

# The $4/\delta$ Bound

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Designing Predictable LLM-Verifier Systems for Formal Method Guarantee  
<https://arxiv.org/abs/2512.02080>

# Probabilistic AI & Formal Verification: Bridging the Gap

We've released a technical report introducing a mathematical foundation with provable guarantees for LLM-based formal verification workflows.

**The Problem:** LLMs show great promise for automating tasks like:

- invariant generation
- proof synthesis
- counterexample-guided repair

# Probabilistic AI & Formal Verification: Bridging the Gap

But today's approaches lack theoretical guarantees. Without convergence bounds, LLM-driven loops may:

- run indefinitely
- waste compute
- require ad-hoc timeouts
- behave unpredictably in CI/CD pipelines

For safety-critical verification, this is unacceptable.

# Probabilistic AI & Formal Verification: Bridging the Gap

## Our Contribution

We present the **LLM-Verifier Convergence Theorem**. To our knowledge, the first formal convergence guarantee for iterative LLM-verification pipelines.

For sequential loops of the form: **propose** → **verify** → **repair** → **retry**

We prove that the system:

- terminates with high probability for any  $\delta > 0$
- converges in expected time  $E[n] \leq 4/\delta$
- satisfies exponential tail bounds

# Probabilistic AI & Formal Verification: Bridging the Gap

Here,  $\delta$  represents the probability that each LLM step makes useful progress, directly linking model quality to runtime guarantees.

## Empirical Validation

- Across 90,000+ trials, observed behavior closely matches theory, showing the bounds are not just asymptotic, they're practical.

# Probabilistic AI & Formal Verification: Bridging the Gap

## Why This Matters for Industry

This turns LLM-assisted verification from guesswork into engineering:

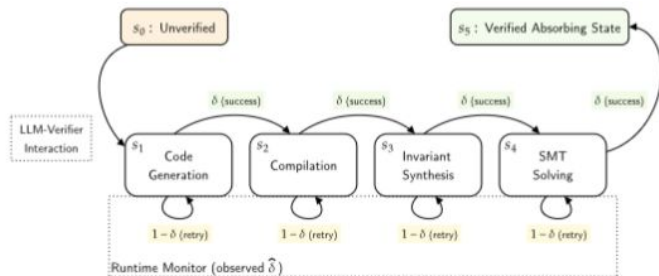
- predictable resource planning for CI/CD
- principled timeout configuration
- performance budgeting based on model capability ( $\delta$ )
- clear deployment regimes: marginal, practical, and high-performance

# Theoretical Background (Absorbing Markov Chain)

The **Theorem** is based on the mathematics of **Absorbing Markov Chains** and the **Geometric Distribution**.

The LLM-Verifier system is a five-state absorbing Markov Chain:

- Transient States (T): Four verification pipeline stages:  $s_1$  (CodeGen),  $s_2$  (Compilation),  $s_3$  (InvariantSynth), and  $s_4$  (SMTSolving).
- Absorbing State (A): One final, irreversible state,  $s_5$  (Verified).



# Theoretical Background (Absorbing Markov Chain)


The Markovian Property holds because the probability of generating a correct verification artifact (the transition probability) depends only on the current state and is independent of the history of previous attempts (memoryless).



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## Requirements for a Geometric Distribution

1. **Binary Outcomes:** Each trial has only two outcomes: success or failure.
2. **Independence:** Each trial is independent of the others.
3. **Constant Probability:** The probability of success ( $p$ ) remains the same for every trial.
4. **Goal:** The experiment stops after the *first* success is achieved. 

# Theoretical Background (Absorbing Markov Chain)

## Convergence Time Derivation (The $4/\delta$ Bound)

The time spent in each transient state,  $M_j$  (the residence time), is modeled by a Geometric random variable.

The **Expected Single-State Time** (time to exit state  $j$ ) is the expected value of the Geometric distribution:  $E[M_j] = 1/\delta$ .

The Total **Expected Convergence Time** is the sum of the independent residence times spent in the four sequential transient states:

$$E[n] = \sum_j (1^4) E[M_j] = 1/\delta + 1/\delta + 1/\delta + 1/\delta = 4/\delta$$

# Theoretical Background (Absorbing Markov Chain)

This theoretical foundation leads to the main result, providing formal guarantees for the system :

**Almost Sure Convergence:** The system reaches the Verified state ( $\tau < \infty$ ) with probability 1, provided  $\delta > 0$ .

**Expected Iteration Bound:** The mean time to convergence is  $E[\tau] = \delta/4$ .

**Tail Bound:** The probability of excessively long verification runs decays exponentially:  $P(\tau > k) \leq \alpha(1 - \delta)^k$ .

# Threats to Validity

## 1. **Assumption of Parameter Stationarity (internal)**

The most critical limitation is the assumption that the LLM's capability ( $\delta$ ) remains constant throughout the entire refinement process.

## 2. **State Space Limitation (internal)**

The model uses a simplified 5-state structure to capture the engineering pipeline.

# Threats to Validity

## **3. Lack of Correlation Between Stages (Advanced Modeling Gap)**

The model assumes independence between the transient states, allowing the total time  $T$  to be a simple sum of the times spent in each state ( $\tau = \sum M_j$ ).

## **4. Generalizability to Real-World Systems (External)**

The core validation relies on a simulation framework rather than real-world LLM deployments.

# Threats to Validity

## **5. Focus on Iteration Count Over Time (Construct)**

The primary metric of convergence is the iteration count ( $E[n]$ ), not the actual time spent.

## **6. Binary Success Definition (Construct Validity Threat)**

The model uses a simple binary definition for progress.

**Theorem 1** (LLM-Verifier Convergence Theorem) We model the LLM-verifier process using a discrete-time Markov Chain, denoted as  $X = \{X_n\}_{n \geq 0}$ . The state space  $S$  consists of a sequential engineering pipeline ( $S = T \cup A$ ). First,  $T = \{s_1, s_2, s_3, s_4\}$  represents the set of transient pipeline stages ( **CodeGen** , **Compilation** , **InvariantSynth** , **SMTSolving** ). Second,  $A = \{s_5\}$  acts as the single absorbing state ( **Verified** ). Assuming a fixed success probability  $\delta \in (0, 1]$  for passing any single stage, we construct the transition matrix  $P = (P_{i,j})$ . The individual entries  $P_{i,j} = \mathbb{P}(X_{n+1} = s_j \mid X_n = s_i)$  are arranged as follows:

(i) **Transient Pipeline States:** For the transient states  $i \in \{1, 2, 3, 4\}$ , the transition probabilities are:

$$P_{i,j} = \begin{cases} \delta & \text{if } j = i + 1 \quad (\text{success: advance to next stage}), \\ 1 - \delta & \text{if } j = i \quad (\text{failure: retry current stage}), \\ 0 & \text{otherwise.} \end{cases}$$

(ii) **Absorbing State:** State  $s_5$  is absorbing, meaning the verification is complete. Hence,  $P_{5,5} = 1$ .

Let  $\tau = \inf\{n \geq 0 : X_n \in A\}$  be the iteration count until verification. The following bounds hold:

1. **Almost Sure Convergence:**

$$\mathbb{P}(\tau < \infty \mid X_0 \in T) = 1$$

2. **Expected Iteration Bound:** The mean time to convergence is given by:

$$\mathbb{E}[\tau \mid X_0 = s_1] = \frac{4}{\delta}$$

3. **Tail Bound:** Consider constants  $\alpha > 0$  and  $\lambda_Q \in (0, 1)$ , where  $\lambda_Q$  is the spectral radius of the transient submatrix  $Q$ . Then, for all  $k \geq 0$ :

$$\mathbb{P}(\tau > k \mid X_0 = s_1) \leq \alpha \lambda_Q^k.$$

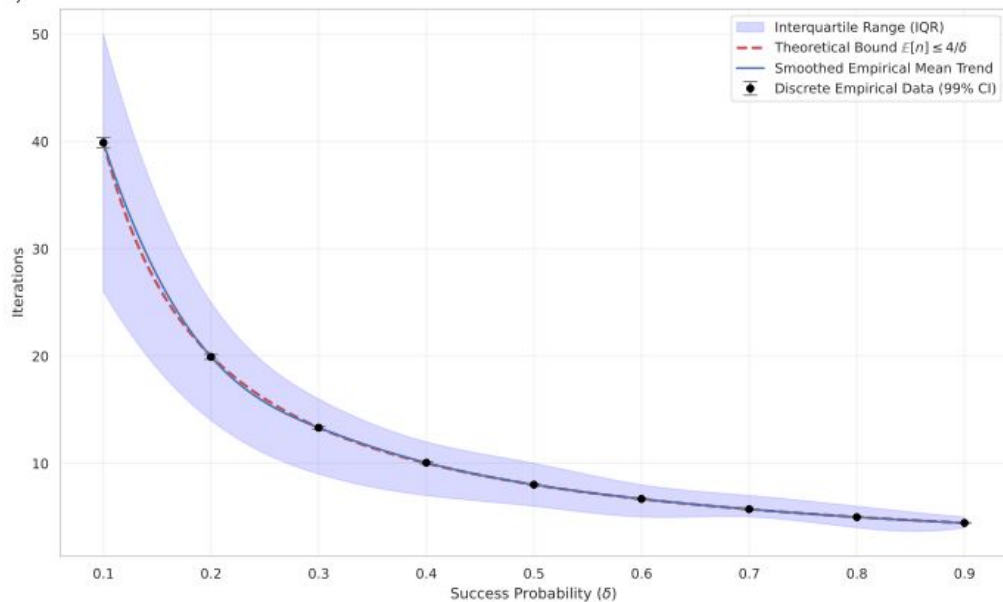
Substituting  $\lambda_Q = 1 - \delta$  results in the following exponential bound:

$$\mathbb{P}(\tau > k \mid X_0 = s_1) \leq \alpha (1 - \delta)^k.$$

*Proof* We analyze the absorbing Markov Chain structure to establish the three guarantees.

# Empirical Validation

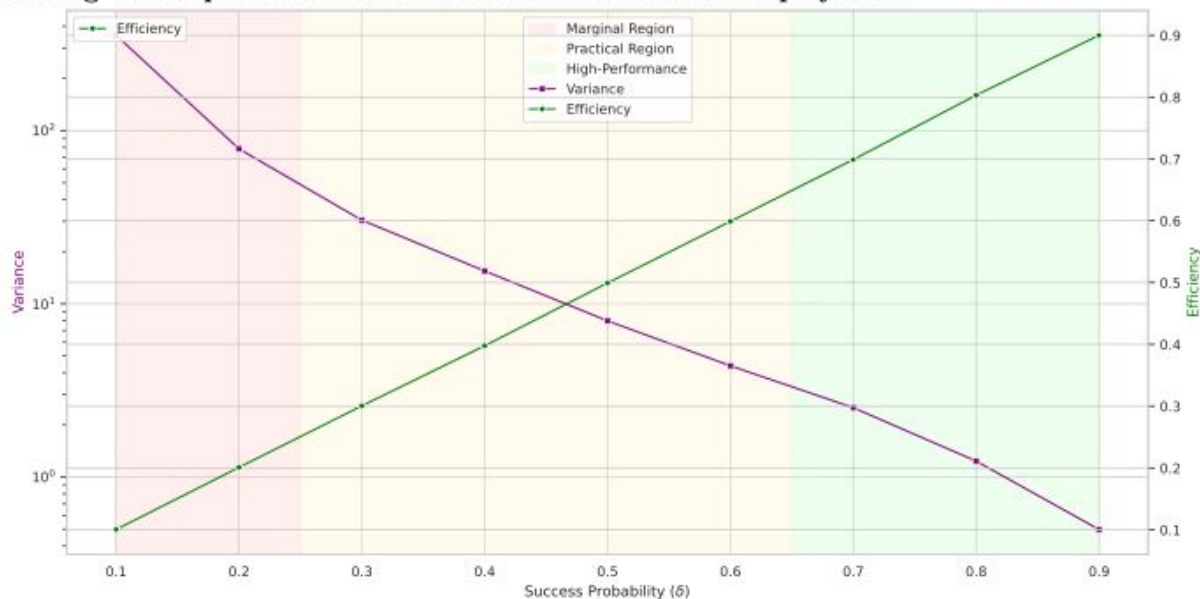
**Fig. 6 Theoretical Bound Alignment and Empirical Convergence Rate.** The figure compares the theoretical expected bound ( $\mathbb{E}[n] \leq 4/\delta$ ) against the empirical mean ( $\mu$ ) across the  $\delta$  spectrum. The extremely close tracking between the empirical curve (blue) and the theoretical curve (red dashed) demonstrates the tight alignment ( $C_f \approx 1.0$ ) of the model prediction. The blue shaded area represents the IQR ( $P_{25} - P_{75}$ ), which visually captures the dramatic decrease in system variance ( $\sigma^2$ ) when transitioning from the marginal region ( $\delta < 0.3$ ) to the high-performance region ( $\delta > 0.6$ )





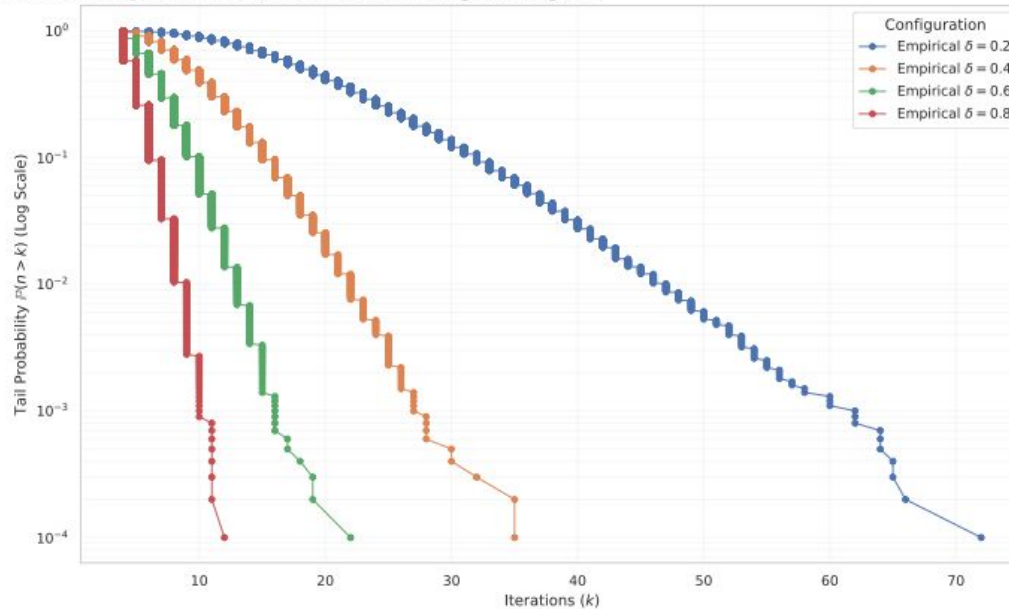
# Empirical Validation

**Fig. 7 Operational Regions Map: Performance and Stability Analysis.** This dual-axis visualization maps the framework's behavior across the  $\delta$  spectrum. The left axis (purple, log scale) tracks the empirical variance ( $\sigma^2$ ), demonstrating the system's predictability. The right axis (green, linear scale) tracks iteration efficiency ( $\eta$ ). The plot clearly illustrates the sharp phase transition in stability: variance collapses rapidly upon entering the practical region ( $\delta \geq 0.3$ ), confirming the empirical boundaries for safe and efficient deployment



# Empirical Validation

**Fig. 8 Tail Probability Analysis: Empirical Validation of Exponential Decay.** The figure plots the Complementary Cumulative Distribution Function (CCDF,  $P(n > k)$ ) on a log-linear scale. The approximately linear decay observed across all tested  $\delta$  values confirms the fundamental assumption that the convergence time follows an exponential tail behavior ( $\mathbb{P}(n > k) \propto e^{-c\delta k}$ ). This validation ensures that the system's worst-case convergence time is reliably bounded and predictable, even in the marginal region



Adaveri

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# Adaveri Architecture

**AdaVeri** is an adaptive, AI-driven framework and real-time control system that enhances the speed, reliability, and safety of formal software verification, especially in safety-critical applications.

It monitors the verification process, adjusting strategies based on success rates, with a mathematical guarantee of convergence for efficient and reliable verification.

# Adaveri Architecture

