# MS&E 233 Game Theory, Data Science and Al Lecture 6

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(by courtesy) Computer Science and Electrical Engineering

Institute for Computational and Mathematical Engineering

#### **Computational Game Theory for Complex Games**

- Basics of game theory and zero-sum games (T)
- Basics of online learning theory (T)
- Solving zero-sum games via online learning (T)
- HW1: implement simple algorithms to solve zero-sum games
- Applications to ML and AI (T+A)
- HW2: implement boosting as solving a zero-sum game
- Basics of extensive-form games
- Solving extensive-form games via online learning (T)
- HW3: implement agents to solve very simple variants of poker
- General games and equilibria (T)

(3)

- Online learning in general games, multi-agent RL (T+A)
- HW4: implement no-regret algorithms that converge to correlated equilibria in general games

#### **Data Science for Auctions and Mechanisms**

- Basics and applications of auction theory (T+A)
- Learning to bid in auctions via online learning (T)
- HW5: implement bandit algorithms to bid in ad auctions

- Optimal auctions and mechanisms (T)
- Simple vs optimal mechanisms (T)
- HW6: calculate equilibria in simple auctions, implement simple and optimal auctions, analyze revenue empirically
- Optimizing mechanisms from samples (T)
- Online optimization of auctions and mechanisms (T)
- HW7: implement procedures to learn approximately optimal auctions from historical samples and in an online manner

#### **Further Topics**

- Econometrics in games and auctions (T+A)
- A/B testing in markets (T+A)
- HW8: implement procedure to estimate values from bids in an auction, empirically analyze inaccuracy of A/B tests in markets

#### **Guest Lectures**

- Mechanism Design for LLMs, Renato Paes Leme, Google Research
- Auto-bidding in Sponsored Search Auctions, Kshipra Bhawalkar, Google Research

## Solving Extensive Form Games via No-Regret Learning

## Recap: No-Regret Learning in Sequence Form

 We have successfully turned imperfect information extensive form zero-sum games into a familiar object

$$\max_{\tilde{x} \in X} \min_{\tilde{y} \in Y} \tilde{x}^{\top} A \tilde{y}$$

• X, Y are convex sets, i.e., sequence-form strategies

- We can invoke minimax theorem to prove existence of equilibria
- We can calculate equilibria via LP duality
- We can calculate equilibria via no-regret learning!

### **Recap:** Regret of FTRL

(FTRL) 
$$x_t = \underset{x \in X}{\operatorname{argmax}} \underbrace{\sum_{\tau < t} \langle x, u_\tau \rangle} - \underbrace{\frac{1}{\eta} \mathcal{R}(x)}_{\text{tunction of } x \text{ that stabilizes the maximizer}}_{\text{Historical performance}}$$

of always choosing

strategy x

Theorem. Assuming the utility function at each period

$$f_t(x) = \langle x, u_t \rangle$$

is L-Lipschitz with respect to some norm  $\|\cdot\|$  and the regularizer is 1-strongly convex with respect to the same norm then

Regret – FTRL(T) 
$$\leq \eta L + \frac{1}{\eta T} \left( \max_{x \in X} \mathcal{R}(x) - \min_{x \in X} \mathcal{R}(x) \right)$$

Average stability induced by regularizer

Average loss distortion caused by regularizer

## **Recap:** Regularizer for the Treeplex Space X

 $\bullet$  The only thing we are missing is a good Regularizer for X

$$U_{t-1} = \sum_{\tau < t} u_{\tau}$$

• **Desiderata.** Be strongly convex in x within X and for the optimization problem to be fast to solve

$$\tilde{x}_{t} = \underset{\tilde{x} \in X}{\operatorname{argmax}} \sum_{\tau < t} \langle \tilde{x}, u_{\tau} \rangle - \frac{1}{\eta} \mathcal{R}(\tilde{x}) = \underset{\tilde{x} \in X}{\operatorname{argmax}} \langle \tilde{x}, U_{t-1} \rangle - \frac{1}{\eta} \mathcal{R}(\tilde{x})$$

• X is no longer a "simplex", so entropy is not a good Regularizer

#### Dilated Entropy

- X is a combination of scaled simplices, i.e.,  $\tilde{x}=\left(\tilde{x}^j\right)_{j\in\mathcal{J}_1}$
- $\tilde{x}^j = (\tilde{x}_a)_{a \in A_j}$ : sequence-form strategies for actions in infoset  $j \in \mathcal{J}_1$   $\tilde{x}^j \in \tilde{x}_{p_j} \cdot \Delta_j \quad \Leftrightarrow \quad \tilde{x}^j / \tilde{x}_{p_j} \in \Delta_j$
- Consider a weighted combination of local negative entropies

$$\mathcal{R}(\tilde{x}) \coloneqq \sum_{j} \beta_{j} \, \tilde{x}_{p_{j}} \, \mathrm{H}\left(\tilde{x}^{j} / \tilde{x}_{p_{j}}\right), \qquad \mathrm{H}(u) = \sum_{l \in \mathbb{N}} u_{l} \log(u_{l})$$
Equivalent to the behavioral strategy  $x^{j}$ 
Negative Entropy

•  $\mathcal{R}(\tilde{x})$  is 1/M strongly convex w.r.t.  $\ell_1$  norm, where  $M = \max_{\tilde{x} \in X} ||\tilde{x}||_1$ , for appropriate choice of  $\beta_i$  based on game tree structure

#### Solving the Optimization Problem

Optimization problem decomposes into local simplex problems

$$\sum_{j \in \mathcal{J}_{1}} \left\langle \tilde{x}^{j}, U_{t-1}^{j} \right\rangle - \left| \frac{1}{\eta} \beta_{j} \right| \tilde{x}_{p_{j}} \operatorname{H} \left( \frac{\tilde{x}^{j}}{\tilde{x}_{p_{j}}} \right) = \sum_{j \in \mathcal{J}_{1}} \tilde{x}_{p_{j}} \left\{ \left\langle \frac{\tilde{x}^{j}}{\tilde{x}_{p_{j}}}, U_{t-1}^{j} \right\rangle - \frac{1}{\eta_{j}} \operatorname{H} \left( \frac{\tilde{x}^{j}}{\tilde{x}_{p_{j}}} \right) \right\}$$

• Quantity  $\frac{\tilde{x}^j}{\tilde{x}_{p_j}}$  is essentially the behavioral strategy  $x^j$  at infoset j

$$\sum_{j \in \mathcal{J}_1} \tilde{x}_{p_j} \left\{ \left\langle x^j, U_{t-1}^j \right\rangle - \frac{1}{\eta_j} H(x^j) \right\}$$

• Quantity  $x^j$  over simplex  $\Delta_j$  is independent of solution  $x_a$  for all ancestral actions and only appears in subsequent infosets

#### Solving the Optimization Problem

• Decomposes in local max over behavioral strategies  $x^j$  solved bottom up

$$V^{j} = \max_{x^{j} \in \Delta_{j}} \left\langle x^{j}, U_{t-1}^{j} \right\rangle - \frac{1}{\eta_{j}} H(x^{j}) \Rightarrow \begin{cases} x^{j} \propto \exp\left(\eta_{j} U_{t-1}^{j}\right) \\ V^{j} = \log \sum_{a \in A_{j}} \exp\left(\eta_{j} U_{t-1}^{a}\right) = \operatorname{softmax}_{\eta_{j}} \left(U_{t-1}^{j}\right) \end{cases}$$

• Value  $V^j$  multiplies  $\tilde{x}_{p_j}$ ; when solving for  $\tilde{x}_{p_j}$  we need to take it into account. If  $p_j \in A_k$ 

$$\max_{\boldsymbol{x}^k \in \Delta_k} \langle \tilde{\boldsymbol{x}}^k, \boldsymbol{U}_{t-1}^k \rangle - \eta_k \; \tilde{\boldsymbol{x}}_{p_k} \; \mathbf{H} \left( \frac{\tilde{\boldsymbol{x}}^k}{\tilde{\boldsymbol{x}}_{p_k}} \right) + \tilde{\boldsymbol{x}}_{p_j} \boldsymbol{V}^j + \cdots$$

• Add  $V^j$  to "cumulative utility"  $Q_{p_j}$  (initialized at  $U_{t-1,p_j}$ ) associated with  $p_j$ 

$$Q_{p_j} \leftarrow Q_{p_j} + V^j$$

#### **Sum:** Nash via FTRL with Dilated Entropy

Each player chooses  $\tilde{x}_t$ ,  $\tilde{y}_t$  based on FTRL with dilated entropy

- For x-player  $u_t = A \tilde{y}_t$  and  $U_t = U_{t-1} + u_t$  and initialize  $Q = U_t$
- Traverse the tree bottom-up; for each infoset  $j \in \mathcal{J}_1$   $x_{t+1}^j \propto \exp(\eta_j Q^j)$ ,  $V^j = \operatorname{softmax}_{\eta_j}(Q^j)$ ,  $Q_{p_j} \leftarrow Q_{p_j} + V^j$
- Define sequence-form strategies top-down:  $\tilde{x}_{t+1}^j = \tilde{x}_{p_j} \cdot x_{t+1}^j$  Similarly, for y player

Return average of sequence-form strategies as equilibrium

#### Interpreting utility vector

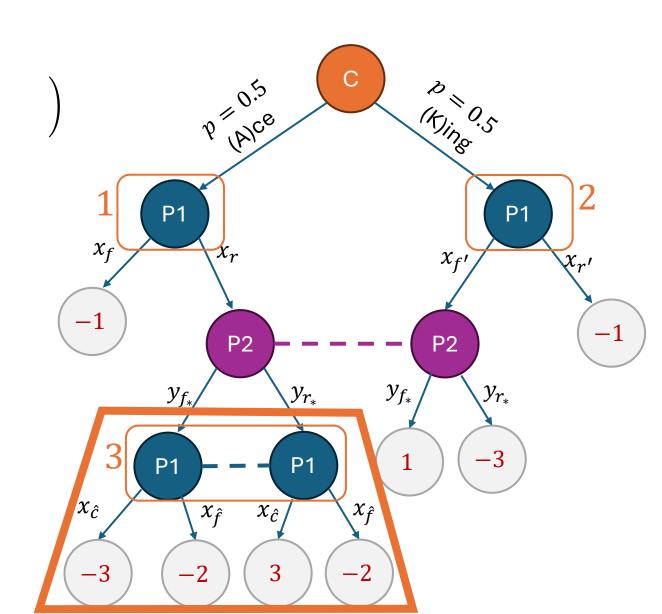
$$u_{t,a} = A\tilde{y}_t = \sum_{a' \in A_{P2}} A_{a,a'} \tilde{y}_{t,a'}$$

 $A_{a.a'}$  is zero if the combination of a, a' does not lead to a leaf node

$$u_{t,a} = \sum_{\substack{a \text{ was last P1 action} \\ a' \text{ was last P2 action}}} u(z) \operatorname{Pr} \begin{pmatrix} \operatorname{Chance chooses} \\ \operatorname{sequence on} \\ \operatorname{path to } z \end{pmatrix} \operatorname{Pr} \begin{pmatrix} \operatorname{P2 plays} \\ \operatorname{sequence} \\ \operatorname{leading to } a' \end{pmatrix}$$

**Interpretation.** If I play with the intend to arrive at action a (i.e.  $\tilde{x}_a = 1$ ) and then don't make any other moves, what is the expected reward that I will collect, in expectation over the choices of my opponent and nature

$$U^3 += \left(u_{\hat{c}}, u_{\hat{f}}\right) = \left(u_{\hat{c}}, u_{\hat{f}}\right) = \left(u_{\hat{c}}, u_{\hat{f}}\right)$$

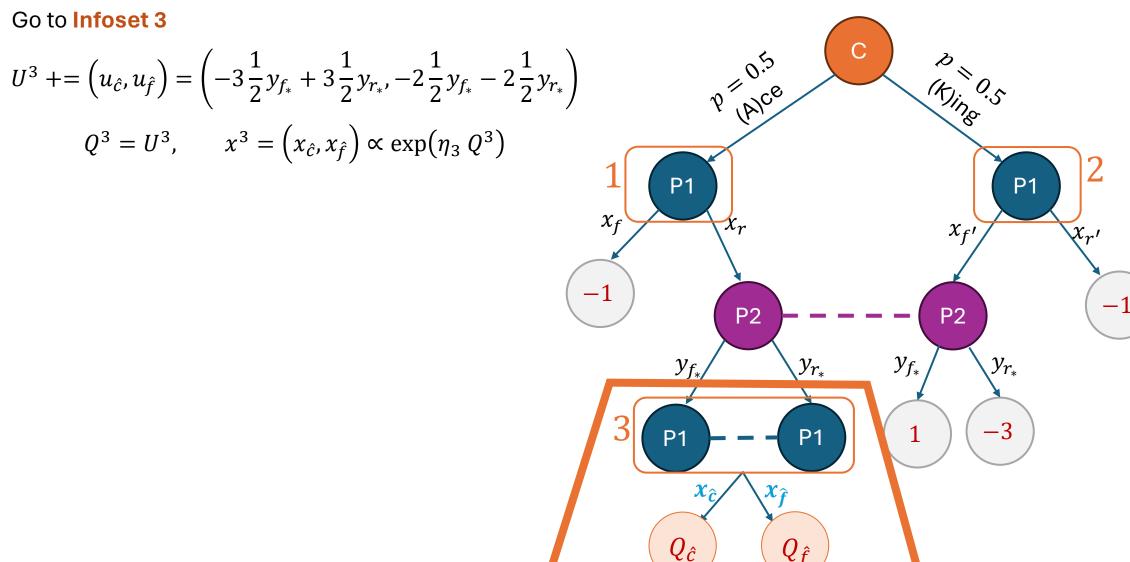


• Go to Infoset 3
$$U^{3} += \left(u_{\hat{c}}, u_{\hat{f}}\right) = \left(-3\frac{1}{2}y_{f_{*}} + 3\frac{1}{2}y_{r_{*}}, -2\frac{1}{2}y_{f_{*}} - 2\frac{1}{2}y_{r_{*}}\right)$$

$$1$$

$$y_{f_{*}}$$

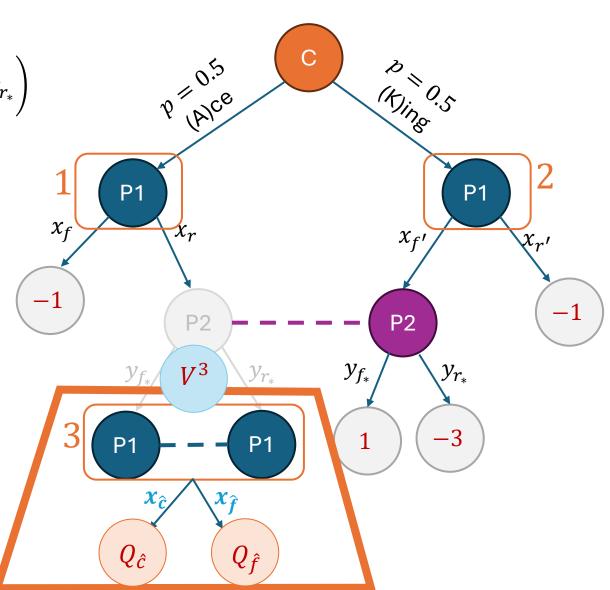
$$y_{f_{*$$



$$U^{3} += \left(u_{\hat{c}}, u_{\hat{f}}\right) = \left(-3\frac{1}{2}y_{f_{*}} + 3\frac{1}{2}y_{r_{*}}, -2\frac{1}{2}y_{f_{*}} - 2\frac{1}{2}y_{r_{*}}\right)$$

$$Q^{3} = U^{3}, \qquad x^{3} = \left(x_{\hat{c}}, x_{\hat{f}}\right) \propto \exp\left(\eta_{3} Q^{3}\right)$$

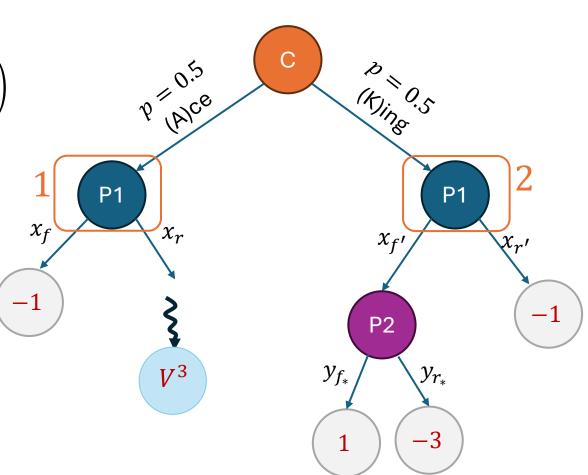
$$V^{3} = \operatorname{softmax}(\eta_{3} Q^{3})$$



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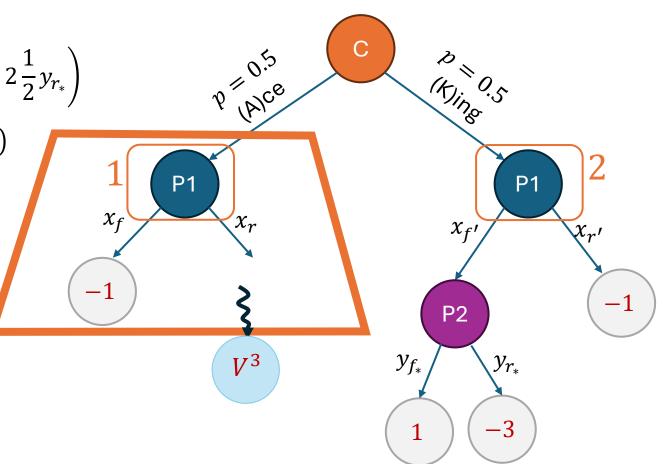
Go to Infoset 3

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$$Q^{3} = U^{3}, \qquad x^{3} = \left(x_{\hat{c}}, x_{\hat{f}}\right) \propto \exp(\eta_{3} Q^{3})$$

$$V^{3} = \operatorname{softmax}(\eta_{3} Q^{3})$$

$$U^1 += (u_f, u_r) = \left( \right)$$



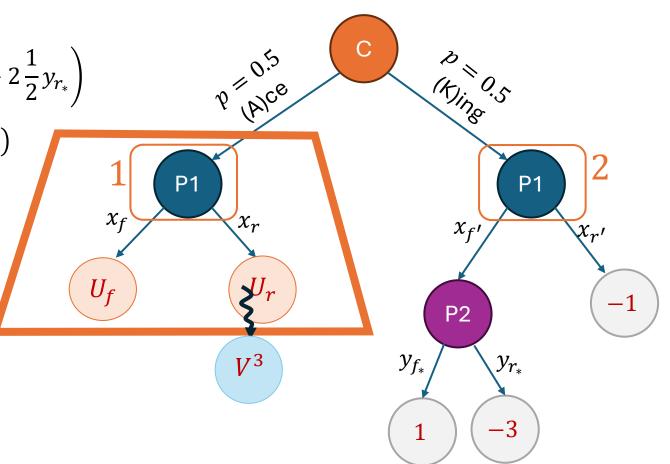
Go to Infoset 3

$$U^{3} += \left(u_{\hat{c}}, u_{\hat{f}}\right) = \left(-3\frac{1}{2}y_{f_{*}} + 3\frac{1}{2}y_{r_{*}}, -2\frac{1}{2}y_{f_{*}} - 2\frac{1}{2}y_{r_{*}}\right)$$

$$Q^{3} = U^{3}, \qquad x^{3} = \left(x_{\hat{c}}, x_{\hat{f}}\right) \propto \exp(\eta_{3} Q^{3})$$

$$V^{3} = \operatorname{softmax}(\eta_{3} Q^{3})$$

$$U^1 += (u_f, u_r) = \left(-1\frac{1}{2}, 0\right)$$



Go to Infoset 3

$$U^{3} += \left(u_{\hat{c}}, u_{\hat{f}}\right) = \left(-3\frac{1}{2}y_{f_{*}} + 3\frac{1}{2}y_{r_{*}}, -2\frac{1}{2}y_{f_{*}} - 2\frac{1}{2}y_{r_{*}}\right)$$

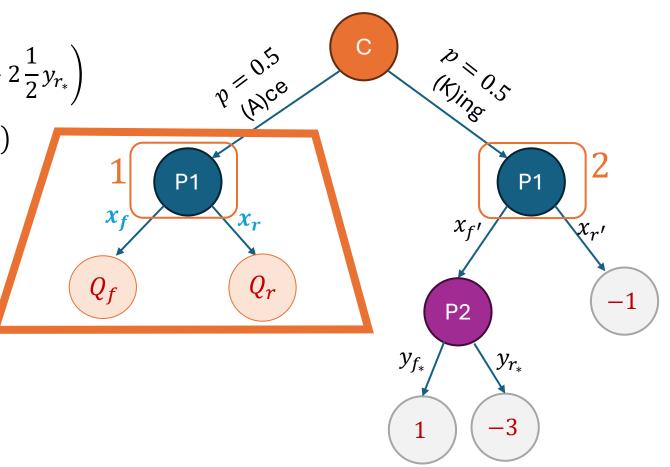
$$Q^{3} = U^{3}, \qquad x^{3} = \left(x_{\hat{c}}, x_{\hat{f}}\right) \propto \exp(\eta_{3} Q^{3})$$

$$V^{3} = \operatorname{softmax}(\eta_{3} Q^{3})$$

$$U^{1} += (u_{f}, u_{r}) = \left(-1\frac{1}{2}, 0\right)$$

$$Q^{1} = U^{1} + (0, V^{3}) = \left(-1\frac{1}{2}, V^{3}\right)$$

$$x^{1} = (x_{f}, x_{r}) \propto \exp(\eta_{1} Q^{1})$$



Go to Infoset 3

$$U^{3} += \left(u_{\hat{c}}, u_{\hat{f}}\right) = \left(-3\frac{1}{2}y_{f_{*}} + 3\frac{1}{2}y_{r_{*}}, -2\frac{1}{2}y_{f_{*}} - 2\frac{1}{2}y_{r_{*}}\right)$$

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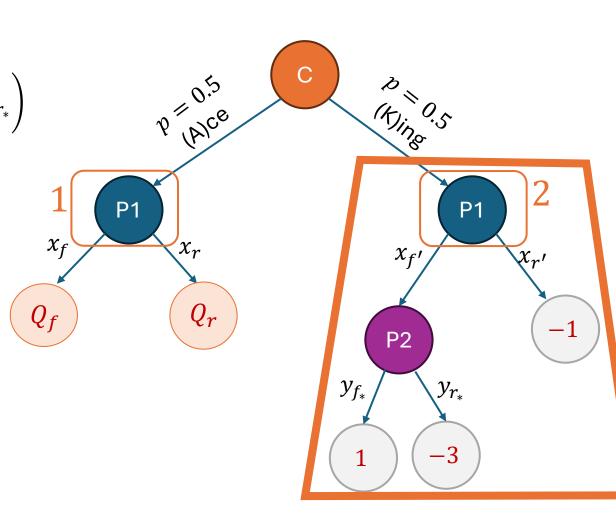
Go to Infoset 1

$$U^{1} += (u_{f}, u_{r}) = \left(-1\frac{1}{2}, 0\right)$$

$$Q^{1} = U^{1} + (0, V^{3}) = \left(-1\frac{1}{2}, V^{3}\right)$$

$$x^{1} = (x_{f}, x_{r}) \propto \exp(\eta_{1} Q^{1})$$

$$U^2 += \left(u_{f'}, u_{r'}\right) = \left( \begin{array}{c} \\ \\ \end{array} \right)$$



Go to Infoset 3

$$U^{3} += \left(u_{\hat{c}}, u_{\hat{f}}\right) = \left(-3\frac{1}{2}y_{f_{*}} + 3\frac{1}{2}y_{r_{*}}, -2\frac{1}{2}y_{f_{*}} - 2\frac{1}{2}y_{r_{*}}\right)$$

$$Q^{3} = U^{3}, \qquad x^{3} = \left(x_{\hat{c}}, x_{\hat{f}}\right) \propto \exp(\eta_{3} Q^{3})$$

$$V^{3} = \operatorname{softmax}(\eta_{3} Q^{3})$$

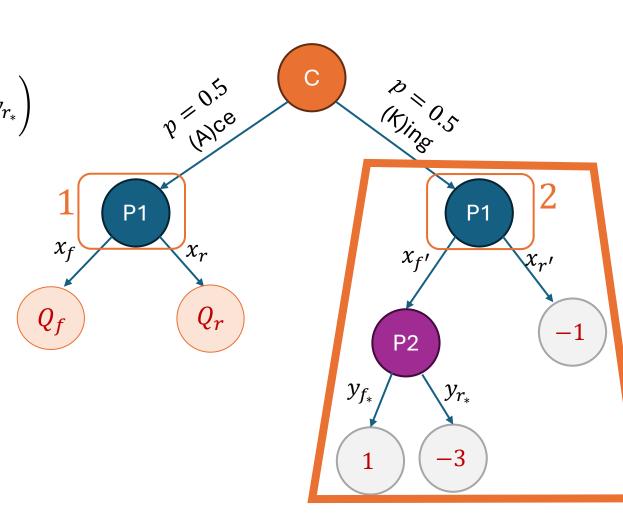
Go to Infoset 1

$$U^{1} += (u_{f}, u_{r}) = \left(-1\frac{1}{2}, 0\right)$$

$$Q^{1} = U^{1} + (0, V^{3}) = \left(-1\frac{1}{2}, V^{3}\right)$$

$$x^{1} = (x_{f}, x_{r}) \propto \exp(\eta_{1} Q^{1})$$

$$U^{2} += \left(u_{f'}, u_{r'}\right) = \left(1\frac{1}{2}y_{f_{*}} - 3\frac{1}{2}y_{r_{*}}, -1\frac{1}{2}\right)$$



Go to Infoset 3

$$U^{3} += \left(u_{\hat{c}}, u_{\hat{f}}\right) = \left(-3\frac{1}{2}y_{f_{*}} + 3\frac{1}{2}y_{r_{*}}, -2\frac{1}{2}y_{f_{*}} - 2\frac{1}{2}y_{r_{*}}\right)$$

$$Q^{3} = U^{3}, \qquad x^{3} = \left(x_{\hat{c}}, x_{\hat{f}}\right) \propto \exp\left(\eta_{3} Q^{3}\right)$$

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Go to Infoset 1

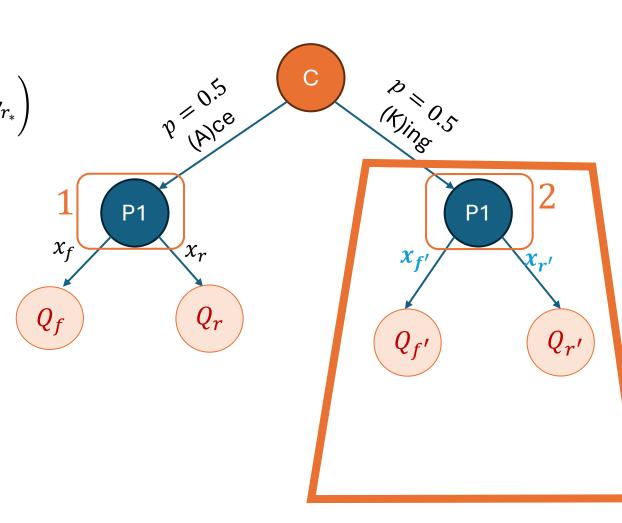
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$$Q^{2} = U^{2}, \qquad x^{2} = \left(x_{f'}, x_{r'}\right) \propto \exp(\eta_{2}Q^{2})$$



#### **Sum:** Nash via FTRL with Dilated Entropy

Each player chooses  $\tilde{x}_t$ ,  $\tilde{y}_t$  based on FTRL with dilated entropy

- For x-player  $u_t = A \tilde{y}_t$  and  $U_t = U_{t-1} + u_t$  and initialize  $Q = U_t$
- Traverse the tree bottom-up; for each infoset  $j \in \mathcal{J}_1$   $x_{t+1}^j \propto \exp(\eta_j Q^j)$ ,  $V^j = \operatorname{softmax}_{\eta_j}(Q^j)$ ,  $Q_{p_j} \leftarrow Q_{p_j} + V^j$
- Define sequence-form strategies top-down:  $\tilde{x}_{t+1}^j = \tilde{x}_{p_j} \cdot x_{t+1}^j$  Similarly, for y player

Return average of sequence-form strategies as equilibrium

#### Fast Rates

**Theorem.** If we use Optimistic FTRL instead of FTRL then we get faster convergence to a Nash equilibrium at rate 1/T instead of  $1/\sqrt{T}$ . Plus, we get last-iterate convergence instead of only average iterate convergence.

#### Monte-Carlo Stochastic Approximation of Utilities

- Calculating utilities on all nodes of the tree can be very expensive
- In linear online learning it suffices that we use an unbiased estimate of the utility vector

$$\tilde{x}_t = \underset{x \in X}{\operatorname{argmax}} \sum_{\tau < t} \langle x, \hat{u}_{\tau} \rangle - \frac{1}{\eta} \mathcal{R}(x), \qquad E[\hat{u}_{\tau} | F_{\tau}] = u_{\tau}$$

All random

before period  $\tau$ 

variables observed

- By standard martingale concentration inequality arguments, the error vanishes with the number of iterations (we will see later)
- In this setting, it suffices that we "sample a path for opponent" and that we "sample chance moves"

- Sample chance moves based on fixed distribution and opponent moves based on  $y_t$ ; Suppose, we sampled A and  $f_{\ast}$
- Go to Infoset 3

$$\widehat{U}^{3} += \left(\widehat{u}_{\hat{c}}, \widehat{u}_{\hat{f}}\right) = (-3, -2)$$

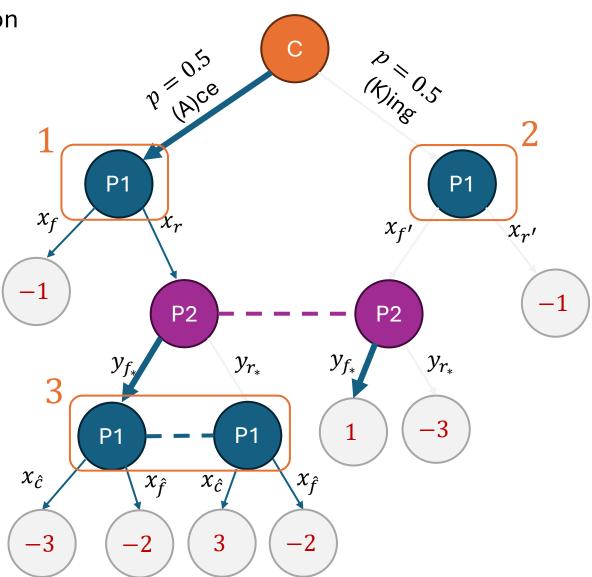
$$\widehat{Q}^{3} = \widehat{U}^{3}, \qquad x^{3} = \left(x_{\hat{c}}, x_{\hat{f}}\right) \propto \exp\left(\eta_{3} \, \widehat{Q}^{3}\right)$$

$$\widehat{V}^{3} = \operatorname{softmax}(\eta_{3} \, \widehat{Q}^{3})$$

$$\widehat{U}^1 += (\widehat{u}_f, \widehat{u}_r) = (-1, 0)$$

$$\widehat{Q}^1 = \widehat{U}^1 + (0, \widehat{V}^3) = (-1, \widehat{V}^3)$$

$$x^1 = (x_f, x_r) \propto \exp(\eta_1 \widehat{Q}^1)$$



- Equivalently top down and evaluate recursively
- Sample chance move (e.g. sampled A)
- Go to Infoset 1

$$\widehat{U}_f += -1, \qquad \widehat{U}_r += 0, \qquad \widehat{Q}_f = \widehat{U}_f, \qquad \widehat{Q}_r = \widehat{U}_r$$

- Recursively go down tree after action r
- Sample P2 move (e.g. sampled  $f_*$ )
- Go down to Infoset 3

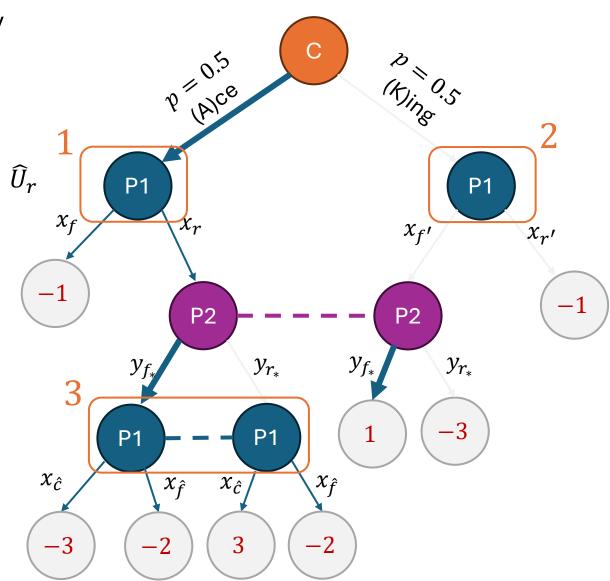
$$\widehat{U}^{3} += (-3, -2)$$

$$\widehat{Q}^{3} = \widehat{U}^{3}, \qquad x^{3} = \left(x_{\hat{c}}, x_{\hat{f}}\right) \propto \exp\left(\eta_{3} \widehat{Q}^{3}\right)$$

$$\widehat{Q}_{r} += \widehat{V}^{3} = \operatorname{softmax}(\eta_{3} \widehat{Q}^{3})$$

Go back up to Infoset 1;

$$x^1 = (x_f, x_r) \propto \exp(\eta_1 Q^1)$$



#### **Local Dynamics**

- These dynamics seem to be doing "local updates" at each node
- They came out of a specific algorithm FTRL with Dilated Entropy
- Is this a general paradigm?
- Can we decompose the no-regret learning problem into local noregret learners at each node?
- What feedback should each node receive from the learners in nodes below?
- What loss should each learner be optimizing?

## Counterfactual Regret Minimization (CRM)

**Interpretation of**  $u_a$ **.** If I play with the intend to arrive at action a (i.e.  $\tilde{x}_a = 1$ ) and then don't make any other moves, what is the expected reward that I will collect, in expectation over the choices of my opponent and nature

What if we now want to express: If I play with the intend to arrive at action a (i.e.  $\tilde{x}_a = 1$ ) and then continue playing based on some behavioral policy x, what is the expected reward that I will collect, in expectation over the choices of my opponent and nature

**Interpretation of u\_a.** If I play with the intend to arrive at action a (i.e.  $\tilde{x}_a = 1$ ) and then don't make any other moves, what is the expected reward that I will collect, in expectation over the choices of my opponent and nature

What if we now want to express: If I play with the intend to arrive at action a (i.e.  $\tilde{x}_{\alpha} =$ 1) and then continue playing based on some behavioral policy x, what is the expected reward that I will collect, in expectation over the choices of my opponent and nature

• Let  $C_a$  be all infosets of the player that are reachable as next infosets after playing a

$$\widetilde{u}_a(x) = \underbrace{\left\{u_a\right\}}_{k \in C_a} + \underbrace{\left\{V^k(x)\right\}}_{pass \ through \ infoset \ k, \ if \ I \ continue}_{playing \ based \ on \ behavioral \ strategy \ x}$$
 this is the last action I play

playing based on behavioral strategy x

**Interpretation of**  $u_a$ **.** If I play with the intend to arrive at action a (i.e.  $\tilde{x}_a = 1$ ) and then don't make any other moves, what is the expected reward that I will collect, in expectation over the choices of my opponent and nature

What if we now want to express: If I play with the intend to arrive at action a (i.e.  $\tilde{x}_a = 1$ ) and then continue playing based on some behavioral policy x, what is the expected reward that I will collect, in expectation over the choices of my opponent and nature

• Let  $C_a$  be all infosets of the player that are reachable as next infosets after playing a

$$\tilde{u}_a(x) = [u_a] + \sum_{k \in C_a} [V^k(x)]$$
 Continuation E[utility] from paths that pass through infoset  $k$ , if I continue playing based on behavioral strategy  $x$  this is the last action I play

• Continuation utility  $V^{j}(x)$  from paths that pass through infoset j recursively defined:

$$V^{j}(x) = \sum_{a \in A^{j}} x_{a} \, \tilde{u}_{a}(x) = \left[\sum_{a \in A^{j}} x_{a} u_{a}\right] + \left[\sum_{a \in A^{j}} x_{a} \left(\sum_{k \in C_{a}} V^{k}(x)\right)\right]$$
"Instantaneous utility", if "Continuation utility", if I this is the last move I make continue playing based on x

• Continuation utility  $V^j(x)$  from paths that pass through j, assuming I play to arrive deterministically at the parent action  $p_j$  (i.e.,  $\tilde{x}_{p_j}=1$ )

$$V^{j}(x) = \sum_{a \in A^{j}} x_{a} \, \tilde{u}_{a}(x) = \sum_{a \in A^{j}} x_{a} \left( u_{a} + \sum_{k \in C_{a}} V^{k}(x) \right)$$

- Obviously  $V^{\text{root}}(x)$  is total expected utility from behavior strategy x
- From equivalence of behavioral and sequence-form strategies

$$V^{\text{root}}(x) = \langle \tilde{x}, u \rangle$$

The same also holds for regrets

$$R^{\text{root}}(x) = \max_{x'} V^{\text{root}}(x') - V^{\text{root}}(x) = \max_{\tilde{x}' \in X} \langle \tilde{x}', u \rangle - \langle \tilde{x}, u \rangle = R(\tilde{x})$$

#### Local Regrets

ullet We can also define infoset regrets based on local utilities  $ilde{u}_a$ 

$$R^{j}(x) = \max_{x'} V^{j}(x') - V^{j}(x) = \max_{x'} \sum_{a} x'_{a} \tilde{u}_{a}(x') - x_{a} \tilde{u}_{a}(x)$$

Right-hand-side can be decomposed as:

$$\max_{x'} \left| \sum_{a} x'_a \tilde{u}_a(x) - x_a \tilde{u}_a(x) \right| + \left| \sum_{a} x'_a \left( \tilde{u}_a(x') - \tilde{u}_a(x) \right) \right|$$

Fix continuation strategy to current strategy and only change the behavioral strategy at the current infoset Weighted average of changes in continuation strategy

#### Local Regrets

ullet We can also define infoset regrets based on local utilities  $ilde{u}_a$ 

$$R^{j}(x) = \max_{x'} V^{j}(x') - V^{j}(x) = \max_{x'} \sum_{a} x'_{a} \tilde{u}_{a}(x') - x_{a} \tilde{u}_{a}(x)$$

Right-hand-side can be decomposed as:

$$\max_{x'} \sum_{a} x'_a \tilde{u}_a(x) - x_a \tilde{u}_a(x) + \sum_{a} x'_a \left( \tilde{u}_a(x') - \tilde{u}_a(x) \right)$$

• Maximum is upper bounded by the decoupled optima

$$\left| \max_{x'} \sum_{a} x'_a \tilde{u}_a(x) - x_a \tilde{u}_a(x) \right| + \sum_{a} \max_{x'} \left( \tilde{u}_a(x') - \tilde{u}_a(x) \right)$$

Local Regret: LR<sup>j</sup>(x)

#### Recursive Bound of Local Regrets

Infoset regrets are bounded by local regret plus continuation terms

$$R^{j}(x) \le LR^{j}(x) + \sum_{a} \max_{x'} (\tilde{u}_{a}(x') - \tilde{u}_{a}(x))$$

• The continuation terms are recursive infoset regrets!

$$\tilde{u}_a(x') - \tilde{u}_a(x) = u_a + \sum_{k \in C_a} V^k(x') - u_a - \sum_{k \in C_a} V^k(x)$$

Deriving the recursive upper bound

$$R^{j}(x) \le LR^{j}(x) + \sum_{a} \sum_{k \in C_{a}} \max_{x'} V^{k}(x') - V^{k}(x)$$
$$\le LR^{j}(x) + \sum_{a} \sum_{k \in C_{a}} R^{k}(x)$$

#### Recursive Bound of Local Regrets

Deriving the recursive upper bound

$$R^{j}(x) \le LR^{j}(x) + \sum_{a} \sum_{k \in C_{a}} R^{k}(x)$$

#### Recursive Bound of Local Regrets

Deriving the recursive upper bound

$$R^{j}(x) \le LR^{j}(x) + \sum_{a} \sum_{k \in C_{a}} R^{k}(x)$$

Theorem. By induction:

$$R^{j}(x) \le LR^{j}(x) + \sum_{k \text{ eventually reachable from } j} LR^{k}(x)$$

#### Local Regrets Upper Bound Total Regret

Deriving the recursive upper bound

$$R^{j}(x) \le LR^{j}(x) + \sum_{a} \sum_{k \in C_{a}} R^{k}(x)$$

Theorem. By induction:

$$R^{j}(x) \le LR^{j}(x) + \sum_{k \text{ eventually reachable from } j} LR^{k}(x)$$

Main Corollary. Regret is upper bounded by sum of local regrets

$$R(\tilde{x}) = R^{\text{root}}(x) \le \sum_{k \in \mathcal{J}_1} LR^k(x)$$

#### Regret over Time

Same inequalities can be followed for the average regret over time

$$R = \max_{\tilde{x}' \in X} \frac{1}{T} \sum_{t} \langle \tilde{x}', u_t \rangle - \langle \tilde{x}_t, u_t \rangle$$

$$LR^{j} = \max_{x^{j}} \frac{1}{T} \sum_{t} \langle x^{j}, \tilde{u}_{t}(x_{t}) \rangle - \langle x_{t}^{j}, \tilde{u}_{t}(x_{t}) \rangle$$

Main CFR Theorem. Regret is upper bounded by local regrets

$$R \le \sum_{j \in \mathcal{L}_1} LR^j$$

## Achieving vanishing Local Regrets

$$LR^{j}(x) = \max_{x^{j}} \frac{1}{T} \sum_{t} \langle x^{j}, \widetilde{u}_{t}(x_{t}) \rangle - \langle x_{t}^{j}, \widetilde{u}_{t}(x_{t}) \rangle$$

#### Counterfactual Regret Minimization

Device local regret algorithms for local regret

$$LR^{j}(x) = \max_{x^{j}} \frac{1}{T} \sum_{t} \langle x^{j}, \tilde{u}_{t}(x_{t}) \rangle - \langle x_{t}^{j}, \tilde{u}_{t}(x_{t}) \rangle$$

• Standard n-action no-regret problem: reward vector at period t is  $\tilde{u}^j(x_t)$  and reward for choice  $x^j$  is  $\langle x^j, \tilde{u}^j(x_t) \rangle$ 

- At period t run bottom-up recursion to calculate  $\tilde{u}^j(x_t)$  for  $j \in \mathcal{J}_1$
- Update probabilities  $x_{t+1}^j$  using reward vectors  $\tilde{u}^j(x_t)$  for  $j\in\mathcal{J}_1$

• Go to Infoset 3 
$$(\tilde{u}_{\hat{c}}, \tilde{u}_{\hat{f}}) = \left(-3\frac{1}{2}y_{f_*} + 3\frac{1}{2}y_{r_*}, -2\frac{1}{2}y_{f_*} - 2\frac{1}{2}y_{r_*}\right)$$

• Go to Infoset 3 
$$\left(\tilde{u}_{\hat{c}}, \tilde{u}_{\hat{f}}\right) = \left(-3\frac{1}{2}y_{f_*} + 3\frac{1}{2}y_{r_*}, -2\frac{1}{2}y_{f_*} - 2\frac{1}{2}y_{r_*}\right)$$

Go to Infoset 3
$$\left(\tilde{u}_{\hat{c}}, \tilde{u}_{\hat{f}}\right) = \left(-3\frac{1}{2}y_{f_*} + 3\frac{1}{2}y_{r_*}, -2\frac{1}{2}y_{f_*} - 2\frac{1}{2}y_{r_*}\right)$$

$$V^3 \leftarrow x_{\hat{c}}\tilde{u}_{\hat{c}} + x_{\hat{f}}\tilde{u}_{\hat{f}}$$

$$V^3 \leftarrow x_{\hat{c}}\tilde{u}_{\hat{c}} + x_{\hat{f}}\tilde{u}_{\hat{f}}$$

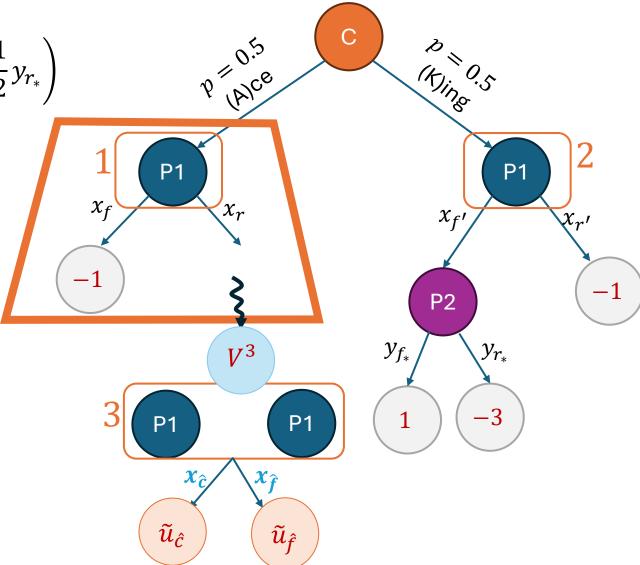
$$V^3 \leftarrow x_{\hat{c}}\tilde{u}_{\hat{c}} + x_{\hat{f}}\tilde{u}_{\hat{f}}$$

Go to Infoset 3

$$\left(\tilde{u}_{\hat{c}}, \tilde{u}_{\hat{f}}\right) = \left(-3\frac{1}{2}y_{f_*} + 3\frac{1}{2}y_{r_*}, -2\frac{1}{2}y_{f_*} - 2\frac{1}{2}y_{r_*}\right)$$

 $V^3 \leftarrow x_{\hat{c}} \tilde{u}_{\hat{c}} + x_{\hat{f}} \tilde{u}_{\hat{f}}$ 

$$\left(\tilde{u}_f, \tilde{u}_r\right) = \left(-1\frac{1}{2}, V^3\right)$$

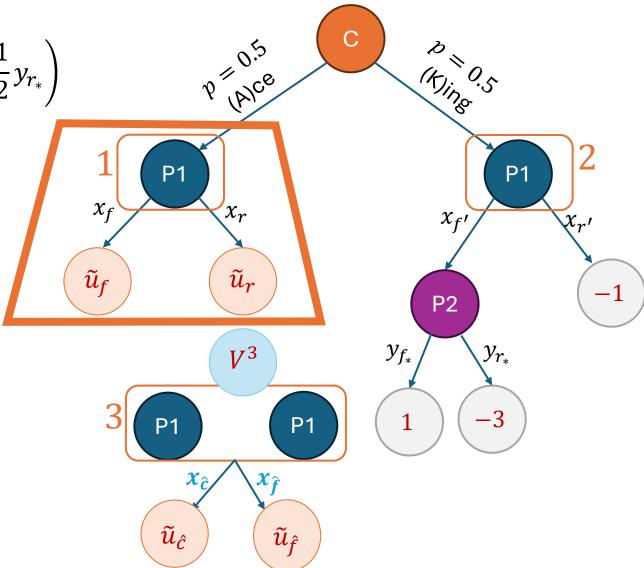


Go to Infoset 3

$$\left(\tilde{u}_{\hat{c}}, \tilde{u}_{\hat{f}}\right) = \left(-3\frac{1}{2}y_{f_*} + 3\frac{1}{2}y_{r_*}, -2\frac{1}{2}y_{f_*} - 2\frac{1}{2}y_{r_*}\right)$$

 $V^3 \leftarrow x_{\hat{c}} \tilde{u}_{\hat{c}} + x_{\hat{f}} \tilde{u}_{\hat{f}}$ 

$$\left(\tilde{u}_f, \tilde{u}_r\right) = \left(-1\frac{1}{2}, V^3\right)$$

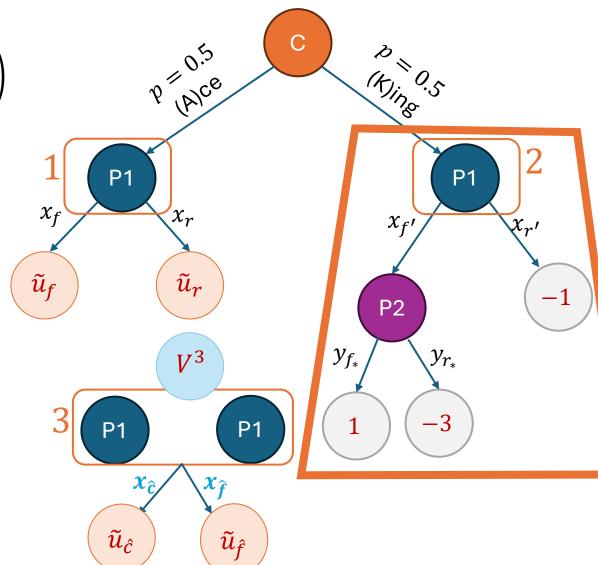


Go to Infoset 3

Go to Infoset 1

$$\left(\tilde{u}_f, \tilde{u}_r\right) = \left(-1\frac{1}{2}, V^3\right)$$

$$\left(\tilde{u}_{f'}, \tilde{u}_{r'}\right) = \left(1\frac{1}{2}y_{f_*} - 3\frac{1}{2}y_{r_*}, -1\frac{1}{2}\right)$$

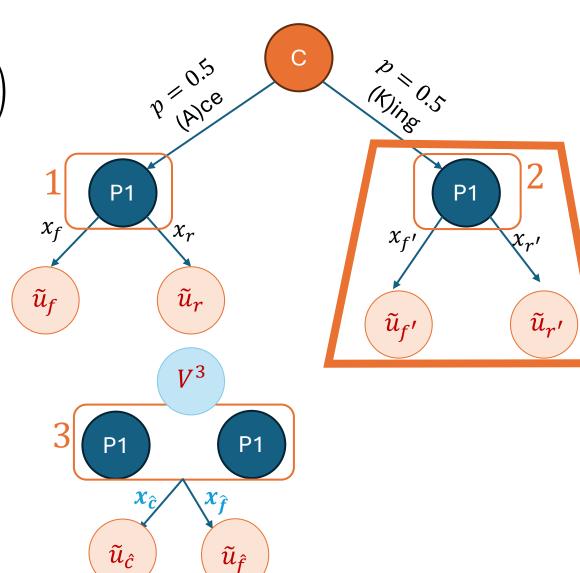


Go to Infoset 3

Go to Infoset 1

$$\left(\tilde{u}_f, \tilde{u}_r\right) = \left(-1\frac{1}{2}, V^3\right)$$

$$\left(\tilde{u}_{f'}, \tilde{u}_{r'}\right) = \left(1\frac{1}{2}y_{f_*} - 3\frac{1}{2}y_{r_*}, -1\frac{1}{2}\right)$$



Go to Infoset 3

Go to Infoset 1

$$\left(\tilde{u}_f, \tilde{u}_r\right) = \left(-1\frac{1}{2}, V^3\right)$$

Go to Infoset 2

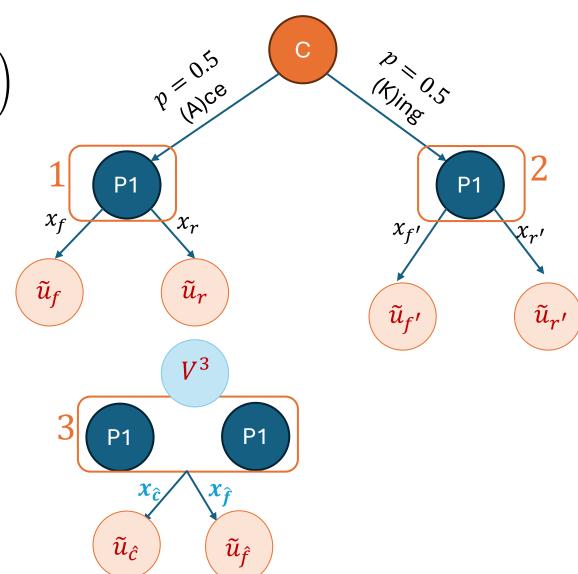
$$\left(\tilde{u}_{f'}, \tilde{u}_{r'}\right) = \left(1\frac{1}{2}y_{f_*} - 3\frac{1}{2}y_{r_*}, -1\frac{1}{2}\right)$$

Update probabilities

$$(x_f, x_r) \leftarrow \text{Update}(\tilde{u}_f, \tilde{u}_r)$$

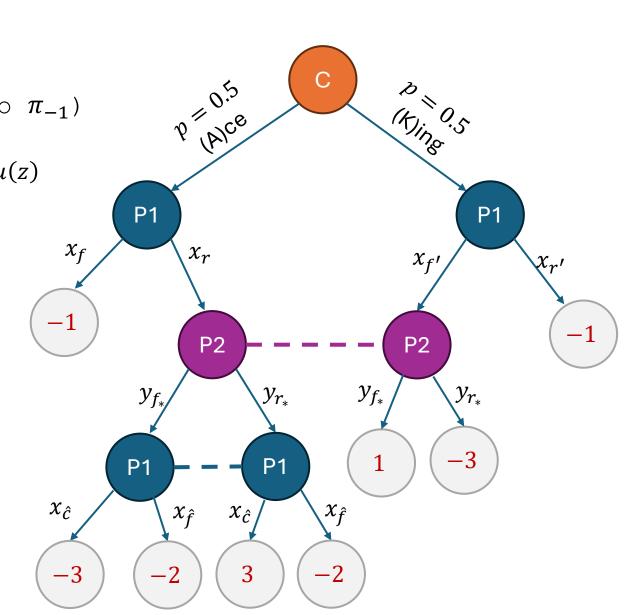
$$(x_{f'}, x_{r'}) \leftarrow \text{Update}(\tilde{u}_{f'}, \tilde{u}_{r'})$$

$$(x_{\hat{c}}, x_{\hat{f}}) \leftarrow \text{Update}(\tilde{u}_{\hat{c}}, \tilde{u}_{\hat{f}})$$



### Recursive Algorithm

```
Value (ActionHistory h, AccOtherProb \pi_{-1})
     Let I be infoset corresponding to h
     If I is terminal node z return \pi_{-1} \cdot u(z)
     If Player(I) = chance
          Return \sum_{a \in A_I} \text{Value}(ha, \pi_{-1}\pi_a^c)
     If Player(I) = 2
          Return \sum_{a \in A_I} Value(ha, \pi_{-1}y_a)
     If Player(I) = 1
          For a \in A_I: \tilde{u}_a += Value(ha, \pi_{-1})
          Return \sum_{a \in A_I} x_a \cdot \text{Value}(ha, \pi_{-1})
Value (\emptyset, 1)
```



### Recursive Algorithm

```
Value (ActionHistory h, AccOtherProb \pi_{-1})
Let I be infoset corresponding to h

If I is terminal node z return \pi_{-1} \cdot u(z)

If Player(I) = chance

Return \sum_{a \in A_I} \text{Value}(ha, \pi_{-1}\pi_a^c)

If Player(I) = 2

Return \sum_{a \in A_I} \text{Value}(ha, \pi_{-1}y_a)

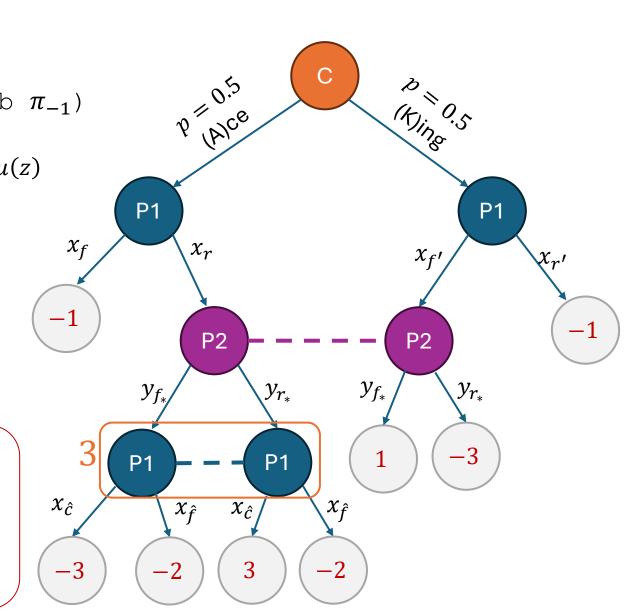
If Player(I) = 1

For a \in A_I: \tilde{u}_a += \text{Value}(ha, \pi_{-1})

Return \sum_{a \in A_I} x_a \cdot \text{Value}(ha, \pi_{-1})
```

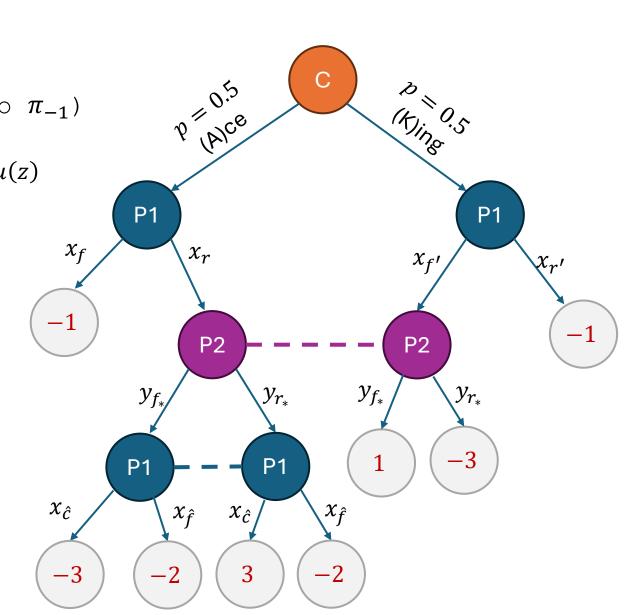
We arrive at the same infoset I multiple times, once for each node in the set;  $\tilde{u}_a$  accumulates continuation utility from taking action a from all these possible "arrival paths".

**Example.** In infoset 3 we arrive once on the left node and add  $-3\frac{1}{2}y_{f_*}$  and once on the right node and add  $3\frac{1}{2}y_{r_*}$  to  $u_{\hat{c}}$ 



### Recursive Algorithm

```
Value (ActionHistory h, AccOtherProb \pi_{-1})
     Let I be infoset corresponding to h
     If I is terminal node z return \pi_{-1} \cdot u(z)
     If Player(I) = chance
          Return \sum_{a \in A_I} \text{Value}(ha, \pi_{-1}\pi_a^c)
     If Player(I) = 2
          Return \sum_{a \in A_I} Value(ha, \pi_{-1}y_a)
     If Player(I) = 1
          For a \in A_I: \tilde{u}_a += Value(ha, \pi_{-1})
          Return \sum_{a \in A_I} x_a \cdot \text{Value}(ha, \pi_{-1})
Value (\emptyset, 1)
```

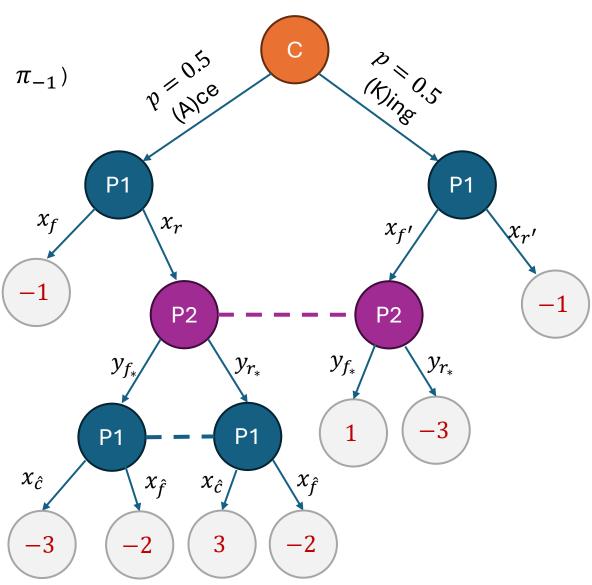


### Equivalent Recursive Algorithm

```
P (A)Ce
CValue (ActionHistory h, AccOtherProb \pi_{-1})
     Let I be infoset corresponding to h
     If I is terminal node z return u(z)
     If Player(I) = chance
           Return \sum_{a \in A_l} \pi_a^c: CValue(ha, \pi_{-1}\pi_a^c)
                                                                           \chi_f
                                                                                          \chi_r
                                                                                                                      x_{f'}
     If Player(I) = 2
           Return \sum_{a \in A} (y_a) CValue(ha, \pi_{-1}y_a)
                                                                         -1
     If Player(I) = 1
           For a \in A_I: \tilde{u}_a + = (\pi_{-1}) \cdot \text{CValue}(ha, \pi_{-1})
                                                                                                                y_{f_*}
                                                                                                                            y_{r_*}
                                                                                                    y_{r_*}
           Return \sum_{a \in A_I} x_a \cdot \text{CValue}(ha, \pi_{-1})
                                                                                                    P1
                                                                                   P1
CValue (\emptyset, 1)
                                                                         \chi_{\hat{c}}
                                                                                        \chi_{\hat{f}}
                                                                                                           \chi_{\hat{f}}
                                                                                                \chi_{\hat{c}}
```

#### The Typical CRM Algorithm Implementation

```
CValue (ActionHistory h, AccOtherProb \pi_{-1})
     Let I be infoset corresponding to h
     If I is terminal node z return u(z)
     If Player(I) = chance
          Return \sum_{a \in A_I} \pi_a^c \cdot \text{CValue}(ha, \pi_{-1} \pi_a^c)
     If Player(I) = 2
          Return \sum_{a \in A_I} y_a \cdot \text{CValue}(ha, \pi_{-1}y_a)
     If Player(I) = 1
          For a \in A_I: \tilde{u}_a += \pi_{-1} \cdot \text{CValue}(ha, \pi_{-1})
          Return \sum_{a \in A_I} x_a \cdot \text{CValue}(ha, \pi_{-1})
CValue (\emptyset, 1)
```



# Recovering Equilibrium from CRM Dynamics

We have run CRM dynamics generating behavioral strategies  $x_t$ ,  $y_t$  for T periods.

How do we calculate the behavioral strategies  $x^*$ ,  $y^*$  that are an approximate Nash equilibrium?

#### Recovering Nash Equilibrium

We need to translate the behavioral strategies into sequence-form

$$\forall a \in A_j \colon \tilde{x}_{t,a} = \tilde{x}_{t,p_j} \mid x_t$$

• Then average the sequence-form strategies

Product of probabilities of actions of player P1 on path to infoset of action *i* 

$$\bar{\tilde{x}} = \frac{1}{T} \sum_{t=1}^{T} \tilde{x}_t$$

• Then translate back to equilibrium behavioral strategies  $x^*$ 

$$\forall a \in A_j \colon x_a^* = \frac{\bar{\tilde{x}}_a}{\bar{\tilde{x}}_{p_j}}$$

#### Recovering Nash Equilibrium

We need to translate the behavioral strategies into sequence-form

$$\forall a \in A_j \colon \tilde{x}_{t,a} = \tilde{x}_{t,p_j} \colon x_t$$

• Then average the sequence-form strategies \to Product of probabilities of actions of

player P1 on path to infoset of action i

$$\bar{\tilde{x}} = \frac{1}{T} \sum_{t=1}^{T} \tilde{x}_t = \frac{1}{T} \sum_{t=1}^{T} \tilde{x}_{t,p_j} \cdot x_t$$

• Then translate back to equilibrium behavioral strategies  $x^*$ 

$$\forall a \in A_j \colon x_a^* = \frac{\bar{\tilde{x}}_a}{\bar{\tilde{x}}_{p_j}} = \frac{\sum_{t=1}^T \tilde{x}_{t,p_j} \cdot x_{t,a}}{\sum_{t=1}^T \tilde{x}_{t,p_j}}$$

#### The Typical CRM Algorithm Implementation

```
CValue (ActionHistory h, AccOtherProb \pi_{-1}, AccProb \pi_{1})
     Let I be infoset corresponding to h
     If I is terminal node z return u(z)
     If Player(I) = chance
          Return \sum_{a \in A_I} \pi_a^c \cdot \text{CValue}(ha, \pi_{-1} \pi_a^c, \pi_1)
     If Player(I) = 2
          Return \sum_{a \in A_1} y_a \cdot \text{CValue}(ha, \pi_{-1}y_a, \pi_1)
     If Player(I) = 1
          For a \in A_I: \tilde{u}_a += \pi_{-1} \cdot \text{CValue}(ha, \pi_{-1}, \pi_1 x_a)
         Set q(I) = \pi_1
          Return \sum_{a \in A_1} x_a \cdot \text{CValue}(ha, \pi_{-1}, \pi_1 x_a)
```

 $CValue(\emptyset, 1)$ 

This is the product of the probabilities of prior actions of player P1 before arriving at infoset I

**Note.** Due to perfect recall this product is the same every time we visit the infoset; irrespective of which node of the infoset we arrived at.

### The Overall Equilibrium Algorithm with CRM

#### After each period *t*:

- With last period behavior strategies  $x_t, y_t$  call  $\text{CValue}(\emptyset, 1, 1)$
- Store  $\tilde{u}_{t,a}$  and  $q_t(I)$  for each action a and infoset I of P1
- Symmetrically, do so for player P2
- Update strategies at all information sets

$$\forall j \in \mathcal{J}_1: x_{t+1}^j \leftarrow \text{Update}\left(\tilde{u}_t^j\right), \quad \forall j \in \mathcal{J}_2: y_{t+1}^j \leftarrow \text{Update}\left(\tilde{u}_t^j\right)$$

$$\forall I \in \mathcal{I}_1 \forall a \in A_I : x_a^* = \frac{\sum_t q_t(I) x_{t,a}}{\sum_t q_t(I)}$$

$$\forall I \in \mathcal{I}_2 \forall a \in A_I : y_a^* = \frac{\sum_t q_t(I) y_{t,a}}{\sum_t q_t(I)}$$

Approximate Equilibrium Strategies

# What algorithm to use for local regret updates?

#### Recursive Value Calculation

#### After each period *t*:

- With last period behavior strategies  $x_t, y_t$  call CValue( $\emptyset, 1, 1$ )
- Store  $\tilde{u}_{t,a}$  and  $q_t(I)$  for each action a and infoset I of P1

• For each infoset 
$$j \in \mathcal{J}_1$$
: 
$$x_{t+1} \leftarrow \operatorname{Update}(\tilde{u}_t) \leftarrow$$

• Symmetrically, do so for player P2

Any no-regret algorithm for the *n*-action no-regret problem can be used, e.g. FTRL, OFTRL, EXP, etc.

What performs well in practice is what is known as Regret Matching!

$$\forall I \in \mathcal{I}_1 \forall a \in A_I : x_a^* = \frac{\sum_t q_t(I) x_{t,a}}{\sum_t q_t(I)}$$

$$\forall I \in \mathcal{I}_2 \forall a \in A_I : y_a^* = \frac{\sum_t q_t(I) y_{t,a}}{\sum_t q_t(I)}$$

Approximate Equilibrium Strategies

### Regret Matching and Regret Matching+

- Consider the n action no-regret learning setting; at each period we choose  $x_t \in \Delta(n)$ , observe utility vector  $u_t$  and get utility  $\langle x_t, u_t \rangle$
- At each period t calculate regret of not playing action a

$$r_{t,a} = u_{t,a} - \langle u_t, x_t \rangle$$

• Calculate cumulative regret of not playing action a

$$R_{t,a} = \sum_{\tau \le t} r_{t,a} = R_{t-1,a} + r_{t,a}$$

Choose next distribution, proportional to positive part of regret

$$x_{t+1,a} \propto \left[ R_{t,a} \right]^+ \coloneqq \max \{ R_{t,a}, 0 \}$$

People typically refer to CFR with RegretMatching as simply "CFR"

## Regret Matching+

- Consider the n action no-regret learning setting; at each period we choose  $x_t \in \Delta(n)$ , observe utility vector  $u_t$  and get utility  $\langle x_t, u_t \rangle$
- At each period t calculate regret of not playing action a

$$r_{t,a} = u_{t,a} - \langle u_t, x_t \rangle$$

ullet Continuously clip above zero, as you accumulate regret of a

$$R_{t,a} = [R_{t-1,a} + r_{t,a}]^{+}$$

• Choose next distribution, proportional to  $R_{t,a}$ 

$$x_{t+1,a} \propto R_{t,a}$$

• Regret Matching and Regret Macthing+ achieve Regret  $\leq \sqrt{n/T}$ 

# Extra Tricks for Empirical Improvement

#### Monte-Carlo Stochastic Approximation of Utilities

- Sample chance move (e.g. sampled A)
- Go to Infoset 1

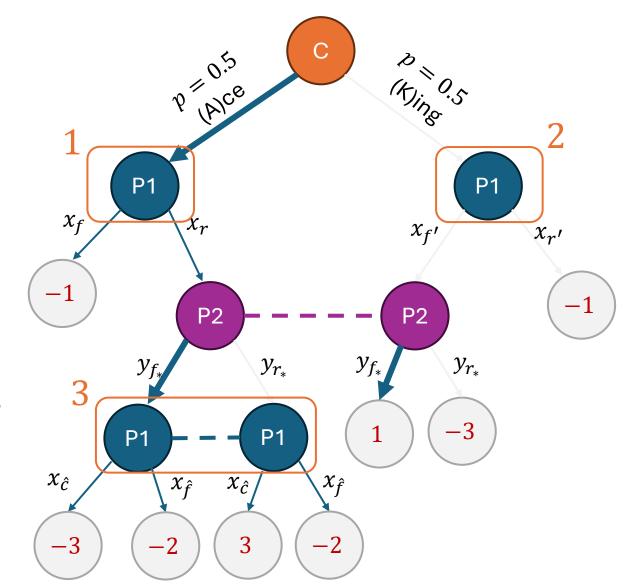
$$\hat{\tilde{u}}_f = -1, \qquad \hat{\tilde{u}}_r = 0$$

- Go down tree the r path
- Sample P2 move (e.g. sampled  $f_*$ )
- Go down to Infoset 3

$$\hat{\tilde{u}}_{\hat{c}} = -3, \qquad \hat{\tilde{u}}_{\hat{f}} = -1$$

$$\hat{\tilde{u}}_r += x_{\hat{c}}\hat{\tilde{u}}_{\hat{c}} + x_{\hat{f}}\hat{\tilde{u}}_{\hat{f}}$$

• Update probabilities of visited infosets  $(x_f, x_r) \leftarrow \text{Update}(\hat{u}_f, \hat{u}_r)$   $(x_{\hat{c}}, x_{\hat{f}}) \leftarrow \text{Update}(\hat{u}_{\hat{c}}, \hat{u}_{\hat{f}})$ 



#### Typical Monte Carlo Algorithm Implementation

```
MCCValue (ActionHistory h, AccProb \pi_1)
    Let I be infoset corresponding to h
     If I is terminal node z return u(z)
     If Player(I) = chance
         Sample a \sim \pi^{C}
         Return MCCValue(ha, \pi_1)
     If Player(I) = 2
         Sample a \sim y^I
         Return MCCValue(ha, \pi_1)
     If Player(I) = 1
         For a \in A_I: \tilde{u}_a += \text{MCCValue}(ha, \pi_1 \cdot x_a)
         Set q(I) = \pi_1
         Return \sum_{a \in A_1} x_a \cdot \text{MCCValue}(ha, \pi_1 \cdot x_a)
Value(\emptyset, 1)
```

#### Can Combine with Update Step in One Pass

```
MCCValue (ActionHistory h, AccProb \pi_1)
     Let I be infoset corresponding to h
     If I is terminal node z return u(z)
     If Player(I) = chance
          Sample a \sim \pi^{C}
          Return MCCValue(ha, \pi_1)
     If Player(I) = 2
          Sample a \sim y^I
          Return MCCValue(ha, \pi_1)
     If Player(I) = 1
          For a \in A_I: \tilde{u}_a += \text{MCCValue}(ha, \pi_1 \cdot x_a)
          Set q(I) = \pi_1
          Update x_{\text{next}}^I \leftarrow \text{Update}(\tilde{u}^I)
          Return \sum_{a \in A_1} x_a \cdot \text{MCCValue}(ha, \pi_1 \cdot x_a)
```

#### Alternation

After each period *t*:

- If t is odd then update the strategy of the x-player
- If t is even then update strategy of the y-player

For most natural algorithms, alternation can only help in terms of reducing the violation of best response constraints!

Can converge faster to equilibrium

## Weighted Averaging

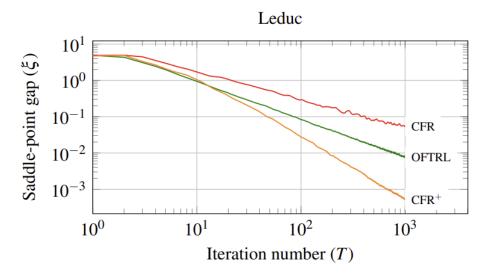
 Instead of uniformly weighting all rounds, put more weight on more recent rounds of play

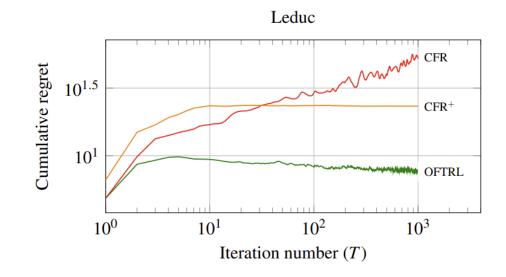
$$\frac{1}{\sum_{t} t^{\alpha}} \sum_{t} t^{\alpha} \tilde{x}_{t}$$

• Typically, one uses linear averaging (i.e.,  $\alpha=1$ )

 The CFR algorithm that uses RegretMatching+, alternation and linear averaging is typically referred to as "CFR+"

#### **Empirical Comparisons**





Violations of best response

 $\left\{ \text{Regret}_{y}(x_{*}, y_{*}) + \text{Regret}_{x}(x_{*}, y_{*}) \right\} = \max_{y} x_{*}^{\mathsf{T}} A y - x_{*}^{\mathsf{T}} A y_{*} + x_{*}^{\mathsf{T}} A y_{*} - \min_{x} x^{\mathsf{T}} A y_{*} = \left\{ \max_{y} x_{*}^{\mathsf{T}} A y - \min_{x} x^{\mathsf{T}} A y - \min_{x} x^{\mathsf{T}} A y_{*} \right\}$ 

$$\begin{bmatrix} R_y + R_x \end{bmatrix} = \max_y \bar{x}^\top A y - \frac{1}{T} \sum_t x_t^\top A y_t + \frac{1}{T} \sum_t x_t^\top A y_t - \min_x x^\top A \bar{y} = \begin{bmatrix} \max_y \bar{x}^\top A y - \min_x x^\top A \bar{y} \end{bmatrix}$$

Sum of learning algorithm regrets

saddle-point gap of average strategies  $\bar{x}$ ,  $\bar{y}$ 

saddle-point gap

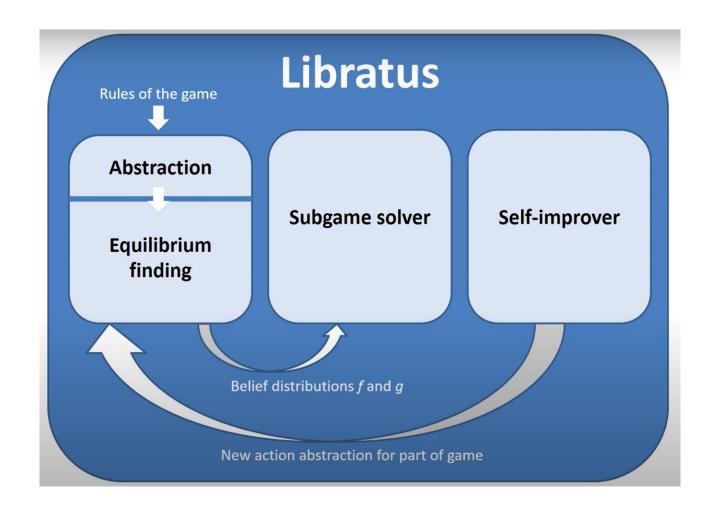
#### Elements of the Libratus Al

• The first agent to achieve superhuman performance in two player No-Limit Texas Hold'em poker ( $10^{161}$  decision points)

• Prior best was Limit Texas Hold'em ( $10^{13}$  decision points); solution is basically "run CFR+"

For No-Limit Texas Hold'em game is too big for this approach!

#### **Elements of Libratus Al**



Credits: Superhuman AI for heads-up no-limit poker: Libratus beats top professionals (youtube.com)