Theory of Generalization

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1 Restriction of Break Point

what 'must be true' when minimum beak point k=2

- N = 1: every $m_H(N) = 2$ by definition
- N = 2: every $m_H(N) < 4$ by definition (so maximum possible = 3)

maximum possible $m_H(N)$ when N=3 and k=2?

maximum possible so far: 4 dichotomies

• N = 3: maximum possible = $4 << 2^3$

Break point k restricts maximum possible $m_H(N)$ a lot for N > k

idea:

$$m_H(N) \le$$
maximum possible $m_H(N)$ given k $\le poly(N)$ (1)

2 Bounding Function: Basic Cases

2.1 Definition

bounding function B(N, k): maximum possible $m_H(N)$ when break point = k

- combinatorial quantity: maximum number of length-N vectors with (o,x) while 'no shatter' any length-k subvectors
- irrelevant of the details of H e.g. B(N,3) bounds both
 - positive intervals (k = 3)
 - 1D perceptrons (k = 3)

new goal: $B(N, k) \leq poly(N)$?

2.2 Table of Bounding Function

3 Bounding Functions: Inductive Cases

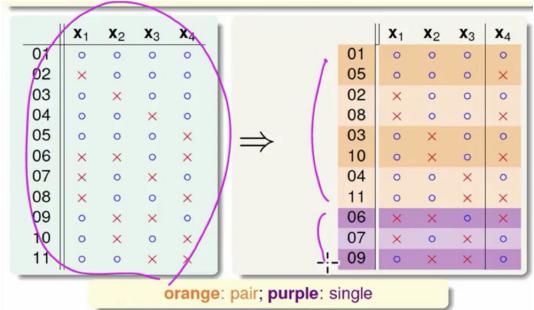
3.1 Estimating B(4, 3)

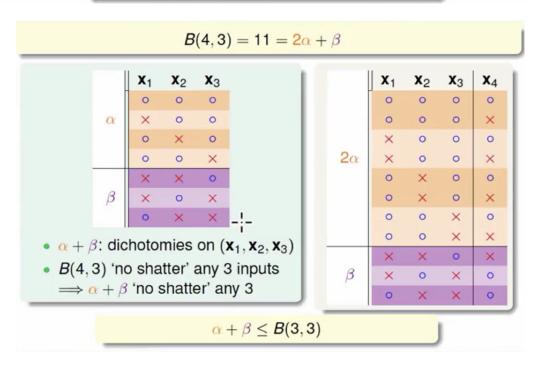
Motivation

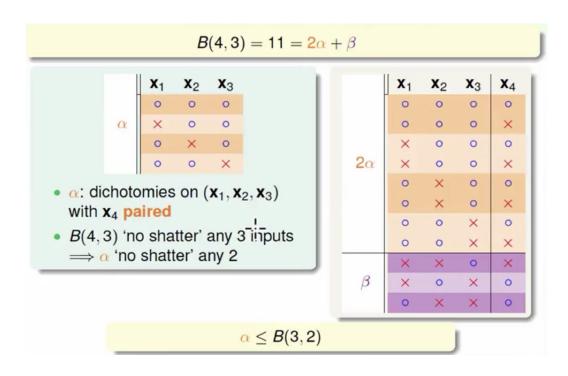
• B(4,3) shall be related to B(3,?) 'adding' one point from B(3,?)

next: reduce B(4,3) to B(3,?)

after checking all 2^16 sets of dichotomies, the winner is...







$$B(N,k) = 2\alpha + \beta$$

$$\alpha + \beta \le B(N-1,k)$$

$$\alpha \le B(N-1,k-1)$$
(2)

$$B(N,k) \le B(N-1,k) + B(N-1,k-1) \tag{3}$$

3.2 The Theorem

$$B(N,k) \le \sum_{i=0}^{k-1} C_N^i N^i \tag{4}$$

- · simple induction using boundary and inductive formula
- for fixed k, B(N, k) upper bounded by poly(N)
- actually \leq can be = $m_H(N)$ is poly(N) if break point exists

4 A Pictorial Proof

Bad Bound for General H,

actually, when N large enough,

$$P[\exists h \in H, s.t. | E_{in}(h) - E_{out}(h)| > \epsilon] \le 2 \cdot 2m_H(2N) \cdot exp(-2 \cdot \frac{1}{16}\epsilon^2 N)$$
(5)