2)

We can rewrite the grammar using the following judgement forms. (e denotes existence within a

set)

Parts of the judgement
$$A(1)$$
 e AObjects, $A(2)$ e AObjects \leftarrow premise, may be left empty for axioms \leftarrow what follows

b)

Show the last grammar is ambiguous.

For example string aba show there are 2 possible parse trees

- c) The grammar forms strings consisting of 1 or more As (A+), or 0 or more Bs (B*), or 1 or more Cs (C+).
- d) 1) $S \Rightarrow AaBb \Rightarrow baab$
 - 2) not in the grammar
 - 3) not in the grammar
 - 4) S => AaBb => AbaBb => bbaBb => bbaab

e)

- 1) $S \Rightarrow aScB \Rightarrow abcB \Rightarrow abcd$
- 2) not in the grammar
- 3) not in the grammar
- 4) not in the grammar
- 5) $s \Rightarrow aScB \Rightarrow aAcB \Rightarrow accB \Rightarrow accA \Rightarrow accc$

- 3)
- a)
- I) This grammar forms strings which may have a large left subtree of operations strung together in a given parse tree. This left tree would then be used in a final evaluation with the original operand. This grammar can also just represent an operand value.
- II) This grammar forms similar expressions, however now the right subtree of the parse tree may be arbitrarily long.
- b) Example from scala REPL:

```
scala> 2-1 << 1
res0: Int2 = 2
– This shows that subtraction has a higher precedence than the binary left shift operator (2-1) << 1 = 2 2-(1 << 1) = 0
```

c) BNF grammar for doubles

```
S ::= fraction \mid fraction \; exp \\ fraction ::= posnum \mid negnum \\ exp: E \; sign \; num \\ posnum = num \; . \; dec \qquad (where "." is a decimal point) \\ negnum = -num \; . \; dec \qquad (where "." is a decimal point) \\ num ::= num \; b \mid a \\ a ::= \{1,2,3,4,5,6,7,8,9\} \\ b ::= \{0,1,2,3,4,5,6,7,8,9\} \\ dec ::= dec \; c \mid c \\ c ::= \{0,1,2,3,4,5,6,7,8,9\}
```