

DAYANANDA SAGAR COLLEGE OF ENGINEERING
 (An Autonomous Institute affiliated to VTU, Approved by AICTE & ISO 9001:2008 Certified)
 Accredited by National Assessment & Accreditation Council [NAAC] with 'A' grade

DEPARTMENT OF TELECOMMUNICATION ENGINEERING
 Accredited by National Board of Accreditation Council (NBA)

SCHEME & SOLUTION: CONTINUOUS INTERNAL ASSESSMENT- 1

Date: 12/10/2020

Course: Cryptography & Network Security Course Code: 17TE7DECS

Semester & Section: VII 'A' 6163 Total no of Pages: 6

Scheme & Solution prepared by: R. VIKRANT

Signature: [Signature]

Q. No.	Description	Marks Distribution
1.	a. iv) Integrity b. ii) Authorization c. i) ENCRYPTED d. ii) Denial of Service e. iii) KLOBV f. ii) Caesar g. ii) Private key h. i) E, T i. ii) cryptography j. ii) other symbol	10 M
2.	Rule for Encryption	2 M

Q. No.	Description	Marks Distribution																									
	<table border="1"> <tr> <td>I</td> <td>N</td> <td>T</td> <td>E</td> <td>R</td> </tr> <tr> <td>S</td> <td>A</td> <td>B</td> <td>C</td> <td>D</td> </tr> <tr> <td>F</td> <td>G</td> <td>H</td> <td>K</td> <td>L</td> </tr> <tr> <td>M</td> <td>O</td> <td>P</td> <td>Q</td> <td>U</td> </tr> <tr> <td>V</td> <td>W</td> <td>X</td> <td>Y</td> <td>Z</td> </tr> </table> <p>P.T = INVESTMENT IN KNOWLEDGE C.T = NT X BI TE NT GENUINERALTY</p>	I	N	T	E	R	S	A	B	C	D	F	G	H	K	L	M	O	P	Q	U	V	W	X	Y	Z	<p>8M } 10</p>
I	N	T	E	R																							
S	A	B	C	D																							
F	G	H	K	L																							
M	O	P	Q	U																							
V	W	X	Y	Z																							
3)	$C = Kp \text{ mod } 26$ <p>P.T = DONT LET TO four del x</p> $\begin{bmatrix} C_1 \\ C_2 \end{bmatrix} = \begin{bmatrix} 11 & 8 \\ 4 & 3 \end{bmatrix} \begin{bmatrix} P_1 \\ P_2 \end{bmatrix} \text{ mod } 26$ $\begin{bmatrix} C_1 \\ C_2 \end{bmatrix} = \begin{bmatrix} 11 & 8 \\ 4 & 3 \end{bmatrix} \begin{bmatrix} 3 \\ 14 \end{bmatrix} \text{ mod } 26 = \begin{bmatrix} 15 \\ 2 \end{bmatrix} = \begin{bmatrix} P \\ C \end{bmatrix}$ $\begin{bmatrix} C_1 \\ C_2 \end{bmatrix} = \begin{bmatrix} 11 & 8 \\ 4 & 3 \end{bmatrix} \begin{bmatrix} 13 \\ 19 \end{bmatrix} \text{ mod } 26 = \begin{bmatrix} 9 \\ 5 \end{bmatrix} = \begin{bmatrix} J \\ F \end{bmatrix}$ $= \begin{bmatrix} 11 & 8 \\ 4 & 3 \end{bmatrix} \begin{bmatrix} 11 \\ 8 \end{bmatrix} \text{ mod } 26 = \begin{bmatrix} 3 \\ 16 \end{bmatrix} = \begin{bmatrix} D \\ U \end{bmatrix}$ <p>C.T = PCTFDUOVICYMTSCXFL</p>	<p>— 5M</p>																									

Q. No.

Description

Marks Distribution

Decryption $\rightarrow P = C K^{-1} \text{ mod } 26$

$$K^{-1} = \frac{\text{adj}(K)}{\text{det}(K)} =$$

$$\text{adj}(K) = \begin{bmatrix} 3 & -8 \\ -4 & 11 \end{bmatrix}$$

$$\text{det}(K) = 7$$

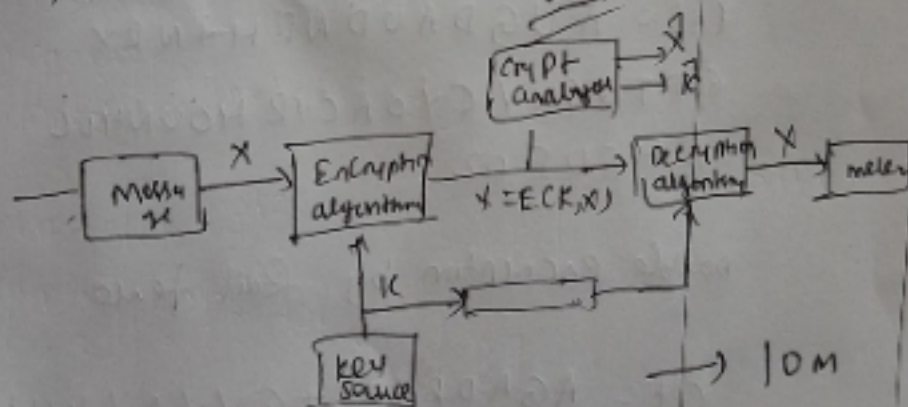
$$K^{-1} = \begin{bmatrix} 3 & 18 \\ 22 & 11 \end{bmatrix}$$

— 5M

$$\begin{pmatrix} P_1 \\ P_2 \end{pmatrix} = \begin{bmatrix} 15 \\ 2 \end{bmatrix} \begin{bmatrix} 3 & 18 \\ 22 & 11 \end{bmatrix} \text{ mod } 26 = \begin{bmatrix} 3 \\ 14 \end{bmatrix} = \begin{bmatrix} D \\ D \end{bmatrix}$$

P.T = DONT USE TO YOURSELF

4)



— 10M

-) Block diagram — 5M
 Explanation — 3M
 Equation — 2M

$$Y = E_K(X)$$

$$X = D_K(Y)$$

(3)

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5)	<table border="1"> <tr> <td>4</td> <td>5</td> <td>1</td> <td>2</td> <td>3</td> <td>6</td> </tr> <tr> <td>C</td> <td>H</td> <td>A</td> <td>N</td> <td>Q</td> <td>E</td> </tr> <tr> <td>Y</td> <td>O</td> <td>U</td> <td>R</td> <td>T</td> <td>H</td> </tr> <tr> <td>O</td> <td>U</td> <td>Q</td> <td>H</td> <td>T</td> <td>S</td> </tr> <tr> <td>A</td> <td>N</td> <td>D</td> <td>Y</td> <td>O</td> <td>U</td> </tr> <tr> <td>C</td> <td>H</td> <td>A</td> <td>N</td> <td>Q</td> <td>E</td> </tr> <tr> <td>Y</td> <td>O</td> <td>U</td> <td>R</td> <td>W</td> <td>O</td> </tr> <tr> <td>R</td> <td>L</td> <td>D</td> <td>X</td> <td>Y</td> <td>Z</td> </tr> </table> <p>C.T = AUGDAUDNRHYNRX GTTQGWYCYOACIRHOUNHOL EHSUEOZ</p> <p>double encryption by Rail fence</p> <p>C.T = AGADRYRGTGYVAYHUHLHEO UDUNHNXTOWCOCRONOES EZ</p>	4	5	1	2	3	6	C	H	A	N	Q	E	Y	O	U	R	T	H	O	U	Q	H	T	S	A	N	D	Y	O	U	C	H	A	N	Q	E	Y	O	U	R	W	O	R	L	D	X	Y	Z	<p>6m</p> <p>4m</p> <p>10</p>
4	5	1	2	3	6																																													
C	H	A	N	Q	E																																													
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6)	<p>P.T = CRYPTOGRAPHY</p> <p><u>Caesar Cipher</u> $K = 7$</p> <p><u>Encryption</u>: $C = (K + P) \bmod 26$ - 1m</p> <p>$= (7 + 2) \bmod 26 = 9 = J$</p> <p>$= (7 + 17) \bmod 26 = 24 = Y$ 4m</p> <p>\vdots</p> <p>$= (7 + 24) \bmod 26 = 5 = F$</p> <p>C.T = JYFNAVN YHWO</p> <p><u>Decryption</u>: $P = (C - K) \bmod 26$ - 1m</p> <p>$= (9 - 7) \bmod 26 = 2 = C$</p> <p>$= (24 - 7) \bmod 26 = 17 = R$ 5m</p> <p>\vdots</p> <p>$= (5 - 7) \bmod 26 = 24 = Y$</p> <p>P.T = <u>Cryptography</u></p>	<p>5m</p> <p>10m</p>

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7)	<p>Take any of the Example and clearly show one to one mapping</p> <p>Ex: p.t = Hello ↓ ↓ key = ABCCD — 6m</p> <p>c.t = HFNNR</p> <p><u>Strengths of monoalphabetic cipher</u></p> <p>① Since we get 4×10^{26} possible combination of key it is difficult to break cipher text by using — 4m</p> <p>Brute force attack</p> <p>② If the cryptanalyst knows the nature of plaintext then he he can exploit the regularities of the language</p>	<p>10m</p>