

9.1

Using that $S = \lambda f.\lambda n.\lambda x.fn(fx)$:

$$\begin{aligned} S(\lambda x.m)(\lambda x.n) &= (\lambda f.\lambda n.\lambda x.fn(fx))(\lambda x.m)(\lambda x.n) \\ &= (\lambda f.\lambda x.f((\lambda x.m)fx))(\lambda x.n) \\ &= \lambda x.(\lambda x.m)x((\lambda x.n)x) \\ &= \lambda x.m((\lambda x.n)x) = \lambda x.mn \end{aligned}$$

9.2

$$\text{xor} = \lambda b. \lambda c. b(\text{not } c)c$$

It is commutative if $\text{xor true false} = \text{xor false true}$

Checking:

$$\text{xor true true} = \beta \text{ true (not true) true} = \beta \text{ false}$$

$$\text{xor true false} = \beta \text{ true (not false) false} = \beta \text{ true}$$

$$\text{xor false true} = \beta \text{ false (not true) true} = \beta \text{ true}$$

$$\text{xor false false} = \beta \text{ false (not false) false} = \beta \text{ not true} = \beta \text{ false}$$

As we can see, it is commutative.

9.3 Problem

$$\text{pred} = \lambda n \lambda f \lambda x. (n(\lambda g(\lambda h. h(gf))))(\lambda u. x)(\lambda u. u)$$

Here we can see that $\lambda u. x$ is const_x .

$\lambda u. u$ is the identity function id .

Looking at this part:

$$\lambda g(\lambda h. h(gf))$$

Say that $\text{evaluator}_y == \lambda h. hy$

$$F = \lambda g. \text{evaluator}_{g(f)}$$

Now, $F^n \text{const}_x = \text{evaluator}_{f^{n-1}(x)}$

$$\text{pred} = \lambda n f x. F^n \text{const}_x \text{id} = \lambda n f x. \text{evaluator}_{f^{n-1}(x)} \text{id} = \lambda n f x. \text{id}(f^{n-1}(x)) = \lambda n f x. f^{n-1}(x)$$

Now, for the proof:

$$\begin{aligned} \text{pred}(\text{succ } m) &= \text{pred}(\lambda f. \lambda n. \lambda x. f^n(fx)m) = \text{pred } (m + 1) = \\ &= \lambda n \lambda f \lambda x. (n(\lambda g(\lambda h. h(gf))))(\lambda u. x) \text{id}(m + 1) = \\ &= \lambda f \lambda x. (m + 1)(\lambda g(\lambda h. h(gf))))(\lambda u. x) \text{id} = \\ &= \lambda f \lambda x. (\lambda h. f(\dots f(x))) \text{id} = m \end{aligned}$$

9.4 Problem

Factorial is defined as follows:

$$\text{fact } n = \text{if (iszro } n) 1 \text{ mult } n (\text{fact (pred } n))$$

$$3 = \lambda sz. sssz$$

Hence,

$$\begin{aligned} \text{fact } 3 &= \text{if (iszro } 3) 1 \text{ mult } 3 (\text{fact (pred } 3)) = \text{mult } 3 (\text{fact (pred } 3)) = \\ &= \text{mult } 3 (\text{fact } 2) = \text{mult } 3 (\text{if (iszro } 2) 1 \text{ mult } 2 (\text{fact (pred } 2))) = \text{if (iszro } 3) 1 \text{ mult } 3 (\text{fact (pred } 3)) = \\ &= \text{mult } 3 (\text{fact (pred } 3)) = \text{mult } 3 (\text{fact } 2) = \text{mult } 3 (\text{mult } 2 (\text{fact (pred } 2))) = \text{mult } 3 (\text{mult } 2 (\text{fact } 1)) = \\ &= \text{mult } 3 (\text{mult } 2 (\text{if (iszro } 1) 1 \text{ mult } 1 (\text{fact (pred } 1)))) = \text{mult } 3 (\text{mult } 2 (\text{mult } 1 (\text{fact (pred } 1)))) = \text{mult } 3 (\text{mult } 2 (\text{fact } 0)) = \\ &= \text{mult } 3 (\text{mult } 2 (\text{mult } 1 (\text{if (iszro } 0) 1 \text{ mult } 0 (\text{fact (pred } 0))))) = \text{mult } 3 (\text{mult } 2 (\text{mult } 1 (1))) = \\ &= \text{mult } 3 (\text{mult } 2 (\text{mult } 1 1)) = \text{mult } 3 (\text{mult } 2 1) = \text{mult } 3 2 = 6 \end{aligned}$$

9.5 Problem

$$\text{mult } m \ n = \lambda m. \lambda n. \lambda f. \lambda x. n(mf)x$$

means the same as "apply s" n times. Thus, when we write m (n s) z, we're adding n to zero m times (i.e., m times).

Now to prove that 1 is the unit of multiplication:

$$\text{mult } 1 \ n = \lambda f. \lambda x. (\lambda f x. (fx))(nf)x = n$$