

UMass Boston Computer Science
CS450 High Level Languages (section 2)
Abstraction

Wednesday, October 11, 2023

AN x64 PROCESSOR IS SCREAMING ALONG AT BILLIONS OF CYCLES PER SECOND TO RUN THE XNU KERNEL, WHICH IS FRANTICALLY WORKING THROUGH ALL THE POSIX-SPECIFIED ABSTRACTION TO CREATE THE DARWIN SYSTEM UNDERLYING OS X, WHICH IN TURN IS STRAINING ITSELF TO RUN FIREFOX AND ITS GECKO RENDERER, WHICH CREATES A FLASH OBJECT WHICH RENDERS DOZENS OF VIDEO FRAMES EVERY SECOND

BECAUSE I WANTED TO SEE A CAT
JUMP INTO A BOX AND FALL OVER.



I AM A GOD.

Logistics

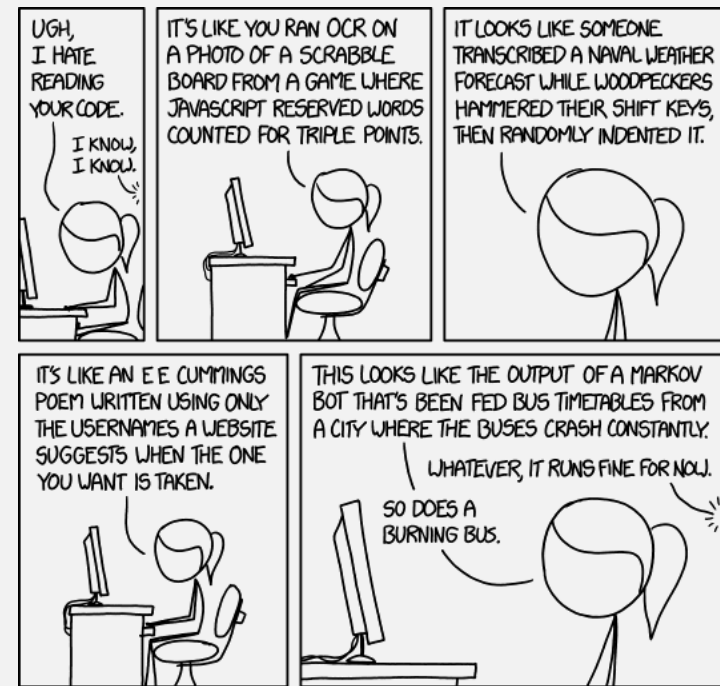
- HW 3 out
 - due: Sun 10/15 11:59 pm EST
- HW 2 (and other) grades returned
 - Use GradeScope re-grade request for complaints or questions



CS 450 so far ...

This class teaches:

- a high-level programming “process”
- i.e., a **design recipe** for creating clean, readable programs
- How to do well: learn and follow the process
- How to not do well: focus only on “getting the code working”

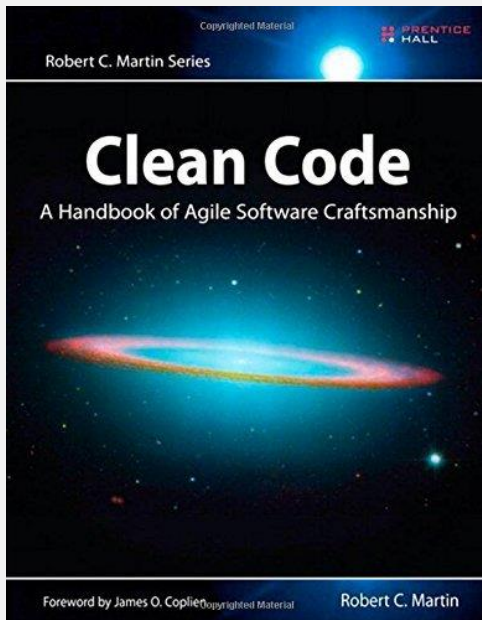


“Perhaps you thought that “getting it working” was the first order of business for a professional developer.

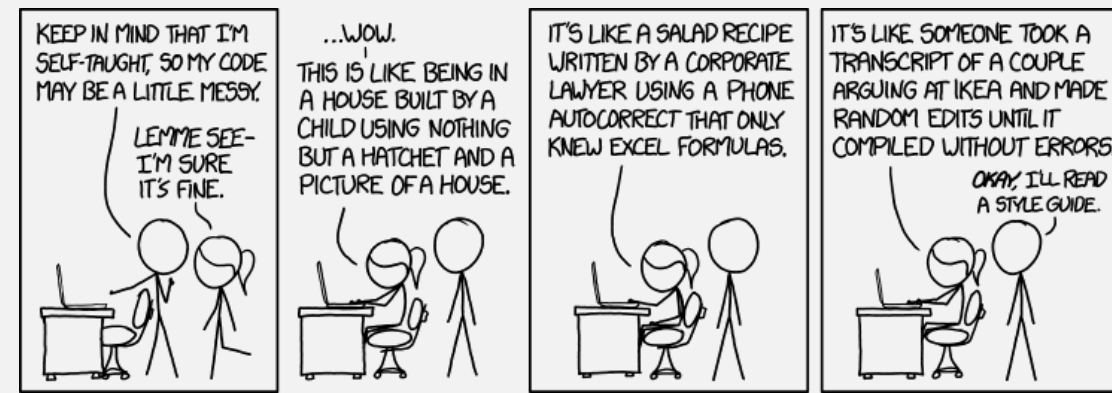
I hope by now, however, that this book has disabused you of that idea.

The functionality that you create today has a good chance of changing in the next release, but the readability of your code will have a profound effect on all the changes that will ever be made.”

— Robert C. Martin,
**Clean Code: A Handbook of Agile
Software Craftsmanship**



HW 2 recap



Many submissions only focused on: “getting the code working”

Many submissions ignored:

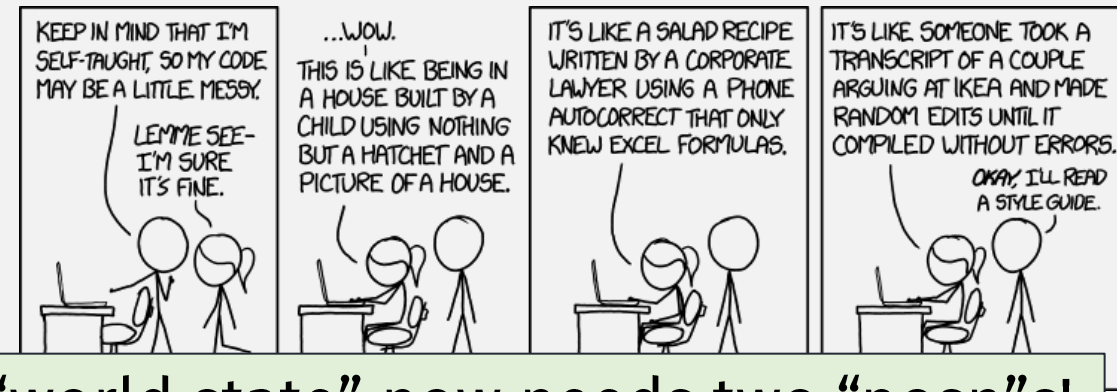
- all other steps of **programming design recipe**
- style guide
- Other instructions in hw

This hw will be graded accordingly:

- correctness (9 pts)
- design recipe (20 pts)
- style (5 pts)
- README (1 pt)

Total: 35 points

HW 2 solution (part 2)



“world state” now needs two “posn”s!

- To practice reading code, save the [big-bang program from lecture 6](#) to a file named `hw2-bigbang.rkt`, and change it so when the user clicks the left-mouse button, a copy of the image is pinned to the canvas. (The original image should continue to move with the cursor.) Subsequent clicks should move the pinned image so that only one image is pinned at any time.

HW 2 recap – key design points

```
;; A WorldState is a:  
(struct world [cursor pinned])
```

```
;; where
```

```
;; cursor : Posn - IMG location that moves with mouse cursor
```

```
;; pinned : MaybePosn - IMG location that is pinned (if there is one)
```

“world state” now needs two “posn”s!

one of them “maybe” will not have a value



```
;; A Posn is a
```

```
(struct posn [x y])
```

```
;; where
```

```
;; x: Integer - represents x coordinate in big-bang animation
```

```
;; y: Integer - represents y coordinate in big-bang animation
```

```
;; A MaybePosn is a
```

```
;; - Posn
```

```
;; - NO-POSN
```

```
;; represents a position, if there is one
```

```
;; constants and predicates for MaybePosn
```

```
(define NO-POSN #false)
```

```
(define (no-posn? x) (equal? x NO-POSN))
```

```
(define (maybe-posn? x) (or (posn? x) (no-posn? x)))
```

Mouse handler

```
;; mouse-handler : WorldState Coordinate Coordinate MouseEvent -> WorldState
;; Returns a WorldState where:
;; - "cursor" = current mouse loc
;; - "pinned" = current mouse loc (if mevt = "button-down")
```

```
;; TEMPLATE fn for MouseEvent : MouseEvent -> ???
(define (mevt-fn mevt)
  (cond
    [(string=? mevt "button-down") ....]
    .... ]))
```

Function splitting rule:

One (data definition
processing) task, one function

When a function has more than
one argument, you choose
which template to use

Mouse handler

```
;; mouse-handler : WorldState Coordinate Coordinate MouseEvent -> WorldState
;; Returns a WorldState where:
;; - "cursor" = current mouse loc
;; - "pinned" = current mouse loc (if mevt = "button-down")
```

```
(define (mouse-handler w x y mevt)
  (cond
    [(string=? mevt "button-down")
     (world (posn x y) (posn x y))]
    [else
     (world (posn x y) (world-pinned w))]))
```

Current mouse loc



```
(define (mouse-handler w x y mevt)
  (define current-mouse-pos (posn x y))
  (cond
    [(string=? mevt "button-down")
     (world current-mouse-pos
             current-mouse-pos)]
    [else
     (world current-mouse-pos
             (world-pinned w))]))
```

This field
always
the same

Mouse handler

```
;; mouse-handler : WorldState  
;; Returns a WorldState w  
;; - "cursor" = current m  
;; - "pinned" = current m
```

```
(define (mouse-handler w x y mevt)  
  (define current-mouse-pos (posn x y))  
  (world  
    current-mouse-pos  
    (cond  
      [(string=? mevt "button-down") current-mouse-pos]  
      [else (world-pinned w)])))
```

WAIT! Function split rule is:

One (data definition
processing) task,
one function

Do we need a
separate (world-
processing) function?

```
(define (mouse-handler w x y mevt)  
  (define current-mouse-pos (posn x y))  
  (cond  
    [(string=? mevt "button-down")  
     (world current-mouse-pos  
             current-mouse-pos)]  
    [else  
     (world current-mouse-pos  
             (world-pinned w))]))
```

On when to create a new function ...

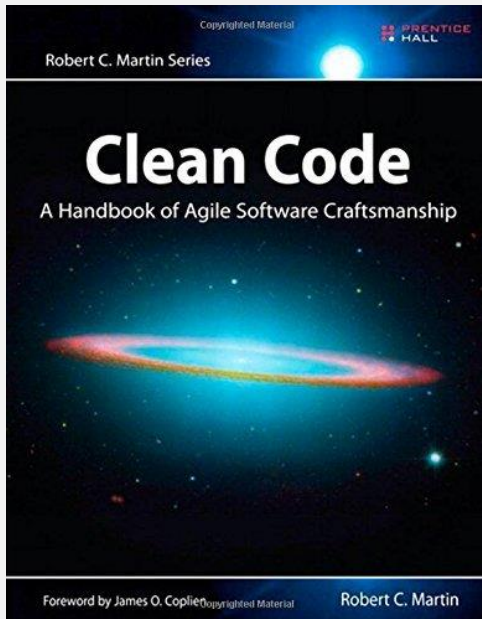
“The first rule of functions is that they should be small.

The second rule of functions is that they should be smaller than that.”

— Robert C. Martin,
**Clean Code: A Handbook of Agile
Software Craftsmanship**

In this class:

create one function per
(data definition processing) task



Mouse handler

```
;; mouse-handler : WorldState  
;; Returns a WorldState w  
;; - "cursor" = current m  
;; - "pinned" = current m
```

```
(define (mouse-handler w x y mevt)  
  (define current-mouse-pos (posn x y))  
  (world  
    current-mouse-pos  
    (cond  
      [(string=? mevt "button-down") current-mouse-pos]  
      [else (world-pinned w)])))
```

WAIT! Function split rule is:

One (data definition
processing) task,

one function

Do we need a
separate (world-
processing) function?

YES. (But in this case,
it would just **get** the
world-pinned value)

Render world

```
;; render-world : WorldState -> Image  
;; draws IMG at the "cursor" and "pinned" posn (if there is one)
```

```
;; TEMPLATE fn for WorldState: WorldState -> ???  
(define (world-fn w)  
  .... (world-cursor w) .... (world-pinned w) ....)
```



These are posns. Should we also extract posn fields here?

NO. They should be handled by a “posn”
function (which follows “posn” template)

create one function per
(data definition processing) task

Render world

```
;; render-world : WorldState -> Image  
;; draws IMG at the "cursor" and "pinned" posn (if there is one)
```

```
;; TEMPLATE fn for WorldState: WorldState -> ???  
(define (world-fn w)  
  .... (world-cursor w) .... (world-pinned w) ....)
```

```
(define (render-world w)  
  (render-posn  
    (world-pinned w)  
    (render-posn  
      (world-cursor w)  
      EMPTY-SCENE))))
```

But this is a "maybe" posn

Wish list:
maybe-render-posn

Render posn

```
;; maybe-render-posn : MaybePosn Image -> Image  
;; draws IMG into given image at given posn;  
;; unless posn is NO-POSN, then return img unchanged
```

```
;; TEMPLATE fn for MaybePosn fn: MaybePosn-> ???  
(define (maybeposn-fn p)  
  (cond  
    [(posn? p) .... (posn-x p) .... (posn-y p) ....])  
    [(no-posn? p) ....]))
```

```
(define (maybe-render-posn p img)  
  (cond  
    [(posn? p) (place-image IMG (posn-x p) (posn-y p) img)]  
    [else img]))
```

Render world

```
;; render-world : WorldState -> Image  
;; draws IMG at the "cursor" and "pinned" posn (if there is one)
```

```
(define (render-world w)  
  (maybe-render-posn  
    (world-pinned w)  
    (maybe-render-posn  
      (world-cursor w)  
      EMPTY-SCENE)))
```

Works for both!

Wish list:

~~maybe-render-posn~~

HW 2 recap - other points

- GitHub repos must be added to cs450f23 organization (not your own account)
 - Otherwise I cant see it
 - (see directions in hw)
- Git commit messages must be meaningful



	COMMENT	DATE
○	CREATED MAIN LOOP & TIMING CONTROL	14 HOURS AGO
○	ENABLED CONFIG FILE PARSING	9 HOURS AGO
○	MISC BUGFIXES	5 HOURS AGO
○	CODE ADDITIONS/EDITS	4 HOURS AGO
○	MORE CODE	4 HOURS AGO
○	HERE HAVE CODE	4 HOURS AGO
○	AAAAA	3 HOURS AGO
○	ADKFJSLKDFJSDKLFJ	3 HOURS AGO
○	MY HANDS ARE TYPING WORDS	2 HOURS AGO
○	HAAAAAANDS	2 HOURS AGO

Do not ...

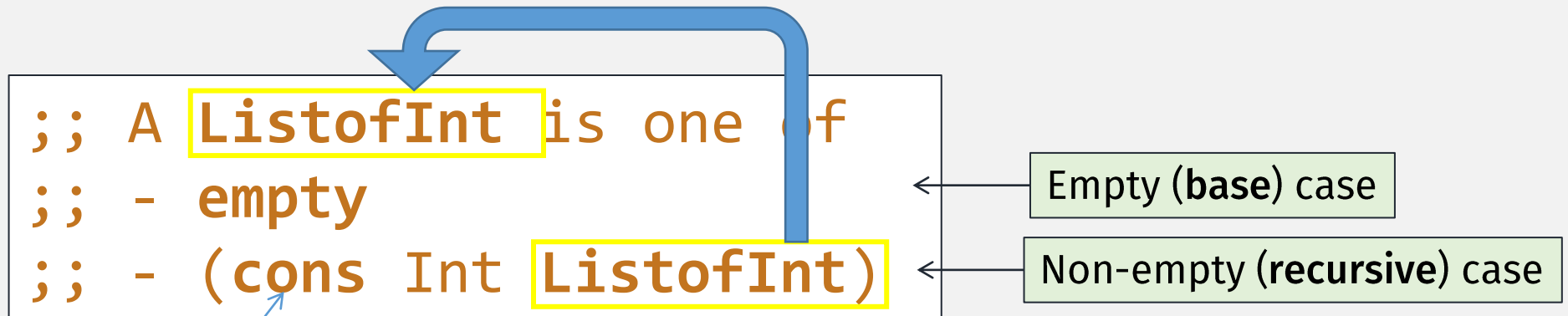
- Add to temp files to repo
- Leave commented out code
- Leave “TODO”s in code

AS A PROJECT DRAGS ON, MY GIT COMMIT MESSAGES GET LESS AND LESS INFORMATIVE.

Lists and List Functions Review

Previously

Racket List Data Definition Example



cons = "node"

Recursive!

(using a definition to define itself)

TEMPLATE??

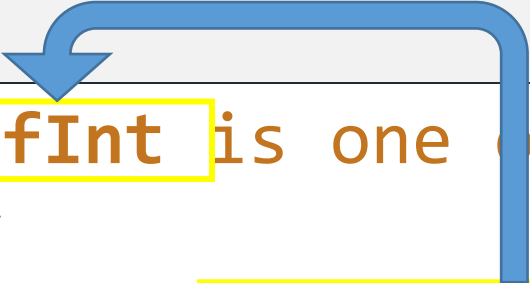
(how can we use a list of ints
to define a list of ints?!?)

Recursion is a valid concept (from math), but only if there is both

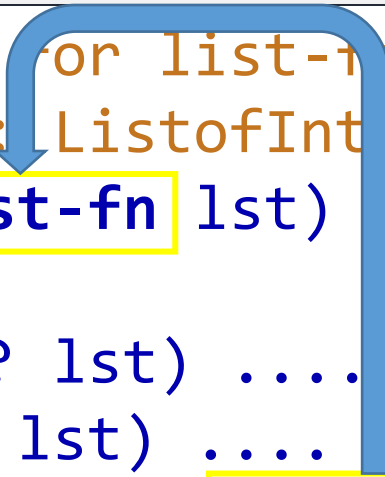
- A **base** case
- A **recursive** case

Previously

Racket Recursive List Fn Template



```
;; A ListofInt is one of  
;; - empty  
;; - (cons Int ListofInt)
```



```
;; TEMPLATE for list-fn  
;; list-fn : ListofInt -> ???  
(define (list-fn lst)  
  (cond  
    [(empty? lst) ....]  
    [(cons? lst) .... (first lst) ....  
      .... (list-fn (rest lst)) ....]))
```

Racket Recursive List Fn: `inc-list`

```
;; TEMPLATE for list-fn
;; list-fn : ListofInt -> ???
(define (list-fn lst)
  (cond
    [(empty? lst) ....]
    [(cons? lst) .... (first lst) ....
      .... (list-fn (rest lst)) ....])))
```

Racket Recursive List Fn: `inc-list`

```
(check-equal?  
  (inc-list (list 1 2 3))  
  (list 2 3 4))
```

```
;; inc-list : ListofInt -> ListofInt  
;; increments each list element by 1  
(define (inc-lst lst)  
  (cond  
    [(empty? lst) ....]  
    [(cons? lst) .... (first lst) ....  
      .... (list-fn (rest lst)) ....]))
```

Racket Recursive List Fn: **inc-list**

```
;; inc-list : ListofInt -> ListofInt
;; increments each list element by 1
(define (inc-lst lst)
  (cond
    [(empty? lst) empty]
    [(cons? lst) .... (first lst) ....
      .... (list-fn (rest lst)) ....]))
```

Racket Recursive List Fn: **inc-list**

```
;; inc-list : ListofInt -> ListofInt
;; increments each list element by 1
(define (inc-lst lst)
  (cond
    [(empty? lst) empty]
    [else (cons (add1 (first lst))
                  (list-fn (rest lst)))]))
```


Previously

Multi-ball Animation

Design a **big-bang** animation that:

- Start: a single ball, moving with random x and y velocity
- On a click: add a ball at random location, with random velocity

;; A WorldState is ... a list of balls!

```
;; A Ball is a
(struct ball [x y xvel yvel] #:transparent)
;; where
;; x: XCoord - represents x coordinate of ball center in animation
;; y: YCoord - represents y coordinate of ball center in animation
;; xvel: Integer - represents x velocity, where
;;               positive = to the right, negative = to the left
;; yvel: Integer - represents y vel, where
;;               positive = down, negative = up
```

```
;; A ListofBall is one of
;; - null
;; - (cons Ball ListofBall)
```

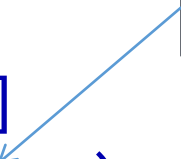
```
;; A WorldState is a ListofBall
```

next-world

List template!

```
;; next-world : WorldState -> WorldState
;; Updates position of all balls by one tick
(define (next-world w)
  (cond
    [(empty? w) ....]
    [else .... (first w) ....
               .... (next-world (rest w)) ....]]))
```

Ball



Create one
function
per “task”

```
(check-equal? (next-world (list (make-ball 0 0 1 1)))
              (list (next-ball (make-ball 0 0 1 1))))
```

next-world

```
;; next-world : WorldState -> WorldState
;; Updates position of all balls by one tick
(define (next-world w)
  (cond
    [(empty? w) empty]
    [else (cons (next-ball (first w))
                  (next-world (rest w)))])
```

next-world

```
:: next-world : ListofBall -> ListofBall  
:: Updates position of all balls by one tick  
(define (next-world w)  
  (cond  
    [(empty? w) empty]  
    [else (cons (next-ball (first w))  
                (next-world (rest w)))]])
```

Comparison

```
;; inc-lst: ListofInt -> ListofInt
;; Returns list with each element incremented
(define (inc-lst lst)
  (cond
    [(empty? lst) empty]
    [else (cons (add1 (first lst))
                 (inc-lst (rest lst)))])])
```

```
;; next-world : ListofBall -> ListofBall
;; Updates position of each ball by one tick
(define (next-world w)
  (cond
    [(empty? w) empty]
    [else (cons (next-ball (first w))
                 (next-world (rest w)))])])
```

Abstraction: Common List Function #1

```
;; lst-fn1: (?? -> ??) Listof?? -> Listof??  
;; Applies the given fn to each element of given lst
```

```
(define (lst-fn1 fn lst)  
  (cond  
    [(empty? w) empty]  
    [else (cons (fn (first w))  
                 (lst-fn1 (rest w)))]))
```

```
(define (inc-lst lst) (lst-fn1 add1 lst))  
(define (next-world lst) (lst-fn1 next-ball lst))
```

Abstraction: Common List Function #1

```
;; lst-fn1: (X -> X) ListofX -> ListofX  
;; Applies the given fn to each element of given lst
```

```
(define (lst-fn1 fn lst)  
  (cond  
    [(empty? w) empty]  
    [else (cons (fn (first w))  
                 (lst-fn1 (rest w)))]))
```

```
(define (inc-lst lst) (lst-fn1 add1 lst))  
(define (next-world lst) (lst-fn1 next-ball lst))
```


Abstraction: Common List Function #1

```
;; lst-fn1: (X -> Y) ListofX -> ListofY  
;; Applies the given fn to each element of given lst
```

```
(define (lst-fn1 fn lst)  
  (cond  
    [(empty? w) empty]  
    [else (cons (fn (first w))  
                 (lst-fn1 (rest w)))]))
```

```
(define (inc-lst lst) (lst-fn1 add1 lst))  
(define (next-world lst) (lst-fn1 next-ball lst))
```

Abstraction Data Definitions

```
;; A ListofInt is one of  
;; - empty  
;; - (cons Int ListofInt)
```

```
;; A ListofBall is one of  
;; - empty  
;; - (cons Ball ListofBall)
```



```
;; A Listof<X> is one of  
;; - empty  
;; - (cons X Listof<X>)
```

NOTE: this shows why our Compound data predicates should be “**shallow**” checks, i.e., list?

Makes abstraction easier

To use this **abstract** data definition, must **instantiate** X with a **concrete** data definition

Listof<Int>

Listof<Ball>

(concrete = opposite of abstract)

Abstract Data Defs common in every PL

```
64 #include<iostream>
65 #include <vector>
66 using namespace std;
67
68 int main()
69 {
70     vector<int> v;
71
72     for (int i = 1; i <= 10; i++)
73     {
74         v.push_back(i);
75     }
76     cout << "Size : " << v.size();
77
78     v.resize(7);
79
80     cout << "\nAfter resizing it becomes : " << v.size();
```

(C++ STL)

Structs define abstract data

Abstract data – “any” x and y allowed

```
;; A Posn is a  
(struct posn [x y])  
;; where  
;; x: Integer - represents x coordinate in big-bang animation  
;; y: Integer - represents y coordinate in big-bang animation
```

(implicit) Instantiation

Common List Function #1

```
;; lst-fn1: (X -> Y) ListofX -> ListofY  
;; Applies the given fn to each element of given lst
```

```
(define (lst-fn1 fn lst)  
  (cond  
    [(empty? w) empty]  
    [else (cons (fn (first w))  
                  (lst-fn1 (rest w)))]))
```

```
(define (inc-lst lst) (lst-fn1 add1 lst))  
(define (next-world lst) (lst-fn1 next-ball lst))
```

Common List Function #1: map

```
;; map: (X -> Y) ListofX -> ListofY  
;; Applies the given fn to each element of given lst
```

```
(define (map fn lst)  
  (cond  
    [(empty? w) empty]  
    [else (cons (fn (first w))  
                  (map (rest w)))])
```

```
(define (inc-lst lst) (map add1 lst))  
(define (next-world lst) (map next-ball lst))
```

Common List Function #1: map

```
(map proc lst ...+) → list?  
proc : procedure?  
lst : list?
```

map: (A B C ... -> Z) Listof<A> Listof Listof<C> ... -> Listof<Z>
;; Applies the given fn to elements (at same index) of given lsts

```
(check-equal? (map + (list 1 2 3) (list 4 5 6)  
                 (list 5 7 9)))
```

Common List Function #2: ???

Previously

Racket Recursive List Fn Example: sum

```
;; TEMPLATE for list-fn
;; list-fn : ListofInt -> ???
(define (list-fn lst)
  (cond
    [(empty? lst) ....]
    [(cons? lst) .... (first lst) ....
      .... (list-fn (rest lst)) ....])))
```

Previously

Racket Recursive List Fn Example: sum

```
;; Returns sum of list of ints
;; sum-lst: ListofInt -> Int
(define (sum-lst lst)
  (cond
    [(empty? lst) 0]
    [else (+ (first lst)
              (sum-lst (rest lst)))]))
```

Render World: ListofBall edition

```
;; render-world : ListofBall -> Image  
;; Draws the given world as an image by overlaying each ball,  
;; at its position, into an initially empty scene
```

```
(define (render-world lst)  
  (cond  
    [(empty? lst) ....]  
    [else .... (first lst).... (render-world (rest lst)) ....])))
```

Render World: ListofBall edition

```
;; render-world : ListofBall -> Image  
;; Draws the given world as an image by overlaying each ball,  
;; at its position, into an initially empty scene
```


```
(define (render-world lst)  
  (cond  
    [(empty? lst) EMPTY-SCENE]  
    [else .... (first lst).... (render-world (rest lst)) ....]))
```

Render World: ListofBall edition

```
;; render-world : ListofBall -> Image  
;; Draws the given world as an image by overlaying each ball,  
;; at its position, into an initially empty scene
```

```
(define (render-world lst)  
  (cond  
    [(empty? lst) EMPTY-SCENE]  
    [else (place-ball (first lst) (render-world (rest lst)))]))
```

Create one
function
per “task”



```
;; place-ball : Ball Image -> Image  
;; Draws a ball, using its pos as the offset, into the given image  
(define (place-ball b scene)  
  (place-image BALLIMG (ball-x b) (ball-y b) scene))
```

Comparison #2

```
:: sum-1st: ListofInt -> Int
(define (sum-1st lst)
  (cond
    [(empty? lst) 0]
    [else (+ (first lst)
              (sum-1st (rest lst)))]))
```

```
:: render-world : ListofBall -> Image
(define (render-world lst)
  (cond
    [(empty? lst) EMPTY-SCENE]
    [else (place-ball (first lst)
                       (render-world (rest lst)))]))
```

Common List Function #2

X = Type of list element

Y = Result Type

```
;; list-fn2 : (X Y -> Y) Y Listof<X> -> Y
```

```
(define (list-fn2 fn initial lst)
  (cond
    [(empty? lst) initial]
    [else (fn (first lst) (list-fn2 fn initial (rest lst)))]))
```

```
;; sum-lst: ListofInt -> Int
(define (sum-lst lst) (list-fn2 + 0 lst))
;; render-world: ListofBall-> Image
(define (render-world lst) (list-fn2 place-ball EMPTY-SCENE lst))
```

Common List Function #2: **foldr** (start at right)

```
:: foldr: (X Y -> Y) Y Listof<X> -> Y
```

```
(define (foldr fn initial lst)
```

```
  (cond
```

Function recurs and builds up fn calls until it gets to the end

```
    [(empty? lst) initial]
```

Then they are evaluated, last one first

```
    [else (fn (first lst) (foldr fn initial (rest lst))))]))
```

```
:: sum-lst: ListofInt -> Int
```

```
(define (sum-lst lst) (foldr + 0 lst))
```

```
:: render-world: ListofBall-> Image
```

```
(define (render-world lst) (foldr place-ball EMPTY-SCENE lst))
```


Common List Function #2: `foldr`

```
;; foldr: (X ... Y -> Y) Y Listof<X> ... -> Y
```

Racket version can also take multiple lists

```
(foldr proc init lst ...+) → any/c  
proc : procedure?  
init : any/c  
lst : list?
```

foldr (start at right): order matters?

For some functions, order doesn't matter, but for others, it does?

`(foldr + 0 (list 1 2 3)) = (1 + (2 + 3))`

`(1 + (2 + 3)) = (1 + 2) + 3`

`(1 - (2 - 3)) = (1 - 2) - 3 ???`



Need List Function #2b: **foldl** (start from left)

Challenge:

- Change **foldr** to **foldl**
- (so that the function is applied to the first element first)

```
(define (foldr fn initial lst)
  (cond
    [(empty? lst) initial]
    [else (fn (first lst) (foldr fn initial (rest lst)))]))
```



```
(define (foldl fn initial lst)
  (cond
    [(empty? lst) ....]
    [else .... (first lst) .... (foldl fn initial (rest lst)) ....]))
```

Next time: Other common list functions

- Filter
- Find
- Reverse
- append

Check-In Quiz 10/11 on gradescope

(due 1 minute before midnight)