

- Set current to first node of list.
- while (current is not empty) do
 - Process current node.
 - Move current to current node's next.
- Stop.