

COMP2201 – Discrete Mathematics
Graph Theory including Planarity and Cycles

1. Prove that K_5 is not planar.

Answer 1.

Revision:

K_n

- the complete graph on n vertices
- the simple graph with n vertices in which there is an edge between every pair of distinct vertices

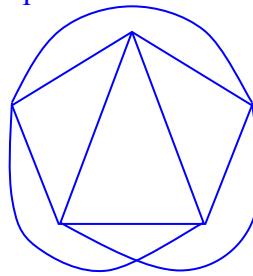
Planar graph

- a graph that can be drawn without its edges crossing

Solution:

- i. Assuming that vertices are numbered from top to bottom proceeding firstly in a left to right direction, we have vertices numbered from 1 (topmost) to 5 (bottom-right most)

By attempting to re-draw graph...



We get edge 3-4 crossing 1-5 or 2-5.

$\therefore K_5$ is not planar

- ii. Where v = number of vertices
 e = number of edges
 f = number of faces, including the exterior face

$$v = 5$$

$$e = 4 + 3 + 2 + 1 = 10$$

$$f = 8$$

If G is *planar* graph, Then

$$v - e + f = 2$$

or

$$f = e - v + 2$$

i.e. $5 - 10 + 8 = 3 \neq 2$

$\therefore K_5$ is not planar

- iii. If G is *planar* graph, Then

$$e \leq 3v - 6$$

$$e = 10, v = 5$$

$$\Rightarrow 10 \leq 3 \times 5 - 6 \quad \text{or} \quad 10 \leq 9$$

$\therefore K_5$ is not planar

- iv. Suppose that K_5 is planar
 Since every cycle has at least 3 edges, each face is bounded by at least 3 edges.
 Thus the number of edges that bound faces is at least $3f$.

In a planar graph, each edge belongs to at most two bounding cycles.

Therefore $2e \geq 3f$

Using $v - e + f = 2$ or $f = e - v + 2$

We find that $2e \geq 3(e - v + 2)$

$$\text{Now } v = 5$$

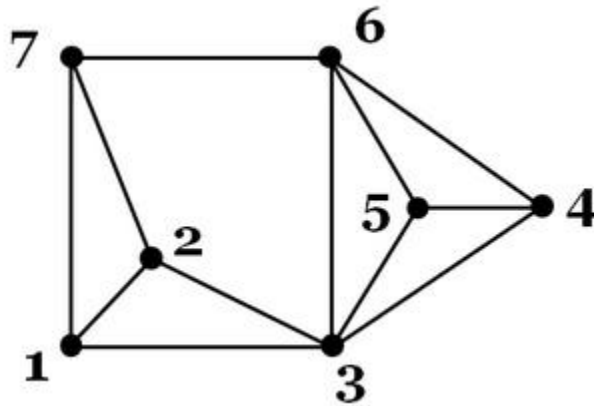
$$e = 4 + 3 + 2 + 1 = 10$$

$$2(10) \geq 3(10 - 5 + 2)$$

$$20 \geq 21$$

$\therefore K_5$ is not planar

2. Given the following undirected graph G with 7 vertices:



- List the clique(s) with 3 or more vertices that are within the graph
- Identify the maximal clique(s)
- Find the maximum clique(s) or state why there is no maximum cliques in the graph G
- Determine the clique number $\omega(G)$

Answer 2.

- i. [A **clique**, C , in an undirected graph $G = (V, E)$ is a subset of the vertices, $C \subseteq V$, such that every two distinct vertices are adjacent.]

The cliques with 3 or more vertices are:

- a. $V = \{1, 2, 3\}$
- b. $V = \{1, 2, 7\}$
- c. $V = \{3, 4, 5\}$
- d. $V = \{3, 5, 7\}$
- e. $V = \{4, 5, 6\}$
- f. $V = \{3, 4, 5, 6\}$

- ii. [A **maximal clique** is a clique that cannot be extended by including one more adjacent vertex, that is, a clique which does not exist exclusively within the vertex set of a larger clique.]

The maximal cliques are:

- a. $V = \{1, 2, 3\}$
- b. $V = \{1, 2, 7\}$
- c. $V = \{3, 4, 5, 6\}$

Cliques $\{3, 4, 5\}$, $\{3, 5, 6\}$ and $\{4, 5, 6\}$ are not maximal as by including vertex 6, 4 and 3 respectively, a larger clique is identified.

- iii. [A **maximum clique** of a graph, G , is a clique, such that there is no clique with more vertices.]

The given graph has one maximum clique. The maximum clique is $\{3, 4, 5, 6\}$

- iv. [The **clique number** $\omega(G)$ of a graph G is the number of vertices in a maximum clique in G .]

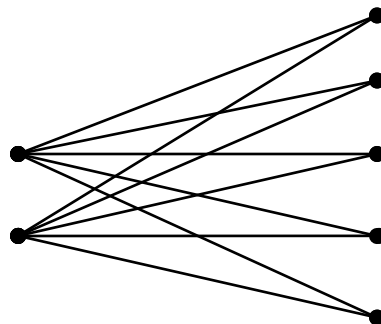
The maximum clique is $\{3, 4, 5, 6\}$.

The number of vertices in the maximum clique is 4.

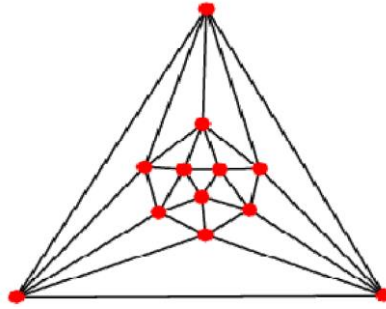
Therefore the clique number $\omega(G)$ is 4.

3. Draw $K_{2,5}$, the complete bi-partite graph on 2 and 5 vertices.

Answer 3.



4. Find a Hamilton cycle in the graph shown below



Revision:

Simple path (from v to w)

- a path from v to w with no repeated vertices

Cycle (or circuit)

- a path of non-zero length from v to v with no repeated edges

Simple cycle (from v to v)

- a cycle with no repeated vertices

Hamiltonian cycle

- A cycle in a graph G that contains each vertex in G exactly once, except for the starting and ending vertex that appears twice.

Hamiltonian path

- A simple path in a graph G that contains each vertex in G exactly once. (A Hamiltonian path begins and ends at different vertices)

Answer 4.

Solution:

Assuming that vertices are numbered from top to bottom proceeding firstly in a left to right direction, we have vertices numbered from 1 (topmost) to 12 (bottom-right most) (1, 2, 3, 4, 5, 6, 9, 7, 8, 10, 12, 11, 1) or many others

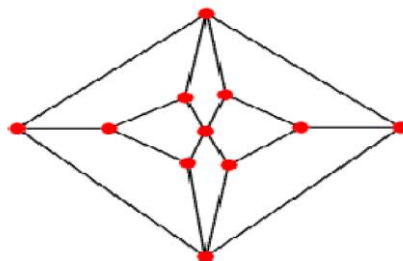
Revision:

Euler Cycle

- a cycle in a graph G that includes all of the edges and all of the vertices of G

Sum of the degrees of all the vertices in a graph is even

5. Consider the graph below



- i. Prove that a Hamiltonian cycle does not exist in this graph

Hamiltonian cycle

- A cycle in a graph G that contains each vertex in G exactly once, except for the starting and ending vertex that appears twice.

Answer 5i.

A Hamiltonian cycle with n vertices should also have n edges. Likewise, each vertex in a Hamiltonian cycle has degree of 2.

Assuming that vertices are numbered from top to bottom proceeding firstly in a left to right direction, we have vertices numbered, we have vertices numbered from 1 (topmost) to 11 (bottommost).

There are 11 vertices and 18 edges.

Suppose that we could eliminate edges from the graph, leaving just a Hamiltonian cycle, as each vertex has degree of 2, we would have to eliminate edges leaving only 2 at each vertex. Consider...

Eliminate two edges from 1, 6, 11 and one edge from 2, 3, 4, 5, 7, 8, 9, 10.

Eliminate

Two edges at vertices 1, 6 and 11:

vertex 1 (edges 1-2, 1-3),

vertex 6 (edges 6-9, 6-10) and

vertex 11 (edges 11-4, 11-8)

resulting in a single edge being eliminated from (possibly)

vertices 2, 3, 9, 10, 4 and 8

One edge at vertices 2, 3, 4, 5, 7, 8, 9, 10

vertex 5 (edge 5-9 OR edge 5-4)

vertex 7 (edge 7-10 OR edge 7-8)

No. of edges eliminated would be $3 \times 2 + 2 \times 1 = 8$

No. of edges remaining = $18 - 8 = 10$

A Hamiltonian cycle with 11 vertices should also have 11 edges.

Therefore the graph does not contain a Hamiltonian cycle.

- ii. If a Hamiltonian path exists, describe one; otherwise, prove that a Hamiltonian path does not exist.

Hamiltonian path

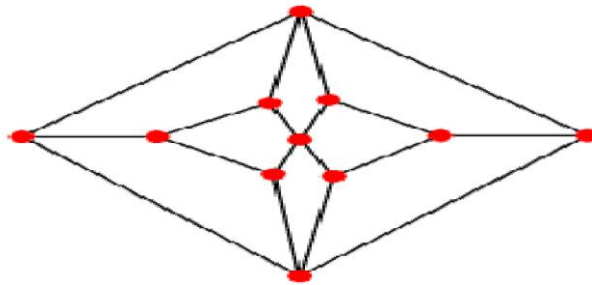
- A simple path in a graph G that contains each vertex in G exactly once. (A Hamiltonian path begins and ends at different vertices)

A Hamiltonian path with n vertices should have $n-1$ edges. Likewise, each vertex in a Hamiltonian cycle has degree of 2 (except for start and end).

Answer 5ii.

(4, 1, 8, 7, 3, 6, 2, 5, 9, 11, 10)

6. Consider the graph below



- If an Euler cycle exists, describe it; otherwise state why one does not exist in this graph
- If an Euler path exists, describe one; otherwise, prove that an Euler path does not exist.

Answer 6.

- [Revision:
Euler Cycle
 - a cycle in a graph G that includes all of the edges and all of the vertices of G
 - Each edge must exist in the cycle only once
 - The degree of each vertex in the Euler cycle is even
 - The sum of the degrees is even]

There are 11 vertices and 18 edges.

If an Euler cycle exists, each vertex has an even degree...

Assuming that vertices are numbered from top to bottom proceeding in a left to right direction in each row, we have vertices numbered from 1 (topmost) to 11 (bottommost) similar to:

1
2, 3
4, 5, 6, 7, 8

9, 10
11

As there exists a vertex which does not have an even degree i.e. any of vertices 2, 3, 4, 5, 7, 8, 9 and 10 all with degree 3, and a graph with an Euler cycle must have every vertex having an even degree

Therefore the graph does not contain an Euler cycle.

ii. [Revision:
Euler Path

- a path in a graph G that includes all of the edges and all of the vertices of G
 - Each edge must exist in the path only once
 - The degree of each vertex in the Euler cycle is even except for the starting and ending vertex
 - The sum of the degrees is even
-]

The sum of the degrees is

$$4+3+3+3+3+4+3+3+3+3+4 = 36$$

However, as there exists more than two vertices with an odd degree i.e. any three of vertices 2, 3, 4, 5, 7, 8, 9 and 10 all with degree 3, and a graph with an Euler path must have exactly two vertices with odd degree

Therefore an Euler path does not exist.