

1 Bin packing.

Bin packing. This is a classical application of McDiarmid's inequality. Let X_1, \dots, X_n be i.i.d. random variables taking values in $[0, 1]$. Each X_i is the size of a package to be shipped. The packages are shipped in bins of size 1, so each bin can hold any set of packages whose sizes sum to at most 1. Let $B_n = f(X_1, \dots, X_n)$ be the minimal number of bins needed to ship the packages with sizes X_1, \dots, X_n . Computing B_n is a hard combinatorial optimization problem; however, we can say something about its mean and tail behavior.

1.1

Let μ be the common mean of the X_i 's. Show that $\mathbb{E}[B_n] \geq n\mu$.

1.2

Prove that for any $\epsilon > 0$,

$$\mathbb{P}(n^{-1}B_n \leq \mu - \epsilon) \leq \exp(-2n\epsilon^2).$$

2 Orthogonal design lasso bits

Let $y_i = x_i^\top \beta^* + \eta_i$ for $i = 1, \dots, n$. Assume that the design matrix $\mathbf{X} = (x_1, \dots, x_n)^\top$ and that $\mathbf{X}^\top \mathbf{X} = nI_p$. Also assume that $\log \mathbb{E}[e^{\lambda \eta_i}] \leq \sigma^2 \lambda^2 / 2$. Suppose that β^* is sparse: $\mathcal{S} = \{j : \beta_j^* \neq 0\}$ and that β^* is such that $\|\beta_{\mathcal{S}^c}^*\|_1 \leq L\|\beta_{\mathcal{S}}^*\|_1$ for some L . Let

$$\hat{\beta} = \underset{\beta}{\operatorname{argmin}} \frac{1}{2n} \|Y - \mathbf{X}\beta\|^2 + \lambda \|\beta\|_1.$$

2.1

Derive the so-called “basic inequality”:

$$\|\mathbf{X}\beta^* - \mathbf{X}\hat{\beta}\| + \lambda \|\beta\|_1 \leq \frac{1}{n} \eta^\top \mathbf{X}(\beta^* - \hat{\beta}) + \lambda \|\beta_{\mathcal{S}}^*\|_1$$

2.2

Let X_j be the j^{th} column of \mathbf{X} . Find a constant M such that $\mathbb{P}(\mathcal{G}) \geq 1 - 2e^{-\delta}$ for

$$\mathcal{G} = \left\{ \max_j |\eta^\top X_j| < M \right\}.$$

3 Empirical Rademacher Complexity

Prove the following result. Let H be a set of hypotheses bounded by $M > 0$ and consider an i.i.d. sample of size n . For any $h \in H$, with probability at least $1 - \delta$

$$R_n(h) \leq \hat{R}_n(h) + \hat{\mathcal{R}}_n(H) + 3M \sqrt{\frac{2 \log(2/\delta)}{2}}.$$