

EXPLORING SEMANTIC HIERARCHIES TO IMPROVE RESOLUTION THEOREM
PROVING ON ONTOLOGIES

by

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ABSTRACT

A resolution-theorem-prover (RTP) evaluates the validity (truthfulness) of conjectures against a set of axioms in a knowledge-base. When given a conjecture, an RTP attempts to resolve the negated conjecture with axioms from the knowledge-base until the prover finds a contradiction. If the RTP finds a contradiction between the axioms and a negated conjecture, it proves the conjecture.

The order in which the axioms within the knowledge-base are evaluated significantly impacts the runtime of the program, as the search-space increases exponentially with the number of axioms.

Ontologies, knowledge bases with semantic (and predominantly hierarchical) structures, describe objects and their relationships to other objects. For example, a 'Car' class might exist in a sample ontology with 'Vehicle' as a parent class and 'Bus' as a sibling class. Currently, any hierarchical structures within an ontology are not taken into account when evaluating the relevance of each axiom. At present, each predicate is automatically assigned a weight based on a heuristic measure (such as the number of terms or the frequency of predicates relevant to the conjecture) and axioms with higher weights are evaluated first. My research aims to intelligently select relevant axioms within a knowledge-base given a structured relationship between predicates. I will use the semantic hierarchy over predicates to assign weights to each predicate passed to a weighting function. The research aims to design heuristics based upon the semantics of the predicates, rather than solely the syntax of the statements.

I plan to develop weighting functions based upon various parameters relevant to the ontological structure of predicates contained in the ontology, such as the size and depth of a hierarchy based upon the structure of the ontology.

I will implement methods to calculate weights for each predicate and thus each axiom in attempts to select relevant axioms when proving a theorem. Then, I will conduct an experimental study to determine if my methods show any improvements over current reasoning methods.

ACKNOWLEDGEMENTS

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List of Figures

1	This simple ontology describes classes (black) and instances (red) in the "automobiles" ontology (inspiration from an example by Natalya F. Noy) [5].	3
2	This shows a resolution tree.	5
3	This shows the input into the prover9 program	5
4	This shows the Prover9 output for the above proof.	6
5	An illustration of my process.	7
6	These are the default weights in the Prover9 program [4].	8
7	These are example weights given to the Prover9 program [4].	8
8	The graphical user interface for Prover9 on macOS [4].	9

List of Tables

1	multidim space voids weights	12
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Contents

ABSTRACT	iii
ACKNOWLEDGEMENTS	iv
LIST OF FIGURES	v
LIST OF TABLES	vi
1 INTRODUCTION	1
2 BACKGROUND AND RELATED WORK	3
2.1 Ontologies	3
2.2 Theorem Proving	4
2.3 Semantic Similarity	5
3 APPROACH	7
3.1 Converting Ontologies	7
3.2 Generating a Complete Heirarchy	7
3.3 Default Weights	8
3.4 Computing Weighting Functions	8
4 EXPERIMENTS	9
4.1 Setup	9
4.2 Results	9
5 CONCLUSIONS	10
REFERENCES	11
A TESTS	12
AUTHOR’S BIOGRAPHY	14

1 INTRODUCTION

The rules of logic enable one to prove theorems from axioms stored in a knowledge-base. Axioms, asserted facts typically expressed in a formal manner, provide a computer program with tools to confirm or refute conjectures without additional user input. Because computers excel at simple and repetitive tasks, one can witness the applications of automated theorem proving in fields which rely heavily on "knowledge acquisition and information retrieval" [sanchez2012ontology]. The ability for machines to deduce logically valid conclusions has applications in artificial intelligence and a variety of scientific domains. Automated theorem proving provides a versatile method for reasoning with a set of facts, and has been used to prove and verify proofs of multiple theorems. The four color map theorem, obtained by an automated proof in 1976 by a general-purpose theorem-proving software in 2005 remains a notable example [2]. Moreover, advances have been made in work on the Kepler conjecture and in finding optimal solutions for a Rubik's Cube. The general-purpose nature of automated theorem proving yields applications to a variety of problems. However, automated theorem proving programs often neglect semantic knowledge embedded in an ontology.

Ontologies provide a "common vocabulary" for researchers to speak about a specific domain by describing entities and the relationships between them [5]. A formal description of a specific environment provides researchers and machines with a shared understanding by aiming to capture the semantics of a domain's concepts and relations. Some relationships between an ontology's terms may be explicitly defined, but many are implicit.

$$\begin{aligned} &isSedan(SubaruLegacy) \\ \forall x \, isSedan(x) &\rightarrow isAutomobile(x) \end{aligned}$$

The former axiom asserts a Subaru Legacy is a sedan and the latter asserts all sedans are automobiles. While one can easily deduce the fact a Subaru Legacy is an automobile, no single axiom explicitly describes such a statement. Fortunately, formal logic defines rules of inference which allow one to transform established facts into new conclusions solely based on the syntax

of these statements. Both humans and computers can clearly distinguish well formed statements $(x + y = 4)$ from those which are not $(x4y+ =)$. Beyond the syntax of statements, ontologies and logics also define the semantics or meaning of sentences (i.e. declaring $x + y = 4$ is true when $x = 1$ and $y = 3$ but false when $x = 0$ and $y = 1$).

Like many taxonomies, one can often form maps of relationships within an ontology which resemble a hierarchy. Currently, knowledge encoded in semantic hierarchies constructed from ontologies is not taken into account when determining which axioms might be most helpful when attempting to prove a specific theorem. This work attempts to improve automated theorem proving with ontologies by identifying relevant facts, and ignoring those less likely to yield a proof.

DESCRIPTION OF THE CONCRETE OBJECTIVE AND APPROACH OF YOUR WORK

2 BACKGROUND AND RELATED WORK

2.1 Ontologies

The word ontology ("study of being") combines Greek *onto-* ("being") and *-logia* ("logical discourse"). The act of studying knowledge and existence in philosophy has given birth to the study of formal logic and automated reasoning in computer science. Researchers often use ontologies to share information among people or computer programs, to enable domain knowledge reuse, to make definitions of a particular domain explicit, or to analyze domain knowledge [5]. While ontologies can often witness reuse, ontology design rarely encounters a single best solution satisfying all goals of the ontology.

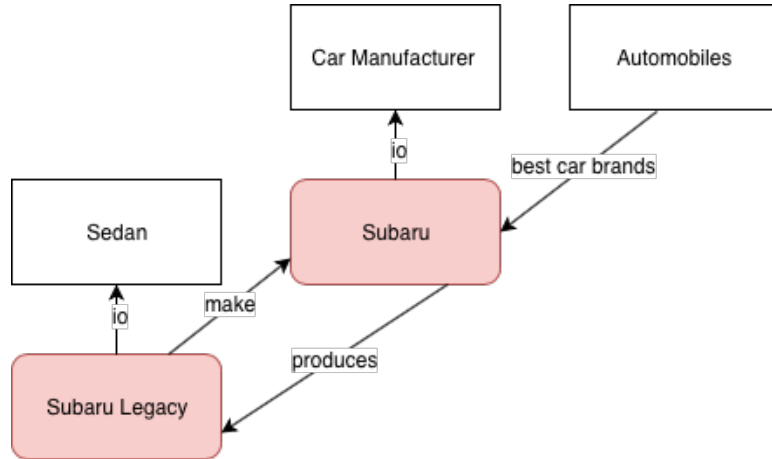


Figure 1: This simple ontology describes classes (black) and instances (red) in the "automobiles" ontology (inspiration from an example by Natalya F. Noy) [5].

In reality, few differences between an ontology and a knowledge base exists. Knowledge engineers must traverse a "fine line where the ontology ends and the knowledge base begins" [5]. At the least, an ontology defines categories (or classes) and relationships among objects. One can think of an ontology as a "vocabulary" used to describe a domain [6, p. 308]. Typically, the classes can be arranged in a hierarchy and the relationships into a separate hierarchy, which are here referred to as semantic hierarchies. When designing an ontology, one must decide the scope and organization

of the knowledge, along with the language used to describe a specific domain. Refer to Figure 1 for an example.

At a simplistic level, semantic hierarchies describing an ontology can be compared to a family tree. Given a pair of two individuals, if one were tasked with determining if two individuals are related, a family tree would prove quite useful.

First order logic exists as one possible language to describe ontologies.

2.2 Theorem Proving

Formal logics like first-order logic defines a structure for statements which can be used to form logical and mathematical proofs. Asserted facts, called axioms, are used to derive facts which logically follow. Consider the following set of facts.

$$\begin{aligned} & isSedan(SubaruLegacy) \\ \forall x \, isSedan(x) & \rightarrow hasFourSeats(x) \end{aligned}$$

The first statement asserts the Subaru Legacy is a sedan. The second statement asserts all sedans have four seats. These two statements do not directly state the Subaru Legacy has four seats. However, one can derive the statement $hasFourSeats(SubaruLegacy)$ by using the inference rule *modus ponens*, defined below.

$$\begin{aligned} & A \\ A & \rightarrow B \\ \therefore & B \end{aligned}$$

One can think of A and B as variables representing statements, and any statement can replace them. An inference rule defines a valid rule for statements. By expressing facts in a formal notation, one makes proofs using such statements mechanical and easily parsed by a computer. Expressing statements in formal logic poses multiple challenges. First, each object and relationship must be explicitly defined.

Resolution is one of many methods for automated theorem proving. Historically, resolution has significance and is widely used [1, p. 51]. In order to use resolution as a proof technique, axioms must

first be expressed in Conjunctive Normal Form. This can be done by following a 7-step procedure of converting the set of facts into a conjunction of disjunctions.

$$isSedan(X) \rightarrow hasFourSeats(X)$$

$$\neg isSedan(X) \vee hasFourSeats(X)$$

One can then resolve the statements.

$$\frac{isSedan(SubaruLegacy), \neg isSedan(X) \vee hasFourSeats(X)}{hasFourSeats(SubaruLegacy)}$$

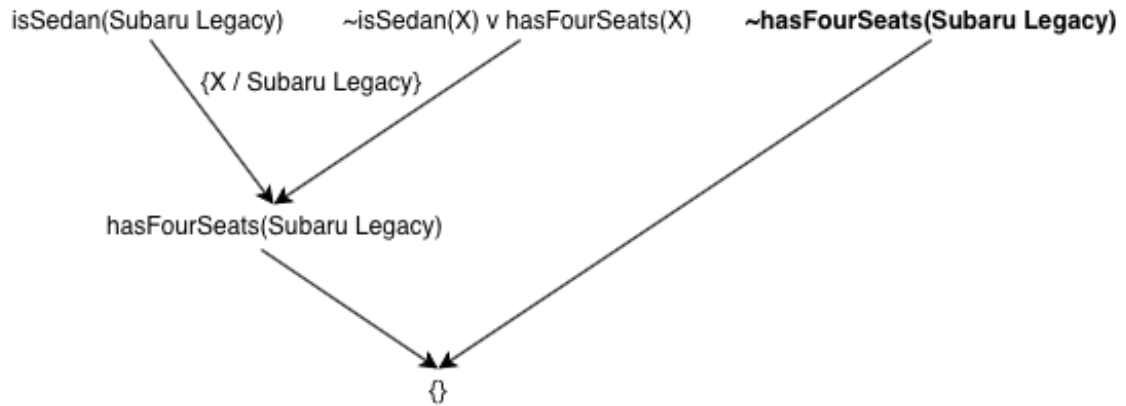


Figure 2: This shows a resolution tree.

Assumptions:	Well Formed?	Clear
isSedan(SubaruLegacy). all x (isSedan(x) -> hasFourSeats(x)).		
<hr/>		
Goals:	Well Formed?	Clear
hasFourSeats(SubaruLegacy).		

Figure 3: This shows the input into the prover9 program

In resolution, the search space increases exponentially with the addition of each new axiom.

2.3 Semantic Similarity

One straightforward method of calculating the semantic similarity between two entities in an ontology.

```

1 (all x (isSedan(x) -> hasFourSeats(x))) # label(non_clause). [assumption].
2 hasFourSeats(SubaruLegacy) # label(non_clause) # label(goal). [goal].
3 -isSedan(x) | hasFourSeats(x). [clausify(1)].
4 isSedan(SubaruLegacy). [assumption].
5 hasFourSeats(SubaruLegacy). [resolve(3,a,4,a)].
6 -hasFourSeats(SubaruLegacy). [deny(2)].
7 $F. [resolve(5,a,6,a)].

```

Figure 4: This shows the Prover9 output for the above proof.

In an undirected graph G defined as a pair (V, E) , where V is a set of vertices, and E is a set of edges between the vertices $E \subseteq (u, v) | u, v \in V$, one can define a path $path(a, b) = l_1, \dots, l_k$ as a set of links connecting a and b in a taxonomy and $|path(a, b)| = k$ as the length of the path [sanchez2012ontology]. One can calculate the semantic distance between a and b using equation 1 [rada1989development].

$$sim(a, b) = \min_{\forall i} |path_i(a, b)| \quad (1)$$

By incorporating depth of the taxonomy into the function, improvement has been seen [wu1994verbs].

$$sim(a, b) = \frac{2 \times N_3}{N_1 + N_2 + 2 \times N_3} \quad (2)$$

3 APPROACH

In efforts to quantitatively evaluate the effectiveness of the proposed weighting functions, a series of experiments were conducted on multiple ontologies. The ontologies selected for testing exist in the COmmon Logic Ontology REpository (COLORE). The objective of COLORE is to serve as a "testbed for ontology evaluation and integration techniques" [3]. Code from the project can be found at <https://code.google.com/p/colore>.

My experiments were conducted using Prover9, written by William McCune [4]. Many tests were conducted using a version of the program which supports a GUI, but a command line version is available for running automated tests. Git was used for version control and a repository containing source code can be found at <https://github.com/stanleysmall/thesis>.

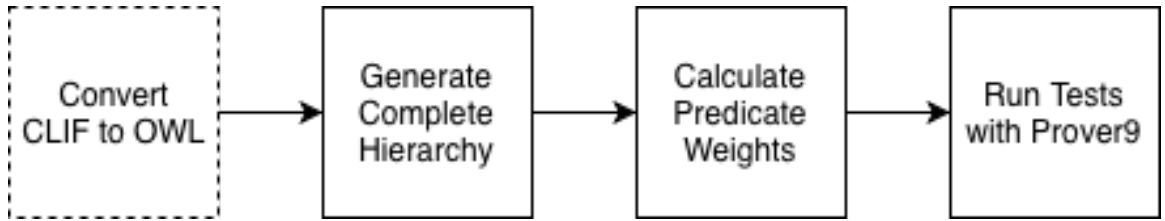


Figure 5: An illustration of my process.

3.1 Converting Ontologies

COLORE

3.2 Generating a Complete Heirarchy

Given an ontology and a conjecture, one can generate weights for classes and properties to reduce the number of clauses generated in a proof. Functions are evaluated by their influence on the number of clauses generated with a proof.

After a hierarchy has been generated, weights can be assigned to each class and subproperty. The same weighting function is applied to both the classes and sub-properties.

The weighting functions are currently applied by hand to the ontologies, with the beginnings of an automated program underway.

3.3 Default Weights

- The default weight of a constant or variable is 1.
- The default weight of a term or atomic formula is one more than the sum of the weights of its arguments.
- The default weight of a literal is the weight of its atomic formula.
- The default weight of a clause is the sum of the weights of its literals.

Figure 6: These are the default weights in the Prover9 program [4].

```
list(weights).  
  weight(a)=3.                % the weight of the constant a is 3  
  weight(f(a,x))=5*weight(x). % weight(f(a,term))=5*weight(term)  
  weight(f(a,_))=-1.          % _ matches any variable  
  weight(x|y)=2+(weight(x)+weight(y)). % add 2 for each "or" symbol  
end_of_list.
```

Figure 7: These are example weights given to the Prover9 program [4].

3.4 Computing Weighting Functions

4 EXPERIMENTS

4.1 Setup

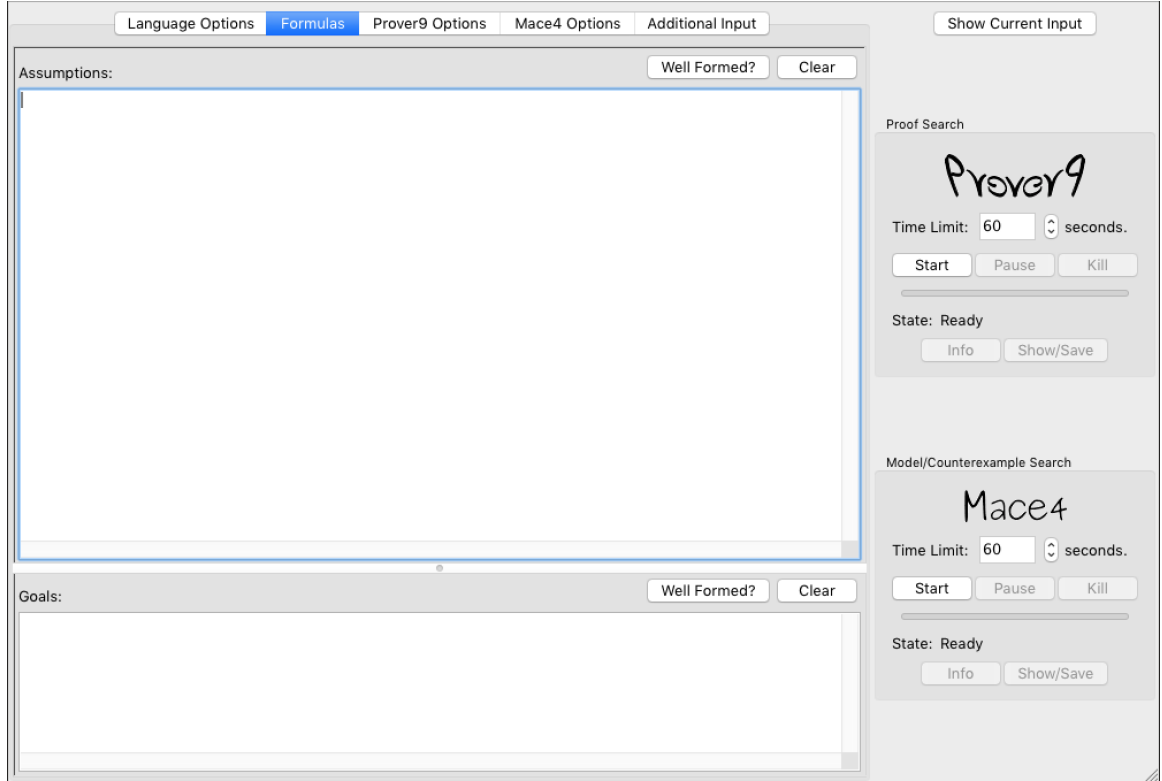


Figure 8: The graphical user interface for Prover9 on macOS [4].

4.2 Results

4.2.1 multidim_space_voids

5 CONCLUSIONS

In many cases, the algorithm increases the number of clauses generated for each test, but does not do so to the point where the tests are unusable. For some very specific ontologies, the number of clauses generated decreases. This can be attributed to, in part, by the small number of ontologies available for testing, along with the specific pattern and hierarchy of each ontology.

Many opportunities for further research include fully automating the search procedure, working with a larger number of ontologies to ensure the weighting functions actually do as they say, developing a new approach towards automatically weighting the predicates.

References

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- [5] Natalya F Noy, Deborah L McGuinness, et al. *Ontology development 101: A guide to creating your first ontology*. 2001.
- [6] Stuart J Russell and Peter Norvig. *Artificial intelligence: a modern approach*. Malaysia; Pearson Education Limited, 2016.

A TESTS

Table 1: multidim space voids weights

Entity	Type	Superclass(es)	1	2	3	4	5	6	7	8	9
CAVITY	Class	owl:Thing									
Closed	Class	owl:Thing									
ComplexV	Class	V									
Con	Class	S									
DPF	Class	F									
F	Class	"PED Ψ RPFforDPF"									
Gap	Class	owl:Thing									
HOL	Class	owl:Thing									
Hole	Class	owl:Thing									
Icon	Class	Con									
M	Class	"PED Ψ mat"									
Max	Class	S									
MaxDim	Class	S									
Min	Class	S									
MinDim	Class	S									
NAPO	Class	POB									
PED	Class	POBorMorF									
POB	Class	PED mat									
POBorMorF	Class	owl:Thing									
POBorMorRPF	Class	owl:Thing									
POBorRPF	Class	owl:Thing									
RPF	Class	F mat									
RPFforDPF	Class	owl:Thing									
S	Class	owl:Thing									
SimpleV	Class	V									
SimpleVorComplexV	Class	owl:Thing									
TUN	Class	owl:Thing									
V	Class	SimpleVorComplexV									
ZEX	Class	S									
mat	Class	POBorMorRPF									
BCont	ObjectProperty										
C	ObjectProperty										
Cont	ObjectProperty										
Covers	ObjectProperty										
DK1	ObjectProperty										
EqDim	ObjectProperty										
ICont	ObjectProperty										
Inc	ObjectProperty										
P	ObjectProperty										
PO	ObjectProperty										
PP	ObjectProperty										
SC	ObjectProperty										
TCont	ObjectProperty										
VS	ObjectProperty										
ch	ObjectProperty										
gt	ObjectProperty										
hosts	ObjectProperty										

Entity	Type	Superclass(es)	1	2	3	4	5	6	7	8	9
hostscavity	ObjectProperty										
hostscavityi	ObjectProperty										
hostscavityt	ObjectProperty										
hostsg	ObjectProperty										
hostsh	ObjectProperty										
hostshollow	ObjectProperty										
hoststunnel	ObjectProperty										
hostsv	ObjectProperty										
hostsve	ObjectProperty										
hostsvi	ObjectProperty										
lt	ObjectProperty										
r	ObjectProperty										
"="	ObjectProperty										
"gt="	ObjectProperty										
"lt="	ObjectProperty										

AUTHOR'S BIOGRAPHY

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