## PL, Jr. Summary: Findler, Felleisen - Contracts for Higher-Order Functions

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- 1. OO languages recognize value of assertion-based contracts (to specify pre/post-conditions) (Eiffel "Design by Contracts", one of most requested Java extensions) (Bigloo Scheme is only functional language with contracts)
- 2. functional languages use type systems to express assertions but many type systems are not expressive enough to express some assertions
- 3. authors present  $\lambda$  calculus + contracts language and prove type soundness
- 4. implementation: contract compiler inserts code to check conditions required by contract
- 5. blame/contract verification: easy with first order contracts caller blamed for input contract violation and function blamed for output contract violation
- 6. not as easy to verify contracts in language with higher order functions, example:

$$g: (\mathbf{int}[> 9] \to \mathbf{int}[0, 99]) \to \mathbf{int}[0, 99]$$
  
val rec  $g = \lambda \ proc. \cdots$ 

It's not enough to monitor applications of proc in g because g may pass proc to another function

- 7. blame is also complicated with higher order functions
- 8. contract syntax example:

```
;; bigger-than-zero?: number \rightarrow boolean
(define bigger-than-zero?(\lambda(x)(\geq x\ 0)))

sqrt: number \rightarrow number
(define/contract sqrt
(bigger-than-zero? \longmapsto bigger-than-zero?)
(\lambda(x)\cdots))
```

Things to note about example:

- contracts can be predicates or function contracts (one predicate for domain and one for range)
- contract can be arbitrary expression that evaluates to a contract use of bigger-than-zero?

9. example of dependent contract – range predicate can depend on function argument

```
sqrt: \mathbf{number} \to \mathbf{number}

(\mathbf{define/contract} \ sqrt

(bigger-than-zero? \overset{d}{\longmapsto})

(\lambda(x).\lambda(res).

(\mathbf{and} \ (bigger-than-zero? \ res)

(\geq (abs \ (-x \ (*res \ res))) \ 0.01))))

(\lambda(x)\cdots))
```

- 10. key to checking higher-order contracts is to postpone enforcement until higher order function is applied delay/save/saved/use example
- 11. contracts also allow for better code modularity because checking for validity of inputs is separated
- 12. assigning blame when first-class functions are involved:
  - if base contract appears to the left of an even number of arrows, then function where first-class function is applied is to blame (covariant)
  - if base contract appears to the left of an odd number of arrows, then calling function is to blame (contravariant)

example:

```
; g : (\mathbf{int} \to \mathbf{int}) \to \mathbf{int}
(\mathbf{define/contract}\ g)
((greater-than-nine? \longmapsto between-zero-and-ninety-nine?)
\longmapsto
between-zero-and-ninety-nine?)
(\lambda(f)(f\ 0))
```

- greater-than-nine? contract is violated and since it appears to the left of two arrows (even), then g is to blame and this is true
- If instead f was applied to 10 and  $(f\ 10) \Rightarrow -10$ , then second between-zero-and-ninety-nine? contract is violated and since it appears to the left of one arrow (odd), then whoever called g is to blame and this is also true
- 13. in Scheme, contracts are first-class
- 14. contracts are useful when dealing with callbacks because the save/use model is frequently encountered (callbacks are first "registered" and then used later)