# CME 307 / MS&E 311: Optimization

## Introduction

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#### **Announcements**

- website: https://stanford-cme-307.github.io/web
- ► Ed for discussion and announcements: https://edstem.org/us/courses/51411/
- fill out course survey (also linked on website): https://forms.gle/7hPniFeC576S12FAA
- ▶ talk to me after class and/or schedule office hours (see website)
- class attendance is required. will post some slides, generally no recordings

#### Intro exercise

Form groups of three people you haven't met before and introduce yourselves.

- name, major, year
- why are you interested in optimization?
- what are you hoping to learn?

# **Goals for today**

- Understand course objectives and expectations
- ▶ Identify several types of optimization problem
- Meet at least two other students you've not met before
- Discuss challenges in a real-world optimization problem

### **Outline**

What is an optimization problem?

Course goals and expectations

#### **Exercise**

## Discuss in groups:

- give an example of a real-world optimization problem
- what is the problem data?
- what are the problem variables?
- how would you write it down?

# (Integer) linear optimization problem

minimize 
$$c^T x$$
  
subject to  $Ax = b$   
 $\ell \le x \le u$   
variable  $x \in \mathbb{Z}^n$ 

- ightharpoonup objective  $c^T x$
- ightharpoonup equality constraints Ax = b
- lower and upper bounds  $\ell \le x \le u$
- integer variable

## problem data:

- $ightharpoonup A \in \mathbf{R}^{m \times n}, \ b \in \mathbf{R}^m$
- $c \in \mathbb{R}^n$
- $\ell \in \mathbb{R}^n$ ,  $u \in \mathbb{R}^n$

# Nonlinear optimization problem

minimize 
$$f_0(x)$$
  
subject to  $f_i(x) \leq 0, \quad i = 1, \dots, m_1$   
 $h_i(x) = 0, \quad i = 1, \dots, m_2$   
variable  $x \in \mathbf{R}^n$ 

- ightharpoonup objective  $f_0$
- ▶ inequality constraints f<sub>i</sub>
- equality constraints h<sub>i</sub>

## problem data:

- ▶ (blackbox) code to evaluate  $f_i$  and  $h_i$  for any  $x \in \mathbf{R}^n$
- ▶ (first order) and to compute gradients
- (second order) and to compute Hessians

# **Optimization problems**

## important optimization problem classes:

- ► linear
- integer
- nonlinear (with linear or nonlinear constraints)
- quadratic
- unconstrained
- ► finite-sum
- conic
- convex
- ▶ black-box with (0, 1, or 2)-order oracle

# **Optimization problems**

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draw a picture relating these

### **Discuss**

Does your problem fit into one of these categories?

- ▶ which?
- ▶ if not, why not?

# Modularity in optimization

#### how to optimize:

- 1. model problem as a mathematical optimization problem
- 2. identify the properties of the problem
- 3. use an appropriate solver (or write a new one)

#### ...and iterate:

- approximate the problem to make it easier
- solve a sequence of approximated problems that converge to solve the original problem
- or initialize ("warm-start") a solver for the original problem with a solution to the approximated problem

### **Outline**

What is an optimization problem?

Course goals and expectations

## **Course goals**

look at goals on course website: <a href="https://stanford-cme-307.github.io/web/">https://stanford-cme-307.github.io/web/</a>

- ► Which goals sound exciting?
- ▶ Which goals don't make sense?
- ▶ What else do you hope to accomplish?

## **Course expectations**

look at grading tab on course website: <a href="https://stanford-cme-307.github.io/web/discuss">https://stanford-cme-307.github.io/web/discuss</a>:

- do expectations make sense given course goals?
- questions about expectations?

#### Course schedule

look at materials tab on course website and at spreadsheet of topics: https://stanford-cme-307.github.io/web/

## **Course project**

read project expectations and ideas. discuss in groups:

- do you want to work on algorithms, applications, or LLM tools?
- pick a problem. what data would you need? what challenges would you foresee in solving it?

#### What next?

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# Quadratic optimization

Nonlinear optimization

Conic optimization

Integer programming

Convex optimization

# **Quadratic optimization**

a quadratic optimization problem is written as

minimize 
$$\frac{1}{2} ||Ax - b||^2 := f_0(x)$$
 variable  $x \in \mathbf{R}^n$ 

#### where

- $ightharpoonup A \in \mathbf{R}^{m \times n}$ : matrix
- ▶  $b \in \mathbf{R}^m$ : vector

how to solve?

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how to solve? take gradient and set to 0:

$$\nabla f_0(x) = A^T (Ax - b) = 0$$

 $\implies$  linear system solvers also solve quadratic optimization problems

matrix  $A \in \mathbf{R}^{m \times n}$ 

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- $ightharpoonup A^T A$  is symmetric positive semidefinite (proof on board)

# Symmetric positive semidefinite matrices

## Definition

a symmetric matrix  $Q \in \mathbf{R}^{n \times n}$  is **positive semidefinite** (psd) if  $x^T Qx \ge 0$  for all  $x \in \mathbf{R}^n$ .

these matrices are so important that there are many ways to write them! for  $Q \in \mathbf{R}^{n \times n}$ ,

$$Q \in \mathbf{S}_{+}^{n} \iff Q \succeq 0 \iff Q = Q^{T}, \ \lambda_{\min}(Q) \geq 0$$

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 $Q \in \mathbf{S}_{+}^{n}$  is symmetric positive definite (spd)  $(Q \succ 0)$  if  $x^{T}Qx > 0$  for all  $x \in \mathbf{R}^{n}$ . why care about psd matrices Q?

- least-squares objective has a psd  $Q = A^T A$
- $\triangleright$  level sets of  $x^T Q x$  are (bounded) ellipsoids
- ▶ the quadratic form  $x^T Qx$  is a metric iff Q > 0
- eigenvalue decomp and svd coincide for psd matrices

# Quadratic program

## a quadratic program is written as

minimize 
$$\frac{1}{2}x^TQx + c^Tx$$
  
subject to  $Ax = b$   
variable  $x \in \mathbf{R}^n$ 

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how to solve? reduce to quadratic optimization problem:

- (explicit) form solution set  $\{x : Ax = b\} = \{x_0 + Vz \mid z \in \mathbf{R}^{n-m}\}$  by computing a solution  $Ax_0 = b$  and a basis V for the null space of A
- ▶ (implicit) use duality to recast problem as larger linear (KKT) system

# Quadratic program: application

## Markowitz portfolio optimization problem:

minimize 
$$\gamma x^T \Sigma x - \mu^T x$$
  
subject to  $\sum_i x_i = 1$   
 $Ax = 0$   
variable  $x \in \mathbf{R}^n$ 

#### where

- $ightharpoonup \Sigma \in \mathbf{R}^{n \times n}$ : asset covariance matrix
- $\blacktriangleright \mu \in \mathbf{R}^n$ : asset return vector
- $ightharpoonup \gamma \in \mathbf{R}$ : risk aversion parameter
- ▶ rows of  $A \in \mathbf{R}^{m \times n}$  correspond to other portfolios
  - ensures new portfolio is independent, e.g., of market returns

## **Outline**

Quadratic optimization

Nonlinear optimization

Conic optimization

Integer programming

Convex optimization

## **Unconstrained smooth optimization**

for  $f: \mathbf{R}^n \to \mathbf{R}$  ctsly differentiable,

minimize 
$$f(x)$$
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how to solve?

## **Unconstrained smooth optimization**

for  $f: \mathbf{R}^n \to \mathbf{R}$  ctsly differentiable,

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how to solve? approximate as a quadratic problem

$$f(x) \approx f(x_0) + \nabla f(x_0)^T (x - x_0) + \frac{1}{2} (x - x_0)^T H(x_0) (x - x_0)$$

and find solution  $x_{\rm quad}$  to the quadratic problem. then set  $x_0 \leftarrow x_{\rm quad}$  and repeat.

#### Finite sum

finite sum optimization problem

minimize 
$$\sum_{i=1}^{m} f_i(x)$$
 variable  $x \in \mathbf{R}^n$ 

**key fact:** can approximate gradient using gradient on **minibatch**  $S \subseteq \{1, \dots, m\}$ :

$$\nabla f(x) \approx \frac{1}{|S|} \sum_{i \in S} \nabla f_i(x)$$

examples:

- statistical learning (logistic regression, SVM)
- deep learning

## **Background: classification**

classification problem: m data points

- feature vector  $a_i \in \mathbf{R}^n$ , i = 1, ..., m
- ▶ label  $b_i \in \{-1, 1\}, i = 1, ..., m$

choose decision boundary  $a^Tx = 0$  to separate data points into two classes

- $ightharpoonup a^T x > 0 \implies \text{predict class } 1$
- $ightharpoonup a^T x < 0 \implies \text{predict class -1}$

classification is correct if  $b_i a^T x > 0$ 

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- projective transformation transforms affine boundary to linear boundary
- classification is invariant to scalar multiplication of x

## **Logistic regression**

(regularized) logistic regression minimizes the finite sum

minimize 
$$\sum_{i=1}^{m} \log(1 + \exp(-b_i a_i^T x)) + r(x)$$
 variable  $x \in \mathbf{R}^n$ 

#### where

- ▶  $b_i \in \{-1, 1\}, a_i \in \mathbb{R}^n$
- ▶  $r: \mathbb{R}^n \to \mathbb{R}$  is a **regularizer**, e.g.,  $\|x\|^2$  or  $\|x\|_1$

support vector machine (SVM) minimizes the finite sum

minimize 
$$\sum_{i=1}^{m} \max(0, 1 - b_i a_i^T x) + \gamma ||x||^2$$
 variable  $x \in \mathbf{R}^n$ 

where  $b_i \in \{-1,1\}$  and  $a_i \in \mathbf{R}^n$ .

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$$\begin{array}{ll} \text{minimize} & \sum_{i=1}^m \max(0, 1 - b_i a_i^T x) + \gamma \|x\|^2 \\ \text{variable} & x \in \mathbf{R}^n \end{array}$$

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where  $b_i \in \{-1,1\}$  and  $a_i \in \mathbb{R}^n$ . not differentiable!

how to solve?

- use subgradient method
- transform to conic form
- solve dual problem instead
- **smooth** the objective

## Nonlinear optimization

optimization problem: nonlinear form

$$\begin{array}{ll} \text{minimize} & f_0(x) \\ \text{subject to} & f_i(x) \leq b_i, \quad i=1,\ldots,m_1 \\ & h(x)=0 \\ \text{variable} & x \in \mathbf{R}^n \end{array}$$

- $ightharpoonup x = (x_1, \dots, x_n)$ : optimization variables
- ▶  $f_0 : \mathbf{R}^n \to \mathbf{R}$ : objective function
- ▶  $f_i : \mathbf{R}^n \to \mathbf{R}$ , i = 1, ..., m: constraint functions

special case: unconstrained optimization

#### **Example: process control**

You are the process engineer for a desalination plant that produces drinking water. The plant has a variety of knobs, collected in vector x, that you can turn to control the process. These control, e.g., how much water is pumped into the plant, how much pressure is used to force the water through filters, and how much of each chemical is added to the water.

- $ightharpoonup f_0(x)$ : cost of water produced
- $ightharpoonup f_i(x)$ : level of each measured impurity in the water
- $\triangleright$   $b_i$ : maximum allowable level of each impurity

Given a setting of the knobs, you can observe the cost of water produced and the levels of impurities.

#### What is the optimal setting of the knobs?

#### **Oracles**

an optimization **oracle** is your interface for accessing the problem data: *e.g.*, an oracle for  $f: \mathbf{R}^n \to \mathbf{R}$  can evaluate for any  $x \in \mathbf{R}^n$ :

- **>** zero-order:  $f_0(x)$
- ▶ **first-order:**  $f_0(x)$  and  $\nabla f_0(x)$
- **second-order:**  $f_0(x)$ ,  $\nabla f_0(x)$ , and  $\nabla^2 f_0(x)$

why oracles?

- can optimize real systems based on observed output (not just models)
- can use and extend old or complex but trusted code (e.g., NASA, PDE simulations, . . . )
- can prove lower bounds on the oracle complexity of a problem class

source: Nesterov 2004 "Introductory Lectures on Convex Optimization"'

#### Nonlinear optimization: how to solve?

#### depends on the oracle:

- first- or second-order: approximate by a sequence of quadratic problems
- zero-order: harder, lots of methods
  - simulated annealing
  - Bayesian optimization
  - pseudo-higher-order methods, e.g., compute approximate gradient

## Solution of an optimization problem

minimize 
$$f(x)$$

for  $f: \mathcal{D} \to \mathbf{R}$ .  $x^*$  is a

- ▶ **local minimizer** if there is a neighborhood  $\mathcal{N}$  around  $x^*$  so that  $f(x) \ge f(x^*)$  for all  $x \in \mathcal{N}$ .
- **plobal minimizer** if  $f(x) \ge f(x^*)$  for all  $x \in \mathcal{D}$ .
- ▶ **strict local minimizer** if there is a neighborhood  $\mathcal{N}$  around  $x^*$  so that  $f(x) > f(x^*)$  for all  $x \in \mathcal{N}$ .
- **isolated local minimizer** if the neighborhood  $\mathcal N$  contains no other local minimizers.
- unique minimizer if it is the only global minimizer.

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#### pictures!

## First order optimality condition

#### Theorem

If  $x^* \in \mathbf{R}^n$  is a local minimizer of a differentiable function  $f : \mathbf{R}^n \to \mathbf{R}$ , then  $\nabla f(x^*) = 0$ .

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**proof:** suppose by contradiction that  $\nabla f(x^*) \neq 0$ . consider points of the form  $x_{\alpha} = x^* - \alpha \nabla f(x^*)$  for  $\alpha > 0$ . by definition of the gradient,

$$\lim_{\alpha \to 0} \frac{f(x_\alpha) - f(x^\star)}{\alpha} = -\nabla f(x^\star)^\top \nabla f(x^\star) = -\|\textit{nablaf}(x^\star)\|^2 < 0$$

so for any sufficiently small  $\alpha > 0$ , we have  $f(x_{\alpha}) < f(x^{*})$ , which contradicts the fact that  $x^{*}$  is a local minimizer.

## Second order optimality condition

#### Theorem

If  $x^* \in \mathbf{R}^n$  is a local minimizer of a twice differentiable function  $f : \mathbf{R}^n \to \mathbf{R}$ , then  $\nabla^2 f(x^*) \succeq 0$ .

# Second order optimality condition

#### Theorem

If  $x^* \in \mathbf{R}^n$  is a local minimizer of a twice differentiable function  $f : \mathbf{R}^n \to \mathbf{R}$ , then  $\nabla^2 f(x^*) \succeq 0$ .

**proof:** similar to the previous proof. use the fact that the second order approximation

$$f(x_{\alpha}) \approx f(x^{\star}) + \nabla f(x^{\star})^{\top} (x_{\alpha} - x^{\star}) + \frac{1}{2} (x_{\alpha} - x^{\star})^{\top} \nabla^{2} f(x^{\star}) (x_{\alpha} - x^{\star})$$

is accurate locally to show a contradiction unless  $\nabla^2 f(x^*) \succeq 0$ : if not, there is a direction v such that  $v^T \nabla^2 f(x^*) v < 0$ . then  $f(x + \alpha v) < f(x^*)$  for  $\alpha$  arbitrarily small, which contradicts the fact that  $x^*$  is a local minimizer.

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#### how to solve?

- use the simplex method
- use a conic solver

#### **Conic form**

**conic form** optimization problem generalizes LP:

minimize 
$$c^T x$$
  
subject to  $b - Ax \in \mathcal{K}$ ,

where K is a **convex cone**:

$$x \in \mathcal{K} \iff rx \in \mathcal{K} \text{ for any } r > 0.$$

#### examples:

- ightharpoonup zero cone  $\mathcal{K}_0 = \{0\}$
- ▶ positive orthant  $\mathcal{K}_+ = \{x : x_i >= 0, i = 1, ..., n\}$
- ▶ second order cone  $\mathcal{K}_{SOC} = \{(x, t) : ||x||_2 \le t\}$
- ▶ positive semidefinite (PSD) cone  $\mathcal{K}_{SDP} = \{X : X = X^T, \ v^T X v \ge 0, \ \forall v \in \mathbf{R}^n\}$
- cartesian products of cones

#### Conic form: how to solve?

Morally, conic problems are solved by reducing to a nonlinear optimization problem

- barrier methods (e.g., interior point methods)
  - add a barrier term to the objective that goes to infinity when constraints are violated
- ▶ penalty methods (e.g., augmented Lagrangian methods, ADMM, ...)
  - add a penalty term to the objective that depends on a dual variable
  - adjust the dual variable to enforce constraints

## Conic form example: nonnegative least squares

$$\begin{array}{ll} \text{minimize} & \|Ax-b\| \\ \text{subject to} & x \geq 0 \\ & & \\ \\ \text{minimize} & t \\ \text{subject to} & x \in \mathcal{K}_+ \\ & (Ax-b,t) \in \mathcal{K}_{\mathsf{SOC}} \end{array}$$

#### Conic form example: SVM

$$\begin{array}{ll} \text{minimize} & \sum_{i=1}^m \max(0,1-b_ia_i^Tx) + \|x\|^2 \\ \text{x in } & \updownarrow \\ \\ \text{minimize} & \sum_i s_i + t \\ \text{subject to} & s \geq \operatorname{diag}(b)Ax - 1 \\ & s \geq 0 \\ & t \geq \|x\|^2 \\ \\ \text{minimize} & \sum_i s_i + t \\ \text{subject to} & s - \operatorname{diag}(b)Ax + 1 \in \mathcal{K}_+ \\ & s \in \mathcal{K}_+ \\ & [t \; x; x^T \; I_n] \in \mathcal{K}_{\mathsf{SDP}} \end{array}$$

## **Schur complement**

Consider the block matrix

$$X = \begin{bmatrix} A & B \\ B^T & C \end{bmatrix}.$$

- ▶ the **Schur complement** of *A* in *X* is  $C B^T A^{-1} B$ .
- ▶  $X \succeq 0$  if and only if  $A \succeq 0$  and  $C B^T A^{-1} B \succeq 0$ . (proof by partial minimization of quadratic form  $(u, v)^T X(u, v)$  over  $u \in \mathbf{R}^m$  for fixed  $v \in \mathbf{R}^n$ )

## Conic form example: semidefinite programming

$$\begin{array}{ll} \text{minimize} & \lambda_{\max}(X) + y^T X^{-1} y \\ \text{subject to} & X \succeq 0 \\ & & \updownarrow \end{array}$$

# Conic form example: semidefinite programming

$$\begin{array}{ll} \text{minimize} & \lambda_{\max}(X) + y^T X^{-1} y \\ \text{subject to} & X \succeq 0 \\ & & \updownarrow \\ \\ \text{minimize} & t_1 + t_2 \\ \text{subject to} & t_1 I - X \in \mathcal{K}_{\text{SDP}} \\ \begin{bmatrix} t_2 & y^T \\ y & X \end{bmatrix} \in \mathcal{K}_{\text{SDP}} \end{array}$$

#### **Outline**

Quadratic optimization

Nonlinear optimization

Conic optimization

Integer programming

Convex optimization

# Integer programming

integer linear programming generalizes linear programming:

minimize 
$$c^T x$$
  
subject to  $b - Ax \ge 0$   
variable  $x \in \mathbf{Z}^n$ 

#### variants:

- ▶ mixed integer linear programming (MILP):  $x \in \mathbf{Z}^{n-m} \cup \mathbf{R}^m$
- ▶ mixed integer nonlinear programming (MINLP):  $x \in \mathbf{Z}^{n-m} \cup \mathbf{R}^m$  and nonlinear objective or constraints

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#### how to solve?

- use Gurobi, CPLEX, . . .
- branch and bound and cut (i.e., a sequence of LPs)
- use duality to decompose into a sequence of simpler LPs

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#### Convex sets

#### Definition

A set  $S \subseteq \mathbf{R}^n$  is convex if it contains every chord: for all  $\theta \in [0,1]$ , w,  $v \in S$ ,

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**Q:** Which of these are convex? ellipsoid, half moon

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Q: Which of these are convex? quadratic, I1, pwl, step, jump, logistic, logistic loss

## **Convex optimization**

an optimization problem is convex if:

- ▶ **Geometrically:** the feasible set and the epigraph of the objective are convex
- ▶ **NLP:** the objective and inequality constraints are convex functions, and the equality constraints are affine
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duality, stopping conditions, ...

- ightharpoonup a function f is concave if -f is convex
- concave maximization results in a convex optimization problem

## Local minima are global for convex functions

### Theorem

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If  $x^*$  is a local minimizer of a convex function f, then  $x^*$  is a global minimizer.

**proof:** suppose by contradiction that another point x' is a global minimizer, with  $f(x') < f(x^*)$ . draw the chord between x' and  $x^*$ . since the chord lies above f, every convex combination  $x = \theta x^* + (1 - \theta)x'$  of x' and  $x^*$  for  $\theta \in (0,1)$  has a value  $f(x) < f(x^*)$ . this is true even for  $x \to x^*$ , contradicting our assumption that  $x^*$  is a local minimizer.

### **Corollary**

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If f is convex and differentiable and  $\nabla f(x^*) = 0$ , then  $x^*$  is a global minimizer.

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Q: Is a global minimizer of a convex function always unique?

A: No. Picture.

#### Modern solvers

- ▶ algebraic modeling languages, *e.g.* 
  - ▶ JuMP facilitates nonlinear and mixed integer optimization
  - ► CVX\* (CVX, CVXPY, Convex.jl, ...) transform a problem into conic form
- and modern solvers

# **Optimization modeling**

- ► Rocket control
- Power systems
- ► AML

#### **Announcements**

- website: https://stanford-cme-307.github.io/web
- ► Ed for discussion and announcements: https://edstem.org/us/courses/51411/
- fill out course survey (also linked on website): https://forms.gle/7hPniFeC576S12FAA
- ▶ talk to me after class and/or schedule office hours (see website)
- class attendance is required. will post some slides, generally no recordings