Diff Trace: Efficient Whole-Program Trace Analysis and Diffing

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ABSTRACT

Abstract to be written

KEYWORDS

diffing, tracing, debugging

ACM Reference Format:

1 INTRODUCTION

When the next version of an HPC software system is created, logical errors often get introduced. To maintain productivity, designers need effective and efficient methods to locate these errors. Given the increasing use of hybrid (MPI + X) codes and library functions, errors may be introduced through a usage contract violation at any one of these interfaces. Therefore, tools that record activities at multiple APIs are necessary. Designs find most of these bugs manually, and the efficacy of a debugging tool is often measured by how well it can highlight the salient differences between the executions of two versions of software. Given the huge number of things that could be different - individual iterative patterns of function calls, groups of functions calls, or even specific instruction types (e.g., non-vectorized versus vectorized floating-point dot vector loops) - designers cannot often afford to rerun the application multiple times to collect each facet of behavior separately. These issues are well summarized in many recent studies [20]

One of the major challenges of HPC debugging is the huge diversity of the applications, which encompass domains such as computational chemistry, molecular dynamics, and climate simulation. In addition, there are many types of possible åĂIJbugsâĂİ or, more precisely, errors. An **error** may be a deadlock or a resource leak.

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These errors may be caused by different **faults**: an unexpected message reordering rule (for a deadlock) or a forgotten free statement (for a resource leak). There exists a collection of scenarios in which a bug can be introduced: when developing a brand new application, optimizing an existing application, upscaling an application, porting to a new platform, changing the compiler, or even changing compiler flags. Unlike traditional software, there are hardly any bug-repositories, collection of trace data or debugging-purpose benchmarks in HPC community. The heterogeneous nature of HPC bugs make developers come up with their own solutions to resolve specific class of bugs on specific architecture or platforms that are not usable on other [20].

When a failure occurs (e.g., deadlock or crash) or the application outputs an unexpected result, it is not economic to rerun the application and consume resources to reproduce the failure. In addition, HPC bugs might not be reproducible due to non-deterministic behavior of most of HPC applications. Also the failure might be caused by a bug present at different APIs, system levels or network, thus multiple reruns might be needed to locate the buggy area. In our previous work[46], we have introduced ParLOT that collects whole program function call traces efficiently using dynamic binary instrumentation. ParLOT captures function calls and returns at different levels (e.g., source-code and library) and incrementally compress them on-the-fly, resulting in low runtime overhead and significantly low required bandwidth, preserving the whole-program dynamic behavior for offline analysis.

In the current work, we introduce DiffTrace, a tool-chain that post-mortem analyze ParLOT traces in order to supply developers with information about dynamic behavior of HPC applications at different levels towards debugging. Topology of HPC tasks on both distributed and shared memory often follow a "symmetric" control flow such as SPMD, master/worker and odd/even where multiple tasks contain similar events in their control flow. HPC bugs often manifest themselves as divergence in the control flow of processes comparing to what was expected. In other words, HPC bugs violates the rule of "symmetric" and "similar" control flow of one or more thread/process in typical HPC applications based on the original topology of the application. We believe that finding the dissimilarities among traces is the essential initial step towards finding the bug manifestation, and consequently the bug root cause. Large-scale HPC application execution would result in thousands of ParLOT trace files due to execution of thousands of processes and threads. Since HPC applications spend most of their time in an outer main loop, every single trace file also may contain million-long sequence of trace entries (i.e., function calls and returns). Finding the bug manifestation (i.e., dissimilarities caused by the bug) among large number of long traces is the problem of finding the needle in the havstack.

Decompressing ParLOT traces collected from long-running largescale HPC applications for offline analysis produce overwhelming amount of data. However, missing any piece of collected data may result in loosing key information about the application behavior. We propose a variation of NLR (Nested Loop Recognition) algorithm [26] that takes a sequence of trace entries as input and by recursively detecting repetitive patterns, re-compresses traces into "iterative sets" in a lossless fashion (intra-trace compression). Analyzing the application execution as a whole is another goal that we are pursuing in this work. By extracting attributes from pre-processed traces, we inject them a concept hierarchy data structure called Concept Lattice [16]. Concept lattices give us the capability of reducing the search space from thousands of instances to just a few equivalent behavior classes of traces by measuring the similarity of traces[3]

** TODO: Highlights of results obtained as a result of the above thinking should be here. THis typically comes before ROADMAP

In summary, here are our main contributions:

- we have a powerful combination of ideas to locate bugs
- A variation of NLR algorithm to compress traces in lossless fashion for easier analysis and detecting (broken) loop structures
- FCA-based clustering of similar behavior traces, efficient,
- Ranking system based on delta-sim
- Visualization diffNLR

The rest of the paper is as follows:

- Sec 2 talks about backgrounds
- Sec 3: Talks about the design of diffTrace
- Sec 4: Evaluates diffNLR by initial recommendation turns into NLR-observable bug and confusion matrix scores how many TP,FP,TN,FN 2.3 Loop structure detection
- Sec 5 major related work will have a detailed discussion of related work
- Sec 6 discusses the potential limitations, future work, and conclusion.

BACKGROUND AND RELATED WORK

There are two major phases in any "Program Understanding" tool: data collection and data analysis. To understand the runtime behavior of applications, an efficient tracing mechanism is required to collect informative data during execution of the application. Par-LOT collects whole program function call traces with the mindset of paying a little upfront and save resource and time cost of reproducing the bug.

- detection of loops, lossless Compression, easier to analyze, information about loop structures, how many iteration each loop executed matters, broken loop matters
- Concept lattices, throw inter-process information, similarity, clustering
- diffNLR visualizes pairs of traces and reflect their differences where they are supposed to be equal.

2.1 Parlot Summary (changeme)

- why binary instrumentation? no need to modify the code to instrument - no need to recompile, unlike most of HPC tools (performance tools like TAU[44], Score-p[27] and debugging tools like STAT[1] and AutomaDeD[30] that requires Dyninst[35] for instrumentation and MRNet[42] and TBON [22])

To overcome the trade-off of comprehensive data collection while adding low time and space overhead, HPC program analysis tools often sacrifice one for the other. However, ParLOT collects wholeprogram function call traces at as low as library level, while incrementally compressing traces on-the-fly and leave majority of the system bandwidth for the application.

Each ParLOT Trace (PT) contain full sequence of function calls and returns for every single thread that running the application code, reflecting the dynamic control flow and call-stack of the application.

2.2 Equivalencing behavior via concept lattices (changeme)

this is how it is cleanly defined (use our context or examples) e.g. high-level: objects to attributes in our case objects can be... and attributes can be how we use them is ...

- Vijay Garg Applications of lattice theory in distributed
- Dimitry Ignatov [?] Concept Lattice Applications in Information Retrieval
- [16] [19] [4] [18] Similarity measurement using FCA [3]
- 2.2.1 Attribute creation... There are many of them (others have not
- 2.2.2 incremental algos. challenges: need incremental algo... end

loop detection has been addressed in xyz challenges in our context are xyz... highlights of what you did (briefly) and why it can help

2.4 diffing (changeme)

what you adapted how does it help [38]

2.5 STAT

Parallel debugger STAT[1]

- STAT gathers stack traces from all processes
- Merge them into prefix tree
- Groups processes that exhibit similar behavior into equivalent classes
- A single representative of each equivalence can then be examined with a full-featured debugger like TotalView or DDT

What STAT does not have?

- FP debugging
- Portability (too many dependencies)
- Domain-specific
- Loop structures and detection

2.6 Program Understanding

- Score-P [27]
- TAU [44]
- ScalaTrace: Scalable compression and replay of communication traces for HPC [40]
- Barrier Matching for Programs with Textually unaligned barriers [49]
- Pivot Tracing: Dynamic causal monitoring for distributed systems Johnathan mace [33]
- Automated Charecterization of parallel application communication patterns [43]
- Problem Diagnosis in Large Scale Computing environments [36]
- Probablistic diagnosis of performance faults in large-scale parallel applications [29]
- detecting patterns in MPI communication traces robert preissl [41]
- D4: Fast concurrency debugging with parallel differntial analysis bozhen liu [32]
- Marmot: An MPI analysis and checking tool bettina krammer [28]
- MPI-checker Static Analysis for MPI Alexandrer droste [15]
- STAT: stack trace analysis for large scale debugging Dorian Arnold [1]
- DMTracker: Finding bugs in large-scale parallel programs by detecting anomaly in data movements [17]
- SyncChecker: Detecting synchronization errors between MPI applications and libraries [11]
- Model Based fault localization in large-scale computing systems Naoya Maruyama [34]
- Synoptic: Studying logged behavior with inferred models ivan beschastnikh [5]
- Mining temporal invariants from partially ordered logs ivan beschastnikh [7]
- Scalable Temporal Order Analysis for Large Scale Debugging Dong Ahn [2]
- Inferring and asserting distributed system invariants ivan beschastnikh stewart grant [21]
- PRODOMETER: Accurate application progress analysis for large-scale parallel debugging - subatra mitra [37]

- Automaded: Automata-based debugging for dissimilar parallel tasks - greg [9]
- Automaded : large scale debugging of parallel tasks with Automaded - ignacio [30]
- Inferring models of concurrent systems from logs of their behavior with CSight - ivan [6]

2.7 Trace Analysis

- Trace File Comparison with a hierarchical Sequence Alignment algorithm [47]
- structural clustering : matthias weber [48]
- building a better backtrace: techniques for postmortem program analysis ben liblit [31]
- automatically charecterizing large scale program behavior timothy sherwood [45]

2.8 Visualizations

- Combing the communication hairball: Visualizing large-scale parallel execution traces using logical time katherine e isaacs [24]
- recovering logical structure from charm++ event traces [23]
- ShiViz Debugging distributed systems [8]

2.9 Concept Lattice and LCA

- Vijay Garg Applications of lattice theory in distributed systems
- Dimitry Ignatov [?] Concept Lattice Applications in Information Retrieval
- [16] [19] [4] [18] [38]

2.10 Repetitive Patterns

• [12] [25] [39] [14] [13]

3 DFFTRACE TOOL ARCHITECTURE

3.1 General Idea

Here is a general overview of DiffTrace and its components

- Motivating example
- Problem statement
- Potential Approaches and Related Work
- Next subsections will explain the components that we have in our framework and the corresponding related work and background

3.2 Fault Injection

3.3 ParLOT

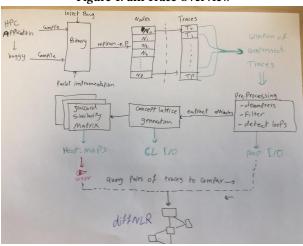
2-3 paragraph explanation about ParLOT and its mechanism [46]

3.4 Filter

Include a table with all filters and their regular expressions

- 3.4.1 General Filters.
 - Returns
 - .plt
 - Memory
 - Network

Figure 1: diffTrace Overview



- Polling
- String
- Customize
- IncludeEverything
- 3.4.2 Target Filters.
 - MPI
 - MPIall
 - MPI_ Collectives
 - MPI_ Send/Recv
 - OMPall
 - OMPcritical
 - OMPmutex

3.5 Nested Loop Recognition

- 3.5.1 Background.
- 3.5.2 Implementation.

3.6 Concept Lattice Analysis

- 3.6.1 Background.
 - FCA (formal concept analysis) background and citations
 - FCA applications in all areas
 - FCA applications in Data Mining and Information Retrieval
 - FCA applications in distributed systems (Garg's work)
 - intro to Concept, Object, Attribute and other definitions

What do we expect to gain by doing so

- Single entity represents the whole execution of HPC application (can be used as signature/model in ML)
- Classifying similar behavior objects(traces)
- Efficient Incremental CL building makes it scalable
- Efficient full pair-wise Jaccard Similarity Matrix extraction
- 3.6.3 CL generation.
 - background

- current approach
- 3.6.4 Jaccard Similarity Matrix.
 - background
 - LCA
 - Benefits

3.7 diffNLR

- motivation
- diff algorithm
- visualization

3.8 FP-Trace

4 EXPERIMENTAL STUDIES

```
5 RESULTS
                                                          //global champion
                                                     47
                                                          MPI AllReduce();
   5.1 Case Study: ILCS-TSP
                                                          // The process with the global champion
   Here is the ILCS framework pseudo-code. User needs to write
                                                          // copy its results to bcast_buffer
   CPU_Init(),CPU_Exec() and CPU_Output().
                                                          if (my_rank == global_champion){
                                                     51
                                                           #pragma omp cirtical
                                                     52
                                                           memcpy(bcast_buffer,local_champ)
                                                     53
   int main(argc, argv){
1
                                                     54
    MPI Init();
2
                                                     55
                                                          //Broadcast the champion
3
    MPI_Comm_size()
                                                     56
                                                          MPI_Bcast(bcast_buffer)
    MPI_Comm_rank(my_rank)
                                                          if(my_rank = = 0)
                                                     57
    // Figuring local number of CPUs
5
                                                           CPU Output (champ)
    MPI_Reduce() // Figuring global number of CPUs
6
7
    CPU_Init();
                                                     60
                                                          MPI_Finalize()
8
    //For storing local champion results
                                                     61
                                                         }
9
    champ[CPUs] = malloc();
                                                     62
10
    MPI_Barrier();
                                                         /* User code for TSP problem */
    #pragma omp parallel num_threads(CPUs+1)
11
                                                     64
12
                                                     65
                                                         CPU Init(){
      rank = omp_get_thread_num()
13
                                                     66
                                                          // Read In data from cities
      if (rank == 0){ //communication thread}
14
                                                          // Calculate distances
                                                     67
15
                                                         // Return data structure to store champion
                                                     68
        //Find and report the thread with
16
                                                     69
                                                         }
        //local champion, global champion
17
                                                     70
        MPI AllReduce();
18
                                                     71
                                                        CPU_Exec(){
19
        //Find and report the process with
                                                         // Find local champions (TSP tours)
                                                     72
        //global champion
20
                                                     73
                                                        }
        MPI_AllReduce();
21
                                                     74
22
        //The process with the global champion
                                                        CPU Output(){
                                                     75
        //copy its results to bcast_buffer
23
                                                     76
                                                          // Output champion
24
        if (my_rank == global_champion){
                                                     77
                                                        }
25
         #pragma omp cirtical
         memcpy(bcast_buffer,local_champ)
26
27
        //Broadcast the champion
28
        MPI_Bcast(bcast_buffer)
29
30
       } while (no_change_threshold);
       cont=0 // signal worker threads to stop
31
      } else{ // worker threads
32
       while (cont) {
33
        // Calculate Seed
34
35
        local result = CPU Exec()
        if (local_result < champ[rank]){</pre>
36
         #pragma omp cirtical
37
         memcpy(champ[rank],local_result)
38
39
        }
40
      }
      }
41
42
    //Find and report the thread with
43
44
    //local champion, global champion
    MPI AllReduce();
45
46
    //Find and report the process with
```

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Table 1 describes the bug that I injected to ILCS-TSP

Table 1: Injected Bugs to ILCS-TSP

ID	Level	Bugs	Description		
1	MPI	allRed1wrgOp-1-all-x	Different operation (MPI_MAX) in only one process (buggyProc = 2) for MPI_ALLREDUCE() in Line 21		
2		allRed1wrgSize-1-all-x	Wrong size in only one process (buggyProc = 2) for MPI_ALLREDUCE() in Line 21		
3		allRed1wrgSize-all-all-x	Wrong Size in all processes for MPI_ALLREDUCE() in Line 21		
4		allRed2wrgOp-1-all-x	Different operation (MPI_MAX) in only one (buggyProc) for first MPI_ALLREDUCE() – L277:ilcsTSP.c		
5		allRed2wrgSize-1-all-x	Wrong size in only one (buggyProc) for first MPI_ALLREDUCE() – L277:ilcsTSP.c		
6		allRed2wrgSize-all-all-x	Wrong Size in all processes for second MPI_ALLREDUCE() – L277:ilcsTSP.c		
7		bcastWrgSize-1-all-x	Wrong Size in only one (buggyProc) of MPI_Bcast() - L290:ilcsTSP.c		
8		bcastWrgSize-all-all-x	Wrong Size n all processes for MPI_Bcast() – L240:ilcsTSP.c		
9		misCrit-1-1-x	Missing Critical Section in buggyProc and buggyThread – L170:ilcsTSP.c		
10		misCrit-all-1-x	Missing Critical Section in buggyThread and all prcoesses – L170:ilcsTSP.c		
11		misCrit-1-all-x	Missing Critical Section in buggyProc and all threads – L170:ilcsTSP.c		
12	OMP	misCrit-all-all-x	Missing Critical Section in all procs and threads – L170:ilcsTSP.c		
13		misCrit2-1-1-x	Missing Critical Section in buggyProc and buggyThread – L230:ilcsTSP.c		
14		misCrit2-all-1-x	Missing Critical Section in buggyThread – L230:ilcsTSP.c		
15		misCrit2-1-all-x	Missing Critical Section in buggyProc and all threads – L230:ilcsTSP.c		
16		misCrit2-all-all-x	Missing Critical Section in all procs and threads – L230:ilcsTSP.c		
17		misCrit3-1-all-x	Missing Critical Section in buggyProc and all threads – L280:ilcsTSP.c		
18		misCrit3-all-all-x	Missing Critical Section in all procs and threads – L280:ilcsTSP.c		
19	General	infLoop-1-1-1	Injected an infinite loop after CPU_EXEC() in buggyProc,buggyThread & buggyIter L164:ilcsTSP.c		

	Filter	Attributes	Top Thread 0	Top Thread 1	Top Thread 2	Top Thread 3	Top Thread 4
8	(14)11.mem.mpicol.ompcrit.cust.0K10	doub.actual	1:(1_0 6_0) 1.00 2:(2_0 6_0) 1.00 3:(5_0 6_0) 1.00	1:(2_1,3_1):0.00 2:(0_1,4_1):0.00 3:(0_1,5_1):0.00	1:(1_2,6_2):0.57 2:(3_2,6_2):0.42 3:(1_2,3_2):0.34	1:(1_3,4_3):0.47 2:(6_3,7_3):0.38 3:(0_3,2_3):0.27	1:(6_4,7_4):0.71 2:(4_4,6_4):0.69 3:(0_4,6_4):0.69
1	(14)11.mem.mpicol.ompcrit.cust.0K10	doub.log10	1:(1_0 6_0) 1.00 2:(2_0 6_0) 1.00 3:(5_0 6_0) 1.00	1:(2_1,3_1):0.00 2:(0_1,4_1):0.00 3:(0_1,5_1):0.00	1:(3_2,6_2):0.36 2:(0_2,2_2):0.30 3:(1_2,3_2):0.29	1:(1_3,6_3):0.62 2:(0_3,5_3):0.29 3:(1_3,4_3):0.26	1:(1_4,5_4):0.62 2:(0_4,5_4):0.39 3:(0_4,4_4):0.33
2	(14)11.mem.mpicol.ompcrit.cust.0K10	doub.orig	1:(1_0,6_0) 1.00 2:(2_0,6_0) 1.00 3:(5_0,6_0) 1.00	1:(2_1,3_1):0.00 2:(0_1,4_1):0.00 3:(0_1,5_1):0.00	1:(3_2,6_2):0.36 2:(0_2,2_2):0.30 3:(1_2,3_2):0.29	1:(1_3,6_3):0.62 2:(0_3,5_3):0.29 3:(1_3,4_3):0.26	1:(1_4,5_4):0.62 2:(0_4,5_4):0.39 3:(0_4,4_4):0.33
3	(14)11.mem.mpicol.ompcrit.cust.0K10	sing.actual	1:(1_0,6_0) 1.00 2:(2_0,6_0) 1.00 3:(5_0,6_0) 1.00	1:(2_1,3_1):0.00 2:(0_1,4_1):0.00 3:(0_1,5_1):0.00	1:(3_2,6_2):0.71 2:(1_2,3_2):0.60 3:(1_2,6_2):0.47	1:(0_3,2_3):0.67 2:(4_3,7_3):0.60 3:(6_3,7_3):0.46	1:(6_4,7_4):0.60 2:(0_4,4_4):0.40 3:(2_4,6_4):0.40
4	(14)11.mem.mpicol.ompcrit.cust.0K10	sing.log10	1:(1_0 6_0) 1.00 2:(2_0 6_0) 1.00 3:(5_0 6_0) 1.00	1:(2_1,3_1):0.00 2:(0_1,4_1):0.00 3:(0_1,5_1):0.00	1:(5_2,7_2):0.67 2:(1_2,3_2):0.67 3:(1_2,6_2):0.50	1:(0_3,2_3):0.40 2:(1_3,6_3):0.40 3:(0_3,5_3):0.27	1:(5_4,6_4):0.50 2:(0_4,5_4):0.50 3:(1_4,5_4):0.40
5	(14)11.mem.mpicol.ompcrit.cust.8K18	sing.orig	1:(1_0,6_0) 1.00 2:(2_0,6_0) 1.00 3:(5_0,6_0) 1.00	1:(2_1,3_1):0.00 2:(0_1,4_1):0.00 3:(0_1,5_1):0.00	1:(3_2,6_2):0.43 2:(0_2,2_2):0.33 3:(1_2,3_2):0.33	1:(0_3,5_3):0.33 2:(0_3,2_3):0.33 3:(1_3,6_3):0.33	1:(5_4,6_4):0.43 2:(0_4,5_4):0.43 3:(1_4,5_4):0.33
6	- Average		1.000	0.000	0.441	0.403	0.488
7	- Mean 1.000		1.000	0.000	0.441	0.403	0.488
8	- Min		1.000	0.000	0.286	0.262	0.333
9	- Max		1.000	0.000	0.714	0.667	0.714

Figure 2: Recommendation Table Bug 1

5.1.1 1: AllReduce wrong op - proc 2. Bug injected within to MPI_AllReduce() in line 21 where the local champions of all processes are about computed and their MIN is going to be distributed to all other processes. However, with injecting the wrong operation (MAX instead of MIN) to one (or more) of the processes, the logic of if statement in line 24 would be violated and the control flow of the code would not take this branch which is the the core part of ILCS (deciding champion to broadcast to other processes).

After collecting traces from bug-free version and this buggy version of ILCS-TSP, and after applying filters and detecting loops, my ranking system, with the highest possible score, recommends to look at figure 3 diffNLR. This diffNLR shows that in the bug free version, we are expecting that at least one of the processes (in this case process 6 to win the championship, copy the result of that process to broadcast_buffer and broadcast it to all other processes. However, in the buggy version, the champion never gets updated (never goes through OMP critical section).

The recommendation system that I designed tries to find the pairs of traces that their similarities changed the most. In other word, this recommendation system wants to suggest top diffNLRs that are candidates to show the impact of the bug based on trace similarities. (table

6 CASE STUDIES

6.0.1 Case Study: ILCS. ILCS [10]

6.0.2 Case Study: MFEM.

7 DISCUSSION AND FUTURE WORK

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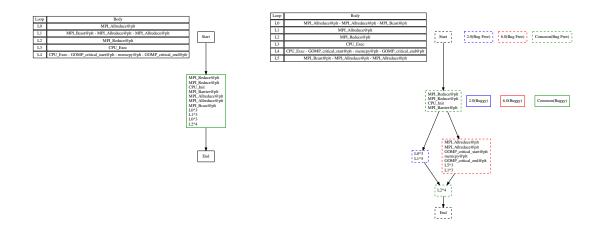


Figure 3: diffNLR of process 2 and process 6 buggy(AllReduce() wrong op) vs. bug-free

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