Q1. For each of the following 6 program fragments, give a Big-Oh analysis of the running time (3 points) -

```
(1)
       sum = 0;
       for (i = 0; i < n; i++)
       ++sum;
       Ans: O(n), input = steps.
(2)
       sum = 0;
       for (i = 0; i < n; i++)
               for (j = 0; j < n; j++)
                       ++sum;
       Ans: O(n<sup>2</sup>)
(3)
       sum = 0;
       for(i = 0; i < n; i++)
               for (j = 0; j < n*m; j++)
                       ++sum;
       Ans: O(n*m)
(4)
       sum = 0;
       for (i = 0; i < n; i++)
               for (j = 0; j < i; j++)
                       ++sum;
       Ans: O(n<sup>2</sup>)
(5)
       sum = 0;
       for (i = 0; i < n; i++)
               for (j = 0; j < i*i; j++)
                       for (k = 0; k < j; k++)
                              ++sum;
       Ans: O(n<sup>5</sup>)
(6)
       sum = 0;
       for (i = 0; i < n; i++)
               for (j = 0; j < i*i; j++)
                       if (j \% i == 0)
                              for (k = 0; k < j; k++)
```

Ans: **O(n⁴)**

Q2. Programs A and B are analyzed and found to have worst-case running times no greater than $150Nlog_2N$ and N^2 , respectively. Answer the following questions (3 points) -

a. Which program has the better guarantee on the running time for large values of N (N > 10,000)?

Ans: Program A

b. Which program has the better guarantee on the running time for small values of N (N < 100)?

Ans: Program B

c. Which program will run faster on average for N = 1000?

Ans: Program B

Q3. Q3. Solve the following recurrence relations using the Master theorem (2 points) -

a. T(n) = 3T(n/2) + n/2

Ans: $O(n) = O(n^{log}_3)$, case 1 of Master Theorem

b. $T(n) = 4T(n/2) + n^{2.5}$

Ans: $O(n) = n^{\log_2(4)}$, case 3 of Master Theorem

Q4. Analyze the run time complexity of the following algorithms (2 points)

a. Given an array (or string), the task is to reverse the array/string.

Algorithm -

- 1) Initialize start and end indexes as start = 0, end = n-1
- 2) In a loop, swap arr[start] with arr[end] and change start and end as follows: start = start +1, end = end -13)

Repeat 2) while start < end

Ans: **O(n)**

Q5. Given an array A[], the task is to segregate even and odd numbers. All even numbers should appear first, followed by odd numbers.

Algorithm -

- 1) Initialize two index variables left and right: left = 0, right = size -1
- 2) Keep incrementing left index until we see an odd number.
- 3) Keep decrementing right index until we see an even number.
- 4) Swap arr[left] and arr[right]
- 5) Repeat 2 4 while left < right

Ans: **O(n)**