

1)

$$8n^2 = 64n \lg n$$

$$n^2 = 8n \lg n$$

$$n = 8 \lg n$$

$$n - 8 \lg n = 0$$

$$n = 43.56$$

insertion sort is faster than merge sort for values $0 < n < 43.56$.

2)

	1 second	1 min	1 hour	1 day	1 month	1 year	1 century
Lg(n)	$2^{(1 \cdot 10^6)}$	$2^{(60 \cdot 10^6)}$	$2^{(3600 \cdot 10^6)}$	$2^{(86400 \cdot 10^6)}$	$2^{(2592000 \cdot 10^6)}$	$2^{(31536000 \cdot 10^6)}$	$2^{(315576000 \cdot 10^6)}$
Sqrt(n)	$(11 \cdot 10^6)^{1/2}$	$(60 \cdot 10^6)^{1/2}$	$(3600 \cdot 10^6)^{1/2}$	$(86400 \cdot 10^6)^{1/2}$	$(2592000 \cdot 10^6)^{1/2}$	$(31536000 \cdot 10^6)^{1/2}$	$(3155760000 \cdot 10^6)^{1/2}$
N	$1 \cdot 10^6$	$60 \cdot 10^6$	$3600 \cdot 10^6$	$86400 \cdot 10^6$	$2592000 \cdot 10^6$	$31536000 \cdot 10^6$	$3155760000 \cdot 10^6$
$N \lg(n)$	62746.1	$2.8 \cdot 10^6$	$1.33 \cdot (10^8)$	$2.75 \cdot (10^9)$	$7.18 \cdot (10^{10})$	$7.97 \cdot (10^{11})$	$6.86 \cdot (10^{13})$
$N^{1/2}$	$(1 \cdot 10^6)^{1/2}$	$(60 \cdot 10^6)^{1/2}$	$(3600 \cdot 10^6)^{1/2}$	$(86400 \cdot 10^6)^{1/2}$	$(2592000 \cdot 10^6)^{1/2}$	$(31536000 \cdot 10^6)^{1/2}$	$(3155760000 \cdot 10^6)^{1/2}$
$N^{1/3}$	$(1 \cdot 10^6)^{1/3}$	$(60 \cdot 10^6)^{1/3}$	$(3600 \cdot 10^6)^{1/3}$	$(86400 \cdot 10^6)^{1/3}$	$(2592000 \cdot 10^6)^{1/3}$	$(31536000 \cdot 10^6)^{1/3}$	$(3155760000 \cdot 10^6)^{1/3}$
2^n	$\lg(1 \cdot 10^6)$	$\lg(60 \cdot 10^6)$	$\lg(3600 \cdot 10^6)$	$\lg(86400 \cdot 10^6)$	$\lg(2592000 \cdot 10^6)$	$\lg(31536000 \cdot 10^6)$	$\lg(3155760000 \cdot 10^6)$
N!	9.45	11.16	12.79	13.9	15.25	16.15	17.757

3)

$$P(k) = T(2^k) = 2^k \lg 2^k$$

$$\text{Base case: } P(1) = T(2^1) = (2^1) \lg(2^1) = 2$$

$$\text{Assume } P(k) \text{ is true, } T(2^k) = 2^k \lg 2^k$$

$$P(k+1) = 2^{k+1} \lg 2^{k+1}$$

Proof for $k+1$:

$$\begin{aligned}
& 2T(2^{k+1}/2) + 2^{k+1} \\
&= (2)T(2^k) + 2^{k+1} \\
&= (2)(2^k \lg 2^k) + 2^{k+1} \\
&= 2^{k+1} \lg 2^k + 2^{k+1} \\
&= 2^{k+1} (\lg 2^k + 1) \\
&= 2^{k+1} (\lg 2^k + \lg 2) \\
&= 2^{k+1} \lg 2^{k+1}
\end{aligned}$$

4)

A: Omega

B: Omega

C: Theta

D: Omega

E: Theta

F: O

G: O

H: O

5)

Getminmax(X[])

Min = X[1]

Max = X[1]

For(i=2 to X length)

 If(X[i] > max)

 Max = X[i]

 If(X[i] < min)

$\text{Min} = X[1]$

Return(min,max)

Worst case this will have a runtime of $2N$

6)

7)