# Programski Prevodioci

#### Druga Laboratorijska Vežba Grupa 8

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Statements → Statements;
Statement | Statement
Statement → Assignment |
WhileStatement
Assignment → ID = Term | ID = Term
+ Term
WhileStatement → repeat ( Term ) {
Statements }
Term → ID | CONST
```

	Statement
(2)	$Statements \rightarrow Statement$
(3)	$Statement \rightarrow Assignment$
(4)	$Statement \rightarrow$
	WhileStatement
(5)	Assignment $\rightarrow$ <b>ID</b> = Term
(6)	Assignment $\rightarrow$ <b>ID</b> = Term +
	Term
(7)	$WhileStatement \rightarrow repeat ($
	Term ) { Statements }
(8)	$Term \rightarrow ID$
(9)	$Term \rightarrow CONST$

(1) Statements → Statements;

(1) 
$$Ss \to Ss$$
;  $S$   
(2)  $Ss \to S$   
(3)  $S \to A$   
(4)  $S \to WS$   
(5)  $A \to ID = T$   
(6)  $A \to ID = T + T$   
(7)  $WS \to \text{repeat}(T) \{ Ss \}$   
(8)  $T \to ID$   
(9)  $T \to \text{CONST}$ 

## 1. Kreiranje kanoničkog skupa LR pravila

$l_0$ :	
	$\underline{Ss' \rightarrow . Ss}$
	$Ss \rightarrow . Ss ; S$
	$Ss \rightarrow . S$
	$S \to A$
	$S \to .WS$
	$\underline{A} \rightarrow . \mathbf{ID} = \underline{T}$
	$\underline{A} \rightarrow . \mathbf{ID} = T + T$
	$\underline{WS} \rightarrow .$ repeat $(T) \{Ss\}$

$$\frac{SS' \to SS}{SS \to SS}$$

 $l_1 = goto(l_0, Ss)$ :

$$l_2 = \text{goto}(l_I, ;):$$

$$\underline{Ss \to SS ; .S}$$

$$\underline{S \to .A}$$

$$\underline{S \to .WS}$$

$$\underline{A \to .ID = T}$$

$$\underline{A \to .ID = T + T}$$

$$\underline{WS \to .repeat(T) \{ Ss \}}$$

$$l_3 = \text{goto}(l_2, S)$$
:
$$\underline{S_S \to SS ; S}.$$
Redukciono stanje za smenu (1)

$$l_4 = \mathrm{goto}(\ l_0, S\ )$$
: 
$$\frac{S_S \to S\ .}{\mathrm{Redukciono\ stanje\ za\ smenu\ (2)}}$$

$$l_5 = \mathrm{goto}(\ l_0, A\ )$$
:
$$\frac{S \to A}{\mathrm{Redukciono\ stanje\ za\ smenu\ (3)}}$$

$$l_6 = goto(\ l_0,\ WS\ )$$
:  $S \to WS$  . Redukciono stanje za smenu (4)

$$l_7 = \text{goto}(\ l_0, \mathbf{ID}\ )$$
:  
 $\underline{A \to \mathbf{ID} \cdot = T}$   
 $\underline{A \to \mathbf{ID} \cdot = T + T}$ 

$$l_8 = goto(l_7, =):$$

$$\underline{A \rightarrow ID = .T}$$

$$\underline{A \rightarrow ID = .T + T}$$

$$\underline{T \rightarrow .ID}$$

$$\underline{T \rightarrow .CONST}$$

 $l_{14} = goto(l_{18}, repeat)$ 

$l_9 = \operatorname{goto}(\ l_8,\ T\ )$ : $A \to \operatorname{ID} = T$ .  Redukciono stanje za smenu (5) $A \to \operatorname{ID} = T \cdot + T$	$l_{10} = \text{goto}(l_9, +):$ $\underline{A \rightarrow \text{ID}} = \underline{T + . T}$ $\underline{T \rightarrow . \text{ID}}$ $\underline{T \rightarrow . \text{CONST}}$	$l_{11} = \text{goto}(\ l_{10},\ T\ )$ : $\frac{A \to \text{ID} = T + T}{\text{Redukciono stanje za smenu (6)}}$
$l_{12} = \text{goto}(l_8, \mathbf{ID})$ : $\underline{T \rightarrow \mathbf{ID}}.$ Redukciono stanje za smenu (8)	$l_{13} = \text{goto}(\ l_8, \mathbf{CONST}\ ):$ $\frac{T \to \mathbf{CONST}.}{\mathbf{Redukciono\ stanje\ za\ smenu\ (9)}}$	$l_{14} = goto(l_0, repeat):$ $\underline{WS} \rightarrow repeat.(T) \{Ss\}$
$l_{15} = \text{goto}(l_{14}, ():$ $\underline{WS} \rightarrow \text{repeat}(.T) \{Ss\}$ $\underline{T} \rightarrow . \text{ID}$ $\underline{T} \rightarrow . \text{CONST}$	$l_{16} = \text{goto}(l_{15}, T)$ :  WS $\rightarrow \text{repeat}(T.) \{Ss\}$	$l_{17} = \text{goto}(l_{16},)$ :  WS $\rightarrow$ repeat $(T) \cdot \{Ss\}$
$l_{18} = \gcd(l_{17}, \{ \}):$ $WS \rightarrow \operatorname{repeat}(T) \{ . Ss \}$ $SS \rightarrow . SS ; S$ $SS \rightarrow . S$ $S \rightarrow . A$ $S \rightarrow . WS$ $A \rightarrow . ID = T$ $A \rightarrow . ID = T + T$ $WS \rightarrow . \operatorname{repeat}(T) \{ . Ss \}$	$l_{19} = \text{goto}(l_{18}, S_S)$ : $WS \rightarrow \text{repeat}(T) \{ S_S . \}$ $S_S \rightarrow S_S . ; S$	$l_{20} = \text{goto}(l_{19}, )$ :  WS $\rightarrow \text{repeat}(T) \{ Ss \}$ .  Redukciono stanje za smenu (7)
$l_2 = goto (l_{19}, ;)$	$l_6 = \text{goto} (l_{18}, WS)$	$l_{13} = \text{goto}(l_{10}, \mathbf{CONST})$
$l_4 = \text{goto} (l_{18}, S)$ $l_5 = \text{goto} (l_2, A)$	$l_7 = \text{goto} (l_2, \mathbf{ID})$ $l_7 = \text{goto}(l_{18}, \mathbf{ID})$	$l_{13} = \text{goto}(\ l_{15}, \mathbf{CONST}\ )$ $l_{14} = \text{goto}(\ l_2, \mathbf{repeat}\ )$

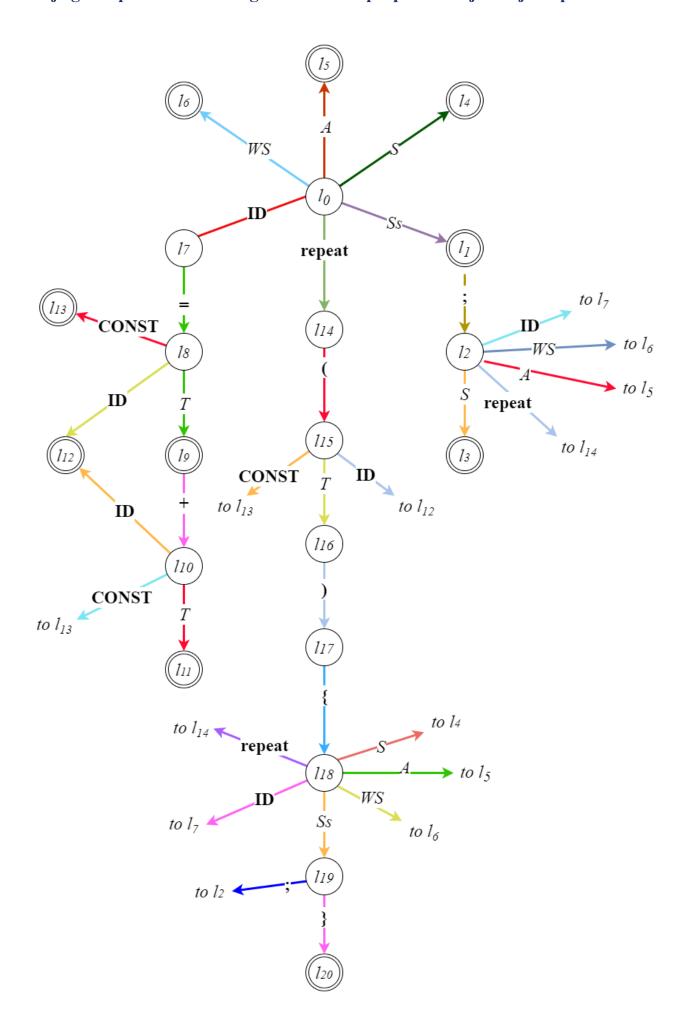
 $l_{12} = \operatorname{goto}(l_{15}, \mathbf{ID})$ 

 $l_{12} = \text{goto}(l_{10}, \mathbf{ID})$ 

 $l_5 = \text{goto}(l_{18}, A)$ 

 $l_6 = goto(l_2, WS)$ 

### 2. Crtanje grafa prelaza konačnog automata za prepoznavanje vidljivih prefiksa



## 3. Popunjavanje LR sintaksne tabele

Pomoćna tabela Follow(A) i First(α) funkcija za konkretne smene										
Rbr. smene	Redukciono stanje	$\mathbf{A}  ightarrow \mathbf{lpha}$	FOLLOW(A)	FIRST(α)						
0.	$l_1$	$Ss' \rightarrow Ss$	#	ID repeat						
1.	$l_3$	$Ss \rightarrow Ss$ ; $S$	#;}	ID repeat						
2.	$l_4$	$S_S \rightarrow S$	#;}	ID repeat						
3.	$l_5$	$S \to A$	#;}	ID						
4.	$l_6$	$S \to WS$	#;}	repeat						
5.	l <sub>9</sub>	$A \to \mathbf{ID} = T$	#;}	ID						
6.	$l_{11}$	$A \to \mathbf{ID} = T + T$	#;}	ID						
7.	$l_{12}$	$T \rightarrow \mathbf{ID}$	#;}+)	ID						
8.	<i>l</i> <sub>13</sub>	$T \rightarrow \mathbf{CONST}$	#;}+)	CONST						
9.	$l_{20}$	$WS \rightarrow \mathbf{repeat} (T) \{Ss\}$	#;}	repeat						

LR sintaksna tabela																
		AKCIJE										PRELAZI				
	;	ID	=	+	CONST	repeat	(	)	{	}	#	Ss	S	A	T	WS
0		<b>s</b> 7				s14						1	4	5		6
1	<b>s2</b>										acc					
2		s7				s14							3	5		6
3	r1									r1	r1					
4	r2									r2	r2					
5	r3									r3	r3					
6	r4									r4	r4					
7			<b>s8</b>													
8		s12			s13										9	
9	r5			<b>s10</b>						r5	r5					
10		s12			s13										11	
11	r6									r6	r6					
12	<b>r</b> 7			<b>r</b> 7				<b>r</b> 7		<b>r</b> 7	<b>r</b> 7					
13	r8			r8				r8		r8	r8					
14							s15									
15		s12			s13										16	
16								s17								
17									<b>s18</b>							
18		<b>s</b> 7				s14						19	4	5		6
19	<b>s2</b>									<b>s20</b>						
20	r9									r9	r9					