CME 307 / MS&E 311: Optimization

Introduction

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Management Science and Engineering Stanford

April 24, 2023

Announcements

- lectures slides on website: https://stanford-cme-307.github.io/website-2023/calendar
- recordings posted on canvas
- office hours posted on website: https://stanford-cme-307.github.io/website-2023/staff

Outline

What topics and why?

Grading

Quadratic optimization

Nonlinear optimization

Conic optimization

Integer programming

Convex optimization

Who's here

What is an optimization problem?

optimization problem: nonlinear form

minimize
$$f_0(x)$$

subject to $f_i(x) \leq 0, \quad i = 1, \dots, m_1$
 $h_i(x) = 0, \quad i = 1, \dots, m_2$
variable $x \in \mathbf{R}^n$

- ightharpoonup objective f_0
- ▶ inequality constraints f_i
- \triangleright equality constraints h_i

Modularity in optimization

how to optimize:

- 1. model problem as a mathematical optimization problem
- 2. identify the properties of the problem
- 3. use an appropriate solver (or write a new one)

...and iterate:

- approximate the problem to make it easier
- solve a sequence of approximated problems that converge to solve the original problem
- or initialize ("warm-start") a solver for the original problem with a solution to the approximated problem

How will you use optimization?

- model problem to match solver capabilities
- use solvers as building blocks in larger algorithms
- write your own solver for a new problem class
- write a specialized solver for a particular problem

What will you need to know to use (or develop) optimization?

- convexity
- geometry
- numerical stability
- complexity
- reliability
- stopping conditions

Optimization courses

compared to

- ► EE364a: more algos, less modeling
- ► EE364b: more nonconvex, different mix of algos

my goal: < 30% overlap with either

Topics

let's look at topics on google sheets: https://docs.google.com/spreadsheets/d/ 1PXv_sFkhz5jNAA765kgHanSPolLm2fHzzRNuOOJzpHM

what other topics should we add?

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Grading

let's look at the course website: https://stanford-cme-307.github.io/website-2023/

- for PhD students, the grading is designed to help (not distract from) your research
- for UG/Masters, this course may push you hard...

Optimization problems

important optimization problem classes:

- quadratic
- unconstrained
- finite-sum
- linear
- conic
- convex
- nonlinear (with linear or nonlinear constraints)
- mixed-integer linear programs
- ▶ black-box with (0, 1, or 2)-order oracle

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Quadratic optimization

a quadratic optimization problem is written as

$$\begin{array}{ll} \text{minimize} & \frac{1}{2}\|Ax - b\|^2 := f_0(x) \\ \text{variable} & x \in \mathbf{R}^n \end{array}$$

where

- ▶ $A \in \mathbf{R}^{m \times n}$: matrix
- ▶ $b \in \mathbf{R}^m$: vector

how to solve?

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how to solve? take gradient and set to 0:

$$\nabla f_0(x) = A^T (Ax - b) = 0$$

⇒ linear system solvers also solve quadratic optimization problems

matrix $A \in \mathbf{R}^{m \times n}$

check matrix calculus results by checking dimensions

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 - ▶ null space of A, **nullspace**(A), is a hyperplane of dimension n-m
 - ▶ solution set is $\{x : Ax = b\} = \{x_0 + Vz\}$ where columns of $V \in \mathbb{R}^{n \times n m}$ span **nullspace**(A)

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- \triangleright A^TA is symmetric positive semidefinite (proof on board)

Symmetric positive semidefinite matrices

Definition

a symmetric matrix $Q \in \mathbf{R}^{n \times n}$ is **positive semidefinite** (psd) if $x^T Q x \ge 0$ for all $x \in \mathbf{R}^n$.

these matrices are so important that there are many ways to write them! for $Q \in \mathbf{R}^{n \times n}$,

$$Q \in \mathbf{S}_{+}^{n} \iff Q \succeq 0 \iff Q = Q^{T}, \ \lambda_{\min}(Q) \geq 0$$

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why care about psd matrices Q?

- least-squares objective has a psd $Q = A^T A$
- \triangleright level sets of $x^T Qx$ are (bounded) ellipsoids
- ▶ the quadratic form $x^T Qx$ is a metric iff Q > 0
- eigenvalue decomp and svd coincide for psd matrices

Quadratic program

a quadratic program is written as

minimize
$$\frac{1}{2}x^TQx + c^Tx$$

subject to $Ax = b$
variable $x \in \mathbf{R}^n$

where

- $\triangleright Q \in \mathbf{R}^{n \times n}$: symmetric positive semidefinite matrix
- $c \in \mathbf{R}^n$: vector

how to solve?

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how to solve? reduce to quadratic optimization problem:

- (explicit) form solution set $\{x : Ax = b\} = \{x_0 + Vz \mid z \in \mathbf{R}^{n-m}\}$ by computing a solution $Ax_0 = b$ and a basis V for the null space of A
- (implicit) use duality to recast problem as larger linear (KKT) system

Quadratic program: application

Markowitz portfolio optimization problem:

minimize
$$\gamma x^T \Sigma x - \mu^T x$$

subject to $\sum_i x_i = 1$
 $Ax = 0$
variable $x \in \mathbf{R}^n$

where

- $\Sigma \in \mathbf{R}^{n \times n}$: asset covariance matrix
- $\blacktriangleright \mu \in \mathbf{R}^n$: asset return vector
- $ightharpoonup \gamma \in \mathbf{R}$: risk aversion parameter
- rows of $A \in \mathbf{R}^{m \times n}$ correspond to other portfolios
 - ensures new portfolio is independent, e.g., of market returns

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Unconstrained smooth optimization

for $f: \mathbf{R}^n \to \mathbf{R}$ ctsly differentiable,

minimize f(x) variable $x \in \mathbf{R}^n$

how to solve?

Unconstrained smooth optimization

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minimize
$$f(x)$$
 variable $x \in \mathbf{R}^n$

how to solve? approximate as a quadratic problem

$$f(x) \approx f(x_0) + \nabla f(x_0)^T (x - x_0) + \frac{1}{2} (x - x_0)^T H(x_0) (x - x_0)$$

and find solution x_{quad} to the quadratic problem. then set $x_0 \leftarrow x_{\text{quad}}$ and repeat.

Finite sum

finite sum optimization problem

minimize
$$\sum_{i=1}^{m} f_i(x)$$
 variable $x \in \mathbf{R}^n$

key fact: can approximate gradient using gradient on **minibatch** $S \subseteq \{1, ..., m\}$:

$$\nabla f(x) \approx \frac{1}{|S|} \sum_{i \in S} \nabla f_i(x)$$

examples:

- statistical learning (logistic regression, SVM)
- deep learning

Background: classification

classification problem: m data points

- feature vector $a_i \mathbf{R}^n$, i = 1, ..., m
- ▶ label $b_i \in \{-1, 1\}, i = 1, ..., m$

choose decision boundary $a^Tx = 0$ to separate data points into two classes

- $ightharpoonup a^T x > 0 \implies \text{predict class } 1$
- $ightharpoonup a^T x < 0 \implies \text{predict class -1}$

classification is correct if $b_i a^T x > 0$

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classification is correct if $b_i a^T x > 0$

- projective transformation transforms affine boundary to linear boundary
- classification is invariant to scalar multiplication of x

Logistic regression

(regularized) logistic regression minimizes the finite sum

minimize
$$\sum_{i=1}^{m} \log(1 + \exp(-b_i a_i^T x)) + r(x)$$
 variable $x \in \mathbf{R}^n$

where

- ▶ $b_i \in \{-1, 1\}, a_i \in \mathbb{R}^n$
- $ightharpoonup r: \mathbf{R}^n \to \mathbf{R}$ is a **regularizer**, e.g., $||x||^2$ or $||x||_1$

support vector machine (SVM) minimizes the finite sum

minimize
$$\sum_{i=1}^{m} \max(0, 1 - b_i a_i^T x) + \gamma ||x||^2$$
 variable $x \in \mathbf{R}^n$

where $b_i \in \{-1,1\}$ and $a_i \in \mathbf{R}^n$.

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how to solve?

- use subgradient method
- transform to conic form
- solve dual problem instead
- **smooth** the objective

Nonlinear optimization

optimization problem: nonlinear form

minimize
$$f_0(x)$$

subject to $f_i(x) \leq b_i, \quad i = 1, \dots, m_1$
 $h(x) = 0$
variable $x \in \mathbf{R}^n$

- $ightharpoonup x = (x_1, \dots, x_n)$: optimization variables
- ▶ $f_0 : \mathbf{R}^n \to \mathbf{R}$: objective function
- ▶ $f_i : \mathbf{R}^n \to \mathbf{R}$, i = 1, ..., m: constraint functions

special case: unconstrained optimization

Example: process control

You are the process engineer for a desalination plant that produces drinking water. The plant has a variety of knobs, collected in vector x, that you can turn to control the process. These control, e.g., how much water is pumped into the plant, how much pressure is used to force the water through filters, and how much of each chemical is added to the water.

- $ightharpoonup f_0(x)$: cost of water produced
- $ightharpoonup f_i(x)$: level of each measured impurity in the water
- ▶ b_i: maximum allowable level of each impurity

Given a setting of the knobs, you can observe the cost of water produced and the levels of impurities.

What is the optimal setting of the knobs?

Oracles

an optimization **oracle** is your interface for accessing the problem data:

e.g., an oracle for $f: \mathbf{R}^n \to \mathbf{R}$ can evaluate for any $x \in \mathbf{R}^n$:

- **>** zero-order: $f_0(x)$
- ▶ first-order: $f_0(x)$ and $\nabla f_0(x)$
- **second-order:** $f_0(x)$, $\nabla f_0(x)$, and $\nabla^2 f_0(x)$

why oracles?

- can optimize real systems based on observed output (not just models)
- can use and extend old or complex but trusted code (e.g., NASA, PDE simulations, . . .)
- can prove lower bounds on the oracle complexity of a problem class

source: Nesterov 2004 "Introductory Lectures on Convex Optimization"

Nonlinear optimization: how to solve?

depends on the oracle:

- first- or second-order: approximate by a sequence of quadratic problems
- zero-order: harder, lots of methods
 - simulated annealing
 - Bayesian optimization
 - pseudo-higher-order methods, e.g., compute approximate gradient

Solution of an optimization problem

minimize
$$f(x)$$

for $f: \mathcal{D} \to \mathbf{R}$. x^* is a

- ▶ **local minimizer** if there is a neighborhood \mathcal{N} around x^* so that $f(x) \geq f(x^*)$ for all $x \in \mathcal{N}$.
- **plobal minimizer** if $f(x) \ge f(x^*)$ for all $x \in \mathcal{D}$.
- ▶ **strict local minimizer** if there is a neighborhood \mathcal{N} around x^* so that $f(x) > f(x^*)$ for all $x \in \mathcal{N}$.
- **isolated local minimizer** if the neighborhood $\mathcal N$ contains no other local minimizers.
- **unique minimizer** if it is the only global minimizer.

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- ▶ global minimizer if $f(x) \ge f(x^*)$ for all $x \in \mathcal{D}$.
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pictures!

First order optimality condition

Theorem

If $x^* \in \mathbb{R}^n$ is a local minimizer of a differentiable function $f : \mathbb{R}^n \to \mathbb{R}$, then $\nabla f(x^*) = 0$.

First order optimality condition

Theorem

If $x^* \in \mathbb{R}^n$ is a local minimizer of a differentiable function $f : \mathbb{R}^n \to \mathbb{R}$, then $\nabla f(x^*) = 0$.

proof: suppose by contradiction that $\nabla f(x^*) \neq 0$. consider points of the form $x_{\alpha} = x^* - \alpha \nabla f(x^*)$ for $\alpha > 0$. by definition of the gradient,

$$\lim_{\alpha \to 0} \frac{f(x_\alpha) - f(x^\star)}{\alpha} = -\nabla f(x^\star)^\top \nabla f(x^\star) = -\|\text{nabla} f(x^\star)\|^2 < 0$$

so for any sufficiently small $\alpha > 0$, we have $f(x_{\alpha}) < f(x^{*})$, which contradicts the fact that x^{*} is a local minimizer.

Second order optimality condition

Theorem

If $x^* \in \mathbb{R}^n$ is a local minimizer of a twice differentiable function $f: \mathbb{R}^n \to \mathbb{R}$, then $\nabla^2 f(x^*) \succeq 0$.

Second order optimality condition

Theorem

If $x^* \in \mathbb{R}^n$ is a local minimizer of a twice differentiable function $f : \mathbb{R}^n \to \mathbb{R}$, then $\nabla^2 f(x^*) \succeq 0$.

proof: similar to the previous proof. use the fact that the second order approximation

$$f(x_{\alpha}) \approx f(x^{\star}) + \nabla f(x^{\star})^{\top} (x_{\alpha} - x^{\star}) + \frac{1}{2} (x_{\alpha} - x^{\star})^{\top} \nabla^{2} f(x^{\star}) (x_{\alpha} - x^{\star})$$

is accurate locally to show a contradiction unless $\nabla^2 f(x^*) \succeq 0$: if not, there is a direction v such that $v^T \nabla^2 f(x^*) v < 0$. then $f(x + \alpha v) < f(x^*)$ for α arbitrarily small, which contradicts the fact that x^* is a local minimizer.

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$$c^T x$$

subject to $b - Ax \ge 0$
variable $x \in \mathbf{R}^n$

where

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how to solve?

- use the simplex method
- use a conic solver

Conic form

conic form optimization problem generalizes LP:

minimize
$$c^T x$$

subject to $b - Ax \in \mathcal{K}$,

where K is a **convex cone**:

$$x \in \mathcal{K} \iff rx \in \mathcal{K} \text{ for any } r > 0.$$

examples:

- ightharpoonup zero cone $\mathcal{K}_0 = \{0\}$
- ▶ positive orthant $\mathcal{K}_+ = \{x : x_i >= 0, i = 1, ..., n\}$
- second order cone $\mathcal{K}_{\mathsf{SOC}} = \{(x, t) : \|x\|_2 \leq t\}$
- ▶ positive semidefinite (PSD) cone $\mathcal{K}_{SDP} = \{X : X = X^T, \ v^T X v \ge 0, \ \forall v \in \mathbf{R}^n \}$
- cartesian products of cones

Conic form: how to solve?

Morally, conic problems are solved by reducing to a nonlinear optimization problem

- barrier methods (e.g., interior point methods)
 - add a barrier term to the objective that goes to infinity when constraints are violated
- penalty methods (e.g., augmented Lagrangian methods, ADMM, ...)
 - add a penalty term to the objective that depends on a dual variable
 - adjust the dual variable to enforce constraints

Conic form example: nonnegative least squares

$$\begin{array}{ll} \text{minimize} & \|Ax-b\| \\ \text{subject to} & x \geq 0 \\ & & \\ \\ \text{minimize} & t \\ \text{subject to} & x \in \mathcal{K}_+ \\ & (Ax-b,t) \in \mathcal{K}_{\mathsf{SOC}} \end{array}$$

Conic form example: SVM

$$\begin{array}{ll} \text{minimize} & \sum_{i=1}^{m} \max(0, 1 - b_i a_i^T x) + \|x\|^2 \\ \text{variable} & x \in \mathbf{R}^n \\ & \updownarrow \\ \\ \text{minimize} & \sum_i s_i + t \\ \text{subject to} & s \geq \operatorname{\mathbf{diag}}(b) Ax - 1 \\ & s \geq 0 \\ & t \geq \|x\|^2 \\ & \updownarrow \\ \\ \text{minimize} & \sum_i s_i + t \\ \text{subject to} & s - \operatorname{\mathbf{diag}}(b) Ax + 1 \in \mathcal{K}_+ \\ & s \in \mathcal{K}_+ \\ & [t \ x; x^T \ I_n] \in \mathcal{K}_{\text{SDP}} \\ \end{array}$$

Schur complement

Consider the block matrix

$$X = \begin{bmatrix} A & B \\ B^T & C \end{bmatrix}.$$

- ▶ the **Schur complement** of *A* in *X* is $C B^T A^{-1} B$.
- ▶ $X \succeq 0$ if and only if $A \succeq 0$ and $C B^T A^{-1} B \succeq 0$. (proof by partial minimization of quadratic form $(u, v)^T X(u, v)$ over $u \in \mathbf{R}^m$ for fixed $v \in \mathbf{R}^n$)

Conic form example: semidefinite programming

$$\begin{array}{ll} \text{minimize} & \lambda_{\max}(X) + y^T X^{-1} y \\ \text{subject to} & X \succeq 0 \\ & \updownarrow \end{array}$$

Conic form example: semidefinite programming

$$\begin{array}{ll} \text{minimize} & \lambda_{\text{max}}(X) + y^T X^{-1} y \\ \text{subject to} & X \succeq 0 \\ & & \\ \\ \text{minimize} & t_1 + t_2 \\ \text{subject to} & t_1 I - X \in \mathcal{K}_{\text{SDP}} \\ & \begin{bmatrix} t_2 & y^T \\ y & X \end{bmatrix} \in \mathcal{K}_{\text{SDP}} \end{array}$$

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Integer programming

integer linear programming generalizes linear programming:

minimize
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subject to $b - Ax \ge 0$
variable $x \in \mathbf{Z}^n$

variants:

- ▶ mixed integer linear programming (MILP): $x \in \mathbf{Z}^{n-m} \cup \mathbf{R}^m$
- ▶ mixed integer nonlinear programming (MINLP): $x \in \mathbf{Z}^{n-m} \cup \mathbf{R}^m$ and nonlinear objective or constraints

how to solve?

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how to solve?

- ▶ use Gurobi, CPLEX, ...
- branch and bound and cut (i.e., a sequence of LPs)
- use duality to decompose into a sequence of simpler LPs

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Convex sets

Definition

A set $S \subseteq \mathbb{R}^n$ is convex if it contains every chord: for all $\theta \in [0,1]$, w, $v \in S$,

$$\theta w + (1 - \theta)v \in S$$

Convex sets

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A set $S \subseteq \mathbf{R}^n$ is convex if it contains every chord: for all $\theta \in [0,1]$, w, $v \in S$,

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Q: Which of these are convex? ellipsoid, half moon

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- a function $f: \mathbf{R}^n \to \mathbf{R}$ is convex iff
 - ▶ **Chords.** it never lies above its chord: $\forall \theta \in [0,1], w, v \in \mathbb{R}^n$

$$f(\theta w + (1 - \theta)v) \le \theta f(w) + (1 - \theta)f(v)$$

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▶ Epigraph. epi $(f) = \{(x, t) : t \ge f(x)\}$ is convex

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- **Epigraph.** epi $(f) = \{(x, t) : t \ge f(x)\}$ is convex
- **First order condition.** if *f* is differentiable,

$$f(v) - f(w) \ge \nabla f(w)^{\top} (v - w) \qquad \forall w, v \in \mathbf{R}^n$$

a function $f: \mathbf{R}^n \to \mathbf{R}$ is convex iff

▶ **Chords.** it never lies above its chord: $\forall \theta \in [0,1], w, v \in \mathbb{R}^n$

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Q: Which of these are convex? quadratic, I1, pwl, step, jump, logistic, logistic loss

Convex optimization

an optimization problem is convex if:

- ▶ Geometrically: the feasible set and the epigraph of the objective are convex
- ▶ **NLP:** the objective and inequality constraints are convex functions, and the equality constraints are affine
- ► Conic: all the cones are convex cones

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duality, stopping conditions, ...

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- ightharpoonup a function f is concave if -f is convex
- concave maximization results in a convex optimization problem

Local minima are global for convex functions

Theorem

If x^* is a local minimizer of a convex function f, then x^* is a global minimizer.

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proof: suppose by contradiction that another point x' is a global minimizer, with $f(x') < f(x^*)$. draw the chord between x' and x^* . since the chord lies above f, every convex combination $x = \theta x^* + (1 - \theta)x'$ of x' and x^* for $\theta \in (0,1)$ has a value $f(x) < f(x^*)$. this is true even for $x \to x^*$, contradicting our assumption that x^* is a local minimizer.

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If f is convex and differentiable and $\nabla f(x^*) = 0$, then x^* is a global minimizer.

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Q: Is a global minimizer of a convex function always unique?

A: No. Picture.

Modern solvers

- ▶ algebraic modeling languages, e.g.
 - ▶ JuMP facilitates nonlinear and mixed integer optimization
 - CVX* (CVX, CVXPY, Convex.jl, ...) transform a problem into conic form
- and modern solvers

Optimization modeling

- ► Rocket control
- ► Power systems
- ► AML

Outline

What topics and why?

Grading

Quadratic optimization

Nonlinear optimization

Conic optimization

Integer programming

Convex optimization

Who's here?

Questions for you

respond at pollev.com/madeleineudell824

- program / major
- year of program
- schedule conflicts
- what other optimization courses have you taken?
- what optimization techniques do you want to learn?
- what's your favorite algorithm (or algorithmic trick)