

CSC421 A3

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Q1.

$$Q2. \quad (p \mid q \mid r) \& ((\neg r \mid q \mid p) \rightarrow ((r \mid q) \& (\neg q \mid \neg p)))$$

Show unsatisfiable

1. (工)

$$(p \vee q \vee r) \wedge (\neg(r \vee q \vee p) \vee ((r \vee q) \wedge \neg q \wedge \neg p))$$

(Negation, distributive)

$$(p \vee q \vee r) \wedge ((r \wedge \neg q \wedge \neg p) \vee ((\neg r \vee q) \wedge \neg r \wedge p))$$

$$(p \mid q \mid r) \wedge ((r \wedge q \wedge \neg p) \mid (r \mid q)) \wedge ((r \wedge q \wedge \neg p) \mid \neg q) \wedge ((r \wedge \neg q \wedge \neg p) \mid \neg p)$$

(Distributive Operators)

$$\textcircled{1} \quad (p \wedge q \wedge r) \wedge$$

$$\textcircled{2} \quad ((r \mid r \mid q) \wedge (\neg q \mid r \mid q) \wedge (\neg p \mid r \mid q)) \wedge$$

$$\textcircled{3} \quad \left\{ (r | \gamma_q) \delta (\gamma_q | \gamma_q) \delta (\gamma_p | \gamma_q) \right\} \delta$$

$$\textcircled{4} \quad ((r | \gamma_p) \wedge (\gamma_q | \gamma_p) \wedge (\gamma_p | \gamma_p))$$

Distributive, remove brackets

$(p \mid q \mid r) \wedge (r \mid r \mid q) \wedge (\neg q \mid r \mid q) \wedge (\neg p \mid r \mid q) \wedge (r \mid \neg q) \wedge (\neg q \mid \neg q) \wedge (\neg p \mid \neg q)$
 $\wedge (r \mid p) \wedge (\neg q \mid \neg p) \wedge (\neg p \mid \neg p)$

↑
clausal form

Derive empty clause using resolution

1. $(p \mid q \mid \neg r)$ Premise
2. $(r \mid \neg r \mid q)$ "
3. $(\neg q \mid \neg r \mid q)$ "
4. $(\neg p \mid \neg r \mid q)$ "
5. $(\neg r \mid \neg q)$ "
6. $(\neg q \mid \neg r \mid \neg q)$ "
7. $(\neg p \mid \neg r \mid \neg q)$ "
8. $(\neg r \mid \neg p)$ "
9. $(\neg q \mid \neg r \mid \neg p)$ "
10. $(\neg p \mid \neg r \mid \neg p)$ "

-
11. $\{\neg p\}$ 1, 5
 12. $\{\neg r, \neg q\}$ 2, Negated conclusion, Idempotent law
 13. $\{\neg r\}$ 4, 12
 14. $\{\}$ 11, 13

Q2.

a) First order logic:

all x all y (Horse(x) & Dog(y) \Rightarrow Faster(x, y)).
exists y (Greyhound(y) & all z (Rabbit(z) \Rightarrow Faster(y, z))).
all y (Greyhound(y) \Rightarrow Dog(y)).
all x all y all z (Faster(x, y) & Faster(y, z) \Rightarrow Faster(x, z)).

Conclusion

all x all y (Horse(x) & Rabbit(y) \rightarrow Faster(x, y)).

b) First order logic:

All x (Hummingbird(x) => bird(x))
All x (Hummingbird(x) => richlyColored(x))
All x (livesOnHoney(x) => ~largeBird(x))
All x (richlyColored(x) => livesOnHoney(x))

Conclusion:

All x (Hummingbird(x) => richlyColored(x)) AND
All x (richlyColored(x) => livesOnHoney(x)) AND
All x (livesOnHoney(x) => ~largeBird(x))
By law of syllogism
All x (Hummingbird(x) => ~largeBird(x))

c) First order logic:

All x (myGardner(x) => worth_listening_military(x))
All x (remember_waterloo(x) => old(x))
All x (worth_listening_military(x) => remember_waterloo(x))

Conclusion:

All x (myGardner(x) => worth_listening_military(x)) AND
All x (worth_listening_military(x) => remember_waterloo(x)) AND
All x (remember_waterloo(x) => old(x))
By law of syllogism
All x (myGardner(x) => old(x))

Q3.

$$P(p_{13} | b_{12}, b_{21}) :$$

$$= \propto \sum_{p_{22}} \sum_{p_{31}} P_r(b_{12} | p_{22}, p_{13}) P_r(b_{21} | p_{22}, p_{31}) P_r(p_{13}) P_r(p_{22}) P_r(p_{31})$$

$$= \propto [P_r(b_{12} | p_{13}, p_{22}) \cdot P(b_{21} | p_{31}, p_{22}) \cdot P(p_{13}) \cdot P(p_{22}) \cdot P(p_{31}) + \\ P(b_{12} | p_{13}, \neg p_{22}) \cdot P(b_{21} | p_{31}, \neg p_{22}) \cdot P(p_{13}) \cdot P(\neg p_{22}) \cdot P(p_{31}) + \\ P(b_{12} | p_{13}, p_{22}) \cdot P(b_{21} | \neg p_{31}, p_{22}) \cdot P(p_{13}) \cdot P(p_{22}) \cdot P(\neg p_{31}) + \\ P(b_{12} | p_{13}, \neg p_{22}) \cdot P(b_{21} | \neg p_{31}, \neg p_{22}) \cdot P(p_{13}) \cdot P(\neg p_{22}) \cdot P(\neg p_{31})]$$

$$= \propto [1 \cdot 1 \cdot 0.01 \cdot 0.01 \cdot 0.01 + \\ 1 \cdot 1 \cdot 0.01 \cdot 0.99 \cdot 0.01 + \\ 1 \cdot 1 \cdot 0.01 \cdot 0.01 \cdot 0.99 + \\ 0 \cdot 0 \cdot 0 \cdot 0 \cdot 0]$$

$$= \propto [0.000001 + \\ 0.000099 + \\ 0.000099 + \\ 0.000000]$$

$$= 0.000199$$

$$P(P_{22} | b_{12}, b_{21}) :$$

$$= \propto \sum_{P_{13}} \sum_{P_{31}} P(b_{12} | P_{13}, P_{22}) P(b_{21} | P_{22}, P_{31}) \cdot P(P_{13}) \cdot P(P_{22}) \cdot P(P_{31})$$

$$= \propto [P(b_{12} | P_{13}, P_{22}) \cdot P(b_{21} | P_{22}, P_{31}) \cdot P(P_{13}) \cdot P(P_{22}) \cdot P(P_{31}) + \\ P(b_{12} | \neg P_{13}, P_{22}) \cdot P(b_{21} | P_{22}, P_{31}) \cdot P(\neg P_{13}) \cdot P(P_{22}) \cdot P(P_{31}) + \\ P(b_{12} | P_{13}, \neg P_{22}) \cdot P(b_{21} | \neg P_{22}, P_{31}) \cdot P(P_{13}) \cdot P(P_{22}) \cdot P(\neg P_{31}) + \\ P(b_{12} | \neg P_{13}, \neg P_{22}) \cdot P(b_{21} | \neg P_{22}, \neg P_{31}) \cdot P(\neg P_{13}) \cdot P(P_{22}) \cdot P(\neg P_{31})]$$

$$= \propto [1 \cdot 1 \cdot 0.01 \cdot 0.01 \cdot 0.01 + \\ 1 \cdot 1 \cdot 0.99 \cdot 0.01 \cdot 0.01 + \\ 1 \cdot 1 \cdot 0.01 \cdot 0.01 \cdot 0.99 + \\ 1 \cdot 1 \cdot 0.99 \cdot 0.01 \cdot 0.99]$$

$$= \propto [0.000001 + \\ 0.000099 + \\ 0.000099 + \\ 0.009801]$$

$$= \propto 0.01$$

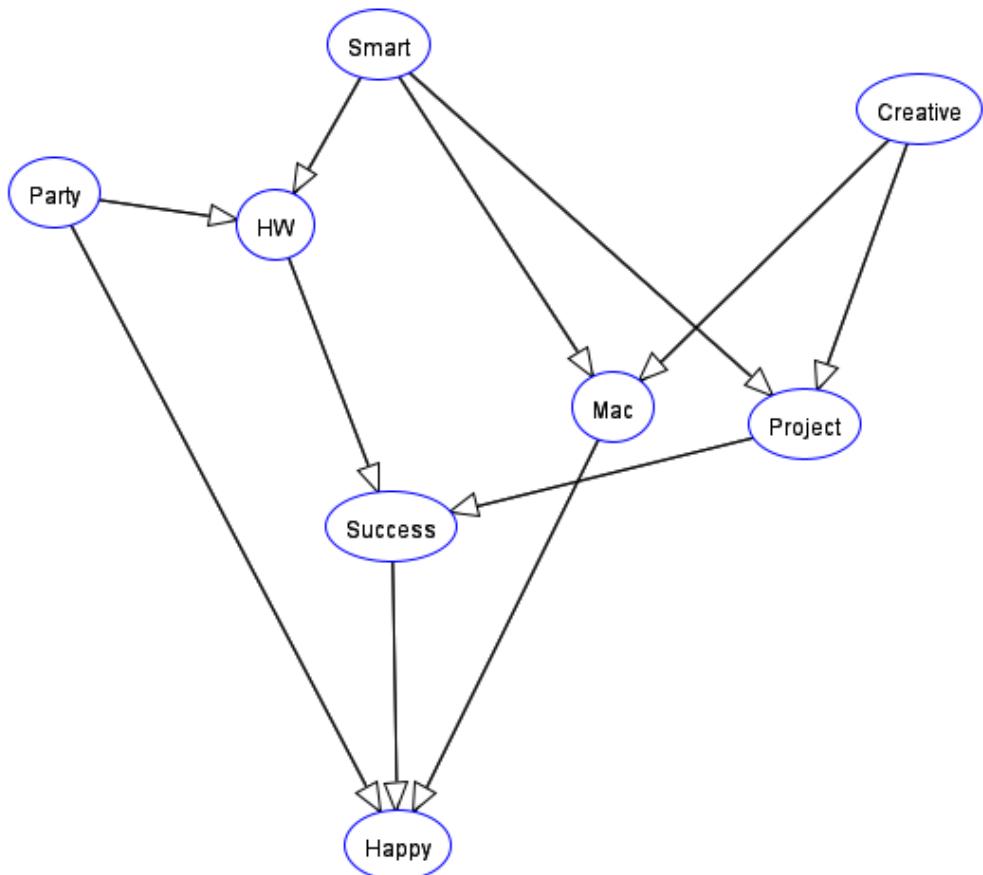
The probability of $P(p_{13} | b_{12}, b_{21})$ ends up to be quite low ($a * 0.000199$) whereas $P(p_{22} | b_{12}, b_{21})$ ends up quite higher ($a * 0.01$).

With a logical agent, all squares look the same, Therefore, the logical agent would choose either one with equal chance ($\frac{1}{3}$) of being chosen. This means the agent will die a third of the time.

With a probabilistic agent, it will never choose to go [2,2], it will always choose [1,3] or [3,1] based on probability.

Q4.

i) Draw the bayesian network



ii) Estimate the probabilities of the conditional probability tables using the data provided

$P(HW | party, smart)$

Party	Smart	cnt1	cnt2	$p(hw party, smart)$
0	0	610	186	0.304918
0	1	1379	1239	0.898477
1	0	866	81	0.093533
1	1	2145	1722	0.802797

P(Mac | smart, creative)

Creative	Smart	cnt1	cnt2	P(Mac creative, smart)		
0	0	436	53	0.12156		
0	1	1067	441	0.413308		
1	0	1040	933	0.897115		
1	1	2457	1685	0.685796		

P(Project | smart, creative)

Creative	Smart	cnt1	cnt2	P(Project creative, smart)		
0	0	436	46	0.105505		
0	1	1067	847	0.793814		
1	0	1040	419	0.402885		
1	1	2457	2224	0.905169		

P (Success | HW, project)

HW	Project	cnt1	cnt2	p(Success HW, project)		
0	0	906	45	0.049669		
0	1	866	179	0.206697		
1	0	558	171	0.306452		
1	1	2670	2394	0.896629		

P(Happy | party, success, mac)

Party	Success	Mac	cnt1	cnt2	P(Happy party, success, mac)		
0	0	0	309	29	0.093851		
0	0	1	483	99	0.204969		
0	1	0	440	135	0.306818		
0	1	1	757	271	0.357992		
1	0	0	507	213	0.420118		
1	0	1	912	449	0.492325		
1	1	0	632	456	0.721519		
1	1	1	960	921	0.959375		

iii)

note* wouldn't let me enter past 5 decimal places for probability tables. IE answers might be slightly off.

3.

$$p(\text{Happy} \mid \text{party, smart, } \neg \text{creative})$$

$$= a \sum_{Hw} \sum_{mac} \sum_{success} \sum_{\text{project}} p(Hw \mid p, s) p(Mac \mid s, \neg c) p(\text{Proj} \mid s, \neg c) \\ P(\text{success} \mid Hw, \text{proj}) \cdot P(\text{Happy} \mid p, \text{success, proj}) \cdot P(p) \cdot P(s) \cdot P(\neg c)$$

$$= a \cdot (\dots)$$

$$p(\neg \text{Happy} \mid p, s, \neg c)$$

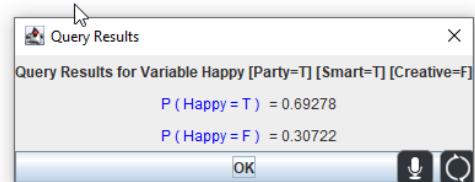
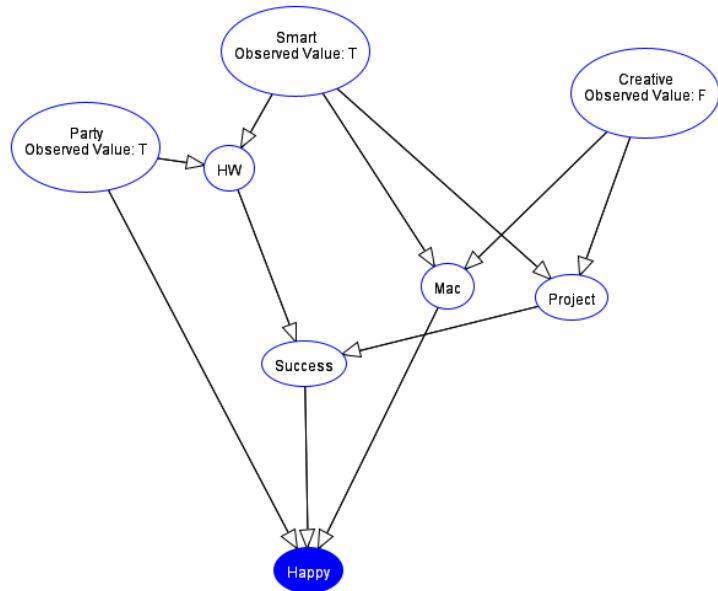
$$= a \sum_{Hw} \sum_{mac} \sum_{success} \sum_{\text{project}} p(Hw \mid p, s) p(Mac \mid s, \neg c) p(\text{Proj} \mid s, \neg c) \\ P(\text{success} \mid Hw, \text{proj}) \cdot P(\text{Happy} \mid p, \text{success, proj}) \cdot P(p) \cdot P(s) \cdot P(\neg c)$$

$$= a \cdot (\dots)$$

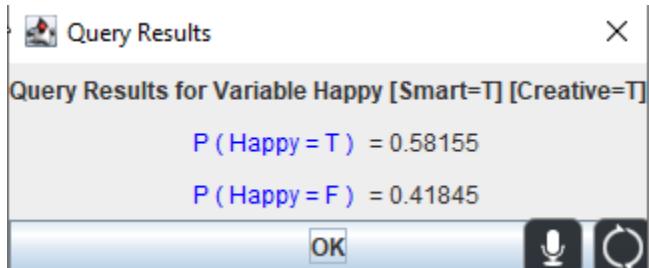
$$\alpha = 1 / (P(h \mid p, s, \neg c) + P(\neg h \mid p, s, \neg c))$$

using Aitool

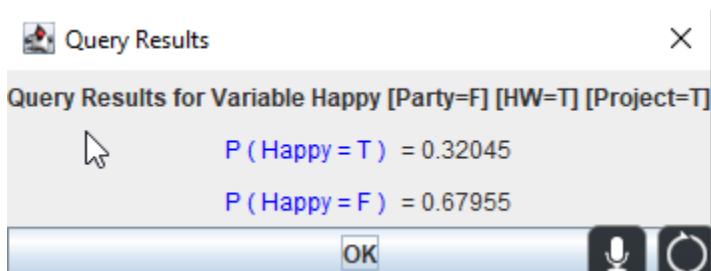
$$P(\text{Happy} \mid p, s, \neg c) = 0.69278$$



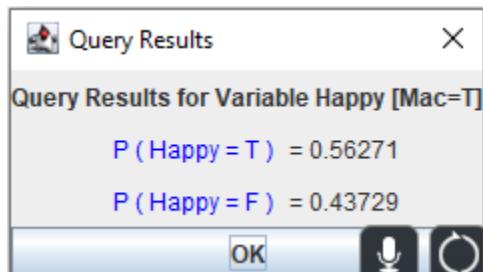
iv) Smart and Creative



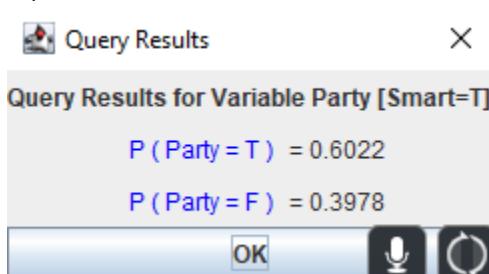
v) Doesn't party, does his homework and project



vi) Own's a mac



vii) Wicked smart



viii)

