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## COM SCI 35L

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## Topics

- **File systems** - organized/organizing collection of data on [persistent] secondary storage
- **Scripting** - the practice of writing “quick and dirty” code, as opposed to more rigorous traditional programming
  - More maintainable, faster to write, but possibly slower, less robust
    - A script will generally be geared towards a single specific task, whereas a program may support many different tasks
  - Scripting languages: Shell, Python, JavaScript (, Lisp)
- **Building & distribution**
- **Version control/configuration management**: maintaining separate versions of a program
  - May distinguish between dev versions vs user versions, or even dev versions vs other dev versions (in the case of collaborative development)
  - Version control typically done using *Git*
    - Users’ view of a repository may differ from the actual internals
- **Low-level debugging** (for OS, using C)
- **Client-server programs** - a form of *distributed application* (operating on >1 computer)
- **CS130**:
  - Programming, data design, integration (i.e. combining different software components), configuration (taking a program for a general purpose, and modifying it to perform a specific task successfully), testing, forensics (retrospective attitude, as opposed to proactive debugging)
    - Configuration - LAUSD payroll system failed due to improper configuration

## Programs & Applications

- Categories of programs:
  - Standalone applications (e.g. machine code) - IoT
  - Applications atop an OS (e.g. terminal scripts, e.g. Emacs)
  - Web applications/PWAs operate atop a browser, atop an OS
    - Utilize a *client-server approach*
- Layers of a computer (lowest to highest):
  - The **hardware** is commanded by instructions from the *kernel* & *applications*
  - The **kernel** (operating space) runs in the kernel space, invisible to *applications*
    - Walls off applications from executing certain instructions on the hardware, allowing only basic safe processes + a limited set of proper system calls
      - Most operations used by applications are just basic load/add/mul
  - **Applications** (*processes*) run atop the kernel, each in their own application space
    - May be built upon the C library, with additional libraries (and eventually the application itself) stacked on top
    - Each application is partly isolated from other applications - can generally only communicate indirectly (e.g. through the kernel (via file system), network, etc.) or through the use of pipes
- Programs vs processes
  - *Program* - a relatively static object (downloaded at some point, doesn't change)
  - *Process* - a program being run by the machine, taking input and producing output
    - Program + internal data + connections to the external world (e.g. IO, file inputs/outputs, network connections, etc.)
    - Can also start/stop its own *subsidiary processes*, pass its own data as input/take the subsidiary's output as input via pipes
    - Output need not be to the screen as text - can also be to a windowing system, which can itself pass input back
      - Cloud systems may forego changes beyond terminal text changes
  - Ex: an Emacs *program* is just a binary file/directory somewhere on the system; an Emacs *process* is an actual running instance of the program.

## Emacs

- Written in C at the bottom level; higher levels at Lisp [Emacs Lisp]
  - C handles low-level IO, signals, Lisp interpreter [~15%]
    - Kept relatively simple; handles any tasks that might take up lots of system resources, CPU instructions, low-level objects
      - Is not properly object-oriented, but mimics the structure
  - Lisp [~85%] as a scripting language - handles most complex decisions, control
    - Does not handle most CPU-intensive tasks (passes them off to the C layer), but makes most decisions
  - Browser low/high-level differentiation: C++ vs JavaScript/Python
- Notes on Emacs
  - Emacs is modeful, i.e. the actions that it takes to a specific character command depend on its mode (the current context)
  - Is capable of reflection-introspection: can modify Emacs from within itself
  - Can be reprogrammed with Lisp code (defun)
- Emacs refcard: <https://www.gnu.org/software/emacs/refcards/pdf/refcard.pdf>
- Commands:
  - Shell: M-X shell RET [M=Meta=Edit/Alt]
    - Quote next character (write to text [translating to plaintext-representable format if needed], e.g. Ctrl-C -> ^C [representation of Ctrl-C as ^C]): C-q
    - Emacs still controls character inputs - simply passes all inputs to shell
  - Enter dired (directory edit) mode: C-x d \*directory name\*
    - Rename: [dired] R
  - Disable [buffer] read-only: C-x C-q (Does not modify system read-only protections)
  - Display character information: C-u C-x =
  - Insert by hexadecimal encoding: C-x 8 RET
  - Replace string: M-x replace-string
- Set key: M-x global-set-key \n \*key\* \*function\*
- Message: message, concat
- Time: current-time-string, format-time-string
- Load file: M-x load-file

## Emacs Internals

- Emacs internal data has two notable types: *buffer* and *string*
- **Buffer** - an Emacs structure (byte sequence) in RAM representing the contents of a file
  - Are easy to edit - cheap insertion
    - Is represented as an array of bytes with a gap [garbage] in the middle, i.e. divided into two halves; modifying fills in the gap with content
      - Is resized as needed when the buffer fills - average  $O(1)$  [amortized] performance for single-character insertions
    - Cursor is always kept at boundary between the first half and gap (*the point*); first/second halves are rebalanced as needed to maintain this (inefficient to move cursor, but cheap to insert)
    - Important location - the mark (user-set)
      - Region of interest (similar to a highlighted area) is the area between the point and the mark (be it an area to be deleted, copied, etc.)
  - Used to read/write to files; given a file, will read the whole file, no matter its size
  - All contents on screen usually represented in data by a buffer/buffers
- **String** [Emacs] - an Emacs structure representing a byte sequence, similar to a C string
  - Are usually not modified (harder to modify, relative to buffers), smaller than buffers

## Character Bit Representation

- C characters represented by 8 bits
  - ASCII characters contained within a 7-bit set (+1 bit for verifying correct format)
    - Contains 127 characters: first 32 control characters (first control character - null), all printable characters, final control character [del]
  - Neither Ctrl nor Meta are actual characters themselves, just modifiers a la Shift
    - Bit representation of a Ctrl-\*char\* is identical to the bit representation of \*char\*, with the exception of one bit (second) turned off relative to \*char\*
      - Meta-\*char\* turns one bit (first) on relative to \*char\*
- Typing Ctrl-A maps to hexadecimal code 1 [Ctrl-A], Ctrl-C to code 3 [Ctrl + \*char\* is a single distinct 7-bit character, not separate Ctrl + \*char\* codes]
  - C-m interpreted as Enter (carriage return = \r), C-i as [horizontal] Tab [= \t]

## Regex

- Regular expressions - used in grep [+ other places], incl:
  - Simple text strings [patterns]
  - Simple text strings + \* [include all]

Regular Expressions	
Pattern	Matches
abc	Matches "abc"
.	Matches any single character
E*	Zero or more Es
\(bcd\)*	Any number of "bcd"s
\(bcd\){5,10}	Between 5 and 10 "bcd"s
\(bcd\){5,}	5 or more "bcd"s
\(bcd\){,10}	Up to 10 "bcd"s
\C (for any special character C, e.g. *)	Matches just C
\ + ordinary character	<i>*Behavior undefined*</i>
\ + number	Matches pattern at the regular expression of that number in the command
^E + *text*	E at line start
*text* + E\$	E at line end
[abcxy]	Matches any character a, b, c, x, y
[A-Za-z0-9]	Equivalent to [ABC...Zabc...z012...9]
[*start char*-end char*]	Matches any character between start, end chars (via ASCII encodings)
[spc~~]	Matches all ASCII characters
[^ABC]	Matches any character not A, B, C
[A^BC]	Matches A, B, C, or ^

<code>[]</code>	Matches <code>[, ]</code> (cannot be other way around)
<code>[*\.]</code>	Matches <code>*</code> , <code>\</code> , or <code>.</code> (not special in <code>[]</code> exps)
<code>[:alpha:]</code>	Matches character class alpha
<code>^[[:ascii:]]</code>	Matches all non-ASCII characters
<code>[A-Za-z_][A-Za-z0-9_]*</code>	Special: Matches C identifiers

- Regular expressions: enclose in `"` [grep]
  - `*` (asterisk), `\` (backslash), `.` (period) not special within bracket expressions
  - Character classes - alpha (alphabetic), digit (English UTF-8 locale), ascii
    - Emacs supports character class `ascii`; grep does not
      - Ascii in grep: `[NULL-DEL]` matches bit patterns 0000000 to 0111111
    - Same as in `<ctype.h>`
- Many other programming languages have their own regex versions (with potentially different syntax)
  - `grep -p` switches to Perl regex
- Option `-E` (extended regular exp. notation = ERE) prior to regex allows for omitting `\`
  - ERE adds certain additional operators: `|` (or), `?` (`=\{0, 1\}`), `+` (`=\{1, \}`)
    - GNU provides `\|` in BRE
    - `ab|cd == (ab)|(cd)` [precedence: counting operators > concatenation > or]
  - Tradeoffs - operator `|` requires backtracking capability (more expensive)
  - Vs basic regular expressions (BRE)
- Regex tools
  - `sed` - stream editor
  - `awk` - a more robust version of `sed` [Perl/Python] - more robust versions of `awk`

## Unicode

- Non-ASCII characters (e.g. Unicode)
  - ASCII only uses 7 of 8 bits in a byte - allows for an additional 128 characters, + ASCII, to be supported in a byte
- Historically - different character sets sprang up adding an additional 128 characters based on the [primarily European] language in question (e.g. French set, Greek set, etc.)
  - ISO-8859-x [x indicating specific char set]
  - Flaws - didn't work for international languages, interpretation of bit patterns relied on knowing which character set was in use
- **Unicode** - international standard assigning character codes to all characters in all supported languages internationally
  - Contains approx.  $2^{24}$  characters
  - Originally stored in 16-bit chunks (general standard for char data types)
    - 8-bits only for C unsigned short
    - Issues: not backward-compatible with pure ASCII files, eventually not big enough to fit entirety of Unicode
  - Next moved to 32-bit chunks
    - Issues: still not backwards-compatible, memory-inefficient
- **UTF-8** - international standard for representing Unicode
  - Designed to be backwards-compatible with ASCII, space efficient
  - Also contains unambiguous character boundaries
  - Encode character string into byte string; decode byte string to character string
  - Has 3 kinds of bytes:
    - 0 + [7-bit sequence] - represents an ASCII character
    - 10 + [6-bit sequence] - trailing (2nd or later) byte in a multi-byte character
    - [Sequence of 1s] + 0 + [sequence] - beginning sequence of 1s encodes # of bytes in character
      - Length encoded in base 1 - max 4 bytes
    - Allows for a character to only use up the minimum number of bytes required to represent that character
- Invalid UTF-8 bytes:
  - Indicated length is incorrect or >4



- Not all 4-byte sequences have defined character representations [Unicode not entirely filled-in - has room for expansion]
- Character is truncated [partial character] (e.g. streaming data from a network)
- An expected trailing character is not a trailing character
- Unexpected [isolated] trailing character
- Specific: Bit pattern: 11000000 (byte 1) + 10111111 (U+003F) == 00111111 (ASCII)
  - ASCII characters can be represented in longer-byte sequences [overlong encoding] - evade regular ASCII regex patterns, e.g.
  - Declared invalid, albeit only discovered 10 years after launch of UTF-8

# The Shell

## Overview

- Emacs IDE/shell/files: persistence, getting work done, writing programs
- Roles of the **shell** (sh/bash):
  - Can be used as a **command language** for running other commands
    - Can be used as a **scripting language**, reading shell commands from a script
    - Is *interactive* - a command is written and immediately executed, all on the fly, without having specified the next line
  - Can be used as a **configuration/setup language**, i.e. as a means of getting other programs configured, set-up, and running
    - In contrast with application-specific languages (ASLs) - the shell eschews performance/efficiency for ease-of-use/integration
    - Shell commands may spawn off other applications (e.g. nano, gedit), which can be both controlled by and controlling the shell
- The shell designed as a small, simple form of getting more robust applications running
  - Designed as a simple way for the user to access/run programs or the OS (i.e. as an interactive layer between programs/OS, and the user)
    - Modularity - programs as modules
- Given a line representing a command: determines what the words (tokens) in the command are, invokes the command given (if found in PATH) and passes it an argument vector representing each word in the command (e.g. `grep word1 word2 -> ["grep", "word1", "word2"]`)\ul>- Words identified by separation by spaces (number of spaces not considered)
- [command name] + [arguments]
- Backslash escapes - Backslash can be used to insert special characters (i.e. `\spc` interpreted as part of a word, not a separator)
  - Two backslashes becomes one backslash
    - Represented by two backslashes in C
  - Backslash + newline eats up the newline
- Quoting:
  - Single quoting: Any characters inside the single quotes interpreted as a single word

- Any special character can be put in as a normal character (incl. newlines), except for a single quote [apostrophe]
- Putting a single quote inside a word (e.g. a'b c'd) works and is interpreted as one word (can be used to insert apostrophe)
- Double quoting: Similar to single quotes, except a few characters are interpreted as special
  - Note: single quote (apostrophe) not special
  - Special characters:
    - \$ replaces a value with the value of the shell variable with the corresponding name (e.g. x = "hi" -> "\$x" becomes "hi")
    - Backslash used to escape
      - Two backslashes produces a single backslash
- Large source of security vulnerabilities - not knowing language quoting rules

## Shell Commands

Command	Function
cat (concatenate)	Echoes all input given in (repeats typed-in text)
ps (process status)	Lists running processes on the current terminal associated with the current user
less	Provides control over display data (similar to nano)
grep	Given a pattern and a directory path, matches all files on path with lines matching the pattern
od	Given a file, outputs any non-ASCII characters as their Unicode character encoding
test	Performs various tests (e.g. checking file existence, string comparisons)
: [single colon]	Always succeeds
kill	Given a process ID, kills said process
wait	Given a process ID, waits for said process to finish
touch	Creates new file
mkdir	Creates a new directory
rm	Given a file, removes said file
ln	Given a file and a link name, creates hard link
ls	Given a directory path, lists contents of directory
find	Search for files recursively within a directory
shuf	Given a file, shuffles and prints the lines of the file
seq	Prints a sequence of numbers

sed	Reads standard input and outputs to standard output, but edits each line before outputting it
awk	A more robust version of sed
chmod	Change file mode bits (notably read/write/execute permissions)
which	Given a shell command, prints binary file associated with said command
tr	Translate or delete characters (tr [SET1] [SET2])
sort	Given a file, sorts lines in file in alphabetical order
comm	Given two files (with lines in sorted order), compares file contents

- cat
  - If given file name(s) as argument, copies file data (in sequence) and output to standard output [terminal]
  - Reads from standard input if no arguments are provided, or an argument is "-"
    - "-" = stdin is convention
  - Mechanism: reads incrementally, e.g. in 1 MB blocks (stores 1 MB in RAM)
- ps
  - -e argument also includes system processes
  - Timestamp indicates amount of CPU time used
- Backslash indicates command continues to next line
- less - control over display of data (similar to nano)
  - /\*search term\* searches
  - Akin to a read-only text editor
  - Shell command | less passes the output of the shell command as input for less
    - Vertical line indicates a pipe (cmd1 | cmd2 -> cmd2 takes cmd1's output as input)
- Grep
  - Grep [pattern] [file] - outputs every line in file matching the pattern [regex]
  - grep -v: print only names of matching files
  - Notably, has two different failure exit statuses:

- 1: No match found, no notable errors
  - 2: Serious error (e.g. file not readable)
- Od [file] - outputs any non-ASCII characters as their Unicode character encoding
  - -c: octal, -t: hexadecimal
  - X1: 1-byte
- test
  - test -r \*file name\*: test if file is readable file
  - test -f \*file name\*: test if file exists
  - test \*string 1\* -eq \*string 2\*: test if strings equal
    - Can be used to compare command exit status (\$?) to exit codes, e.g.
    - Equal to [str1 -eq str2]
- sed "s/\*search\*/\*replace\*/2" replaces 2nd instance of search with replace in every line
  - /g for all; /3g for 3 and up; "1,3 s/..." for lines 1,3; -n only replaced lines
  - '5,\$d' deletes 5th, last line; '\*pattern\*/d' deletes pattern matches
- Chmod bits
  - \*file type\*rwxrwx...; file type can be 'f', 'd', 'l', etc.
- Tr [SET1] [SET2]; -c = "complement"
- true, false
- set a b c: sets \$1, \$2, \$3 to a, b, c
- shift: Shifts \$1, \$2, etc. (arguments) right by 1
- Shortcuts
  - Control-D/^D/C-D indicates EOF if at start of line, from a terminal
  - Control-C interrupts the current process
- Features
  - Write multiple commands on one line - separate with semicolons
  - Appending & to the end of a command will run it in the background as a separate process (as a descendant of the current shell)
- /dev/null: always-empty file (will always return nothing when read from)

## Shell Keywords

- Some command names [words] reserved by the shell (designated keywords) for compound commands

- Keywords only considered special [reserved] at the start of a new command
- *Shell keywords:*
  - **exit** - exits the shell
  - **if, then, else, fi** - parts of a compound command for conditional statements
    - if, then, else, fi (closing word); also: elif
    - if, then, else, fi only keywords at start of command
    - Ex: if test -r foo \n then cat foo \n else echo "foo missing" \n fi
      - Can also put sequences of commands, nest commands
        - Nesting conditionals - use else, fi to exit a nest
    - If a; then \n b \n else \n c \n fi
    - Exit status is the exit status of the last command of the then or the else (whichever was executed)
      - Omitting else - evaluates as true
  - **case, in, esac** - parts of a compound command for case statements
    - esac - closing word for a case chain
    - Ex: case \$? in \n 0) echo str;; \n 1) echo str2;; \n 2) echo str3;; \n esac
      - Default case: \*)
      - Cases can also be patterns
        - Use globbing patterns (same as used by the shell for file names)
          - Ex: \*.c;;, \*.ao)
      - Multiple cases: separate with vertical bar (e.g. 0|1)
      - Would work if all written on one line
    - Exit status is the exit status of the last command of the then or the else (whichever was executed)
      - No case match - evaluates as true
  - Keyword { + } (group a series of commands; run in sequence)
    - Exit status = exit status of last command executed
      - If an earlier command fails, shell will continue to run subsequent commands [ignores failure]
  - Keyword ! (negates command[’s exit status])
    - Is considered its own word [need space between words]
      - Is a command, not an operator

- Common extension (present in Bash; not universal):
      - `!!` = text of previous command
        - `!! *arguments*` add the arguments to the end of previous command
      - `!*`  = arguments of previous command
      - `!$` = last argument of previous command
    - Exit status: 0 or 1
  - ***while, do, done*** - parts of a compound command for [while] loops
    - Also: until (alternative), break, continue
      - No do-while loop defined
        - Alternative: `while *cmds* \n do : \n done`
    - Ex: `while *cmd* \n do *cmd* \n done`
  - ***for, in, do, done*** - parts of a compound command for [for] loops
    - Ex: `for i in a b 27 c \n do *body* \n done`
      - Runs body 4 times, where `$i` is equal to a, b, c, 27, and c for the 1st, 2nd, 3rd, 4th iteration, respectively
      - Similar to python for loops
      - Alt: `for i in *.c`
- *Extra shell tokens:*
    - `&&` - if the left-hand command succeeds, run the right-hand command
      - Returns exit status of last command run
    - `||` - if the left-hand command fails, run the right-hand command
      - Returns exit status of last command run
    - `()` - act similar to `{}`, but runs contained commands as a subprocess
      - Subprocess - akin to running commands in separate shell process [a descendant of the original]
        - Any changes to the shell itself (e.g. changes to shell variables) done in a subprocess do not affect the main shell [are erased when the subprocess exits]
      - Curly braces do not guarantee commands are run as a subprocess
  - *Shell variables: `*name*=*value*`*
    - Can only contain strings (incl. strings of digits)
      - Using special characters (e.g. spc) necessitates quoting



- Can have multi-character names
- Expands right-hand side (e.g. references to shell variables with \$ are replaced by the value of the respective variable)
- Spaces around the = break the assignment
- Can perform multiple assignments on one line
- Used for two purposes - internal use (local to shell), export to subsidiary processes
  - Exporting:
    - \*assignment\* \*command\*
      - Exports assigned variable as environment variable of descendant process, for descendant process only
        - Parent shell unaffected; variables unchanged
    - export \*assignment\*
      - Affects both shell, subsidiary processes
        - Changes to exported environment variables also affect the exported values [even if not explicitly re-exported]
- Built-in variables
  - \$? - exit status of most recent command
    - Convention: 0 = success, nonzero = failure
    - Max: 128 or 256
  - \$! - process id of [most recent] background command
  - \$0, \$1, \$2, ... correspond to the arguments of the process
    - \$0 corresponds to current process's name
  - \$\* concatenates all arguments \$1, \$2, etc. (except \$0)
  - @\$ acts like \$\*, except when quoted
    - When quoted ("@\$"), separates each argument into its own word
  - \$(*\*command name\**) runs the command as if it were a shell command (run as subshell), then saves the output
    - Is not a proper variable
    - Can replace \$(, ) with ` [not common]
- **Stages of shell command interpretation (done strictly in order):**
  - Delimit the command (default delimiting tokens: ; | \n)

- Any commands not properly delimited (e.g. within a `$()` expression) may result in unintended output
- Read in command as a series of words, and expand:
  - **Tilde expansion:** `~` expanded into `$HOME` path
    - `~, ~/*` expanded into `$HOME`
    - `~*user*` expanded into user's `$HOME`
    - Only expanded if at the start of a word
  - **Variable expansion:** `$x` evaluates to the value of shell variable `x`
  - **Command substitution:** `$(*cmd*)` is expanded into output of the command
  - **Arithmetic expansion:** `$((x + 5))` evaluates to value of shell variable `$x`, plus 5
    - Arithmetic expansion takes priority over `$(*cmd*)`
      - Must add spaces between parentheses for latter expansion
  - **Field splitting:** some words in a command may split into multiple words once interpreted
    - Ex: `x='a b'; cat $x` splits `x` into two words: `'a'` and `'b'`
    - `$IFS` [inter-field separator] = tab, space, newline
      - Can be modified by modifying `$IFS`
- **Pathname expansion/globbing:**
  - Characters:
    - `*`: matches any sequence of characters [incl. empty]
      - `*.c` returns all files ending with `.c`, e.g.
    - `?`: matches one char
    - `[*chars*]`: matches one char in a set
      - Negation: `[!*chars*]`
  - `*, ?` do not match file names starting with `.` [hidden files]
- **I/O redirection:** some words are treated as directives to the shell in regards to what to do with input/output [rather than simple arguments]
  - Characters:
    - `>*file*`: create/truncate file [discarding all existing data], write standard output to the file
      - `>>*file*` appends to a preexisting file, rather than truncating
    - `<*file*`: write file to standard input

- Will only be effective if the command has been told to read from standard input
- <<EOF [+other commands] RET \*lines\*\n EOF: effectively creates a temporary file to pass as stdin
  - EOF can be any word - << will keep reading until it finds a match
  - Contents in EOF expressions are expanded unless the EOF at start is quoted
- Done left to right, independently
  - cat <x <y evaluates to just cat <y, e.g.
- Can be done without expressions
  - >x >y just creates two empty files, e.g.
- Note: can read from and output to the same file, possibly causing infinite loop
- (#>/#<): 0=stdin, 1=stdout, 2=stderr
  - & after >/<: 2>&1 changes the contents of 2 to be written to 1, e.g.
  - 2>&-: closes 2 [writing to stderr causes error]
- <> opens a file for both input, output
- Expansions do not occur in single-quoted strings
  - Only variable, command, arithmetic [\$] expansions occur in double-quoted strings

## Shell Scripting

- Shell scripting
  - A shell script is essentially a sequence of terminal commands written within a file
  - A shell script “expanded”, command-by-command, into a simpler language during execution
- Traditionally, a shell script is split into multiple files (as needed) rather than combined into single big scripts with functions (unlike higher-level scripting languages, e.g. Lisp/Python)
  - Pros of split - forgoes | [pipe] functionality, split files do not all need to be shell scripts [can be C++ scripts, e.g.], more modular
  - Pros of all in one - ease of integration, viewability, package/distribute, more [space/memory] efficient [less overhead - 1 rather than 2 processes]

## The File System

- Linux file systems structured as rooted trees, starting at the root
  - Directories as nodes, non-directory files as leaves
    - Non-directory files: regular files, symbolic links, special files
    - Files, directories, symlinks distinguished by first bit: '-' vs 'd' vs 'l'
- A file system is akin to a simple hash table: a [non-null] string (e.g. /usr/local/bin) maps to a directory/file
  - Consequence: a single string cannot map to multiple files (files cannot share the same full path, including name)
  - String is composed of file name components, separated by slashes (e.g. "usr", "local", "bin", directories, etc.)
- A single non-directory file can be present in several different directories at once
  - **Hard links:** In `*location1/name1* *location2/name2*` makes it such that the file with name `location1/name1` appears also present as `location2/name2` [files only]
    - Is the same file - modifying one will also be reflected in the other
      - Essentially creates another "name" for the file
      - rm-ing a file only removes a name/directory entry, not a file itself
        - Similar to a pointer falling out of scope in C++
    - Are identified by unique internal IDs (shown by `ls -li`)
      - Also do not share metadata
    - Akin to a reference in C++
  - Consequence: a single Linux file can have multiple names
    - Link count akin to the degree of the file in the file system graph (# of directories pointing to it)
- **Symbolic links:** In `-s *file* *link_name*` creates a symbolic link pointing to the original file
  - Will point to a file, even if the file does not exist [dangling link]
    - Can store anything in a symbolic link, incl. directories, itself
  - During file name expansion, will replace the link with its contents
  - Akin to a pointer in C++

## Lisp Scripting

- Lisp [LISt Processing]
  - Relatively old [1950s] + simple (root language - inspired other languages)
  - Originated as an AI language
    - Was originally slow, theoretical - made more practical over time
  - Variants: Lisp 1/1.5, Scheme, Emacs Lisp [Elisp]
    - Emacs Lisp is an **extension language** for Emacs - is not meant to be the primary driver of functionality (generally forgoes a main function, e.g.), used instead for small additions/changes to the main application [Emacs]
      - Similar to Chrome extensions, for Chrome
- Lisp scripting uses an object-oriented approach:
  - Lisp objects represented by [64-bit] pointers to values in memory
    - Done so that all sorts of values (i.e. of any size  $\geq 1$  machine word) can be accommodated by a single syntax
    - Uses *tag bits* to denote the type of variable specified
      - Primitives - denoted by 3 tag bits at the end of the pointer
        - Put at the end of a pointer since [x86-64] addresses are aligned & byte-addressable (end 3 bits [8] are not used)
        - Exception: fixnums are represented by their actual values, rather than through pointers [end 3 tag bits indicate such]
          - More memory-efficient
      - Other objects - denoted by tag bits at the end of the pointer (denoting a non-primitive) + tag bits at the head of the value (denoting the actual type)
    - Note: Javascript uses a similar approach
- Lisp object types:
  - Primitives: Numbers, strings, etc.
    - Numbers: fixnums (small integers - fit in a single machine word), bignums (larger integers), floats
    - Chars are essentially just an alternate notation for integers (?a returns ASCII integer encoding 97), rather than a distinct type
  - User interaction: Buffers, windows, frames, terminals

- Shell: Processes
- Atom (symbol)
  - Represents a single data point, equal [eq] only to itself
    - Any given name can only be associated with exactly zero or one atoms; thus, any two references to a given atom name will always be to the same object in memory
      - Implemented internally via hash table
    - Can be used to represent “concepts” in the context of AI, e.g.
  - Ex: t, nil
- Lists constructed via cons [construct]: (cons val1 val2) -> (val1 , val2)
  - car, cdr - return first and second values of a pair, respectively
    - If the second value of a pair is a pointer to remaining list elements, cdr will return all remaining list elements
    - Naming - historical IBM convention
  - nil stands for empty list
  - Lists can then be made out of value-ptr pairs (a la a linked list), nesting cons
    - nil as the pointer stands for end of list/no value
  - Can also be constructed via (list \*values)
- Other notable types:
  - Hash tables (constructed as an empty table, then filled - keys of any type)
  - Char table (hash table that only take chars as keys - more efficient)
  - Markers (moveable pointer into buffer, moves along with contents of buffer when changes are made)
    - E.g. “Hello \*mrkr\*world” -> “Hello, \*mrkr\*world” [marker moves when string changes]
    - Can also integer index into buffers - does not move
  - Buffer - contains copy of the contents of a file
  - Frame, window - specifies display (frame - overall region, window - subregion)
  - Terminal
  - Process
  - Function: ( defun function\_name (args) body )
    - Function call: ( function\_name args )

- Lambda expressions: (lambda (x) (x+1)) defines function -> call: (lambda (x) (x+1)) 27 ) [functions do not require names]
- Lisp arithmetic operations, comparisons, etc. utilize prefix notation
  - Ex.: (+ 10 5), (= 3 8)
    - Negative/positive infinity: (/ 1 -0.0), (/ 1 +0.0)
      - -0.0, +0.0 are numerically distinct values
    - NaN (/ 0.0 0.0) does not numerically equal any number, not even itself
  - = compares object numerical values; eq compares object bit patterns (as opposed to object values)
    - eq fails to compare numerical equality for pointer-represented numbers (i.e. only works for fixnums)
      - (eq \*immediate\* \*immediate\*) will construct and compare two different objects
    - “Equal” - general object value comparison
  - (setq \*variable name\* \*value\*) sets a variable
  - t [true] vs nil [false]
- Control structures [Emacs: special forms]
  - Are also called special forms - resemble functions, but are not so
    - Does not need to explicitly evaluate all the provided arguments before executing, e.g.

## Lisp control structures

- Conditionals: (if (= a b) [then] (cons a b) [else] (cdr b))
  - Else part can be written as a sequence of expressions, to be executed in sequence, or zero expressions [then part must be a single expression]
  - Shorthand [no else]: (when/unless (condition) (do))
    - do can be a sequence of commands
  - Executing a sequence of commands - value returned is value of last expression
  - (cond ((condition1) expression) ((condition2) expression)...((true) expression)
- Not: (not (expression))
- And: (and (expression1) (expression2)...(expressionN))
  - Stops prematurely if an expression evaluates to False, similar to C++ &&
- Or: (or (expression1) (expression2)...(expressionN))



- Stops prematurely if an expression evaluates to True, similar to C++ ||
- While loops, recursion

### Lisp variables:

- (setq \*varname\* \*value\*) assigns variables
- (let ((\*varname\* \*value\*)) declares a local variable
  - Can declare multiple local variables simultaneously: (let ((\*var1\* \*val1\*)) (\*var2\* \*val2\*)), e.g.
- Functions can have local variables

### Emacs programs

- Emacs programs are represented as lists
  - Hence: same syntax for lists, functions + args
  - Function calls represented as lists also, though not all lists are valid function calls
  - Apostrophe [quoting] - an apostrophe before a list keeps list as data, rather than evaluating as function call
    - Lists: (let ((x '(3 4 5)))
      - Shorthand for (quote (3 4 5))
    - Quoting a function call, arithmetic expression, variable reference, etc. will store the function name, arithmetic expression, variable reference, etc. as a string
- C-] exits Lisp debugger
- Emacs scratch buffer - used to hold temporary values (C-x b \*scratch\*)
  - Can also be used to evaluate Lisp expressions (C-j; C-x C-e for general buffers)

## Python

### History of Python

- Python initially motivated by a desire to improve computer science education
- Historically [1980s] - BASIC (introduced as teaching language in 1962) was the entry language for programmers, taught in high school
  - Now largely dead, aside from Microsoft
  - Was originally simpler than FORTRAN, but became more complex over time + was difficult to teach [relied on go-tos, e.g.]
    - Practices in BASIC had to be un-taught to teach other languages + proper programming practice (e.g. indentation)
    - Involved teaching a large number of low-level functions
- ABC introduced as replacement to BASIC as a teaching language
  - Made indentation built-into the language
  - Came with an IDE, built-in library for higher-level functions (e.g. sort, hash tables)
  - Failed because it could not be used to get jobs (not enough users)
- Reminder: little languages: mix-and-match different “little languages” (e.g. sh, sed, awk, grep) to accomplish a task
  - Effective, but difficult to teach - requires knowing many languages
  - Attempts made to replace/unify it with a “big language”
    - Perl [scripting language, made by linguist Larry Wall] - “there’s more than one way to do it”
      - Inspired by English also being very flexible
- Python made as a reaction to Perl, descendant to ABC
  - Unified “little languages”, but tried to be more organized than Perl
    - “There is one best way to do it, and Python has it”

## Overview

- Python not generally used as an extension language, a la Lisp
  - Python may be used as a “glue language” - used for converting data in one form to another, as cleanly/conveniently as possible
    - Relies on calls to higher level functions, which may or may not be implemented in other languages at a lower level
    - Emphasizes OOP very strongly
- Python syntax:
  - Requires indentation
  - Single quotes/double quotes evaluated identically
    - Three single quotes allows for strings to cross line breaks
  - Every entity is an *object*: has an **identity**, a **type**, and a **value**
    - **Identity** roughly corresponds to C++ addresses, but no pointers/references
      - `id(object)` returns an address as an integer, but can't be used for anything else (i.e. no pointers)
    - **Type, value** stored together at some address
      - `type(object)` returns another object [of type Object]
    - Identity, type are immutable (address, typefield won't change)
      - Value mutability depends on object type
- Everything in Python is an object, or rather a reference to an object
  - [Object] copying is just pointer reassignment
- Came with an original 7 built-in types: None [NoneType] (akin to nullptr), int/IntType, Float, Complex, Boolean, List [array], String [array of chars], Tuple [immutable list], Buffer [mutable String, a la buffer vs string in Elisp]
  - `1+2J` returns complex number [J instead of i for readability]
    - Number types (int, float, complex, boolean) can be converted between each other, even bool to complex
      - No integer overflows in Python
  - Also contains its own internal types [Code]
- CPython converts Python to bytecode before running, allowing for optimizations

## Sequences (Python)

- Sequence types (**List + Tuple, Buffer + String**) akin to C++ arrays
  - List and buffer are mutable; tuple and string are not
  - Indexed by integers from 0 to n-1 ( $n = \text{len}(\text{sequence})$ )

### General Sequences

Sequence Operations (Python)	
Syntax	Result
<code>arr[i]</code>	Returns value at index i
<code>arr[-i]</code>	Returns value at index ( $\text{len}(\text{arr}) - i + 1$ )
<code>len(arr)</code>	Returns length of arr
<code>min(arr)</code>	Find minimum value in arr
<code>max(arr)</code>	Finds maximum value in arr
<code>list(arr)</code>	Constructs new list from sequence elements

### Mutable Sequences (List, Buffer)

Mutable Sequence Operations (Python)	
Syntax	Result
<code>arr[i] = v</code>	Sets value at index i to v
<code>arr[i:] = arr2</code>	Splices in arr2 at index i
<code>del arr [i]</code>	Deletes value at index i
<code>del arr[i:j]</code>	Deletes value at indices i through j (incl. i)
<code>arr.append(val)</code>	Adds copy of val to end of the array
<code>arr.extend(seq)</code>	Appends elements from a sequence
<code>arr.count(x)</code>	Counts instances of x in arr
<code>arr.index(x)</code>	Finds first instance of x in arr
<code>arr.insert(i, v)</code>	Inserts v before i <sup>th</sup> element of arr

<code>arr.pop(i)</code>	Removes $i^{\text{th}}$ element and returns its value
<code>arr.pop()</code>	Equivalent to <code>arr.pop(-1)</code>
<code>arr.remove(x)</code>	Equivalent to <code>arr.pop(arr.index(x))</code>
<code>arr.reverse()</code>	Reverses the elements of <code>arr</code>
<code>arr.sort()</code>	Sorts the elements of <code>arr</code> in-place

## Lists

- Lists implemented as an object with values: pointer to vector head in memory, length, alloc
  - *Length* counts current # of elements; *alloc* counts total allocated size of vector
    - Appending increments length; when `length == alloc` [list is full], will copy the list to another allocated vector with double the size
  - List operation performance
    - Append:  $O(1)$ , amortized [averages out to  $O(1)$ ]
      - Math of amortization: half the list cost 1 operation, a quarter cost 2 operations (copied once), an eighth 3, etc.  $((1 + \frac{1}{2} + \frac{1}{4} + \frac{1}{8} + \dots) = 2n)$
- Python List akin to a stack (push=append, pop)
- `arr.sort()`: Numbers [non-complex] compared directly, strings via string comparison, custom types via custom operator<

## Strings

String Operations (Python)	
Syntax	Result
<code>str = "this is a string"</code>	Creates a new string
<code>s1.join([s2, s3, ...])</code>	Returns a string with value: <code>s2 + s1 + s3 + s1 + ...</code>
<code>str.split(separator)</code>	Divides <code>str</code> into an array, using <code>separator</code> as the delimiter
<code>str.replace(old, new)</code>	Replaces every instance of <code>old</code> in <code>str</code> , with <code>new</code>

- Strings are immutable - mutable operations don't apply
  - Operations work by constructing new strings, rather than overwriting the old one

## Mappings (Python)

- Python mappings (notably Dictionary/dict) implemented internally via hash tables
  - `dict[key]` performs a hash lookup in dict with key
    - Key must be immutable (no lists, dictionaries, e.g.)
    - Using variables as keys will use the variable value as key
- Notes on dictionary operations
  - Using `[]` to index/delete on a nonexistent key returns a `KeyError`
  - `dict.update(dict2)` will overwrite any preexisting item in dict overlapping with dict2

Dictionary Operations (Python)	
Command	Result
<code>dict = {key: value, ...}</code>	Creates new dictionary, mapping keys to values
<code>dict[key]</code>	Returns value associated with key in dict
<code>dict.get(key)</code>	Returns <code>dict[key]</code> if key is valid, <code>None</code> if not
<code>dict.get(key, val)</code>	Returns <code>dict[key]</code> if key is valid, <code>val</code> if not
<code>dict[key] = val</code>	Maps key to val in dict (overwriting any prior mapping with key)
<code>dict.has_key(key)</code>	Looks up if key is in dict
<code>del dict[key]</code>	Deletes key-value pair from dict
<code>len(dict)</code>	Returns number of key-value pair
<code>dict.keys()</code>	Returns list of keys in dict
<code>dict.values()</code>	Returns list of values in dict
<code>dict.items()</code>	Returns list of key-value pairs as an array of tuples
<code>dict.clear()</code>	Removes all items from dict
<code>dict.update(other_dict)</code>	Copies every key-value pair in other_dict, to dict
<code>dict.popitem()</code>	Deletes some item and returns the corresponding key-item pair

## Callables (Python)

- Callables: Function, Method [class function], Class [calling Class akin to C++ constructor]
- Functions
  - Can be called with arguments in any order by specifying the argument called (e.g. `f(arg2=val, arg1=val2, ...)`)
  - Can take in trailing arguments (e.g. `f(a, b, *c)` will return any trailing arguments not a, b in a tuple as c)
  - Can take in keyword arguments (e.g. `f(a, b, **d)`, then calling `f(a, b, key1=val1, key2=val2)` will return `d={key1: val1, key2=val2}`)
    - Can be used in conjunction with trailing arguments
- Classes
  - Syntax: `class(*parent classes*)`
    - Allows multiple inheritance
      - Conflicts: Inherits closest (depth-first, left-to-right)
      - Later parent classes “fills in” gaps of first parent class [effectively superclass]
  - Can put any number of methods, variables in definition
    - Is dynamic - can be changed during runtime
    - `myClass.__dict__`; instance variable matches names [of functions/methods] associated with class, to values
      - Can be modified, e.g. for use in *meta-programming* (building upon classes made by other people/libraries)
        - Only works in Python2; Python3 changed `__dict__` to read-only MappingProxy (only string keys -> faster)
    - Defining methods with underscores defines built-in methods/operations
      - `myClass.__init__`: executed on object creation
      - `myClass.__str__` defines the return of casting obj to string
      - `myClass.__cmp__(self, other)` comparison operator
        - -1 for <, 0 for =, 1 for >
        - Used for `list.sort()`, e.g.
        - `myClass.__lt__` [less than]
      - `myClass.__nonzero__`

- Used for if object: checks
  - `myClass.__add__(self, other)`
- Class methods - first argument [self] is generally the instance for which the method is being called
  - Does not need to be called self
- Can rename classes (e.g. `class c; d = c`)
- Immutable classes can be made dictionary keys
  - `def __hash__(self):` define hash function returning integer

misc

- Documentation puts square brackets around optional arguments
- Python string processing
  - `Line = "str,10,10.2"; types = [str, int, float]; fields = [t(v) for t,v in zip(types, line.split(',')) -> [(str, "str"), (int, '10'), (float, '10.2')] -> fields = ['str', 10, 10.2]`



## Client-Server/JavaScript

- Node/React - built on JS [+POSIX+OS[+sh as a configuration language]]
  - Quoting issues, configuration in React
- Client-server system: clients and servers [separate computers] communicate over/through a network [e.g. the cloud]
  - Clients send queries to the server; server sends information back
  - (Other models exist for distributed computing)
    - Ex: Multiple clients, one server
      - Clients communicate only via the one server [may cause scaling issues with more clients]
    - Multiple clients, multiple servers
      - Servers are closely connected, can spread request load
    - Peer-to-peer: each peer [computer] maintains part of the state of the system, communicate with other peers via network
      - No single server; spreads server responsibilities across peers
    - Divide server into primary and secondary nodes: secondary nodes receive questions, send to primary node (primary node returns information)
      - Used in machine learning, e.g.
- Performance issues in distributed computing
  - *Throughput* - how much work [queries per second] is the overall system doing?
    - Primarily a server-side issue
  - *Latency* - how long is the delay between a request and a response?
    - Primarily a client-side issue
  - Balancing throughput vs latency - can improve one at the cost of the other
    - Batch requests (out-of-order execution) - can increase throughput by collecting requests until a number of sufficient related requests [querying same information] are received, then performing them all at once
      - Increases efficiency, since the information is in the RAM for each request; but hurts latency
    - Caching - Can decrease latency by client-side caching and reusing information received from past queries
      - Forgoes having to send a request and wait for a server response

- Stale caches - what happens when the information in the cache is out-of-date?
  - Cache validation - looking for a means of checking the validity of a cache, cheaply
- Preloading/prefetching - caching information a client has not requested, in the expectation they will request it soon
  - Aggressive caching

## The Internet

- The Internet & alternatives:
  - Internet precursor: telephone system [landlines] operated on circuit switching
    - Many wires [individual phones] are connected to central offices; central offices communicate via long series of switches
      - Switch paths can traverse long distances geographically
      - A single set of wires is limited to a single connection (wires are reserved)
    - Cons: inefficient, no fallback in the event of failure between switches
      - 1960s [Cold War] - US military concerned about the possibility of attacks on US communication system
- Paul Baran [RAND corporation, ~1961] proposed **packet switching**: conversations broken up into small pieces [packets, e.g. 1 KiB], each packet routed dynamically between central offices [rather than following a preplanned route]
  - Any two packets with the same origin & destination may take entirely different paths – no reservation of wires
  - Packet composed of two parts - **header** (a la metadata) and **payload**
  - More robust in the presence of an unreliable network, more efficient [multiple connections can use the same wires]
    - Cons: worse latency [dynamic routing takes time, since each switch must determine next switch], data may not always be received in order of transmission
  - Used by Zoom, modern phones, e.g.
  - 3 problems for packet switching:
    - Packets may be received out-of-order
    - Packets can be lost during transmission
    - Packets can be duplicated
      - Usually due to network issues (e.g. a switch thinks a send failed, and sends again)
- Packet switching issues dealt with using **protocols** - communication between switches to handle the transmission of packets
  - Can be complicated

- Complexity dealt with using layers of protocols
- *Protocol layers [high-level to low-level]*
  - **Application layer** - application-dependent series of protocols
    - Ex: World Wide Web [WWW] allows for asking for, receiving webpages; File Transfer Protocol [FTP] for transferring large files
  - **Transport layer** - views packets as a stream of data from origin to destination
    - Has protocols attempting to resolve 3 issues
    - Ignores low-level implementation
    - Intended to replicate circuit switching functionality
  - **Internet layer** - switches [routers] determine pathing from origin to destination through individual switches
    - Based off Internet protocol [IP]: packet header contains information about packet, what to do with it
  - **Link layer** - low-level [hardware-level], point-to-point communication between switches
    - Only deals with packets - not packet switching/routing
- Versions of Internet Protocol (IP):
  - **IPv4 (1983)**
    - Team incl. J. Pastel (UCLA)
    - Specified contents of packet and packet header
      - Packet header contains:
        - Length of packet
        - Protocol number, indicating # of high-level protocol to use
          - Intended to use IP as base layer for higher layers
        - Source, destination addresses
          - Represented as 32-bit ints [4 decimal bytes each]
          - Source address can be used to communicate transmission failure, e.g.
        - Checksum [16-bit] - used to check for bit errors in transmission
        - Time-to-live (TTL) [8-bit] - tells the router to drop the packet after a certain number of routers are traversed
          - Decrement for each router traversed

- Used to prevent infinite router loops
  - Abuses/problems:
    - Source address can be faked by sender
    - Nonexistent destination addresses
  - **IPv6 (1998):**
    - Made addresses larger + other changes
  - Many networks may involve both IPv4 and IPv6 machines, use gateways to convert IPv6 addresses to IPv4
    - Loses IPv6 features
- 2 major categories of transport layer protocols: **User Datagram Protocol (UDP)**, **Transmission Control Protocol (TCP, TCP/IP)**
  - **UDP** (David Reed, MIT) forms a very thin layer atop IP
    - Easy to implement, but vulnerable to packet loss
      - App is responsible for dealing with packet loss, duplication, etc.
  - **TCP** (Vint Cerf, [UCLA alum] Stanford; Bob Kahn, Princeton) provides reliable, ordered, error-checked stream of information
    - Divides data stream into packets [negotiating packet size with hardware]
    - Also handles retransmission, in the event of dropped packets; reassembly [reordering], in the event of out-of-order packet transmission
    - Handles flow control - avoids problem of greedy senders [senders shipping too much data to local central office, to be handled by global network]
- Application layer: WWW uses **HyperText Transfer Protocol (HTTP)**, **Real-Time Protocol (RTP)**
  - Also uses other protocols (e.g. video protocols)
  - Defines protocols for classes of similar/related applications (e.g. browsers)
  - **WWW** and **HTTP** developed by Tim Berners-Lee, CERN (1991) to handle displaying of, and linking between [footnotes], digital copies of physics papers
    - Implemented via combination of **TCP-like** protocol (handling transfers) + **SGML** (data protocol) + new linking protocol
      - **Standard Generalized Markup Language (SGML)** developed as standard interchange language between publishers of books/papers
        - Had HTML-like syntax
- **HTTP** (original main Internet protocol) written as a client-server protocol for transfers, similar to TCP

- Operates in terms of requests and responses between client, server
  - Example request: GET /index.html HTTP/1.1 \n\n
  - Server response: Header + \*webpage contents\*
    - Header: 200 OK + \*webpage metadata\* \n
    - Metadata: length, content type field (e.g. HTML; used to aid in client interpretation), content transfer encoding (e.g. UTF-8)
- HTTP 1: client setup, one request, one response, connection drops [TCP teardown]
  - Eventually evolved to handle multiple request/responses (HTTP 1.1)
- HTTP/2 (2015): made various changes for efficiency
  - Compressed headers (client-server had to agree on compression format)
    - Before HTTP/2: headers not compressed
  - Allowed for *server pushing* - server can send a response, even if client has not sent a request
  - *Pipelining* - connections kept alive after server response (no teardown) + client can issue multiple requests at once without waiting for a response
    - Issues: browser now has to deal with out-of-order responses, e.g. by labeling requests and responses
  - Multiplexing - using a single TCP connection for multiple conversations
    - Used in the case that a single client wants multiple sessions, e.g.
      - Requests include host/domain name/session of interest
- HTTP/individual TCP connections layered atop data streams [TCP] atop packets [IP]
  - Results in “*head of the line problem*” - TCP trying to preserve chronological order of sent packets means, in the case of packets not being sent/dropping, any later packets will not be received by the browser
  - HTTP/3 (2022): replaces TCP with QUIC protocol to resolve head-of-the-line
    - QUIC (“TCP v2 + UDP”) - attempts to preserve chronological order of sent packets (a la TCP), but has ways to deal with packet dropping
    - Allowed for further multiplexing

## HTML

- HTML derived from publishing format SGML

- Philosophy [SGML, sorta HTML] - the language should contain only the contents of the book, not any form of program to actually run them/display them
  - “Declarative, not imperative” - does not tell the browser what to do (e.g. styling), only what they should display (content)
  - Makes porting easier, since each browser may have its own command set
- SGML philosophy inherited (to some extent) by HTML
- Terminology of HTML
  - HTML elements consist of an opening tag (enclosed in `<*tag name*>`) and closing tag (enclosed in `</*tag name*>`), with text in between
    - Some *void elements* consist of only an opening tag (e.g. `<br/>`)
    - Opening tags may have attributes: `<*tag name* *attributes*>`
      - Attributes consist of name-value pairs of text
        - Names are distinct identifiers; values can be any string
      - Attributes may describe to a browser how an element should be displayed, or denote relationships between HTML elements
    - Elements can generally nest inside each other, in a hierarchical structure
      - Exception: raw text elements contain only text, paragraph elements cannot nest directly inside other paragraph elements
      - Elements are nodes of the tree; internal text are the leaves
  - HTML documents (also called resources) can point to other documents via links
    - `<a>` [anchor] has `href` [HTML reference] attribute, holds a URL (uniform resource locator) to another document
      - `#id*` at end of URL - move user cursor to element with that ID
        - In the absence of a domain name (e.g. just `href="#intro"`), defaults to URL of current webpage
- HTML problems:
  - Possible errors in HTML:
    - Element lacks a closing tag
      - General mantra: be strict in what is generated, but be generous in what you accept
      - If adding a closing tag would fix error, add a closing tag
        - Omitting closing tags on the developer side can be done to reduce overhead, if there is no ambiguity

- Element has an extra closing tag: things break
- Possible upgrades for HTML
  - Extending/changing HTML (e.g. by adding an attribute) to allow for additional functionality
    - Consequence: servers and clients may update at different times and potentially have slightly different versions of HTML
      - Servers on newer versions may send HTML that clients on older versions cannot parse
      - Methods for resolving:
        - Clients, upon seeing HTML they cannot parse, will make something up
        - Clients tell server the kind of HTML they can handle
          - Server tries to use specified HTML format; get more complicated
        - Server tells client the version of HTML in response; client figures out what to do
          - Clients have to support more HTML versions
            - HTML elements may change, be removed between versions (e.g.)
          - Used in the early web
  - HTML version denoted by Document Type Definition (DTD/Doctype)
    - Taken from SGML - different types of documents, publications may be shared, each having different needs, specified by Doctype
      - Doctype lists all elements + their allowed attributes, subelements, void-ness, raw text-ness, etc.
    - HTML Doctype also specifies whether closing tags are optional
    - Doctype standardization effort ultimately failed
      - Web evolved too rapidly - Doctype standards came too late
      - HTML became too complicated to express in DTD
  - HTML/5 [current version] - uses evolving standard in place of a static DTD
    - Commits are made to a GitHub repository
  - XML (eXtensible Meta Language) - barebones for HTML syntax, with no semantics
    - Simple way to send any form of tree-structured data



- Leverages browser infrastructure for parsing HTML, as a means of transmitting other data (e.g. payroll)
  - Element tag names, attr names, etc. all ignored
    - Insists on closing brackets
  - XHTML - “strict” HTML
    - Is valid XML - insists on closing brackets, e.g.
- **Document object model (DOM)** - HTML ships elements within an tree data structure, where nodes are elements (w/ attributes), leaves are text
  - **DOM** - object-oriented model for examining/manipulating said trees
    - Primarily used for JavaScript; theoretically intended for any OO language
    - Attempt at making HTML more imperative (rather than declarative)
      - DOM itself is not imperative, but using JavaScript makes it so
- **Cascading Style Sheets (CSS)** - separate HTML declarations into two categories: content (HTML) and presentation (appearance on the screen; CSS)
  - Presentation can be specified by user preferences, the element involved (e.g. paragraph attributes), ancestor elements, etc.
    - CSS provides a standard for combining preferences/styles
- Declarative philosophy is limited in its ability to handle more complicated interactions
  - Done using imperative technologies, namely code in JavaScript
- **Javascript** - programming language designed specifically for use in browsers
  - Is directly embedded into and run by web browsers
  - Can be hooked into HTML via the <script> tag
    - Can either import external files or write Javascript directly
  - Can modify elements in the DOM

- Programming paradigms - classifications of programming languages based on their features/functionalities
  - Methodologies/strategies languages follow, and the capabilities they have (can have multiple in one)
  - Can be classified by organization of code, style, syntax, grammar
- Web app vs website: web app more complex (interactive, integrated, authentication)
  - Interaction - self-explanatory
  - Integration - bringing together different components
  - Authentication - self-explanatory
- Layers:
  - Frontend - user-side interface
  - Backend - services & presentation logic; returns responses
  - Backend (database) - store & serve data

## Version Control

- Configuration control - ensuring deployed software is properly setup (ops)
- Version control - (dev)
  - Tracks version history (history - versions)
    - Allows for reversion, version disambiguation, compatibility, ability to see past changes
  - Allows for working in different directions + changes in different areas of project
  - Bug reports
  - Merging projects
  - Project reviewers - reviewers can look at changes made, e.g.
- Similarities - record/develop variants of a single program
- Git - developed for Linux kernel
  - Contains object database recording variants
    - Commit object - records project state + metadata (time, author, etc.)
  - Indexes - records planned changes for next commit, etc.
    - Used in the context of day-to-day changes
  - Is a distributed database - multiple copies of the same project database can exist, and can differ; none are strictly "in charge"
- Info
  - Remote: # of objects received = # of objects in database (roughly)
  - Change history stored as a set of changes from the previous version
    - During initialization - must compute all changes from oldest to newest version
- Commands
  - Git init - initializes empty git repo [.git]
  - Git clone \*repository link\*
    - Creates copy of source files + copy of repository history (in .git) (40%/60% of space, possibly)
  - Git status - tells status of current repository, working project
    - Indicates current branch, sync status with upstream branch, staged & unstaged & untracked changes [for commit]
    - Finding current project state

- Git log - displays metainfo
  - Commit ID (40 hexadecimal digits), author, date, commit message
    - Commit ID associated with only one commit in the entirety of git
      - Computed as a cryptographically secure checksum of commit contents [SHA-1] - difficult to find a match with same checksum
    - Commit message human-generated: important for reviewers, other developers, etc.
      - Format convention: 1st line - key, 2nd line - empty, other lines - describe motivations for change
        - Other lines also include API changes
    - Can be identified within repositories using a prefix set of digits (initial 12, e.g.)
  - Git log \*commit ID\* - view status at certain commit
  - Git log \*commitID1...commit ID2\* - view entire history from commit1 (exclusive) to commit2 (inclusive)
  - Git log -pretty=\*option\* changes formatting
  - Git log \*commit\*@[num] - state of repository num changes before commit
- Git ls-files - lists all files under git control [working files] in current version (can be grep'd, e.g.)
  - Git grep \*pattern\* effectively greps through files in git ls-files
- Git diff - compares current state of working files vs state of files in repository
  - Git diff compares current working files vs plans for commit
    - Adding files via git add -> will no longer appear in diff
    - Option: -cached compares files in index to most recent commit
  - Git diff \*commit1\* \*commit2\* - compares files in different repository versions (omit commit2 -> compare to current state of working files)
  - Options: -w (ignore whitespace-only changes)
  - Git apply \*patch\*
- Git add \*FILES\* - stages FILES for next commit [adds to index]
- Git checkout
  - Can checkout individual files from current version - git checkout \*path/filename\*

- Git commit - enters editor for writing commit message; creates commit upon exiting editor
  - Copies index into current repository; does not necessarily change files
  - Projects can setup style rules for commits (e.g. message 1st line must contain colon [for describing project component, e.g.])
  - Options -a (add all files), -m (specify commit message within shell command)
- Git configurations
  - Options: -l (list all command variables)
    - Comes from/stored into: current git installation, project's repo [.git/config], personal preferences
  - Must be properly formatted
- Git show \*commit\* - similar to git log, for just the single commit
  - Git show \*commit:FILE\*
- Git reset reverts to HEAD, discarding all changes
  - Git reset --hard \*commit\* - reverts to said commit
- Git blame FILE - shows who last edited each line of the file
  - Git blame COMMIT -- FILE
  - Struggles with nonlinear history (e.g. in the case of merges)
- Git bisect - binary search to find specific bug-introducing version
  - Git bisect start
    - Git bisect bad; git checkout \*good\*; git bisect good
      - Git bisect bad \*commit ID\*
    - Git bisect start \*bad\* \*good\*
    - Git bisect run \*any shell command\*
      - Git bisect skip
  - Difficult for multiple branches
    - Git wants both maximal node (most # of ancestors), and minimal node (farthest away from HEAD) -> picks compromise
      - Has option for randomization
      - $N = \# \text{ ancestors}; M = \text{largest } N \text{ in graph}; \text{bisect: } \max(N, M-N)$
    - Can mark bad nodes to skip (e.g. broken versions)

- Sporadic test – usually report good in the case of a working program/bad in the case of a non-working program, but sometimes fails to report the correct result
    - May result from performance testing/race conditions (tests dependent on many external variables)
- Git checkout
  - Git checkout COMMIT - changes working files to match commit
- Git switch - change branch [alt: git checkout -b]
- Git cat-file -p \*hash\* – view contents associated with hash
  - -p [print contents], -t [print type]
  - HEAD^{tree} – tree associated with head [e.g.]
- Git hash -stdin – generate unique ID based on particular string
  - Can take stdin, file as input
  - ID generated will be identical for the same string
  - Is mostly cryptographically secure checksum
    - Was originally used for verifying successful transmission of a file
    - Relatively difficult to generate two strings making the same checksum (roughly proportional to  $2^{\text{key length}}$  (in bits)  $\rightarrow 2^{160}$ )
    - Is meant to be generally collision-free
  - Git hash -w – write object into repository (.git/objects), as a blob
- Git update-index – builds new tree from planned changes
  - Update-index -add -cacheinfo “permissions,hash,name” creates tree with blob in index
  - Git write-tree – copies index into repository [gives checksum for new tree]
- Git commands - remote servers
  - Git clone \*URL\*
    - Repository hosted on a git server; copied to local machine
      - Cloning local repository - .git directory of newer clone will use hard links to .git directory of older clone [saves space]
  - Git fetch - downloads data from remote repository without making changes to working files
    - Local repository stays on same commit, same index

- Git pull downloads from remote repository (a la git fetch), but merges changes into commits - changes working files
- Git remote - lists name of all remote, fetchable repositories
  - Default (upon cloning): origin
  - Git remote -v
- Branches allow for multiple versions of a project
  - Different teams may be working on different features independently
    - Ex: main (development) + feature branches
  - May have different tolerances for new features, depending on purpose
    - Ex: maint (bug fixes to previous versions) & old (older version of maint - only security fixes)
  - Commands
    - Merging two branches creates a merge commit - commit with two parent commits (combines changes in both branches)
    - Git checkout \*branch\* switches to a new branch - attaches commits to new branch
- Other
  - .gitignore - file distinguishing between files to be version-controlled, files not to be version-controlled
    - Contains set/list of patterns (similar to shell globbing patterns) to ignore
      - Pattern starting with slash - only matches files in current directory
      - Pattern starting without slash - matches in working directory & subdirectories
      - Pattern ending with slash - match only directories
    - Can be used to ignore build files, private files, e.g.
  - Commit arithmetic:
    - HEAD - latest commit
    - ^ after commit ID = preceding commit (HEAD^^ = commit two commits before latest commit)
- Git tags – named Git object pointing to a commit
  - Syntax: git tag -a *tagname* *commit* [optional: -m *tagmsg*]
    - git tag [show tags], git tag -l '*globbingpattern*'
  - Can be used to mark important commits or signify milestones, e.g.

- Are shared when the repo is cloned
  - Are designed to be immutable – cannot be moved to a different commit, e.g.
- Using a tag to a commit in git commands equivalent to using the commit hash
- Git branches –
  - Are essentially just lightweight, movable pointers to a commit
    - Will move as development occurs, commits are made
  - Making commits on a branch will sprout off a separate branch of a commit tree
    - Branches can sprout others branches; can have many branches
  - Motivations for branches:
    - Can have different feature branches – lets different features be developed in isolation until the feature is ready
      - Can change the same things other people are changing
    - Can have maintenance branches – separate branches from “mainline” development branches, usually older, more stable versions
      - Receive only bug fixes, i.e.
      - Security branches – even older, receive only security fixes
      - Drawback: may have to copy the same fix across multiple branches
    - Can fork – separate a piece of software into two separate versions, developed entirely independently
      - No attempts made to sync forks
      - May be used in the case of disagreements between developers regarding future directions, e.g.
  - Syntax: `git branch b commithash` creates branch at that commit
    - Alt: `git checkout -b branchname commithash`
      - `git branch --track branchname upstream`
      - `git checkout -f branchname` (switch branch, discarding current changes to working files)
      - Moves HEAD to the commit where branch was started
    - `git branch` [lists all branches]; `git tag -l 'globbingpattern'`
    - `git branch -d branchname` [deletes a branch]
      - Deletes only the pointer to branch commits; actual commit objects are not initially affected [but may become orphaned]
        - `-D` – force delete [no warnings]



- Git garbage collector will eventually kill orphans (`git gc`)
  - `git branch -m branchold branchnew` [renames a branch]
    - Can rename branches to have different names to avoid conflicts with upstream branches, e.g.
  - `git diff branch1 branch2` [compare two branches]
- Git
  - Merging – move the changes made in one branch into another (“combine”)
    - Creates a merge commit in the merging-into branch – commit with multiple parents (such that both branches will point to the same commit)
      - Can merge multiple branches simultaneously
    - If the two branches have colliding changes (i.e. made their own different changes to the same lines), cause merge conflicts
      - Can be resolved manually, or automated via Linux *diff3* (compares three files/versions - runs diff twice)
        - Diff3: compares files B, A, C [in that order]
          - A - oldest file (common ancestor of two branches), B and C - branched-off alternatives
          - Alt: `diff3 MYFILE OLDFILE YOURFILE`
        - In the case that two versions are in agreement and one disagrees, diff3 picks disagreeing branch [wants to capture changes]
          - `Diff3 -m` [show merged results]
        - diff smart enough to determine what regions have changed, what regions have not [even if differing regions are of different lengths]
          - Developer of diff later worked for comparing DNA for Human Genomics Project
      - `git merge` essentially uses diff3 internally
    - `git merge` may result in a merged version with the wrong behavior even in cases where changes do not directly collide (do not change the same lines)
      - May want to compare both parent commits to new merge commit to see if the changes made sense (`git diff branch1 merged`; `git diff branch2 merged`; [`git diff ancestor merged`])

- Rebasing – essentially “copying” the changes made in one branch (from the common ancestor) to another branch, such that the rebasing branch now appears to have been branched off of *otherbranch*/HEAD
  - Theoretically results in the same changes as a merge, but results in different-looking git history
  - (In rebasing branch) git rebase *branchname* [rebases changes made between common ancestor to head of rebasing branch, to the head of the branch named in the rebase command]
    - git rebase -i *branchname* [selectively choose which commits to apply to rebased branch, in what order, etc. - may risk collisions]
      - Can even rebase a branch on itself (from an older commit) to throw away future commits as desired
    - git rebase -ignore-date [changes dates of rebased commits to appear to have occurred on the date on which the rebase occurred, rather than when the original commits were made]
      - Otherwise [default]: rebase will keep original commit dates, even if that results in what appear to be out-of-order commits (i.e. rebased commits older than main commits)
    - Rebasing other onto main: essentially changes other/HEAD, from branching off an older commit in main, to appear to branch off main/HEAD instead [and have all the changes made in main in the interim period]
  - Sort of “rewrites history”
    - Motivation:
      - Simplifies git history (makes history appear to be a single line of sequential commits, rather than many merges)
        - May want to rebase a branch against current master HEAD to show only the changes actually made by commits in the branch (i.e. won't also show “differences” caused by commits in master)
- Git stash – temporarily saves current set of changes to working files in the index
  - git stash push (saves current changes to index)
    - git stash list [show stashes]

- `git stash apply` [apply stashed changes]
- Can create multiple stashes
- Can be used in the case of context switching (working on multiple branches)
- Alternative: `git diff >changes.diff` (stores differences in file), `patch -p1 <changes.diff` (re-performs changes based on file)
- Version control philosophies
  - Git philosophy - only important files should be committed
    - Automatically computed files ignored, e.g.
  - Alternative - save everything
- Aside on `git bisect` - help bisect out
  - Make changes small and self-contained – break large commits (touching multiple source commits) into smaller ones before publishing
    - Makes changes more easily analyzed for bugs via `git bisect`, e.g.
    - Don't make any commits that break things (break the build)
- Aside on `git status`
  - Generally clean, but can become messy in the case of a merge/rebase/bisect, e.g.
- Git submodules – ways of combining different projects
  - Can make projects depend on one another to build/run, e.g.
    - Places dependency source code as submodule within depending directory; allows depending directory to access dependency's source code
  - Syntax: `git submodule add *link to repo*` [from within depending git directory]
    - Adds new type of object to depending directory's `.git` – submodule object, pointing to commit in dependency directory [dependency version]
      - Dependency's `.git` contents unaffected
  - Case: dependency project gets an update
    - Can choose to do nothing (continue on old version; insulated from updates)
    - Can alternatively update to newer version
      - Syntax: `git submodule foreach git pull origin/main` [e.g.]
        - `foreach` – run command for every submodule
        - Can just run `git pull` inside dependency directory instead
      - Need to then add submodule update to git commit: `git add submodule_name`
        - Updates only submodule object of depending directory

- Cloning a directory with submodules will not clone submodules automatically
  - Would need to explicitly initialize submodules: `git submodule init`
- Drawbacks – relies on close dependence on other people's repositories/projects
  - Cannot make own changes to submodules; only rely on others for fixes
  - May prefer to use higher-level importing tools (e.g. PyPi) instead, which assume the dependency is installed elsewhere
  - Submodules may be used to avoid needs for install dependency elsewhere (i.e. putting burden of dependency management on dev, not the customer)
    - Used more commonly in embedded systems without robust package managers, e.g.

## Git Internals

- Git internals: porcelain [Git commands] vs plumbing [internals]
  - Cloning only involves copying a `.git` folder (faster); updating files involves resolving changes to get files to their current version (slower)
  - Contents of `.git` subdirectory:
    - `.git/branches/` – obsolete, just kept for backwards compatibility
      - Formerly – listed branches [had small files describing branches]
    - `.git/config` – file describing repository configuration
      - Describes repository format, various settings
        - Bare – bare repository does not contain working files, only the `.git` repository [git ignores source files]
          - Only used for cloning/viewing history
      - Contains remote origin (where repo was cloned from, where to fetch from) + info about local branches + corresponding remotes, where to merge from
        - Transfer - determines how much checking git should be on pushes/pulls to/from upstream
      - Can be modified directly via text editor or `git config` command
        - `git config *variable name* *value*`
    - `.git/description` – obsolete, would be used by other applications (e.g. GitWeb) to describe repository; replaced by READMEs

- `.git/HEAD` – lists current location
  - Gives a branch, or a checksum if not on a branch
- `.git/hooks/` – contains scripts + sample scripts written by developers
  - Hooks - scripts executed by scripts at various points of operation
    - `.sample` files [sample scripts] ignored by git, just indicate nature of corresponding hooks
    - Maintained manually by developers (e.g. created in scripts), not automatically by Git
    - Ex: script to verify commit message formatting
      - Commit messages, pre-commit, pre-push
- `.git/index` – data structure [binary file] representing index (“future plans”)
  - Filled in by git when files are added to index by git add, e.g.
- `.git/info/` –
  - `info/exclude` – file under Git control, detailing files to ignore
- `.git/logs/` – keep track of changes to branches
  - `/logs/HEAD` – file containing history of all commits/changes to current repository
    - Is a metahistory: Tracks what was changed, by who, when, via what command
  - `/logs/reflogs` - temporary auxiliary file tracking changes by the user to the local git repository
    - Expires after 90 days [changeable]; not shared with others
- `.git/objects/` – stores all Git objects
  - Git repo state stored via object-oriented database; each object associated with unique ID (40-digit hexadecimal)
    - Stored in directories: `.git/objects/*first 2 hex digits of ID*/*remaining digits of ID*`
- `.git/refs/` – stores pointers to commit objects (branches, tags, remotes, etc.)
  - Stores branch heads
- **Git objects** – content for database
  - Objects named by their contents via hash; same contents -> same hash
    - Object naming divorced from the file system file [node] #s; means knowing contents of a git object -> knowing name/hash of git object

- File #s map names to contents, but not vice versa
- Ensures atomic updates – when pulling from git, multiple files are changed “simultaneously”, e.g. [no intermediate state with half old, half new]
  - In a regular file system - would require two copies of the directory [old and new], re-pointing of symbolic link from old to new version
- Generally meant to be immutable
- *Types of git objects:*
  - **Blob** - chunk of binary data, storing file contents/data [generally a file]
    - Does not store attributes, file name (just contents)
    - Is entirely defined by its data - same hashed contents create same blob, same SHA hash
    - Blob permissions (six digits): first two digits = type [directory vs file, e.g.], last three digits = three digits of chmod [user, group, other]
    - Blob - low-level representation of a git object
      - Git representation - contains “blob”, null byte, contents of blob, newline; compressed via gzip algorithm into a blob
        - Gzip \*stdin\* compresses stdin; gzip -d decompresses
      - Source code generally fairly compressible
      - Compression used as a space-saver (takes some CPU time for compression/decompression, but saves space on disk + time for reading/writing from/to disk)
  - **Tree** - directory-like (references other trees, blobs)
    - Stores to pointers to blobs, other trees
      - Typically used to represent contents of a directory/subdirectory
    - Maps names to IDs [other git objects - blobs, trees]
  - **Commit** - points to a single tree [represents project state at a certain point in time]
    - Commit object - contains hash of tree [top-level directory of commit - of working files/directories], hashes of parent commits, author, committer, commit message
      - Newer commits point back to older commits (linked list)
      - Author/committer contain names, email addresses, possibly authentication string
    - Does not commit changes from previous commit to current commit

- Tag - marker for a single commit

## Compression

- Compression
  - Speed vs compression quality trade-offs
    - More RAM also generally leads to better compression
  - Git compression algorithms: Huffman coding vs dictionary coding
    - Huffman coding - basic, cheap, well-understood
      - Basic problem: how to send a symbol string to a recipient
        - Symbols encoded in bytes; transmitted as bit stream
          - Common symbols given short bit encodings, shorter than the original byte-length symbol
          - All other symbols become longer (e.g. 0, -> 100, -> 10101; as long as no bit encoding is a prefix of another, can later be decoded)
            - Symbol bit encodings transmitted separately as a table of representation [i.e. of encodings]
    - Encoding process performed via binary tree
      - Create binary tree, with leaf nodes being symbols [containing the symbol + its rate of occurrence]
      - While nodes are not all connected: connect two rarest symbols with “supernode” (e.g. “j”, “q” -> parent node “j or q”)
      - Decoding can then traverse binary tree for each new symbol
    - Huffman – proved tree is optimal under assumptions (finite set of symbols, encoding into bits, known rates of occurrence)
    - May be used as a subroutine by higher-level algorithms
    - Dynamic Huffman coding – approach for Huffman coding in the case that rates of occurrence are not known/not assumed
      - Both compressor, decompressor begin with same, perfectly balanced Huffman tree (assume all characters occur equally, character bit representation = byte representation)
        - Result: initial characters uncompressed

- As compressor sends characters, both compressor and decompressor begin to fill in a Huffman tree on the fly based on the rates of occurrence of symbols in sent bits
  - Compressor and decompressor keep their own copies of the tree, but should be the same
  - Huffman trees stored in RAM
- Less space-efficient than static coding, but more flexible
- Dictionary coding -
  - Look for common strings of symbols (e.g. words), and find shorter encodings for those strings
    - Assigns integer encodings to each word
      - May be able to get away with 16 bits ( $2^{16}$  words) rather than writing out all 48 bits, e.g.
    - Requires more memory for dictionary (word-encoding table), but better compression
  - Dynamic dictionary coding – analogous to dynamic Huffman coding
    - zlib uses a window over the input - sender uses last region before currently-being-compressed region, as a dictionary
      - If a word in the currently-being-compressed region matches one in the window, express as the offset from start of window to old word (integers)
      - Window moved forward as compression progresses
- zlib essentially Huffman-encodes output of dictionary coding

## History of Git

- Motivation of version control – want to use small repository to record a large history
  - Cheaper than making full copies
- Principle of version control: save just one copy of each unchanged line in history, rather than making a new copy for each version
  - Later versions can obtain the line by looking at previous versions
  - Is an application-specific compression method (application: version control)
    - Uses a bit of CPU time to reduce memory consumption



- Approach: Files stored as an interleaved history – stores all lines that have ever appeared in the lines of the source code
  - When referring to individual copies, auxiliary table can simply refer to whichever lines in the interleaved history, were present in that version (contains line #s)
    - Can use line ranges to save space
    - Retrieving a version simply fetches the line numbers associated with that version in the auxiliary table
  - Downside: cost of version retrieval increases as the size of the history grows
    - Commits relatively expensive – must make new copies of the entire history
  - Used by SECS (proprietary AT&T Unix version control program)
- Revision Control System (RCS) – files + past versions stored in memory as a full copy of the latest version, plus reverse deltas to later versions
  - Cheap to retrieve latest version
  - Versions of a file in other branches accessed via reverse deltas to the common ancestor, then forward deltas to the branch head
    - Deltas stored as deletions/insertions at certain line numbers
      - Two-way deltas exist, but are more expensive
  - Developed by Walty Tichy @ Purdue
    - Drawbacks: couldn't handle multiple files at once
      - CVS: could handle multiple files
- CVS succeeded RCS; SVN (free) succeeded CVS; Bitkeeper (proprietary) combined elements of SVN, SECS
  - Git developed as an open-source alternative to Bitkeeper
- Git – operates on two-way deltas

## Low-Level Programming

- Importance of low-level programming: ensures performance, reliability
- C - language designed to be compiled + run across platforms, providing low-level access to memory + language constructs readily analogous to machine instructions, all with minimal runtime support
- C to C++ loses:
  - Data types: Classes [OOP, inheritance], structs
  - Language features: Namespace control, overloads, cin/cout
  - Exception handling [C exception library - comparatively clumsy]
  - Memory allocation on heap (just malloc, free)
    - No double frees, just null tests [(ptr)]
      - free(0) - true
- Malloc [ptr=malloc(idk)], void\* [generic pointer]
  - malloc(# bytes) vs calloc(num elements)
  - Issues: using before writing, not freeing [leak], double freeing, using after free
  - Malloc: SIZE\_MAX, -1, 0 [some object size 0 - allocates 1 byte]
    - realloc(ptr, n) -> reallocates contents to new container size n
      - Truncates if needed
- Misc: C library, kernel: meta programs
  - time \*pgm\*: real/user/system [CPU] time
  - strace \*pgm\*: outputs any program system calls [-o \*output file\*]
  - valgrind \*pgm\*: looks for suspicious activity (e.g. loads from far up the stack, memory leaks, bad memory accesses, etc.)
    - Runs as an interpreted program, sort of
  - gdb
  - top/ps - lists processes (esp. resource-intensive)

## GCC Flags

- Security: **-fstack-protector** [stack canaries]/-fstack-noprotector
  - **-mshstk**: Intel x86-64 protector (new)
    - Shadow stack - partial copy of return addresses, stored in \*shadow area\* of memory [only writable by certain things]

- Security: interpreting random bytes as jump instructions
    - mshstk: allows only jumps to endbranch instructions [new inst]
- Performance: **-O0**, **-O1**, ..., **-O3** [optimize: faster, slower compilation, larger memory footprint, less human-readable code, may reveal bugs] (higher # -> more optimization; -O0 = none; -O2: general)
  - **\_\_builtin\_unreachable()** [C; C++]: uncallable function (results in undefined behavior; compiler may do anything)
    - Not an instruction in and of itself; indicates to the compiler unintended behavior
  - **-fllto** - enables link-time optimization/whole-program optimization (GCC performs another optimization pass during the linking process)
    - Combination of -o [optimization], -f [may change instructions] flags
    - Optimizes each individual compiled file, knowing the contents of the other compiled files (but before linking)
      - May inline calls to short functions in another file, e.g.
    - Compiler may also find more bugs
  - Many optimization techniques are  $O(n^2)$
- Performance attributes:
  - **\_\_attribute\_\_((\_\_aligned\_\_(4096)))** - ensures alignment of data with 4096-size pages
    - Formatted for inclusion at start of C program [e.g. # !\_\_GNUC\_\_; #define \_\_attribute\_\_(x); #endif; #\_attribute\_\_((\_\_aligned...]
  - **function()\_\_attribute\_\_((cold))** - tells compiler a function is rarely called
    - Stores function in memory/RAM - does not bring instructions into instruction cache unless called
      - Instructions [in instruction cache] jump to cold section [somewhere in memory] rarely, as needed
    - In the case of sequential code (i.e. no loops, conditionals) with a cold function, the entire code section will be evaluated as code
  - **function()\_\_attribute\_\_((hot))** - tells compiler a function is often called
- Misc misc
  - ptrdiff\_t (C) - type denoting differences between pointers

## Debugging

- Debugging (>50% of time):
  - **Performance debugging** - making a program smaller/faster
  - **Correctness debugging** - fixing incorrect behavior
    - May include performance debugging (e.g. in real-time applications)

## Static Checking

- **Static checking** - running a program to examine the program-to-test for correctness, without running the program-to-test [i.e. just looking at the code]
- Has many related compiler options, depending on the needs/types of concerns of a project
  - Notable: **-Wall** [W -> warning, static checks many common cases], **-Wextra** [-Wall + more checks]
    - **-Wall: -Wcomment** [checks for possible mistakes in comments],  
**-Wparentheses** [checks for possible mistakes with operator precedence],  
**-Waddress** [always-false pointer comparisons], + more
      - **-Wstrict\_aliasing** [banned in Linux kernel] - warns against casting between pointers of different types
      - **-Wmaybe-uninitialized** - warns against accesses to possibly uninitialized variables
      - Arithmetic operators higher precedence than shifts; && > ||
    - **-Wextra: -Wtype-limits** [comparisons exceeding max size of data types (i.e. impossible comparisons) - may not be “impossible” for all system]
    - **-Wno-*\*warning-type\**** removes certain warning types
  - Recent flag: **-fanalyzer** (like -Wmaybe-uninitialized, more checks + interprocedural)
    - Catches more errors, but expensive ( $O(n^3)$ ) [should not combine w/ -fllto]
- **C static checking:**
  - C: **static\_assert(*\*some boolean\**)** will only allow a program to be compiled if the boolean evaluates as true
    - Does not generate instructions, but must be able to be evaluated at compile time [constants only]
  - C keywords:

- **`_Noreturn`** declares a function that does not return [to the caller function] (e.g. an exit function) [ex: `_Noreturn void exit(int status)`]
  - Performed as a jump, rather than call instruction [faster]
  - Can also be used by compiler to improve checks elsewhere (e.g. a branch calling a noreturn means the branch condition will always be false anywhere else)
- **`__attribute__((pure, access(readonly, parameter#, #bytesaccessed)))`** - calling a function twice with same arguments, same memory (i.e. same state in the #bytesaccessed after the parameter# [ptr]) will return the same value
  - Will not modify storage [memory]
  - Returned value depends only on current state of the program
  - Lets compiler reorder calls to pure functions if the program state is not changed between calls [allows for optimization] + improves static checking, knowing memory will not change
  - [[reproducible]] in C23
    - [[unsequenced]] – value independent of program state
- **`attribute((malloc(other_func,1)))`** – prints warnings if data is ever allocated by the function, but never freed [other\_func never called]
  - Has limits [static analysis] – if another function is called and takes a pointer [possibly freeing it], compiler may just give up
    - Rust “does better”
- **`#ifdef`** – if a compiler does not see a line of code, it may throw warnings that would have been solved by said line
  - [[maybe\_unused]] before a function parameter
- Static checking can be employed when going for especially refined programs (as opposed to quick, just functional programs)
  - Static analysis for performance –

## Dynamic Analysis

- **Profiling** - running a program on “typical” input and tracking which parts of the program (e.g. which instructions/branches/functions) are called + how often
  - Can be used to determine hot vs cold functions - can be fed back into the compiler

- Can also be used to determine areas of the code not covered in test cases
- GCC coverage options [**--coverage**] - generate profiles, creating a histogram of times run for each instruction address
- Issues: program runs more slowly when being profiled, “typicality” of input based on developer guesses [atypical -> potentially bad advice], takes time
- Static analysis for correctness
  - const on pointer (const char\*, e.g.) – tells compiler program will not store through the pointer [allows for more compiler checks]
    - Caution – make sure annotations match actual intended functionality
- **Dynamic analysis** - changing the behavior of a program, to dynamically inspect itself for bugs even while it runs
  - Can detect a wider variety of bugs than static analysis
  - Downsides: incurs runtime performance cost, only detects errors that occur in a particular run [still not 100%]
- **Runtime checking**
  - Behavior to multiplication overflow undefined for C, C++
    - Can be used for attacks – better to crash than to continue
  - GCC options (**-fsanitize**):
    - **-fsanitize=undefined**: crash on undefined behavior
    - **-fsanitize=address**: crash on bad pointer references
      - Catches subscript errors, dereferences to freed storage
      - Catches most, not all cases: low-level pointer manipulation, e.g.
      - Mutually exclusive with =undefined [difficult for GCC to do both]
    - **-fsanitize=leak**: check for memory leaks
      - Faster + more accurate than valgrind, since fsanitize can see the source code
    - **-fsanitize=thread**: dynamically detect race conditions in multithreaded [parallel] applications
    - **-fwrapv**: reliable wrap-around on integer overflow (i.e. will always wrap around - no undefined overflow behavior)
      - Used in the Linux kernel for low-level manipulation
      - Side effect: cannot check for integer overflows

- **Portability checking** – the process of checking to ensure that a single program can run on many different systems
  - Can work across different OSes/ISAs, e.g.
  - GCC options:
    - `-m32`: generates a 32-bit executable [ex: 4-byte pointers]
      - If this fails, likely indicates portability bug (e.g. implicit assumption of 8-byte pointers)
- *Avoiding debugging*:
  - Test cases
    - **Test-driven development** – write tests first, then the code
      - “Every bug fix should come with a test case”, e.g.
      - No commits without passing every test
    - **Reliability checking** [regression testing]
  - **Defensive programming** – assume other modules in the program may or may not be buggy, and program extra cautiously [with extra protections] as a result
    - **Exception handling** – try/catches, asserts, e.g.
    - Traces and logs – logging program activity in the event of a crash, e.g.
  - **Barricades** – separate app into organized, bug-free safe space, potentially dangerous “outside world” (bad data, e.g.)
    - Outside portion rigorously checks passed-in data via runtime checks; only allows approved data into safe space [no checks – more efficient]
  - **Interpreters** – programs that run other programs
    - Used in browsers (JavaScript interpreter), e.g.
    - Usually contain many runtime checks
    - Virtual machines – hardware support for hardware-level interpretation
- General debug advice:
  - Don’t guess where bugs are; find them
  - Stabilize the failure first (make it reproducible)
    - Upon seeing the symptoms, locate the failure’s case
- Debugger contains program execution history, exploration tool
  - Runs program under debugger control
  - Debugger exerts significant control over subsidiary programs/processes
  - `Gcc -g3`: generates extra info in executable to help debugger (e.g. gdb)

- Avoid -O2 with -g3 [O2 may reorder code]
    - -Og – optimize for debugging (avoid reordering, e.g.)
  - Gdb: attach \*process ID\*
- *Gdb commands:*
  - Detach - let program run freely [cede control]
    - Cannot attach multiple gdb's to the same program
  - Exit - exit gdb
    - Lets program run freely [if already existing]
  - Run \*arguments\* - runs underlying program (with arguments)
    - Effectively suspends gdb until something happens (gdb waits)
  - C-c - suspends program, returning to gdb control
    - Program crashes place program into suspended state, allowing backtrace
  - Bt - gets a backtrace
    - Examines suspended program's state, namely contents of the stack + machine addresses
      - Prints all functions down to main
  - P \*address\* – prints contents of memory at address
    - Options: p/x = print hexadecimal, e.g.
  - Info registers – print registers
  - P \*variable\* = \*value\* – set variable value
    - P \$\*register\* = \*value\* – set register value
    - P \*variable\* = \*function\* – calls function, sets variable value
    - P \*function\* – calls function
- Gdb - “program execution history explorer”
  - Breakpoints - pre-set locations to automatically suspend a program when encountered
    - Gdb b \*address\*; then run the program
      - Gdb b \*function name\*
    - Continue – continue after breakpoint
      - Continue 10 – continue for the next 10 breakpoints
        - Alt: step (single source code line), stepi (single machine instruction), next (like step, but upon function call – keeps running until function returns)



- In the case of interweaved source code lines in machine instructions, step will stop at every change in source code line
  - Fin – finish program
  - U \*address\* – execute until the line is reached, or program finishes (sets temporary breakpoint)
- Reverse continue [rc] – runs program backward until previous breakpoint
  - Requires starting gdb with special option
- Watch \*variable\* – stops program whenever the variable is modified
  - Sets a breakpoint on data, essentially
- Checkpoint – sets checkpoint at current line (e.g. known “good” location)
  - Restart # – restarts at checkpoint number
- Gdb actually modifies memory for the duration of the program
  - Breakpoints - temporarily adds an instruction before a breakpoint [in machine code] that crashes program; gdb then re-adds original instruction after crash upon continue
- Disas – disassembles machine code
- Debugging gdb
  - GDB – call old, unbuggy version of gdb on itself [new, buggy version]

## Makefiles

- Command ``make`` used for incremental builds (e.g. recompiling only one file in a project, out of many files)
  - Build configuration specified in Makefiles
    - Use combination of shell, “make” notation
- Makefile syntax:
  - *Make rule*: `*target file name*: *included/dependency files* \n *shell commands*`
    - Will only rebuild target files if at least one included file [dependency] is newer than the target file
      - Does not automatically build dependencies, if they do not exist; can be written such that ``make`` automatically builds all dependencies
  - Can declare, use variables (similar to shell variables, but no quotes in declaration; called via `$(varname)` rather than `$varname`)
- Command ``make``
  - Builds first target, by default [no arguments]; with command-line arguments, can specify the make target (e.g. ``make foo.o``)
  - ``make clean`` removes all target files (can be used for complete rebuild, e.g.)
  - Missing dependencies can cause issues; extra dependencies lead to inefficiencies
    - Dependency loops [circular dependencies] not forbidden (will run one of two ways), but discouraged
  - Flags: `-j` [build in parallel; `-j10` = at most 10 jobs, e.g.]; `-j` = no limits
    - `make` will automatically handle waiting for a target’s dependencies to build before building the target
- *Makefiles for larger systems*:
  - **Approach 1**: centralized Makefile at root, importing subsidiary Makefiles in the various subdirectories [include `sub/sub1.mk`]
  - **Approach 2**: Makefile at root that `cd`’s into and runs ``make`` within subdirectories (based on those subdirectories’ individual Makefiles)
    - Primarily used in earlier versions of `make`

## Backups

- Motivation: ~100 ZB of data is produced/year by the Internet; ~2/3 of it is copied
  - Devs use git; ops staff use backups
- *Backup strategies:*
  - Use **failure models** to determine what constitutes a failed backup
    - Potential models: hardware failure (e.g. disk drive failure) [partial vs whole-drive failure], accidental (or purposeful/malicious) file deletion/modification
    - Judging risk:
      - TBW denotes manufactured-rated, warranted amounts of safe writes to a flash drive from ordinary wear
      - Disk drives have AFR (annualized failure rate – generally sub-1%)
  - May not need to back up everything: can omit OS installation files, temporary files (e.g. /tmp/ – generally stored in RAM)
    - Do back up: most files, incl. data [file contents], metadata
  - **Break backup into different units**
    - Break data, metadata into discrete byte blocks of a specified size – faster for a backup system to run at a lower level, not need to worry about distinct file boundaries
    - Alt: files as units (simpler to users, but less efficient)
  - **Time units** - how often to backup
    - Compromise between backing up every change, never backing up: back up on important changes/at important times
  - **Garbage collection** - reclaiming storage from unneeded backups
- Can save data via reducing backup frequency, reducing storage of unimportant data, more aggressive garbage collection
  - Can use cheaper, slower drives; remote backup services
  - Automation of cheaper backups:
    - **Data grooming** – removal of unnecessary data
    - **Deduplication** – looking for duplicates of stored data, and removing them (replaces them with pointers to the original, essentially)

- **Incremental backup** – later backups copy only changed data relative to the previous backup (rather than copying all data)
  - Same philosophy as git (inspired git)
  - Ex: if backing up by file, back up only recently changed files
  - Inherently more fragile – losing a single past backup prevents retracing
- Can **compress** file data for smaller backups
- Other techniques
  - Can **encrypt data** for greater security
  - Can stage **different levels** of copies on different storage devices
- **Backup recovery** – backups are useless unless they can be recovered from
  - Require testing to make sure backups are sufficient for restarting from even in the case of a complete data loss, e.g.
- Strategies for file-system-controlled backups
  - **Versioning file systems** – similar to git, but run on a filesystem (see: TENEX)
    - Slightly worse than git, but automatic – no user input needed
    - Keeps copies of older versions of files
      - May be kept to important versions only to save space
      - File version added as a suffix to the file name (lowest number -> most recent backup; no suffix -> latest version)
        - Multiple versions of the file will be seen upon ls -l
        - Saving the file bumps version number
    - Possible downsides: changes to multiple files, system overhead from storing/updating copies
      - Losing a past version makes it difficult to reclaim the rest
      - Hard to understand for users
  - **Snapshotting file systems** – taking pictures of the state of the entire file system at single points in time
    - File system states stored as read-only copies of the entire file system tree
    - Snapshotting operation made cheaper via copy on write (CoW) – full copies of files are only performed when the file has been modified
      - File systems modeled via trees of addressed block-structured objects; a new snapshot is just a pointer to a tree

- Updates to a file create minimal amounts of new objects (typically just copies of old ones) to be created
- More popular (used on SEASnet, e.g. – creates .snapshot subdirectories)
  - Note: not listed by `ls -al` (but still `cd`-into-able)
- Drawbacks: still susceptible to losing parts of history, incurs overhead

## “Cloud Computing”

- The name “cloud computing” originated as a joke (“the server is just somewhere off in space – the cloud”); to some extent a marketing term
- [Technical aspects of] **Cloud computing** – an on-demand, self-service method for clients to gain access to servers with a low barrier of entry
  - Intended to be rapidly scalable (elastic) for users – clients should be able to spin up servers quickly as needed
    - Broad network access – should be (relatively) location independent
  - Multitenancy – shared resources (server pool) distributed across many tenants
  - Downsides: service providers must keep track of client server usage (incurs overhead)
- Possible implementations:
  - **Infrastructure-as-a-service (IaaS)** – service provider provides actual physical hardware + networking infrastructure; customer is responsible for actually configuring + running it
    - Customer provides + maintains OS, apps, etc.
    - Provided by AWS, e.g.
  - **Platform-as-a-service (PaaS)** – service provider provides infrastructure + middleware: preconfigured OS, standard apps (e.g. Python)
    - Customer provides only custom software
    - Creates some hassle for the service provider: might need to tailor software versions to client requests, e.g.
    - Provided by Google App Engine + AWS, e.g.
  - **Software-as-a-service (SaaS)** – service provider provides platform + entire software stack (i.e. programs to be run)
    - Customer responsible only for configuring the provided software (e.g. tweaking settings); does not need to write any software
    - Provided by Salesforce, e.g.
- **Virtual machine** – “software pretending to be hardware”
  - VMs may emulate other hardware
    - Can be written such that the emulated environment is entirely insulated from the outside environment (provides a barrier)

- Drawback: Incurs large amounts of overhead (performance)
  - VMs may instead take a hardware-supported approach to emulation: simulated machine/architecture is limited to match that of the actual hardware
    - Allows for utilization of hardware support for most common virtualized instructions (faster than straight emulation)
      - Some commands (e.g. halt) may not be hardware-supported
    - Can be used to emulate programs running on the same ISA/hardware (e.g. Linux atop macOS)
      - Drawback: still has some overhead, relative to running natively
  - **Containers** – VMs specialized for virtualizing the same OS as the physical machine
    - Can share hardware instructions + shared OS files (e.g. /lib, /usr/bin)
      - Stores only one copy of kernel, executables
      - Executables stored in read-only memory
    - Containers (notably Docker) – standard approach for most cloud computing platforms, IaaS excepted
      - VMs insulated – client cannot see what others are doing, e.g.
      - Alternative to giving each client a process on the same server, e.g.
- Docker plumbing:
  - Objects:
    - Containers – environments to run applications
    - Images – templates for building containers
    - Services – let containers scale across multiple Docker domains
      - Can coordinate large numbers of containers into a single service
  - Registries – repositories for Docker images
    - Can pull, push images
  - Dockerd[daemon] – Docker daemon responsible for containers
    - Containers implemented as Linux processes controlled by the Docker daemon (running at a root level), essentially
    - Docker shell commands operate through requests to the daemon
- Docker porcelain:
  - Docker maintenance/management programs responsible for load balancing, managing resource constraints, logging/monitoring systems and system errors, security (avoiding break-ins to containers), updates to a running system

- Updates to a running system: blue/green upgrades (installing updates on a new set of containers separate from the running containers, then switching over when the new version is ready)
  - Programs: Kubernetes, cloud-specific programs (e.g. AWS: Elastic Container Service/ECS)
- Alternate model to containerization: “serverless” computing
  - Serverless computing: writing code ignorant of the server to be run on, then allowing the infrastructure to route transactions to the appropriate server
    - Servers maintained by service provider
    - Allow for simpler code – can ignore server state (down vs up?), server location, e.g.
    - Downside: not knowing which server the code will be run on makes caching data on a single server more difficult (performance downside)



## Software & The Law

- Law & ethics (law vs ethics)
- Types of law: commercial law (to do with contracts, business disputes), criminal law (crimes, injury), international law (international agreements), admiralty law (ships at sea)
  - Have differences (unanimous consent vs majority opinion in juries, e.g.)
- Commercial law: most software falls under purview of copyright law (protecting creative works), patent law (inventions), trade secrets, trademarks (logos, short descriptions of the product), personal data (modern field), future: AI law
  - Different expirations: copyright (life of author + 75 years), patents (20 years), trade secrets (indefinite)
  - Enforcement: sue [for \$\$] under commercial law for using intellectual property without permission, e.g.
    - Alternative: technical protections – preventing trade secrets from leaking via simply limiting access to them, e.g.
- License – legal permits [grants of permission] to use/copy/etc. software
  - Free software/open-source software – use licenses to further the goals of the organization, rather than for commercial purposes
  - Approaches:
    - Public domain – no rules on usage; given away (openest approach)
    - Academic license – requires citations of sources [BSD, e.g.]
    - Reciprocal license – similar to academic license; any changes built upon the original must also be redistributed under reciprocal license [GPL, e.g.]
    - Forced contribution – any changes must be given back to, owned by the original authors (Apple, Microsoft, etc.)